

Lattice

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Ordered Sets

Definition 1.1. **Partially ordered set** is a system $\mathcal{P} = (P, \leq)$ where P is a nonempty set and \leq is a binary relation on P satisfying, for all $x, y, z \in P$,

1. $x \leq x$, (reflexivity)
2. if $x \leq y$ and $y \leq x$, then $x = y$, (antisymmetry)
3. if $x \leq y$ and $y \leq z$, then $x \leq z$. (transitivity)

Definition 1.2. \mathcal{C} is a **chain** if for every $x, y \in \mathcal{C}$, either $x \leq y$ or $y \leq x$.

chain 上的元素都可以相互比较, 所以它是 totally ordered.

Definition 1.3. We say that x is **covered** by y in \mathcal{P} , written $x \prec y$, if $x \leq y$ and there is no $z \in P$ with $x \leq z \leq y$.

Definition 1.4. **Hasse diagram** for a finite partially order set \mathcal{P} : the elements of P are represented by points in the plane, and a line is drawn from a up to b precisely when $a \prec b$.



Definition 1.5. Given a partially order set, f is a **order preserving map** satisfying the condition $x \leq y$ implies $f(x) \leq f(y)$.

Definition 1.6. Given two posets (P, \leq_P) and (Q, \leq_Q) , an **order isomorphism** from (P, \leq_P) to (Q, \leq_Q) is a bijective order preserving map.

Definition 1.7. Given two posets (P, \leq_P) and (Q, \leq_Q) , an **order embedding** from (P, \leq_P) to (Q, \leq_Q) is a both order-preserving and order-reflecting map that $x \leq y \iff f(x) \leq f(y)$.

相比 order isomorphism 而言稍微弱一点, 不需要是一个 surjective.

Definition 1.8. An **ideal** I of a partially ordered set \mathcal{P} is a subset of the elements of P which satisfy the property that if $x \in \mathcal{P}$ and exists $y \in I$ with $x \leq y$, then $x \in I$.

衍生自 the ideal of ring, 后面我们将会看见 the ideal of lattice.

Definition 1.9. Given an ordered set $\mathcal{P} = (P, \leq)$. The **dual of P** is another poset $\mathcal{P}^d = (P, \leq^d)$ with the order relation defined by $x \leq^d y \iff y \leq x$.

Definition 1.10. The dual notion of an ideal is called a **filter** that F is a subset of P such $x \geq y \in F$ implies $x \in F$

类似的还有 principle ideal 和 principle filter. 就是通过一个元素生成的.

Definition 1.11. The poset \mathcal{P} has a **maximum**(element) if there exists $x \in P$ such that $y \leq x$ for all $x \in P$.

An element $x \in P$ is **maximal** if there is no element $y \in P$ with $x \leq y$ and $x \neq y$.

maximum 是一个名词表示最大值 (greatest), maximal 是一个形容词表示极大的意思. 在 poset 中可能不只有一个 maximal element.

Lemma 1.12. The following are equivalent for an poset \mathcal{P} .

1. Every nonempty subset $S \subseteq P$ contains an element minimal in S .
2. \mathcal{P} contains no infinite descending chain

$$a_0 > a_1 > a_2 > \cdots$$

这里去掉等号是指 $a_0 \neq a_1 \neq a_2 \neq \cdots$

3. If

$$a_0 \geq a_1 \geq a_2 \geq \cdots$$

in \mathcal{P} , then there exists k such that $a_n = a_k$ for all $n \geq k$.

这个 lemma 被称为 descending chain condition(DCC). 对偶地也有 ascending chain condition(ACC). original 'a partially ordered set \mathcal{P} requires that all decreasing sequences in \mathcal{P} become eventually constant'.

证明. (2) \Rightarrow (3) 前提只存在 finite descending chain. 假设 (3) 不成立, 且 $a_0 \geq a_1 \geq a_2 \geq \cdots$ 是 infinite chain. 则对于任意的 k , 都能找到 $n \geq k$ 使得 $a_n \neq a_k$ 且 $a_k \geq a_n$, 那么 $a_k > a_n$. 这样从 $k = 0, 1, 2, \cdots$ 开始我们每次都可以找到 $a_{n_0} > a_{n_1} > \cdots$. 这样我们实际构造了一个 infinite descending chain, 这是和前提矛盾的. 若 $a_0 \geq a_1 \geq a_2 \geq \cdots$ 是一个 finite chain, 它的最后一个元素显然是满足 (3), 这和假设是矛盾的.

(3) \Rightarrow (2) 也是分 infinite chain 和 finite chain 来讨论, finite 是显然的, infinite 的时候可以把它变成 finite.

(1) \Rightarrow (2) (1) 前提满足下, 假设 (2) 不成立, 即 \mathcal{P} 存在 infinite descending chain. 把这个 chain 上的元素取出来组成一个 subset S , 那么任取 a_k 都有 $a_{k+1} \leq a_k$. 即找不到 minimal.

(2) \Rightarrow (1) (2) 前提满足下, 假设 (1) 不成立. 这里需要用一下[选择公理](#)了, 定义 S 上一个选择函数 $f: S \rightarrow T$, 其中 $T \subseteq S$. 让 $a_0 = f(S)$, 递归地定义对任意的 $i \in \omega$ 有 $a_{i+1} = f(\{s \in S \mid s < a_i\})$. 接下来让这个 definition make sense, (2) 前提下 S 是没有 minimal, 所以 $\{s \in S \mid s \leq a_i\}$ 不是 empty set. 这样就找到了一个 infinite descending chain, 与假设矛盾.

(1) \Rightarrow (2) \Rightarrow (3)

(3) \Rightarrow (2) \Rightarrow (1)

done well!

□

Lemma 1.13. Let \mathcal{P} be an poset satisfyint the DCC. If $\varphi(x)$ is statement such that

1. $\varphi(x)$ holds for all minimal elements of P , and
2. whenever $\varphi(y)$ holds for all $y < x$, then $\varphi(x)$ holds,

then $\varphi(x)$ is true for every element of P .

这个 lemma 有点意思, 如果对 P 上所有的 minimal element m 都有命题 $\varphi(m)$ 成立, 且 \mathcal{P} 满足 DCC. 那么再加上一个条件: 只要对任意元素 $x \in P$, 满足 $y < x$ 都有 $\varphi(y)$ 成立. 则对任意元素 $x \in P$ 都有 $\varphi(x)$ 成立.

证明. 其实 (1) 是 (2) 的一个 special case. 在 (1)(2)hold 的情况下, 我们试想一下 $\varphi(x)$ 没有被 hold 住的是哪些元素呢? 即对于某个 x , 存在 $y < x$ 使得 $\varphi(y)$ 没有被 hold. 递归地, 我们再去考虑这个 y . 那么这里就存在一条 descending chain 在这里, 由于 \mathcal{P} 是满足 DCC, 所以这个 descending chain 是 infinite 的. 这条 chain 的结尾显然是一个 minimal element, 但是它是满足 $\varphi(x)$. 所以实际上是不存在这里的 x 不满足 $\varphi(x)$. □

Definition 1.14. Let \mathcal{P} be poset. Two elements a and b of \mathcal{P} are called **comparable** if $a \leq b$ or $a \geq b$. Otherwise, they are called **incomparable**.

[可比性](#).

Definition 1.15. An **antichain** in \mathcal{P} is a subset A of \mathcal{P} in which each pair of different element are incomparable.

Definition 1.16. Define the **width** of an poset \mathcal{P} by

$$w(\mathcal{P}) = \sup\{|A| \mid A \text{ is an antichain in } \mathcal{P}\}$$

where $|A|$ denotes the cardinality(集合的势) of A .

Definition 1.17. We define the **chain-covering-number** CCN $c(\mathcal{P})$ to be the least cardinal number k , such that P is a union of k chains(finite) of P , means $P = \bigcup C_i$

另一种 covering number, 有趣.

Lemma 1.18. Suppose $P = \bigcup C_i$ where $i \in I$, then $w(\mathcal{P}) \leq |I|$.

证明. 因为 $|A \cap C_i| \leq 1$ for $i \in I$. 也就是说你把 A 里面的元素分开塞到 C_i 上, 每次都只能塞一个. 那么最多你可以每个 C_i 上都塞一个. \square

Theorem 1.19. (Dilworth, 1950) Let \mathcal{P} be a finite poset. $w(\mathcal{P})$ is width. Then \mathcal{P} is a union of $w(\mathcal{P})$ -chains.

证明. TODO. \square

Semilattices, Lattices and Complete Lattices

Semilattice

Definition 2.1. A **semilattice** is an algebra $\mathcal{S} = (S, *)$ satisfying, for all $x, y, z \in S$,

1. $x * x = x$,
2. $x * y = y * x$,
3. $x * (y * z) = x * (y * z)$.

where $*$ is binary operator. 换句话说 **semilattice** 就是一个 **idempotent commutative semigroup**(幂等交换半群).

Theorem 2.2. In a semilattice \mathcal{S} , define $x \leq y$ if $x * y = x$. Then (S, \leq) is a poset in which every pair of elements has a greater lower bound.

Conversely, given an poset P with that property, define $x * y = g.l.b(x, y)$. Then $(P, *)$ is a semilattice.

semilattice 上弄了一个特殊的 **poset** 出来, 它最好的性质就是任意两个元素都有一个下确界. 把 $*$ 换成 \cap , 然后把 \leq 换成 \subseteq , 可能就很熟悉了. A semilattice with the above ordering is usually called **meet semilattice**.

证明. 先证明这个是一个 poset.

1. $x * x = x$ implies $x \leq x$,
2. if $x \leq y$ and $x \geq y$, then $x = x * y = y * x = y$,
3. if $x \leq y$ and $y \leq z$. then $x * z = (x * y) * z = x * (y * z) = x * y = x$, so $x \leq z$.

这个 greater lower bound 就是 $x * y$. 首先证明它是一个 lower bound, $x * (x * y) = x * y$ and $y * (x * y) = x * y$, 所以 $x * y$ 是一个 lower bound. 再来证明所有的 lower bound 都比它小, 假设 $z \leq x$ 和 $z \leq y$, 即 z 是 $\{x, y\}$ 的一个 lower bound. 那么 $z * (x * y) = z * y = z$, 所以 $z \leq (x * y)$. 最后 $x * y$ 的一个 greater lower bound.

□

对偶地, 使得 $x \geq y \iff x * y = x$, 则称 \mathcal{S} 为是一个 **join semilattice**. 自然地在 (S, \leq) 下任意的 pair 都有一个 least upper bound $x \vee y$.

Definition 2.3. A **homomorphism** between two semilattice is a map $f: \mathcal{S} \rightarrow \mathcal{T}$ with the property that $f(x * y) = f(x) * f(y)$. An **isomorphism** is a homomorphism that injective and surjective.

nothing new...

Theorem 2.4. Let \mathcal{S} be a meet semilattice. Define $\phi: \mathcal{S} \rightarrow \mathcal{O}(\mathcal{S})$ by

$$\phi(x) = \{ y \in \mathcal{S} \mid y \leq x \}$$

where $\mathcal{O}(\mathcal{S})$ is collection of all order ideals of \mathcal{S} . Then \mathcal{S} is isomorphic to $(\mathcal{O}(\mathcal{S}), \cap)$ (注意这里是 \mathcal{S} 的 image).

怎么感觉这些 ideal 都是 principle ideal.

证明. \cap 表示 set inclusion, ϕ 是 order-preserving 和 order-reflecting 还是比较 obvious. 所以 ϕ 是一个 order embedding of \mathcal{S} into $\mathcal{O}(\mathcal{S})$. Moreover $\phi(x \wedge y) = \phi(x) \cap \phi(y)$ because $x \wedge y$ is the greatest lower bound of $\{x, y\}$, so that $z \leq x \wedge y$ if and only if $z \leq x$ and $z \leq y$. \square

Lattice

Definition 2.5. A **lattice** is an algebra $\mathcal{L} = (L, \wedge, \vee)$ satisfying, for all $x, y, z \in L$,

1. $x \wedge x = x$ and $x \vee x = x$,
2. $x \wedge y = y \wedge x$ and $x \vee y = y \vee x$,
3. $x \wedge (y \wedge z) = (x \wedge y) \wedge z$ and $x \vee (y \vee z) = (x \vee y) \vee z$,
4. $x \wedge (x \vee y) = x$ and $x \vee (x \wedge y) = x$.

就第四个在我们眼里似乎没有那么自然, 它叫 absorption laws(吸收律), 它在这里可以保证后面 \wedge 和 \vee 定义了相同的 partial order(虽然是 dual). 前三个我们知道 lattice 同时在两种 binary operator 都是 semilattice, 所以我们只要在 lattice 上定义前面合适的 partial order, 它就是 both meet and join semilattice.

Theorem 2.6. In a lattice \mathcal{L} , define $x \leq y$ if and only if $x \wedge y = x$. Then (L, \leq) is a poset in which every pair of elements has a greatest lower bound and a least upper bound.

证明. 给定一个 pair (x, y) . 前面已经证明了 $x \wedge y$ 是它的一个 greater lower bound. 再根据 lattice definition 的第四条的第一个式子, $x \vee y$ 是它的一个 upper bound, 第二式子说明当 $x \geq y$ 时, 有 $x \vee y = x$, 对偶地 $x \vee y$ 是 least upper bound.

这里若 $x \wedge y = x$, 则 $x \vee y = (x \wedge y) \vee y = y$. 类似地 $x \vee y = y$, 则 $x \wedge y = x \wedge (x \vee y) = x$. 所以有一个很重要的结论就是 $x \wedge y = x \iff x \vee y = y$. \square

类似的我们可以通过一个 poset 构造 lattice.

Theorem 2.7. Given an poset \mathcal{P} with that above property, define $x \wedge y = \sup\{x, y\}$ and $x \vee y = \inf\{x, y\}$. Then (P, \wedge, \vee) is a lattice.

所以实际上 lattice 可以有两种定义第一种是前面的代数定义, 第二种就是在 poset 上定义 join 和 meet 操作, 这一点要清楚.

the definitions of sublattice, homomorphism and isomorphism).

Complete Lattice

Definition 2.8. For a subset A of a poset P , let A^u denote the set of all upper bounds of A ,

$$\begin{aligned} A^u &= \{x \in P \mid x \geq a \text{ for all } a \in A\} \\ &= \bigcap_{a \in A} \uparrow a \end{aligned}$$

where $\uparrow a = \{x \in P \mid x \geq a\}$. Dually, A^l is the set of all lower bounds of A ,

$$\begin{aligned} A^l &= \{x \in P \mid x \leq a \text{ for all } a \in A\} \\ &= \bigcap_{a \in A} \downarrow a \end{aligned}$$

where $\downarrow a = \{x \in P \mid x \leq a\}$.

思考一个问题 poset P 的一个 subset A 什么时候 least upper bound? 很显然 A^u 一定不是空的, 更确切地说 A^u 有一个 greatest lower upper z , 而且 $z \in A^u$, 根据 z 的 definition 它是 A 的 least upper bound. 这种情况下我们就说 the join of A exists, and write $z = \bigvee A$. 对偶地, 考虑 A 的 greatest lower bound, 则 A^l 一定不为空, 那么 A^l 里面是有一个 lower upper bound 的 w , 根据 w 的 definition 它是 A 的 greatest lower bound. 这种情况下我们就说 the meet of A exists, and write $w = \bigwedge A$.

这样我们 define 两个特殊的 meet 和 join 作用在一个集合上.

Theorem 2.9. Let \mathcal{S} be a finite meet semilattice with greatest element 1. Then \mathcal{S} is a lattice with join operation defined by

$$x \vee y = \bigwedge \{x, y\}^u = \bigwedge (\uparrow x \cap \uparrow y).$$

证明. \mathcal{S} 有 greatest element, 则 A^u 肯定不是空了, 至少这个 greatest element 里面. $\bigwedge A^u$ 就是要找 A^u 的 lower upper bound. 由于 \mathcal{S} 是一个 finite lattice, 所以 A^u 也是 finite. A^u 里面的元素做有限次 meet 操作得到就是一个 lower upper bound, 但是你还得说明它在 A^u 里面. 这是很显然的, $x \wedge z_1 \wedge \cdots \wedge z_k = x$ 其中 $z_i \in A^u$, 所以 $\bigwedge A^u$ 是它的一个 upper bound.

还得 proof 一下它是一个 lattice, 上面只是证明了这个东西是 well behaved. Lattice definition 中前三条还是比较明显的.

$$x \wedge (x \vee y) = x$$

这也很显然, 因为 $x \vee y \in \{x, y\}^u$.

$$x \vee (x \wedge y) = x$$

因为 $x \wedge y$ 是 $\{x, y\}$ 的一个 greatest lower bound, 有 $x \geq x \wedge y$, 那么 $\inf(x, x \wedge y) = x$. □

这个 theorem 告诉我们: if a finite poset P has a greatest element and every pair of elements has a meet, then P is a lattice.

Theorem 2.10. Every finite subset of a lattice has a greatest lower bound and a least upper bound.

证明. \mathcal{L} 是 finite, 则它的 subset 也是 finite. 前面我们知道 lattice 中任意一个 pair 都有 greatest lower bound 和 least upper bound, 这是 meet 和 join 定义下的 partial order 所带来的性质. 在 finite subset 里面先挑两个出来做 meet 或者 join 可以得到 inf 和 sup 它们也是属于 L 的, 再从剩下的 subset 里面再挑一个出来做同样的操作, 这个操作只会做有限多次, 所以最终我可以得到这个 subset 的 greatest lower bound 和 least upper bound. \square

这个性质在 infinite lattice 下可能就无法成立. 例如 infinite subset 上述操作可能根本就停不下. 由此我们定义另外一个概念.

Definition 2.11. Given poset \mathcal{L} . If every subset A of \mathcal{L} has a greatest lower bound $\bigwedge A$ and a least upper bound $\bigvee A$, then \mathcal{L} is called complete lattice.

subset 其中就包含了 pair, 所以它是一个 lattice 还是比较明显的. 此外, finite lattice 是 complete 的, 并且所有 complete lattice 都包含 greatest element 和 least element.

Definition 2.12. a complete meet semilattice is an poset \mathcal{S} with greatest element and the property that every nonempty subset A of S has a greatest lower bound $\bigwedge A$.

下面这个 theorem 让我们抛弃了更强的 finite, 只需要 complete 就可以在 meet semilattice 上构造一个 lattice 出来.

Theorem 2.13. If \mathcal{L} is a complete meet semilattice, then \mathcal{L} is a complete lattice with the join operation defined by

$$\bigvee A = \bigwedge A^u = \bigwedge (\bigcap_{a \in A} \uparrow a).$$

证明. 和前面在 finite meet semilattice 上构造 lattice 类似, 这里 finite 换成了 complete. 这里我们直接就可以知道 $\bigwedge A^u$ 是有意义的, 前面也证明了它还是一个 A 的 upper bound. 那么 $\bigwedge A$ 的 definition 是满足 A 的 least upper bound. \square

Closure System

Definition 2.14. A **closure system** on a set X is a collection \mathcal{C} of subsets of X that is closed under arbitrary intersections(任意的交). The sets in \mathcal{C} are called closed set.

Example 2.15. 有一些 closure system 的例子

1. closed subsets of topological space,
2. subgroups of group,
3. subspace of vector space.
4. convex subsets of euclidean space \mathbb{R}^n ,
5. order ideals of an poset.

By convention, 把 $\bigcap \emptyset = X$ 的话, closure system 配合集合操作就是一个 complete lattice.

Definition 2.16. A **closure operator** on a set X is a map $\Gamma: \mathfrak{P}(X) \rightarrow \mathfrak{P}(X)$ satisfying, for all $A, B \subseteq X$,

1. $A \subseteq \Gamma(A)$,
2. $A \subseteq B$ implies $\Gamma(A) \subseteq \Gamma(B)$,
3. $\Gamma(\Gamma(A)) = \Gamma(A)$.

closure 的 general definition 有点意思. 简单地来说, 就是 (1) 集合的闭包包含集合本身; (2) 如果两个集合有包含关系, 则它们的闭包也有相同的包含关系; (3) 闭包的闭包是其本身.

如何在一个在 X 上利用已知的 closure system 构造一个 closure operator, 即让 X 上的任意一个子集对应一个 closure?

Theorem 2.17. If \mathcal{C} is a closure system on a set X , then the map $\Gamma_{\mathcal{C}}: \mathfrak{P}(X) \rightarrow \mathfrak{P}(X)$ defined by

$$\Gamma_{\mathcal{C}}(A) = \bigcap \{ D \in \mathcal{C} \mid A \subseteq D \}$$

is a closure operator. Moreover $\Gamma_{\mathcal{C}}(A) = A$ if and only if $A \in \mathcal{C}$

所有包含这个集合的 closed set 的交是这个集合的 closure. 在 topology 里面 closure 是包含这个集合最小的 closed set.

Definition 2.18. A set of **closure rules** on a set X is a collection \sum of properties $\varphi(S)$ of subsets of X . where each $\varphi(S)$ has one of the forms

$$x \in S$$

or

$$Y \subseteq S \Rightarrow z \in S$$

with $x, z \in X$ and $Y \subseteq X$. A subset D of X is said to be closed with respect to these rules if $\varphi(D)$ is true for each $\varphi \in \sum$.

你看到这里一定会感觉非常的困惑，closure rules 到底是个啥东西？

Example 2.19. 对应前面列举到的 closure system.

1. In topological space, all rules $Y \subseteq S \Rightarrow z \in S$ where z is an accumulation point of Y .
2. In subgroup, the rule $1 \in S$ and all rules

$$x \in S \Rightarrow x^{-1} \in S \{x, y\} \in S \Rightarrow xy \in S$$

3. In vector space, $0 \in S$ and all rules $\{x, y\} \subseteq S \Rightarrow ax + by \in S$ with a, b scalars.

closure rules 就是一系列判断 closed set 的命题.

Theorem 2.20. If Γ is a closure operator on a set X , \sum_{Γ} be the set of (1)all rules where $c \in \Gamma(\emptyset)$, and (2)all rules

$$Y \subseteq S \Rightarrow z \in S$$

with $z \in \Gamma(Y)$. Then a set $D \subseteq X$ satisfies all the rules of \sum_{Γ} if and only if $\Gamma(D) = D$.

证明. (1) 是在说 \emptyset 的 image 非空? (2) 是在说取 S 上任意子集 Y , 都有 $\Gamma(Y) \subseteq S$. 直觉上就说这群规则是一个 closure rule, 满足它的只有 closed set, 自然地在 closure operator 上也是一个 closed set. 可这条件都 TM 都 abstract nonsense 了!!!

我尝试用 closure operator 的 definition 来推一下, $Y \subset S$, 结合前面那么有

$$Y \subseteq \Gamma(Y) \subseteq S.$$

特殊点, 把 Y 换成 S , 有 $S \subseteq \Gamma(S) \subseteq S$, 所以 $\Gamma(S) = S$. 但是 (1) 在这里有啥用啊? 保证 S 非空?

反过来若 $\Gamma(D) = D$. 自然地, 当 $Y \subseteq D$, 则有 $Y \subseteq \Gamma(Y) \subseteq \Gamma(D) = D$. □

说实话上面这个命题有点 abstract, 当有一个 closure operator 之后, 我们知道 closed set 在它的作用下是它本身, 但是能不能找到一个 closure rule 来和它对应呢?

Theorem 2.21. If Σ is a set of closure rules on set X , let \mathcal{C}_Σ be the collection all subsets of X that satisfy all the rules of Σ . Then a set \mathcal{C}_Σ is a closure system.

证明. 这个定理可以更形象地去理解 closure rule 到底是什么? 假设 A, B 是满足 Σ 里面所有 rules 的两个集合. 我们看它们的交, 对于第一类规则 $x \in S$, 很显然在交下是保持的, 因为 $x \in A$ 和 $x \in B$, 则 $x \in A \cap B$. 对于第二类的规则, 若 $C \subseteq A \cap B$, 且它是某个规则里面对应的 Y , 那么 $C \subseteq A$ 和 $C \subseteq B$, 对应地有某个 $z \in A$ 和 $z \in B$, 所以 $z \in A \cap B$. 综上 $A \cap B$ 也是属于 \mathcal{C}_Σ . \square

在这里我们才终于认识到这样 closure rules 这样抽象的东西, 它确实可以刻画一堆 closed set 组成了一个 closure system.