

Categorical Semantics for STLC

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First Try

Category of Baby Type Theory

Definition 1.1. A *baby type system* with only atomic types and where typing context are singletons.

Definition 1.2. A interpretation of baby type system is category \mathbf{C}_{bT} such that objects are interpretations of types and morphism are interpretations of term-in-context(sequent).

- *objects*: $\llbracket A \rrbracket$ where A is atomic type;
- *morphism*: $\llbracket \Gamma \vdash E : A \rrbracket : \llbracket \Gamma \rrbracket \rightarrow \llbracket A \rrbracket$, abbreviate it to $\llbracket E \rrbracket : \llbracket \Gamma \rrbracket \rightarrow \llbracket A \rrbracket$;
- *identity*: $\llbracket x : A \vdash x : A \rrbracket = \text{id}(\llbracket A \rrbracket) : \llbracket A \rrbracket \rightarrow \llbracket A \rrbracket$, it corresponds to

$$\overline{x : A \vdash x : A} \text{ VAR}$$

- *composition*: $\llbracket E_2[y \rightarrow E_1] \rrbracket = \llbracket E_2 \rrbracket \circ \llbracket E_1 \rrbracket : \llbracket A \rrbracket \rightarrow \llbracket C \rrbracket$, it corresponds to

$$\frac{x : A \vdash E_1 : B \quad y : B \vdash E_2 : C}{x : A \vdash E_2[y \rightarrow E_1] : C} \text{ Sub}$$

- *unit law*:

$$\begin{aligned} \llbracket x : A \vdash E[x \rightarrow x] \rrbracket &= \llbracket x : A \vdash E : B \rrbracket \circ \llbracket x : A \vdash x : A \rrbracket \\ \llbracket x : A \vdash E[y \rightarrow E] \rrbracket &= \llbracket y : B \vdash y : B \rrbracket \circ \llbracket x : A \vdash E : B \rrbracket \end{aligned}$$

- *associative law*: Given $x : A \vdash E_1 : B, y : B \vdash E_2 : C, z : C \vdash E_3 : D$, we have

$$\llbracket E_3[z \rightarrow E_2[y \rightarrow E_1]] \rrbracket = \llbracket E_3[z \rightarrow E_2][y \rightarrow E_1] \rrbracket : \llbracket A \rrbracket \rightarrow \llbracket D \rrbracket$$

参考文献

- [1] Edward Morehouse. Basic Category Theory. OPLSS, 2016.