# 211: Computer Architecture Fall 2019

#### Topic:

- C programming
- Data representation

# **Debugging C**

```
#include <stdlib.h>
int main(int argc, char* argv[])
  int* p = malloc(10 * sizeof(int));
  for (int i = 0; i <= 10; i++)
     p[i] = i;
```

# **Debugging C**

- printf("Reached here\n");
- -Wall
- -Werror
- -fsanitize=address
- -g
- valgrind
- gdb

# **Valgrind**

- Run as "valgrind ./a.out" (where a.out is your binary)
- Finds invalid reads
- Finds invalid writes
- Finds leaks
- Any program you submit should have 0 of these!

# **Checking output**

- Piping: take output from one program and give it as input to another – "./a.out | wc -l"
- Redirect output to a file: "./a.out > out"
- Compare two files: "diff myOutput correctOutput"
- See exact file contents: "hexdump -c myFile"

## GDB – The GNU debugger

- Run as "gdb a.out"
- Then "run arg1 arg2 arg3", where the args are commandline arguments for your program
- To stop before the program finishes, set a breakpoint: "break main"
- Once you hit a breakpoint, use "next" to step one line of code at a time
- Use "print x" to see the contents of a variable x
- Use "backtrace" to see the current list of functions called
- Use "frame X" to select frame X from this list

# **Classic Memory Bugs**

Memory management is one of the biggest differences between C and Java

Here are some classic bugs that might afflict you

The classic scanf bug

scanf("%d", val);

#### Reading Uninitialized Memory

Assuming that heap data is initialized to zero

```
/* return y = Ax */
int *matvec(int **A, int *x) {
   int *y = malloc(N*sizeof(int));
   int i, j;
   for (i=0; i<N; i++)
      for (j=0; j<N; j++)
      y[i] += A[i][j]*x[j];
   return y;
}</pre>
```

#### **Overwriting Memory**

Allocating the (possibly) wrong sized object

```
int **p;
p = malloc(N*sizeof(int));
for (i=0; i<N; i++) {
   p[i] = malloc(M*sizeof(int));
}</pre>
```

#### **Overwriting Memory**

Off-by-one error

```
int **p;
p = malloc(N*sizeof(int *));
for (i=0; i<=N; i++) {
   p[i] = malloc(M*sizeof(int));
}</pre>
```

#### **Overwriting Memory**

Misunderstanding pointer arithmetic

```
int *search(int *p, int val) {
   while (*p && *p != val)
      p += sizeof(int);
   return p;
}
```

#### Referencing Nonexistent Variables

Forgetting that local variables disappear when a function returns

```
int *foo () {
  int val;
  return &val;
}
```

#### Freeing Blocks Multiple Times

```
x = malloc(N*sizeof(int));
<manipulate x>
free(x);

y = malloc(M*sizeof(int));
<manipulate y>
free(x);
```

#### Referencing Freed Blocks

```
x = malloc(N*sizeof(int));
<manipulate x>
free(x);
...
y = malloc(M*sizeof(int));
for (i=0; i<M; i++)
  y[i] = x[i]++;</pre>
```

#### Failing to Free Blocks (Memory Leaks)

Slow, long-term killer

```
foo() {
  int *x = malloc(N*sizeof(int));
  ...
  return;
}
```

#### Failing to Free Blocks (Memory Leaks)

Freeing only part of a data structure

```
struct list { int val; struct list *next;};
foo() {
   struct list *head = malloc(sizeof(struct list));
   head->val = 0;
   head->next = NULL;
   <create and manipulate the rest of the list>
   ...
   free(head);
   return;
}
```

# What Do Computers Do?

Manipulate stored information

Information is data

How is it represented?

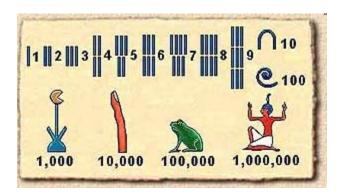
Basic information: numbers

Human beings have represented numbers throughout history

- Egyptian number system
- Roman numeral

Typically decimal

Natural for humans



Discoveregypt.com

I	VI	XI XII	-	XXXI	
-	## 1	XIII	***	XXXIII	
II	VII	XIV		XXXIV	
		XV		XXXV	
II	ATT	XVI		XXXVI	
II.	100.00	XVII		XXXVII	
TIT	VIII	XVIII		XXXVII	
				XXXIX XL	
III	4 111	XXI		XLI	-
III	HH III	XXII		XLII	
***		XXIII		XLIII	
IV	IX	XXIV		XLIV	==
III	HH 1111	XXV		XLV	==
<b>X</b> 7		XXVI		XINI	==
		XXVII		XLVII	==
Y	X	XXVIII		XLVIII	
###	****	XXX		L	
L *****	X ***	MMV		мм	
C	$\overline{L}_{-x+v+v+v}$	MMIV		MCMXCIX	
D	C 1-1	ммш		MCMXCVIII	
M	D receive	ммп		MCMXCVII	
V MANAGERIA	<u>M</u> >	MMI		MCMXCVI	

# **Number System**

#### Comprises of

- Set of numbers or elements
- Operations on them
- Rules that define properties of operations

#### Need to assign value to numbers

#### Let us take decimal

- Base 10
- Numbers are written as d<sub>n</sub>...d<sub>2</sub>d<sub>1</sub>d<sub>0</sub>
- Each digit is in [0-9]
- Value of a number is interpreted as  $\sum_{i=0}^{n} d_i \times 10^i$

# **Binary Numbers**

Base  $2 \Rightarrow$  each digit is 0 or 1

Numbers are written as d<sub>n</sub>...d<sub>2</sub>d<sub>1</sub>d<sub>0</sub>

Value of number is computed as  $\sum_{i=0}^{n} d_i \times 2^i$ 

Binary representation is used in computers

- Easy to represent by switches (on/off)
- Manipulation by digital logic in hardware

Written form is long and inconvenient to deal with

#### **Hexadecimal Numbers**

Base 16

Each digit can be one of 16 different values

 $\blacksquare$  Symbols = {0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F}

First 10 symbols (0 through 9) are same as decimal

■ A=10,B=11,C=12, D=13, E=14, F=15

Numbers are written as d<sub>n</sub>...d<sub>2</sub>d<sub>1</sub>d<sub>0</sub>

Value = 
$$\sum_{i=0}^{n} d_i \times 16^i$$

#### **Octal Numbers**

Base 8  $\Rightarrow$  each digit is in [0-7] Numbers are written as  $d_n...d_2d_1d_0$ Value of number is computed as  $\sum_{i=0}^n d_i \times 8^i$ 

# **Converting Hex to Binary**

Each hexadecimal digit can be represented by 4 binary digits

Why?

0x2A8C (hex) = 0b0010101010001100 (binary)

- $0xC = 12 \times 16^{0} = 8 + 4 = (1 \times 2^{3}) + (1 \times 2^{2}) = 0b1100$
- $0 \times 80 = 8 \times 16^1 = 2^3 \times 2^4 = 2^7 = 0b \ 10000000$
- And so on ...

So, to convert hex to binary, just convert each digit and concatenate

What about octal to binary?

# **Converting Binary to Hex**

#### Do the reverse

- Group each set of 4 digits and change to corresponding digit in hex
- Go from right to left

Example 1011011110011100

 $\blacksquare$  0b1011011110011100 = 0xB79C

What about binary to octal?

## **Decimal to Binary**

#### What's the largest p, q, r ... such that

- $n = 2^p + r_1$ , where  $r_1 < 2^p$
- n  $2^p = 2^q + r_2$ , where  $r_2 < 2^q$
- $n (2^p + 2^q) = 2^r + r_3$ , where  $r_3 < 2^r$
- ...

#### The above means that

- $n = (1 \times 2^p) + (1 \times 2^q) + (1 \times 2^r) + ... + r_m$ , where  $r_m = n \% 2$
- Can you see why this now allows n to be easily written in binary form?

#### Example: convert 21 to binary

$$\blacksquare$$
 21 = 2<sup>4</sup> + 5, 5 = 2<sup>2</sup> + 1 ⇒ 21 = 0b10101

# **Decimal to Binary and Back**

How to do the conversion algorithmically?

What about binary to decimal?

What about decimal to hex? Hex to decimal?

Decimal to octal? Octal to decimal?

Hex to octal? Octal to Hex?

# **Decimal and Binary fractions**

In decimal, digits to the right of radix point have value 1/10<sup>i</sup> for each digit in the i<sup>th</sup> place

-0.25 is 2/10 + 5/100

Similarly, in binary, digits to the right of radix point have value  $1/2^i$  for each i<sup>th</sup> place

Just the base is different

8.625 is 1000.101

 $\blacksquare$  .625 = 6/10 + 2/100 + 5/1000 = 1/2 + 1/8

How to convert?

# **Decimal to Binary Example**

```
Algorithm

Number = decimalFraction
while (number > 0) {
    number = number*2
    if (number >=1) {
        Output 1;
        number = number-1
    }
    else {
        Output 0
    }
}
```

Why does it work?

Example: 0.625 to binary

- ANS: 0.101
  - $\bullet$  0.625\*2 = 1.25
  - output 1
  - $\bullet$  0.25\*2 = 0.5
  - output 0
  - $\bullet$  0.5\*2 = 1
  - output 1
  - Exit

# C shift operations

- x << n shift x left by n bits</li>
- x >> n shift x right by n bits
- A left shift = multiplying by 2
- A right shift = dividing by 2 (and discarding remainders)

- 0x01 << 1 == 0x02
- 25 << 2 == 100</li>
- 25 >> 1 == 12

# C mask operations

- x & m do a bitwise AND operation
- x | m do a bitwise OR operation
- x ^ m do a bitwise XOR operation
- ~x do a bitwise NOT operation

- 0x01 & 0x01 = 0x01
- $0x11 \mid 0x44 = 0x55$
- $0x11 \mid 0x14 = 0x15$
- $0x55 ^ 0x53 = 0x06$

#### C shift and mask

• Given the integer 0x11223344, grab the 3<sup>rd</sup> byte (0x33)

- int x = 0x11223344;
- int y = (x >> 8) & 0xff;