DUBLIN CITY UNIVERSITY

ELECTRONIC AND COMPUTER ENGINEERING

EE562 - Network Stack Implementation

Assignment 2



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1 Introduction

The Red-Black Tree is an implementation of the Binary Search Tree which maintains balance. Nodes of a Red-Black tree are either coloured red or black, hence the name. By ensuring that there are never two directly connected red nodes, rotate operations ensure that the tree is always balanced.

The concept for this type of binary search tree was instroduced by Rudolf Bayer in 1972[1]. They are particularly useful in real-time applications, and as such were included in the implementation of the "Completely Fair Scheduler" in the Linux Kernel in version 2.6.23.

The main advantage of using Red-Black trees is the time complexity. Red-Black trees have a time complexity of O(logN) for insertion and deletion operations, improving upon the typical complexity of O(N) for more basic Binary Search Tree implementations[1].

For each node in a tree there is a set colour, up to one parent node, and up to two child nodes. The root node if a red-black tree is always coloured black, along with any null nodes (i.e. empty/leaf nodes). Two red nodes can not be directly connected[2]. Each left child node in a tree must have a value less than the parent node, and each right child node in a tree must have a value greater than the parent node. Maintaining balance within the tree means that the case of the root node being the lowest value in a constantly incrementing tree, i.e. all child nodes are in the right branch, cannot occur. This balancing means that the time complexity of traversing the tree does not increase past O(logN).

2 Red-Black Tree Operation

The main feature of a red-black tree is it's ability to maintain balance. It does this through a number of operations. On adding a new value to to tree, an Insertion opera-

tion is used to define where within the tree the new value will be placed. On deleting a value from the tree, the tree must be restructured in order to allow any child node of the deleted node to remain connected to the rest of the tree. Recoloring is a technique which is unique to Red-Black Trees, and allows for a tree to be rebalanced without using a rotation operation. If recoloring is not successful, rotations are used following these operations to ensure that balance is maintained within the tree.

2.1 Insertion

Insertions are done based on the value of the incoming node. The incoming node value is compared agains the root node value. If the incoming node value is greater than the root node value, the comparison is repeated with the root nodes right child node. If the incoming node value is less than the root nodes value, the comparison is repeated with the root nodes left child node[3]. This operation is typically called recursively in order to traverse the tree until the point that there is no left or right child node with which to compare, at which point the incoming node can be placed in the tree.

Following the insertion of the incoming node to the tree, the tree must be rebalanced. This is to maintain the traversal speed of the tree. Rebalancing is done via either recoloring or rotation, as described below.

2.2 Deletion

Deletion within a Red-Black tree involves removing the specified value from the tree. The specified value is compared to the root node value, and, if greater than the root node value, a comparison is made with the right child node value, and if less than the root node value, a comparison is made with the left child node value. This operation is called recursively until the specified node value is found. This node is then deleted from the tree.

If the specified node value is a leaf node, or a node with one child, then this deletion

is relatively simple. If the node has one child, then the child node replaces the deleted node within the tree and a recolouring or rotation must be performed. If the node has no children, i.e. is a leaf node, then the node is simply deleted[4].

For any node deletion in which the specified node is not a leaf node, or a node with only one child, i.e. the specified node is "internal", the order of all child nodes must be retained, and after deletion must be reordered into the existing tree[4]. This is, therefore, a much more complex operation.

2.3 Tree Balancing

Tree balancing is required in order to maintain the traversal time complexity within a red-black tree. The two techniques used in tree balancing are "Recoloring" and "Rotation". Recoloring involves changing the colours associated with the nodes in the tree, whereas rotation involves moving the nodes of the tree, replacing the existing root node.

2.3.1 Recoloring

There are a number of cases for recoloring of nodes within a red-black tree. As the root node is required to be black, and two red nodes cannot be adjacent, the recoloring must be done in such a way as to maintain these rules.

Any new node being inserted to the tree is initially coloured red. The parent node colour is then checked. If this node is red, the rule stating that two red nodes cannot be directly connected is being violated. The parent node color is changed to black, and the colour of the parents parent (i.e. the grandparent), is changed to red. Following this, rotations may be required to make a valid red-black tree, for example if the grandparent node is the root of the tree (which must be black)[3].

2.3.2 Rotation

Rotations are used within the red-black tree to rebalance the tree without modifying the overall order of the elements in the tree. When a node is inserted, recoloring is first performed, however, if this does not result in a valid, balaced, red-black tree, a rotation operation must be completed.

Following recolouring, if the root node is red, a left or right rotation must occur. If the newly added node is in the left branch, a right rotation occurs, and if in the right branch, a left rotation occurs[3]. These rotations may be required to be done recursively in order to acquire a valid and balanced red-black tree.

3 Red-Black Tree Code

The aim of the code portion of this assignment is to implement a Red-Black Tree which can store paired "finish number" and "VPN" values. The Red-Black Tree is required to sort the pairs by finish number. In order to ensure that the code is working as intended, 100 input pairs must be inserted to the tree, with the original input pairs and the tree following all insertions being printed to verify its operation.

Provided for the purposes of the coding section of this assignment were a selection of random number generating functions, along with an implementation of a red-black tree[5] which needs to be modified. The red-black tree implementation takes integer values for its elements, however, the output from the random number generator is of type long. The required implementation must also be able to perform insertions of two values.

In order to store both the "finish key" and the "vpn" values, a C structure named "node_entry" was defined in the "node_entry.h" file. In this header file, the finish key is defined as an unsigned long, while the vpn is defined as an unsigned int. A function to display the values of the struct is defined within this header file, and is implemented

within the "node_entry.c" file.

```
struct node_entry
{
    unsigned long finish_key;
    unsigned int vpn;
};
```

Figure 1: Node Entry Structure

Within the "redblack.h" file, the type definition for "ElementType" must be changed from type int to type "struct node_entry". This allows for the finish key/vpn pairs to be passed in as arguements to the functions of the red-black tree, such as insertion etc.

```
while (X->Element.finish_key != Item.finish_key)
/* Descend down the tree */
{
    GGP = GP;
    GP = P;
    P = X;
    if (Item.finish_key < X->Element.finish_key)
        X = X->Left;
    else
        X = X->Right;
    if (X->Left->Color == Red && X->Right->Color == Red)
        HandleReorient(Item, T);
}
```

Figure 2: Example of Modifications to the Red-Black Tree Code

Within the "redblack.c" file, all references to "Element", which is the value in the tree, must be replaced with "Element.finish_key". This allows for all sorting within the tree to be done based on the finish key, rather than on the vpn. The return type of the "Retrieve" function must also be changed to "long", as it is no longer returning an

integer value, rather the value of the finish key, which is long.

Within the main function, the tree and random number generator are both initialized. The tree is initialized as an empty red-black tree, and the random number generator is initialised with the seeds (0, 0, 1541, 0402). Within the "init" function, the seeds are set to (1, 2, 3, 4). An empty array of type node_entry is defined, with a size of 100. A loop iterates from 0 to 100, populating the finish key values with random long values from 0-999, using the "fairdie" function. The vpn values are populated with the iteration value, i.e. 0-99. Within this loop, each new node_entry item is inserted into the tree using the "Insert" function. The "PrintTree" function is then used in order to print the sorted contents of the tree, which verifies the correct operation of the tree.

4 Conclusion

Red-Black Trees offer traversal time complexity in the order O(logN). While the implementation may be complex, the benefits of the red-black tree are clear, and its benefits in real-time applications are why it was included in the Linux Kernel's "Completely Fair Scheduler".

The red-black tree code shows how the tree can be used to sort values, while maintaining balance, and the time complexity of the tree. An issue was experienced in the code, when using the randomly generated numbers the tree would overwrite each number in the tree with the last input number, resulting in only one number (the final input number) in the completed tree. As no changes were made to the structure of the tree's code, this issue was unable to be resolved.

References

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