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USING GRAPHIC NETWORK SIMULATOR 3 FOR DDOS ATTACKS SIMULATION

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Abstract: Distributed Denial of Service (DDoS) attacks are still one of the major cybersecurity threats and the focus of much research on developing DDoS attack mitigation and detection techniques. Being able to model DDoS attacks can help researchers develop effective countermeasures. Modeling DDoS attacks, however, is not an easy task because modern DDoS attacks are huge and simulating them would be impossible in most cases. That's why researchers use tools like network simulators for modeling DDoS attacks. Simulation is a widely used technique in networking research, but it has suffered a loss of credibility in recent years because of doubts about its reliability. In our previous works we used discrete event simulators to simulate DDoS attacks, but our results were often different from real results. In this paper, we apply our approach and use Graphical Network Simulator-3(GNS3) to simulate an HTTP server's performance in a typical enterprise network under DDoS attack. Also, we provide references to related work. Copyright © Research Institute for Intelligent Computer Systems, 2017. All rights reserved.

Keywords: GNS3;DDoS attack; network simulator.

1. INTRODUCTION

Despite over a decade of research into DDoS attack detection ([1], [2], [3]), mitigation ([4], [5], [6]), and advanced source detection ([7], [8], [9]), these attacks are still one of the most dangerous threats to computer networks. Modern DDoS attacks can vary in size from several PCs to huge botnets consisting of tens of thousands of PCs from all over the world. The DDoS attack on Russian banks in 2016 was carried out by a huge botnet. Being able to model DDoS attacks is helpful in developing new techniques for mitigating them. Modeling DDoS attacks [10]-[12] in real life is not an easy task. For one thing, one must select the approach for modeling attacks. In our previous work [15] we surveyed the main approaches in this area. One can model DDoS attacks using either a specialized testbed or network simulator software. In this paper we will concentrate on the last and the most affordable option. The rest of this paper is organized as follows: in Section II we provide an overview of related work, in Section III we justify our choice of network simulator,

Section IV describes the simulation, and conclusion is in Section V.

2. RELATED WORK

A performance comparison of network simulators can be found in [17]. In [17], the authors focus on the open source simulators NS2, NS3, OMNeT++, JiST, and SimPy, and compare their performance by implementing the same model on each simulator. Performance comparison is done using two performance metrics: effective simulation runtime and memory usage. In conclusion, the authors states that ns-3, OMNeT++ and JiST are all capable of carrying out large-scale network simulations. Overall, ns-3 demonstrated the best overall performance. A detailed comparison of network simulators was done in [18], which focused on the network simulators NS2, NS3, QualNet, GloMoSim, NetSim, OMNeT++, OPNET, TOSSIM, J-SIM, NCTUns, DRMSim, SSFNet, GrooveNet, and TraNS. The paper [18] contains information about the main features, advantages, limitations, supported

OS, hardware requirements etc. of all the above mentioned simulators, also it includes comparison tables listing license types, languages, GUI types, document availability, etc. Authors [13] analyze the accuracy of NS2 and the OPNET Modeler comparing the test bed results for CBR and FTP traffic with simulated results from ns-2 and OPNET Modeler, and concluding that significant effort was required to match the simulators with the test bed. In [14] the authors compared wireless network simulators (NS2, Qualnet, and OPNET) to a real testbed. The authors of [19] collate the results obtained from running NS2, Matlab, Opnet and Graphical Network Simulator-3 (GNS3) with the results obtained from a real network made up of Cisco routers. At first the authors [19] used a very simple network containing one IP routing device and measured the delay for single ICMP packets across the device, later they repeated the procedure in a more complex network similar to what can be found in a typical IP network. In order to compare the results from the simulations and real network results the authors [19] used Wireshark, and the results of OPNET were different from the real network results in the first scenario. It was not possible to run the second scenario because of the lack of parameters for traffic control. The results of the GNS3 simulation matched the results obtained from the Cisco network, and the authors [19] concluded that the only way of getting accurate simulation results about real networks is to use a mathematical model and implement it in Matlab or to create an application. In [20] the authors use datasets of actual

attack traffic to create simulations in ns-2 simulator.

3. THE SIMULATOR CHOICE

According to the information in related works, there is no universal network simulator which can be used for creating any of the simulations. Each simulator has its advantages and disadvantages. That is why, it is very important to make a list of the research requirements when selecting a tool for simulation. Having studied the most commonly used network simulators we decided to use GNS3 simulator in our research. While using network simulators the researchers should compare the simulation results with the real network results. Comparing them we can see that many of the parameters (like application server settings), which can significantly affect the results, are missing in most of network simulators. This causes difficulties while comparing the simulated results with the real network results. In our previous work [16] we used Riverbed Opnet modeler for simulating a DDoS attack. Even if we were able to set traffic parameters, network links speed and server applications, more important parameters would be missing. That is why we've searched for an alternative. One of them is Graphical Network Simulator-3. The GNS3 is a free network software emulator first released in 2008. GNS3 provides a user friendly graphical interface displayed in Fig. 1, which allows us to create simulated topology without spending too much time.

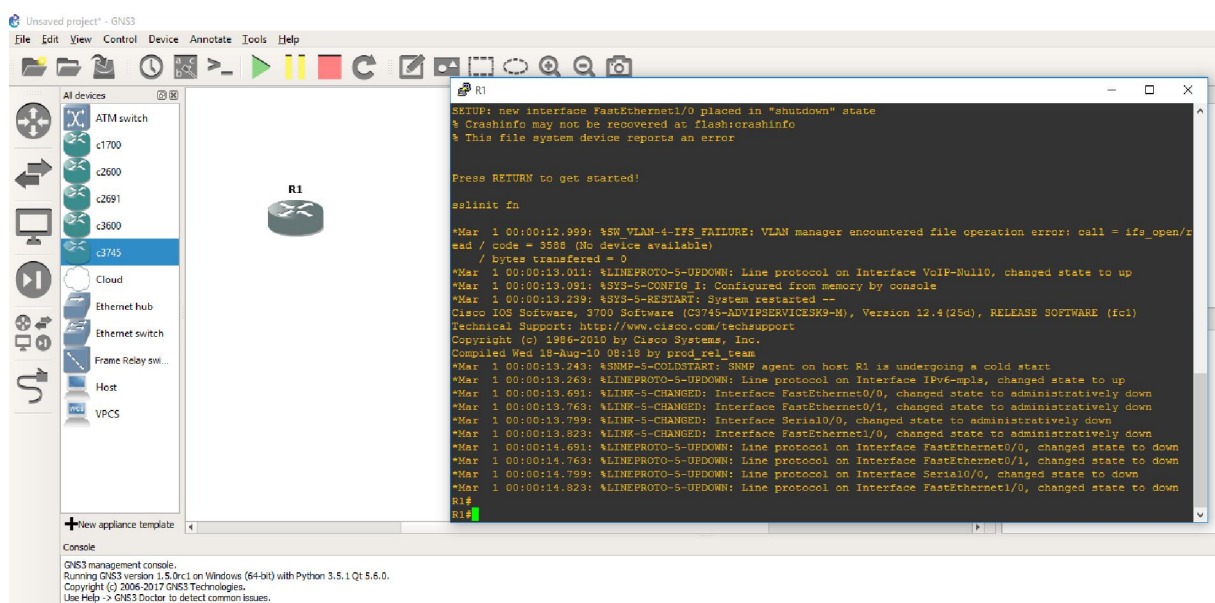


Fig. 1 – GNS3 GUI.

With GNS3 the combination of virtual and real devices can be made and used to simulate complex

networks. It uses Dynamips emulation software to simulate Cisco IOS, it also supports devices from

other network vendors like Juniper and others. If a network device IOS image is introduced into GNS3 then we may select allocated hardware resources, a number of network interfaces and their type. When the simulated device is added into the topology we can access it with a ssh remote control as it is shown in Fig. 1. One of GNS3 important advantages is the possibility to connect the simulated network topology to the real network environment. This can be done using the cloud virtual device from the device list in Fig. 1. We may select there a real or virtual network interface available on PC running GNS3. GNS3 is used by many large companies including Exxon, Walmart, AT&T and NASA, and is also popular while preparing for network professional certification exams.

4. THE SIMULATION

A model of computer network was created including a web server, 3 PCs of regular users and one attacker host. The network is served by Ethernet switches and Cisco routers. Then, we simulated a DOS attack from attacker host to see how it affects the work of web server and its accessibility for regular users. After that we try out some approaches for mitigating this attack. In Fig. 2 you can see what our topology looks like.

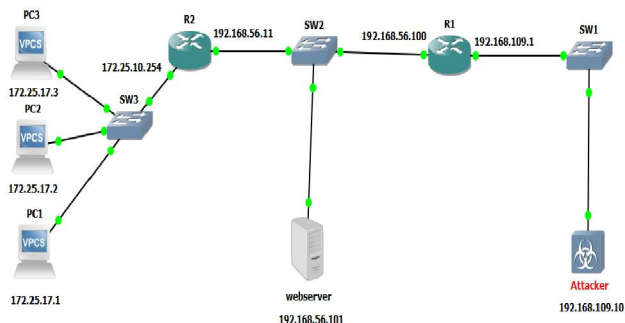


Fig. 2 – Simulated network topology.

In this topology the following devices are used:

- 1) The simulation host – OS: windows 10, CPU: core i5 6600 CPU, 16 GB RAM, HDD: 250 GB Samsung EVO 850;
- 2) Webserver – Fedora core 22 64 bit Linux system running apache 2.4.12 web server and mariadb 10.0.17 database server in default configuration. On the web server we have a default Wordpress 4.7 CMS installed. The web server OS is running in Oracle Virtual Box with 1 CPU core and 2 GB RAM;
- 3) Attacker – Kali Linux 4.6.0 OS running in virtual box with 1 CPU core and 2 GB RAM;
- 4) R1 and R2 routers are Cisco 3745 routers with 256 Mb RAM,
- 5) SW1, SW2, SW3 are GNS3 generic Ethernet

switches;

6) PC1, PC2, PC3 are GNS3 Virtual PC Simulator.

All links in this simulation are set to 100Mbit/s speed. Virtual PC Simulator can be used to simulate end host in the network topology in Gns3 and run simple reachability tests like ping and traceroute. Though there are other alternatives available like Qemu and Virtual box, however, they are CPU intensive. Virtual PC Simulator is integrated with Windows and Linux machine and is very CPU light. GNS3 generic Ethernet switches are virtual devices created by GNS3 that do provide virtual connections between devices with much less resource usage compared to Cisco devices.

A. Scenario 1

Virtual PC Simulator allows making TCP ping by specifying destination port and protocol parameters. In the internal network we have 3 PCs which are in a separate LAN 172.25.10.0/24 and can access webserver through Cisco 3745 router. As it can be seen from Fig. 3, Fig. 4 and Fig. 5 we've launched tcp ping towards 3 on our PC's webserver to simulate regular users accessing web server.

Fig. 3 – PC1 tcp ping towards web server before attack.

Fig. 4 – PC2 tcp ping towards web server before attack.


```

PC3
Connect 80@192.168.56.101 seq=137 ttl=63 time=35.572 ms
SendData 80@192.168.56.101 seq=137 ttl=63 time=19.319 ms
Close 80@192.168.56.101 seq=137 ttl=63 time=21.032 ms
Connect 80@192.168.56.101 seq=138 ttl=63 time=45.570 ms
SendData 80@192.168.56.101 seq=138 ttl=63 time=23.562 ms
Close 80@192.168.56.101 seq=138 ttl=63 time=13.994 ms
Connect 80@192.168.56.101 seq=139 ttl=63 time=37.025 ms
SendData 80@192.168.56.101 seq=139 ttl=63 time=20.324 ms
Close 80@192.168.56.101 seq=139 ttl=63 time=19.572 ms
Connect 80@192.168.56.101 seq=140 ttl=63 time=36.480 ms
SendData 80@192.168.56.101 seq=140 ttl=63 time=20.111 ms
Close 80@192.168.56.101 seq=140 ttl=63 time=20.076 ms
Connect 80@192.168.56.101 seq=141 ttl=63 time=34.676 ms
SendData 80@192.168.56.101 seq=141 ttl=63 time=20.905 ms
Close 80@192.168.56.101 seq=141 ttl=63 time=22.132 ms
Connect 80@192.168.56.101 seq=142 ttl=63 time=42.465 ms
SendData 80@192.168.56.101 seq=142 ttl=63 time=16.383 ms
Close 80@192.168.56.101 seq=142 ttl=63 time=16.601 ms
Connect 80@192.168.56.101 seq=143 ttl=63 time=34.180 ms

```

Fig. 5 – PC3 tcp ping towards web server before attack.

No.	Time	Source	Destination	Protocol	Length	Info
909	3.787072	192.168.109.10	192.168.56.101	TCP	66	44762 → 80 [ACK] Seq=1 Ack=27513 Win=1020 Len=0 TSval=1562007 TSecr=6263057
910	3.790592	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
911	3.790592	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
912	3.790592	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
913	3.790592	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
914	3.790592	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
915	3.790592	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
916	3.791083	192.168.56.101	192.168.109.10	TCP	128	[TCP segment of a reassembled PDU]
917	3.794090	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
918	3.794090	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
919	3.797098	172.25.10.3	192.168.56.101	TCP	66	45463 → 80 [RST, ACK] Seq=4124089901 Ack=1 Win=5840 Len=0 TSval=1489850920 TSecr=0
920	3.797599	192.168.109.10	192.168.56.101	TCP	74	44894 → 80 [SYN] Seq=0 Win=2920 Len=0 MSS=1460 SACK_PERM=1 TSval=1562008 TSecr=0 WS=128
921	3.797599	172.25.10.3	192.168.56.101	TCP	74	[TCP Port numbers reused] 45463 → 80 [SYN] Seq=0 Win=2920 Len=0 MSS=1460 TSval=1489850920 TSecr=0 WS=2
922	3.797599	192.168.56.101	192.168.109.10	TCP	74	80 → 44894 [SYN, ACK] Seq=0 Ack=1 Win=28960 Len=0 MSS=1460 SACK_PERM=1 TSval=6265535 TSecr=1562008 WS=128
923	3.797599	192.168.56.101	172.25.10.3	TCP	74	[TCP Port numbers reused] 80 → 45463 [SYN, ACK] Seq=0 Ack=1 Win=28960 Len=0 MSS=1460 TSval=6265535 TSecr=1489850920 WS=128
924	3.801610	192.168.56.101	192.168.109.10	TCP	1514	[TCP segment of a reassembled PDU]
925	3.807626	192.168.109.10	192.168.56.101	TCP	66	44892 → 80 [ACK] Seq=1 Ack=1 Win=29312 Len=0 TSval=1562009 TSecr=6263172
926	3.818157	172.25.10.3	192.168.56.101	TCP	66	45463 → 80 [ACK] Seq=1 Ack=1 Win=2920 Len=0 TSval=1489850920 TSecr=0
927	3.818157	192.168.109.10	192.168.56.101	TCP	66	44820 → 80 [ACK] Seq=1 Ack=5793 Win=410 Len=0 TSval=1562012 TSecr=6263184
928	3.818157	192.168.56.101	192.168.109.10	TCP	1514	80 → 44820 [ACK] Seq=15929 Ack=1 Win=235 Len=1448 TSval=6265555 TSecr=1562012
929	3.818157	192.168.56.101	192.168.109.10	TCP	1514	80 → 44820 [ACK] Seq=17377 Ack=1 Win=235 Len=1448 TSval=6265555 TSecr=1562012
930	3.828685	192.168.109.10	192.168.56.101	TCP	66	44820 → 80 [ACK] Seq=1 Ack=7241 Win=432 Len=0 TSval=1562015 TSecr=6263184
931	3.829187	192.168.56.101	192.168.109.10	TCP	1514	80 → 44820 [ACK] Seq=18825 Ack=1 Win=235 Len=1448 TSval=6265566 TSecr=1562015
932	3.829187	192.168.56.101	192.168.109.10	TCP	1514	80 → 44820 [ACK] Seq=20273 Ack=1 Win=235 Len=1448 TSval=6265566 TSecr=1562015

Fig. 6 – Webserver switch traffic flow statistics.

As we can see in Fig. 6 a large number of connections is generated by IP address 192.168.109.10 which is the attacker host and only few by 172.25.10.1 and 172.25.10.3 which are regular user hosts. Since the web server is running a default configuration, then after running the script multiple server instances are created and web server quickly goes out of memory and stops responding. Also since the server runs Wordpress CMS, it makes database connections on each page request, after launching the attack the maximum connection limit is overreached. In Fig. 7, Fig. 8, Fig. 9 we can see that ping statistics of legitimate users' hosts simulated by Virtual PC Simulator indicates that server stopped responding to users requests.

```

PC1
SendData 80@192.168.56.101 seq=128 ttl=63 time=18.551 ms
Close 80@192.168.56.101 seq=128 ttl=63 time=16.524 ms
Connect 80@192.168.56.101 seq=129 ttl=63 time=42.617 ms
SendData 80@192.168.56.101 seq=129 ttl=63 time=16.555 ms
Close 80@192.168.56.101 seq=129 ttl=63 time=17.037 ms
Connect 80@192.168.56.101 seq=130 ttl=63 time=42.133 ms
SendData 80@192.168.56.101 seq=130 ttl=63 time=15.512 ms
Close 80@192.168.56.101 seq=130 ttl=63 time=26.562 ms
Connect 80@192.168.56.101 seq=131 ttl=63 time=35.625 ms
SendData 80@192.168.56.101 seq=131 ttl=63 time=20.025 ms
Close 80@192.168.56.101 seq=131 ttl=63 time=18.931 ms
Connect 80@192.168.56.101 seq=132 ttl=63 time=35.909 ms
SendData 80@192.168.56.101 seq=132 ttl=63 time=20.035 ms
Close 80@192.168.56.101 seq=132 ttl=63 time=502.627 ms
Connect 80@192.168.56.101 seq=133 ttl=63 time=12.161 ms
SendData 80@192.168.56.101 seq=133 ttl=63 time=14.222 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=134 ttl=63 time=37.608 ms
SendData 80@192.168.56.101 seq=134 ttl=63 time=19.024 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=135 ttl=63 time=40.420 ms
SendData 80@192.168.56.101 seq=135 ttl=63 time=17.550 ms
Close 80@192.168.56.101 timeout

```

Fig. 7 – PC1 tcp ping towards web server during attack.


```

PC2
Connect 80@192.168.56.101 seq=131 ttl=63 time=53.195 ms
SendData 80@192.168.56.101 seq=131 ttl=63 time=20.493 ms
Close 80@192.168.56.101 seq=131 ttl=63 time=21.590 ms
Connect 80@192.168.56.101 seq=132 ttl=63 time=41.855 ms
SendData 80@192.168.56.101 seq=132 ttl=63 time=16.656 ms
Close 80@192.168.56.101 seq=132 ttl=63 time=15.628 ms
Connect 80@192.168.56.101 seq=133 ttl=63 time=35.177 ms
SendData 80@192.168.56.101 seq=133 ttl=63 time=20.291 ms
Close 80@192.168.56.101 seq=133 ttl=63 time=20.594 ms
Connect 80@192.168.56.101 seq=134 ttl=63 time=41.304 ms
SendData 80@192.168.56.101 seq=134 ttl=63 time=16.541 ms
Close 80@192.168.56.101 seq=134 ttl=63 time=18.054 ms
Connect 80@192.168.56.101 seq=135 ttl=63 time=42.106 ms
SendData 80@192.168.56.101 seq=135 ttl=63 time=16.919 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=136 ttl=63 time=40.634 ms
SendData 80@192.168.56.101 seq=136 ttl=63 time=17.124 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=137 ttl=63 time=44.623 ms
SendData 80@192.168.56.101 seq=137 ttl=63 time=14.537 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=138 ttl=63 time=39.642 ms
SendData 80@192.168.56.101 seq=138 ttl=63 time=19.541 ms

```

Fig. 8 – PC2 tcp ping towards web server during attack.

```

PC3
SendData 80@192.168.56.101 seq=127 ttl=63 time=16.856 ms
Close 80@192.168.56.101 seq=127 ttl=63 time=16.623 ms
Connect 80@192.168.56.101 seq=128 ttl=63 time=44.131 ms
SendData 80@192.168.56.101 seq=128 ttl=63 time=26.780 ms
Close 80@192.168.56.101 seq=128 ttl=63 time=15.276 ms
Connect 80@192.168.56.101 seq=129 ttl=63 time=42.076 ms
SendData 80@192.168.56.101 seq=129 ttl=63 time=15.597 ms
Close 80@192.168.56.101 seq=129 ttl=63 time=16.091 ms
Connect 80@192.168.56.101 seq=130 ttl=63 time=42.921 ms
SendData 80@192.168.56.101 seq=130 ttl=63 time=16.714 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=131 ttl=63 time=41.557 ms
SendData 80@192.168.56.101 seq=131 ttl=63 time=16.531 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=132 ttl=63 time=34.184 ms
SendData 80@192.168.56.101 seq=132 ttl=63 time=22.422 ms
Close 80@192.168.56.101 timeout
Connect 80@192.168.56.101 seq=133 ttl=63 time=42.806 ms
SendData 80@192.168.56.101 seq=133 ttl=63 time=15.373 ms

```

Fig. 9 – PC3 tcp ping towards web server during attack.

```

3982 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3983 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3984 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3987 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3988 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3989 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3910 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3913 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3914 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3924 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3925 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3926 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3927 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3928 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3929 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3935 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3941 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3944 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3948 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3950 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3951 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3955 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
3957 ? S 0:00 /usr/sbin/httpd -DFOREGROUND
1110 tty1 R+ 0:00 grep --color=auto httpd
root@localhost httpd#

```

Fig. 10 – List of apache child processes when server is under attack.

In Fig. 10 we can see the list of apache child processes running on web server, at that time server stopped responding to legitimate users' requests.

B. Scenario 2

We installed the mod_evasive on the web server. It is an evasive maneuvers module for Apache that provides evasive action in the event of an HTTP DoS attack or brute force attack. It is also designed to be a detection and network management tool, and can be easily configured to talk to ipchains, firewalls, routers, and more. The mod_evasive presently reports abuse via email and syslog facilities. The mod_evasive enables to set the threshold for the number of requests for the same page (or URI) per page interval. Once the threshold for that interval has been exceeded, the IP address of the client will be added to the blocking list. Then we launch an attack again. As a result the attacking script, produced the output, is shown in Fig. 11.

```

Kali [Running] - Oracle VM VirtualBox
File Machine View Input Devices Help
Applications Places Terminal Sun 13:06
root@kali: ~/Desktop
File Edit View Search Terminal Help
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
HTTP/1.0 200 OK
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HTTP/1.1 403 Forbidden
HTTP/1.1 403 Forbidden
HTTP/1.1 403 Forbidden
HTTP/1.1 403 Forbidden
HTTP/1.1 403 Forbidden

```

Fig. 11 – Attacking script output.

This means that after a short period of time an attacker IP was blacklisted by the web server. And instead of serving attacker's requests, the server started to respond with 403 Forbidden. This prevents server from making database connections, decreases server load and allows web server to be accessible for legitimate users.

5. CONCLUSION

In this paper we've shown some possibilities which GNS3 simulator can provide for scientists in the area of DoS and DDoS attacks simulation. The proposed simulation describes one of the DoS mitigation methods. However in real networks this method alone won't stand a chance against full scale DoS or DDoS attack. The aim of this simulation was not to present the best DDoS mitigation solution but to demonstrate a variety of parameters which can be simulated using GNS3. As we can see, such parameters as web server settings and defense modules settings can be used in GNS3 simulations. These parameters influence on performance of the server under attack and are unavailable in popular simulators like OPNET NS3 and others. GNS3 provides a very realistic approach to creation of the network simulations allowing setting a full variety of parameters which are available in the real computer networks. However, using of GNS3 compared to other network simulators has also some disadvantages. Because it employs hardware resources to simulate the work of all devices and a scalability is limited inside its topology. Another disadvantage is that GNS3 currently supports a limited amount of simulated hardware. Creation of more advanced simulations with the comparison of their results with real networks results should be a topic for the future studies.

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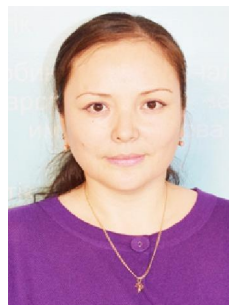
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