## Computer Architecture

Lecture-05

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## Reference

- "Computer Organization and Architecture" by William Stallings; 8<sup>th</sup> Edition (Chapter-03).
  - Any later edition is fine.

#### Interconnection Structure

- The collection of paths connecting the various modules is called the interconnection structure.
- The interconnection structure must support the following types of transfers:
  - Memory to processor: The processor reads an instruction or a unit of data from memory.
  - Processor to memory: The processor writes a unit of data to memory
  - I/O to processor: The processor reads data from an I/O device via an I/O module.
  - Processor to I/O: The processor sends data to the I/O device
  - I/O to or from memory: For these two cases, an I/O module is allowed to exchange data directly with memory, without going through the processor, using direct memory access (DMA).

#### Bus Interconnection

- A bus is a communication pathway connecting two or more devices.
  - It is a shared transmission medium.
  - Consists of multiple communication pathways, or lines, each line is capable of transmitting signals representing binary 1 and binary 0.
- Multiple devices connect to the bus, and a signal transmitted by any one device is available for reception by all other devices attached to the bus.
- If two devices transmit during the same time period, their signals will overlap and become garbled.
  - Only one device at a time can successfully transmit.
- A bus that connects major computer components (processor, memory, I/O) is called a system bus.

#### Bus Structure

- A system bus consists, typically, of from about 50 to hundreds of separate lines.
- Each line is assigned a particular meaning or function.
- Any bus lines can be classified into three functional groups: data, address, and control lines.
- There may be power distribution lines that supply power to the attached modules.

#### Data Bus

- The data lines provide a path for moving data among system modules.
  - These lines, collectively, are called the data bus.
- The data bus may consist of 32, 64, 128, or even more separate lines, the number of lines being referred to as the width of the data bus.
- Because each line can carry only 1 bit at a time, the number of lines determines how many bits can be transferred at a time.
- Width: If the data bus is 32 bits wide and each instruction is 64 bits long, then the processor must access the memory module twice during each instruction cycle.

### Address Bus

- The address lines are used to designate the source or destination of the data on the data bus.
- For example, if the processor wishes to read a word (8, 16, or 32 bits) of data from memory, it puts the address of the desired word on the address lines.
- The width of the address bus determines the maximum possible memory capacity of the system.
- Typically, the higher-order bits are used to select a particular module on the bus, and the lower-order bits select a memory location or I/O port within the module.
- For example, on an 8-bit address bus, address 01111111 and below might reference locations in a memory module (module 0) with 128 words of memory, and address 10000000 and above refer to devices attached to an I/O module (module 1).

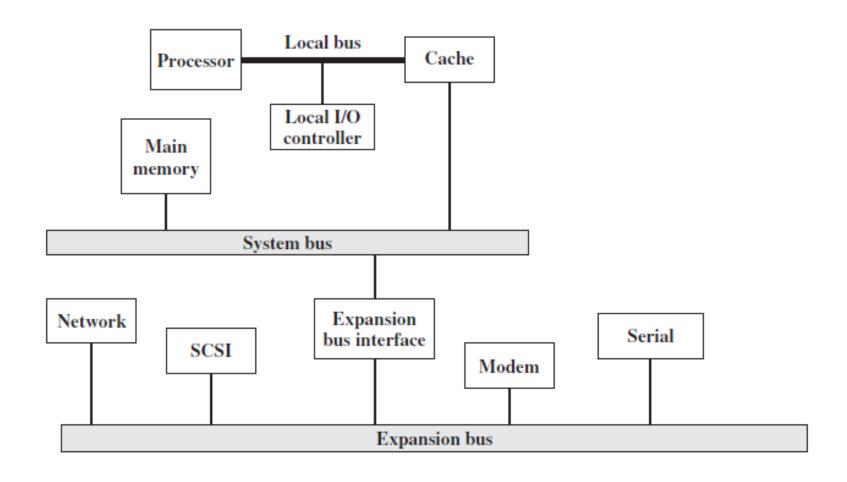
#### Control Bus

- The control lines are used to control the access to and the use of the data and address lines.
- Control signals transmit both command and timing information among system modules.
  - Timing signals indicate the validity of data and address information.
  - Command signals specify operations to be performed.
- Typical control lines include:
  - Memory write, Memory read, I/O write, I/O read, Transfer ACK, Bus request, Bus grant, Interrupt request, Interrupt ACK, Clock, Reset.

## Multiple-Bus Hierarchies

- Motivated by two major concerns:
- In general, the more devices attached to the bus, the greater the bus length and hence the greater the propagation delay.
- The bus may become a bottleneck as the aggregate data transfer demand approaches the capacity of the bus.
- Most computer systems use multiple buses, generally laid out in a hierarchy.

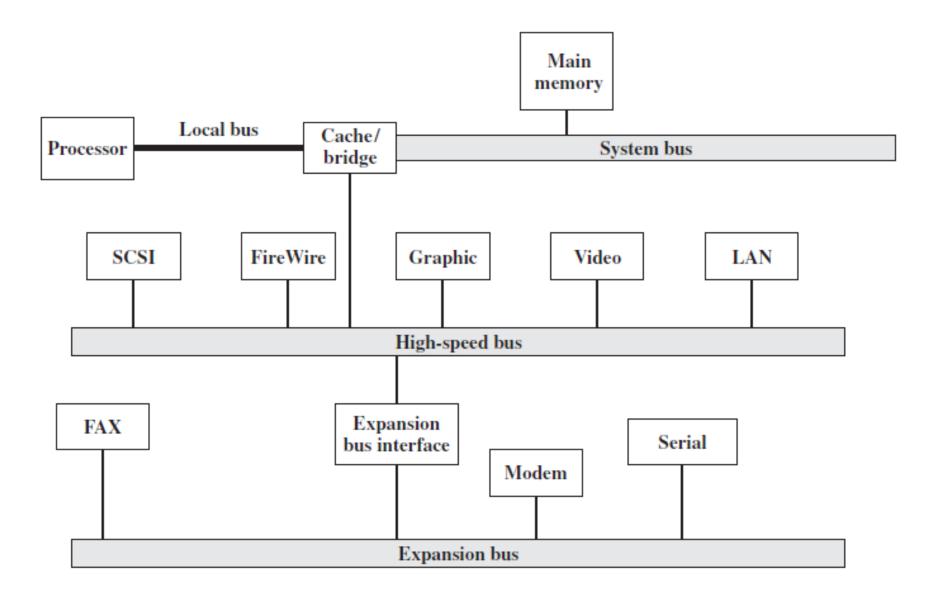
## Traditional BUS Architecture



## Multiple-Bus Hierarchies

- Network connections include local area networks (LANs) such as a 10-Mbps Ethernet and connections to wide area networks (WANs) such as a packet-switching network.
- SCSI (small computer system interface) is itself a type of bus used to support local disk drives and other peripherals.
- A serial port could be used to support a printer or scanner.

## High-Performance Architecture



#### Continue...

- There is a local bus that connects the processor to a cache controller, which is in turn connected to a system bus that supports main memory.
- The cache controller is integrated into a bridge, or buffering device, that connects to the high-speed bus.
- This bus supports connections to high-speed LANs, such as Fast Ethernet at 100 Mbps, video and graphics workstation controllers, as well as interface controllers to local peripheral buses, including SCSI and FireWire.
- The latter is a high-speed bus arrangement specifically designed to support high-capacity I/O devices.
- Lower-speed devices are still supported off an expansion bus, with an interface buffering traffic between the expansion bus and the high-speed bus.

#### Continue...

- The advantage of this arrangement is that the high-speed bus brings high demand devices into closer integration with the processor and at the same time is independent of the processor.
- Thus, differences in processor and high-speed bus speeds and signal line definitions are tolerated.
- Changes in processor architecture do not affect the high-speed bus, and vice versa.

## Elements of Bus Design

Type Bus Width

Dedicated Address

Multiplexed Data

Method of Arbitration Data Transfer Type

Centralized Read

Distributed Write

Timing Read-modify-write

Synchronous Read-after-write

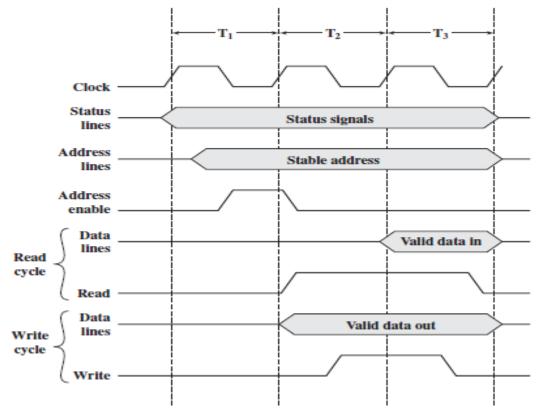
Asynchronous Block

#### Method of Arbitration

- Roughly Two categories:
  - Centralized
  - Distributed
- In a centralized scheme, a single hardware device, referred to as a bus controller or arbiter, is responsible for allocating time on the bus.
- The device may be a separate module or part of the processor.
- In a distributed scheme, there is no central controller.
- Rather, each module contains access control logic and the modules act together to share the bus.
- With both methods of arbitration, the purpose is to designate one device, either the processor or an I/O module, as master.

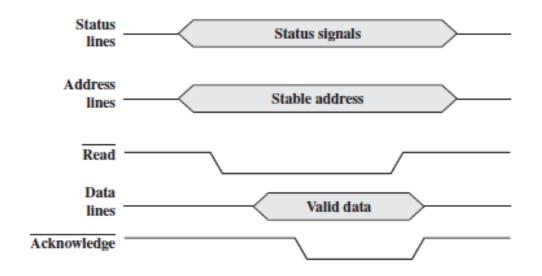
## Timing

- With **synchronous timing**, the occurrence of events on the bus is determined by a clock.
- The bus includes a clock line upon which a clock transmits a regular sequence of alternating 1s and 0s of equal duration.



## Timing

• With asynchronous timing, the occurrence of one event on a bus follows and depends on the occurrence of a previous event.



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