



***INSTITUTE OF INFORMATION TECHNOLOGY***  
***JAHANGIRNAGAR UNIVERSITY***

**Number of Assignment : 01**

**Course Tittle : Data Structure**

**Course Code : ICT – 2101**

**Submission Date : 23/02/2021**

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Roll – 2023

2<sup>nd</sup> year 1<sup>st</sup> Semester

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## **Answer to the Question no – 01**

If we want time complexity for the problem as  $O(n \log n)$ , we can then simply apply merge sort and add the last  $k$  elements of the array.

### **Pseudocode:**

- Declare left and right variables.
- Left will be assigned to 0 and right will be assigned to  $n-1$ .
- Find  $\text{mid} = (\text{left} + \text{right} / 2)$
- Call merge sort on  $(\text{left} . \text{mid})$  and  $(\text{mid} + 1 . \text{rear})$
- Then we will merge on the 2 sub problems.

[Sorting completed]

- Using  $\text{Sum} = 0$ , we will simply add the last ' $k$ ' element.

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## **Answer to the Question no - 02**

We can use an AVL tree.

- ❖ `new()`: create a new AVL tree, which takes  $O(1)$  time
- ❖ `insert(x)`: use AVL insertion, which takes  $O(\log n)$  time
- ❖ `member(x)`: use BST membership testing, which takes  $O(\log n)$  time
- ❖ `increaseBy(x)`: use an inorder traversal to iterate through all nodes of the tree. For each node, add  $x$  to the value. Note that this preserves the relative order of all elements so will not break the AVL invariant. It takes  $O(n)$  time because each node of the tree is visited once.

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**THE END**