



***INSTITUTE OF INFORMATION TECHNOLOGY***  
***JAHANGIRNAGAR UNIVERSITY***

**Number of Assignment : 01**

**Submission Date : 14/07/2021**

**Course Title : Discrete Mathematics**

**Course Code : ICT – 1207**

**Submitted To**

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Class Roll – 2023

Exam Roll – 192340

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Answer to the question no - Q.1

(a) If  $R$  is reflexive,

Then, for every  $x$ ,

Since  $xRx$ ,

we have  $xSx$ .

$\therefore S$  is reflexive.

[Proved]

(b) Given that

$$xSy \Rightarrow xRy \text{ and } yRx$$

$$\Rightarrow xRy \wedge yRx$$

$$\Rightarrow yRx \wedge xRy$$

$$\Rightarrow ySx$$

$\therefore S$  is symmetric

[Proved]

(c) Suppose  $R$  is transitive.

Then,

$$\begin{aligned}xSy \wedge ySz &\Rightarrow (xRy \wedge yRx) \wedge (yRz \wedge zRy) \\&\Rightarrow (xRy \wedge yRz) \wedge (zRy \wedge yRx) \\&\Rightarrow xRz \wedge zRx \\&\Rightarrow xSz\end{aligned}$$

$\therefore$  If  $R$  is transitive,  $S$  is transitive. [Proved]

(d) If  $R$  is antisymmetric

$$\begin{aligned}\text{Then, } xSy \wedge ySx &\Rightarrow xRy \wedge yRx \\&\Rightarrow x = y\end{aligned}$$

In the same way,  $S$  is antisymmetric.

[Proved]

(e) By a, b, and c

we can say  $S$  is an equivalence relation.  
[Proved]

(f) By a, b, and d

we can say  $S$  is an Partial Order [Proved]

Answer to the Question no- Q2

(a) Given that we said

For any  $x$ ,

$xRx$  nor  $xRx$  is false

So,  $xSx$  is also false.

$\therefore S$  is irreflexive [Proved]

(b) Given that

$$\neg Sx \Rightarrow \neg Ry \text{ nor } yRx$$

$$\Rightarrow yRx \text{ nor } \neg Ry$$

$$\Rightarrow ySx$$

$$\therefore \neg Sx \Rightarrow ySx$$

$\therefore S$  is symmetric.

(c) Let  $R$  be the subset relation as before.

Then,  $\{1\} S \{1, 2\}$  and  $\{1, 2\} S \{2\}$  can happen.

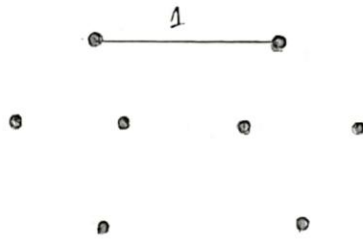
But  $\{1\} S \{2\}$  is false.

$\therefore$  If  $R$  is transitive,  $S$  is not necessarily transitive. [Proved]

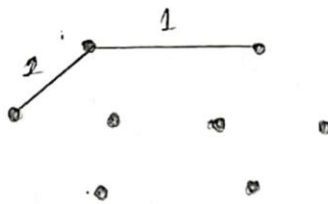
Answer to the question no-Q3

1. using Prim's algorithm

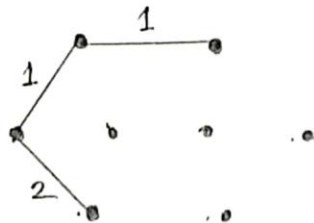
Firstly we choose edge a which has the lowest weight 1.



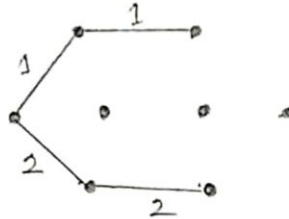
Then, we add the adjacent edge with lowest weight b, weight 1.



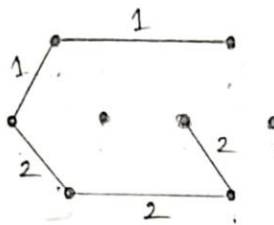
Then we add i,



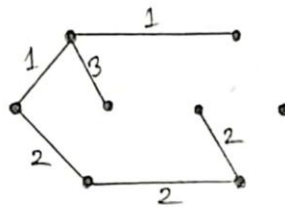
Then add n,



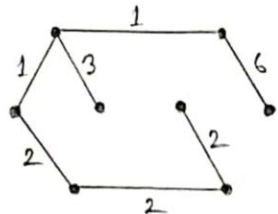
then add k,



Then add c,



Then add e,



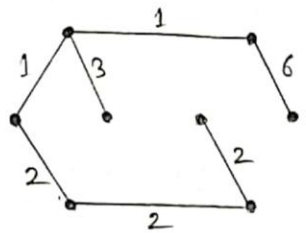
$$\begin{aligned} \text{The cost} &= 1+1+2+2+2+3+6 \\ &= 17 \end{aligned}$$

$$\text{Order} = a \rightarrow b \rightarrow i \rightarrow n \rightarrow k \rightarrow c \rightarrow e.$$

## 2. Using Kruskal's Algorithm

The edge with weights are sorted here.

Edge	Weight
a	1
b	1
i	2
k	2
n	2
c	3
d	3
f	4
j	4
g	5
e	6
h	6
m	7



Order =  $a \rightarrow b \rightarrow i \rightarrow n \rightarrow k \rightarrow c \rightarrow e$ .



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**THE END**