

Problems Project

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Problem I: Statement

Construct an NPDA that accepts the language generated by the productions $S \rightarrow aSa/bSb/c$. Show an instantaneous description of this string $abcba$ for this problem. [WBUT 2007]

[Introduction to Automata Theory, Formal Language and Computation, Shyamalendu Kandar - Page 421]

Problem I: Solution

The production rules are not in GNF. So, we need to first convert it into GNF. The production rules are

$$S \rightarrow aSa \mid bSb \mid c$$

Let us introduce two new productions $C_a \rightarrow a$, $C_b \rightarrow b$. The new production rules become

$$S \rightarrow aSC_a$$

$$S \rightarrow bSC_b$$

$$S \rightarrow c$$

$$C_a \rightarrow a$$

$$C_b \rightarrow b$$

Problem I: Solution, Cont.

Now, all the productions are in GNF. Now, from these productions, a PDA can be easily constructed. First, the start symbol S is pushed into the stack by the following production

$$\delta(q_0, \epsilon, z_0) \rightarrow (q_1, Sz_0)$$

Problem I: Solution, Cont.

For the production $S \rightarrow aSC_a$, the transitional function is

$$\delta(q_1, a, S) \rightarrow (q_1, SC_a)$$

For the production $S \rightarrow bSC_b$, the transitional function is

$$\delta(q_1, b, S) \rightarrow (q_1, SC_b)$$

For the production $S \rightarrow c$, the transitional function is

$$\delta(q_1, c, S) \rightarrow (q_1, \lambda)$$

For the production $C_a \rightarrow a$, the transitional function is

$$\delta(q_1, a, C_a) \rightarrow (q_1, Sz_0)$$

For the production $C_b \rightarrow b$, the transitional function is

$$\delta(q_1, b, C_b) \rightarrow (q_1, \lambda)$$

For acceptance, the transitional function is

$$\delta(q_1, \lambda, z_0) \rightarrow (q_f, z_0) \text{ // accepted by the final state}$$

$$\delta(q_1, \lambda, z_0) \rightarrow (q_1, \lambda) \text{ // accepted by the empty stack}$$

Problem I: Solution, Cont.

ID for the String 'abcba'

$\delta(q_0, \underline{e}abcba, z_0) \rightarrow \delta(q_1, \underline{a}bcba, Sz_0) \rightarrow \delta(q_1, \underline{ab}cba, SC_a z_0) \rightarrow \delta(q_1, \underline{abc}ba, SC_b C_a z_0) \rightarrow \delta(q_1, \underline{abcba}, C_b C_a z_0) \rightarrow \delta(q_1, \underline{abcba}B, z_0) \rightarrow (q_f, z_0)$ (Acceptance by FS).

Problem II: Statement

Construct a PDA, A , equivalent to the following context-free grammar

$$S \rightarrow 0BB, B \rightarrow 0S \mid 1S \mid 0$$

Test whether 0104 is in $N(A)$.

[Introduction to Automata Theory, Formal Language and Computation,
Shyamalendu Kandar - Page 422]

Problem II: Solution

Solution: The CFG is $S \rightarrow 0BB$, $B \rightarrow 0S \mid 1S \mid 0$

All the production rules of the grammar are in GNF. First, the start symbol S is pushed into the stack by the following production

$$\delta(q_0, \epsilon, z_0) \rightarrow (q_1, Sz_0)$$

Problem II: Solution, Cont.

For the production $S \rightarrow 0BB$, the transitional function is

$$\delta(q_1, 0, S) \rightarrow (q_1, BB)$$

For the production $B \rightarrow 0S$, the transitional function is

$$\delta(q_1, 0, B) \rightarrow (q_1, S)$$

For the production $B \rightarrow 1S$, the transitional function is

$$\delta(q_1, 1, B) \rightarrow (q_1, S)$$

For the production $B \rightarrow 0$, the transitional function is

$$\delta(q_1, 0, B) \rightarrow (q_1, \lambda)$$

For acceptance, the transitional functions are

$$\delta(q_1, \lambda, z_0) \rightarrow (q_f, z_0) \text{ // accepted by the final state}$$

$$\delta(q_1, \lambda, z_0) \rightarrow (q_1, \lambda) \text{ // accepted by the empty stack}$$

Problem II: Solution, Cont.

The ID for the String 010000

$(q_0, \underline{\epsilon}10000, z_0) \rightarrow (q_1, \underline{0}10000, Sz_0) \rightarrow (q_1, 0\underline{1}0000, BBz_0) \rightarrow (q_1, 01\underline{0}000, SBz_0) \rightarrow (q_1, 010\underline{0}00, BBBz_0) \rightarrow (q_1, 0100\underline{0}0, BBz_0) \rightarrow (q_1, 01000\underline{0}, Bz_0) \rightarrow (q_1, 010000\underline{\epsilon}, z_0) \rightarrow (q_f, 010000\underline{\epsilon}, z_0)$ (Accepted by the final state).



An Introduction to Automata Theory, Formal Language and Computation,
Shyamalendu Kandar - Pages 421-422

The End