

## 1 Problem 11.2

The corresponding HMM model for this Problem is illustrated in Figure1. The states where the dealer is tossing the biased coin are represented by  $B_1, \dots, B_{10}$ , the fair coin is represented by the states  $F_1, \dots, F_{10}$ . From the start state we go into state  $B_1$  or  $F_1$  with equal probability. After step  $i$ , with  $i \leq 10$  we end in state  $B_i$  if we went into  $B_1$  after the start and we end in state  $F_i$  if we went into  $F_1$  after the start. After the 10th toss we change from  $B_{10}$  to  $F_1$  (if we went into  $F_1$  from the start) or from  $F_{10}$  to  $B_1$  (if we went into  $F_1$  from the start) both with probability  $\frac{1}{10}$  or we stay in the same state  $B_{10}$  or  $F_{10}$  with probability  $\frac{9}{10}$ . If we change into  $B_1$  ( $F_1$ ) we again traverse  $B_2, B_3, \dots, B_{10}$  ( $F_2, F_3, \dots, F_{10}$ ) before we can change into  $B_1$  ( $F_1$ ) again. By always traversing  $B_1, B_2, B_3, \dots, B_{10}$  ( $F_1, F_2, F_3, \dots, F_{10}$ ) before we can go into  $F_1$  ( $B_1$ ) and because we stay in  $B_{10}$  ( $F_{10}$ ) with probability  $\frac{9}{10}$  the case is represented that the dealer keeps each coin for at least 10 tosses.

For the decoding algorithm we now have to deal with 20 states per step (see Figure 2). Let  $B_{i,k}$  and  $F_{i,k}$  be the probability for the most probable sequence of states for emitting  $x_1, \dots, x_k$  and ending at state  $B_i$  or  $F_i$ . For any given state  $i$  with  $i \leq 9$  and step  $k$  we can compute  $B_{i,k}$  as  $B_{i-1,k-1} * e_{B_i}(x_k)$  and likewise for  $F_{i,k}$ , because we can only come from one previous state. When  $k \geq 20$  we can come from two previous states to get to  $B_{10,k}$  and  $F_{10,k}$  and thus have to calculate two values and chose a maximum: to get  $B_{10,k}$  we have to get the maximum of  $F_{1,k-1} * \frac{1}{10} * e_{B_{10}}(x_k)$  and  $B_{10,k-1} * \frac{9}{10} * e_{B_{10}}(x_k)$ . When  $k < 20$  we can also just come from one previous state to get to  $B_{10,k}$  and  $F_{10,k}$ .

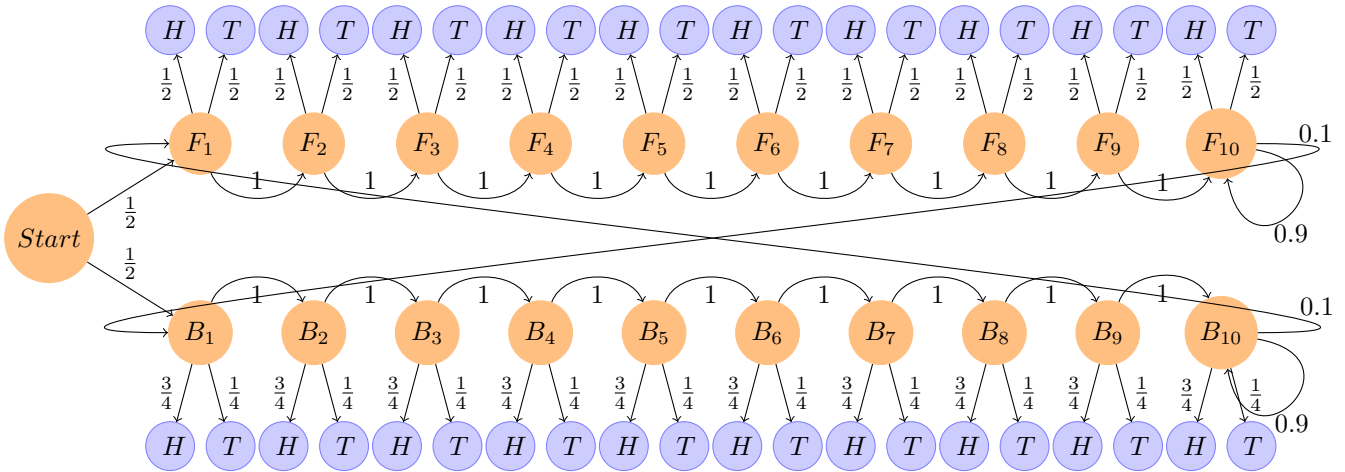


Figure 1: A HMM for the Fair Bet Casino where the dealer keeps every coin for at least ten tosses.

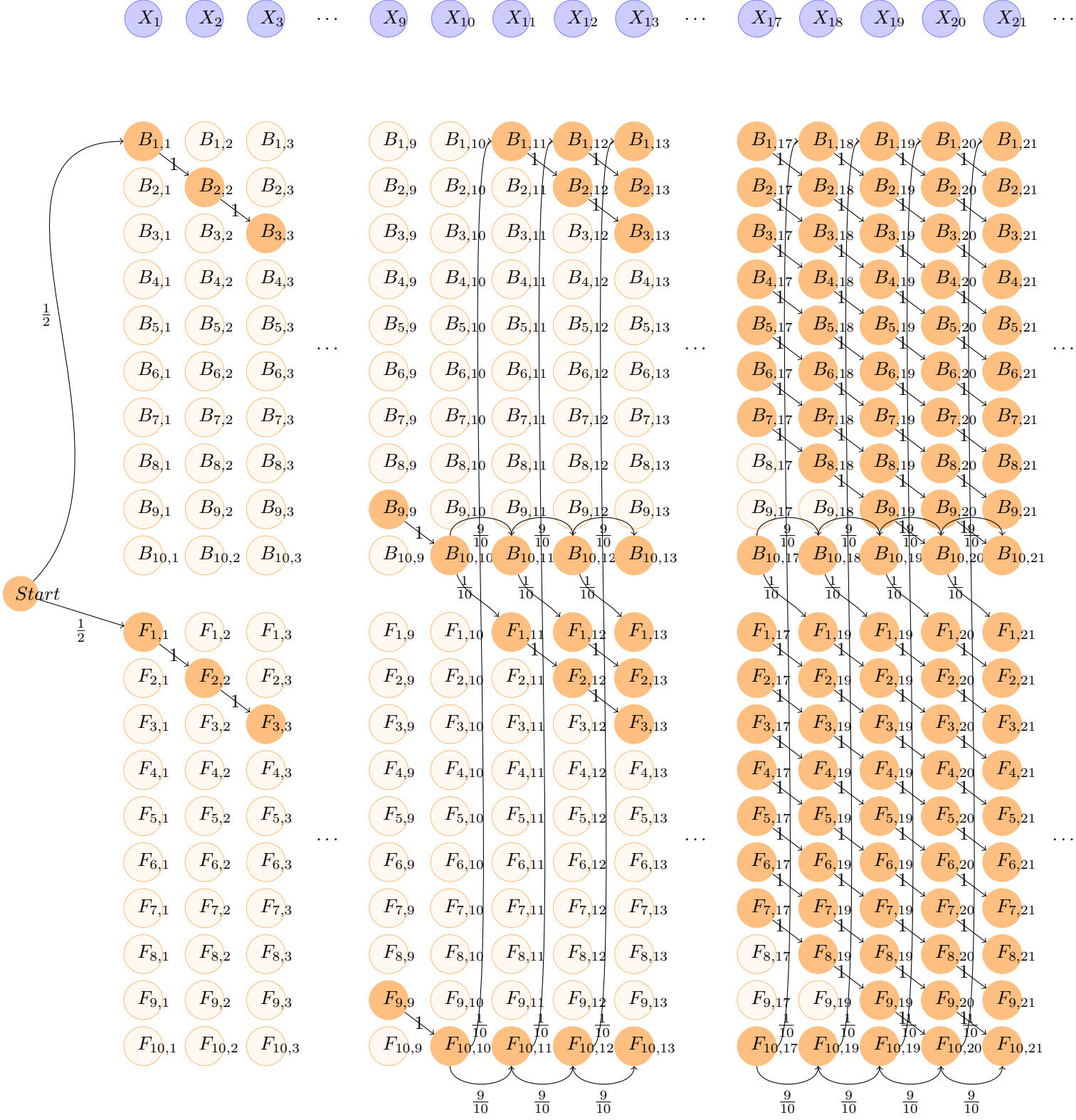


Figure 2: Directed Acyclic Graph for the possible sequence of states that emit the observation  $x_1x_2x_3 \dots x_n$  of the Fair Bet Casino. The light drawn nodes  $B_{i,k}$  and  $F_{i,k}$  indicate that we can not emit  $x_1, \dots, x_k$  and reach state  $i$ .

## 2 Problem 11.4

Figure 3 illustrates running the decoding algorithm for this Problem. As the most probable sequence of states we get  $\beta\beta\beta\beta$ .

## 3 Problem 11.5

We compute the probabilities  $P(x|fair\ coin)$  and  $P(x|biased\ coin)$  using log probabilities:

- $P(x|fair\ coin) = 17 * \log(\frac{1}{2}) = -17$
- $P(x|biased\ coin) = 7 * \log(\frac{1}{4}) + 10 * \log(\frac{3}{4}) = -14 + 10 * (\log(3) + \log(\frac{1}{4})) = -14 + 10 * (1.585 - 2) = -14 + 15.85 - 20 = -18.15$

We get a higher probability for  $P(x|fair\ coin)$  thus the sequence was more likely generated by the fair coin.

## 4 Problem 11.6

A HMM for this problem is given by:  $\mathcal{M} = (\Sigma, Q, A, E)$  with

- $\Sigma = \{1, 2, 3\}$
- $Q = \{Begin, D_1, D_2, End\}$
- $A = \{a_{begin, D_1} = \frac{1}{2}, a_{begin, D_2} = \frac{1}{2}, a_{D_1, D_1} = \frac{1}{2}, a_{D_1, D_2} = \frac{1}{4}, a_{D_1, D_{End}} = \frac{1}{4}, a_{D_2, D_2} = \frac{1}{2}, a_{D_2, D_1} = \frac{1}{4}, a_{D_2, D_{End}} = \frac{1}{4}\}$
- $E = \{e_{D_1}(1) = \frac{1}{2}, e_{D_1}(2) = \frac{1}{4}, e_{D_1}(3) = \frac{1}{4}, e_{D_2}(1) = \frac{1}{4}, e_{D_2}(2) = \frac{1}{2}, e_{D_2}(3) = \frac{1}{4}\}$

Figure 4 shows the graph of this HMM. We observed the sequence 112122end. We find a sequence of states which best explains this observation by running the decoding algorithm according to emission probabilities  $E$  and transition probabilities  $A$ . The computed values  $D_{xy}$  as log probabilities are listed in the following table:

	-	1	1	2	1	2	2	end
Start	0							
$D_1$	-2	$\max \begin{cases} -2-1-1 \\ -3-2-1 \end{cases}$ = -4 coming from $D_{11}$	$\max \begin{cases} -4-1-2 \\ -6-2-2 \end{cases}$ = -7 coming from $D_{12}$	$\max \begin{cases} -7-1-1 \\ -7-2-1 \end{cases}$ = -9 coming from $D_{13}$	$\max \begin{cases} -9-1-2 \\ -10-2-2 \end{cases}$ = -12 coming from $D_{14}$	$\max \begin{cases} -12-1-2 \\ -12-2-2 \end{cases}$ = -15 coming from $D_{15}$		
$D_2$	-3	$\max \begin{cases} -3-1-2 \\ -2-2-2 \end{cases}$ = -6 coming from $D_{11}, D_{21}$	$\max \begin{cases} -6-1-1 \\ -4-2-1 \end{cases}$ = -7 coming from $D_{12}$	$\max \begin{cases} -7-1-2 \\ -7-2-2 \end{cases}$ = -10 coming from $D_{23}$	$\max \begin{cases} -10-1-1 \\ -9-2-1 \end{cases}$ = -12 coming from $D_{24}, D_{14}$	$\max \begin{cases} -12-1-1 \\ -12-2-1 \end{cases}$ = -14 coming from $D_{25}$		
End							$\max \begin{cases} -15-2 \\ -14-2 \end{cases}$ = -16 from $D_{26}$	

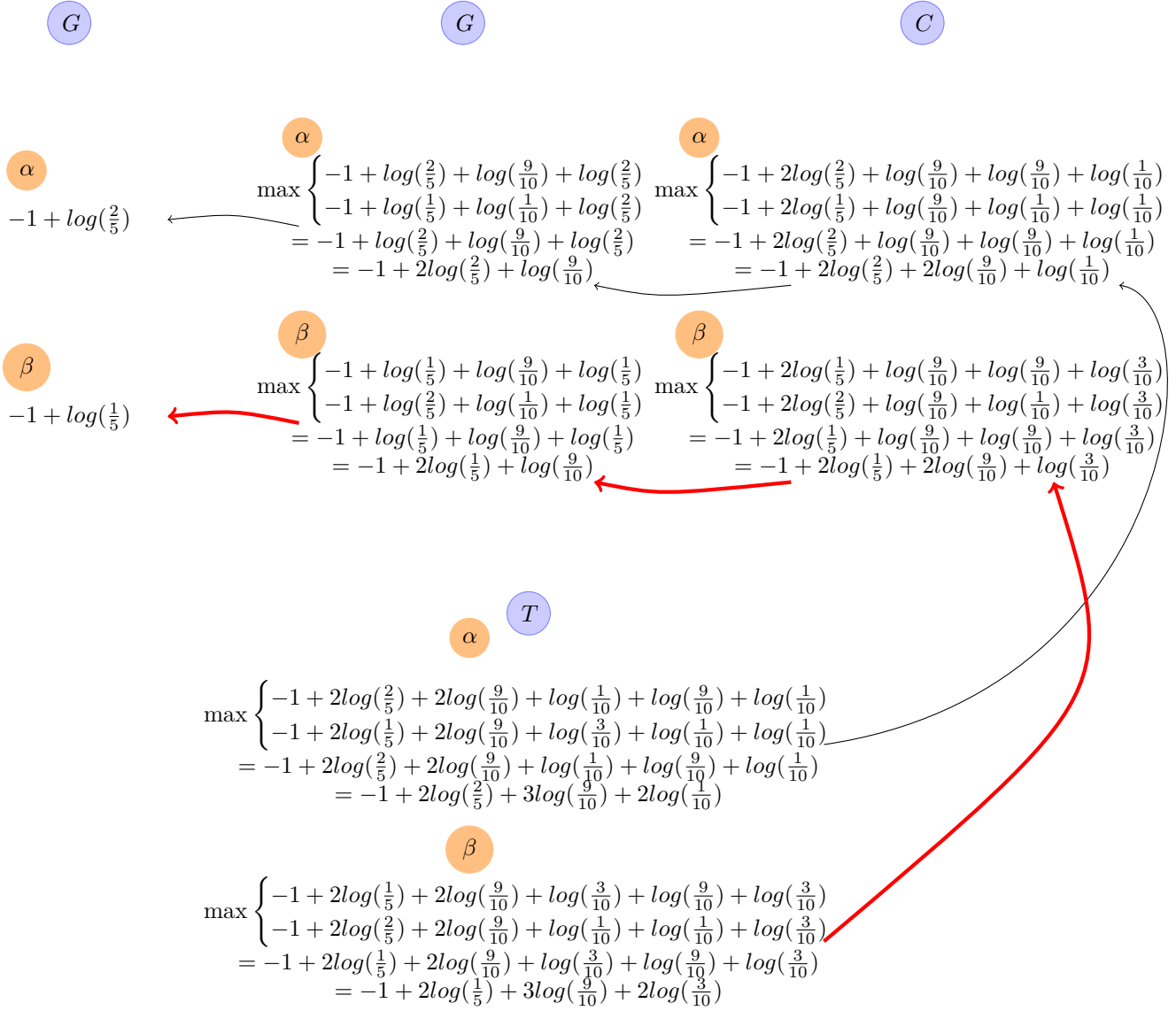


Figure 3: decoding algorithm for the HMM of Problem 11.4. The upper max value always represents the  $\alpha$  state and the lower max value represents the  $\beta$  state. After emitting the final T we get a higher probability in state  $\beta$  and therefore backtrack from there.

As the two most probable sequence of states by backtracking through the table (illustrated in Figure 5) we get:

- $Start-D_1-D_1-D_1-D_2-D_2-End$
- $Start-D_1-D_1-D_2-D_2-D_2-End$

These are the most likely sequences of states that emitted our observation. As a log probability in the final state we get -16 thus the probability of these paths, given the observation 112122end is  $\frac{1}{2^{16}} = \frac{1}{65536}$

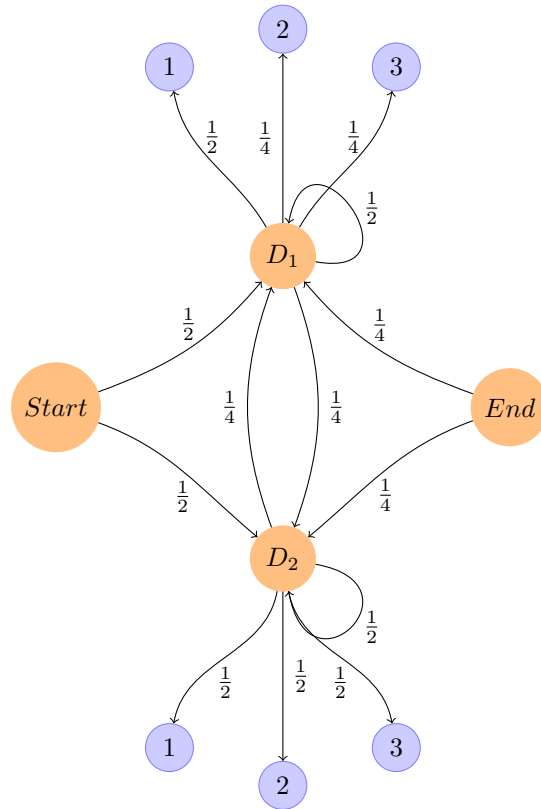


Figure 4: HMM for rolling the three sided dice of Problem 11.6

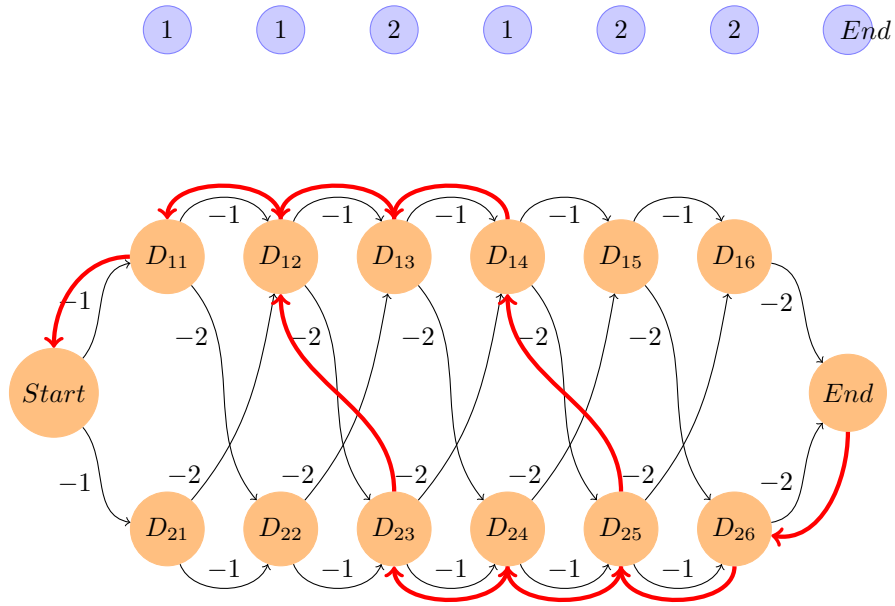


Figure 5: Running the decoding algorithm for the sequence of states 1121221end.  $D_{xy}$  is the probability of the most likely path of  $x_1, \dots, x_y$  that ends at state  $D_x$ . The values between  $D_{xi}$  and  $D_{yi+1}$  describe the log-probability of changing from state  $D_x$  to  $D_y$  from time step  $i$  to  $i + 1$ . The red arrows indicate the backtracking paths.