

giaco.ml

an interpreted EDSL written in OCAML

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1 Introduction

giaco.ml is an imperative (but also functional) EDSL interpreter written in OCAML.

The interpreter is capable of handling expressions, commands and declarations all within a program.

A program can be parsed into valid code from a file/string thanks to reflection.

Static taint analysis can be performed, guards have been also added to the `Reflect` command as a form of protection, similar to perl.

The main execution functions are: `interpret` (or `interpret'` which returns also the declarations' output), `reflect` and `taint_analysis` which all operate on a `Prog(ds,cs)` construct.

2 Design Choices

The interpreter works with 3 domains:

generic this domain is shared amongst the syntax (external) and semantic (internal) domains

syntax this domain contains all the domains accessible to the user. it features 3 sub-domains

expressions these are entities which may be mapped directly to a value

commands these are entities which allow modifying the global state of the program, though allocating new memory is not allowed

declarations these are entities which allow allocation of new memory in the global state

semantic this domain is used for representing the internal state of the interpreter.

evaluable values these are values that may be directly evaluated from an expressions. they represent the base internal type for the interpreter. Basic allowed types are Integer, Boolean, Float, String and Subprograms.

environment this is a mapping representing internal non-mutable state. It maps variables (identifiers) to values.

store this is a mapping representing internal mutable state. It allows aliasing of 2 identifiers for one value. It maps addresses (pointers or locations) to values.

denotable values these are values contained in the environment. Basic allowed types are Integer, Boolean, Float, String and Locations.

mutable memory values these are values contained in the store. Basic allowed types are Integer, Boolean, Float, String.

The interpreter itself is implemented by 3 main functions:

eval this function basically maps expressions to evaluable values

cval this function maps commands to a changed internal state (store)

dval this function maps declarations to a new internal state (environment + store)

Some important choices:

- memory is viewed as a mapping. it is therefore impossible to modify in ocaml. changing the type signature of the interpreter would allow this
- because memory is a mapping, new locations are simulated by randomly selecting a new location based on given store size
- regular expressions are used in reflection, so the Str module must be available
- fixed points are calculated through ocaml's `function` and `let rec` operations
- to ease debugging, the interpreter is not compiled but interpreted in an ocaml shell with the `#use "file.ml"` construct
- since the program is interpreted, `let and` between multiple files is not available, this has been fixed by using mutable state pointers to future functions. These pointers will be updated when the function is created.
- many utility functions have been created to ease debugging. most notable are `emptyenv emptystore tenv' tstore' tnew'` which simplify dealing with memory
- due to the non-recursive nature of `Procedure` and `Call`, they have been omitted from reflection and taint-analysis. The syntax restriction is due to the difficulty of implementing a parser for the complete OCAML list syntax, and the static analysis of commands that cannot be recursively unpacked (one lies in the *expressions*, the other in the *commands*)

3 Usage

the interpreter can be loaded up with

```
>> #use "giaco.ml";;
```

and subsequent testing may be performed:

```
>> (* have a functional sample: *)
>> Plus(Int(3), Int(4)) in eval e emptyenv emptystore;;
- : eval = EInt 7

>> (* now let's try a complete program *)
>> let e = Lambda("x", Multiply(Val("global"), Var("x"))) in
>> let d = New("global", Str("yes")) in
>> let c = Assign("global", Apply(e, Int(5))) in
>> let program = Prog(d,c) in
>> let result_environment, result_store = interpret' program emptyenv emptystore in
>> eval (Val("global")) result_environment result_store;;
- : eval = EStr "yesyesyesyesyes"
```

The most useful interpreter functions are:

`eval: expr -> env -> store -> eval` evaluates expressions

`cval: com -> env -> store -> store` converts commands into a modified mutable memory

`dval: dec -> env -> store -> env*store` allows extending mutable memory

`interpret: prog -> env -> store -> store` combines all of the above. cannot return expression values like `eval`, though, as we do not have print functionality

interpret': `prog -> env -> store -> env*store` just like **interpret** but return the last evaluated environment as well, which allows for further analysis of the output

emptyenv and **emptystore** already initialized empty environment and store

env' and **store'** allow extending environments and stores outside of syntax. **new'** combines this and allows doing something like the **New(...)** command outside of syntax.

ereflect and **creflect** and **dreflect** and **reflect** reflection of expressions, commands, declarations, full blown programs.

etaint and **ctaint** and **dtaint** and **taint_anlaysia** static taint analysis of expressions, commands, declarations and full blown programs.

3.1 EXPRESSIONS

EXAMPLE	DESCRIPTION
<code>Int(3)</code>	basic integer
<code>Str("hello world")</code>	basic ASCII string
<code>Bool(true)</code>	basic boolean
<code>Float(4.5)</code>	basic float
<code>Lambda("x", <exp containing x>)</code>	typical function
<code>RecLambda("f", "x", <exp containing f and x>)</code>	typical recursive function
<code>Rec("f", Lambda(...))</code>	just another way to define recursive lambdas
<code>Proc(["x";"y";"z";...], Block(...))</code>	this is a procedure, check the commands section
<code>IfThenElse(Bool(true), .., ..)</code>	control flow element
<code>Var("x")</code>	this is a way to retrieve an immutable variable's content
<code>LetIn("x", e1, e2)</code>	this is a way to nest functional blocks and scopes
<code>Val("x")</code>	this is a way to retrieve a mutable variable's content
<code>Plus(e1, e2)</code>	plus function, applies to: Int, Str, Float
<code>Multiply(e1, e2)</code>	multiply function, appliest to: Int, Str, Float
<code>Apply(e1, e2)</code>	typical function application, e1 is of type: Lambda, RecLambda, Rec
<code>Equals(e1, e2)</code>	like C's ==
<code>Greater(e1, e2)</code>	like C's >
<code>Not(e)</code>	like C's !
<code>Or(e1, e2)</code>	like C's
<code>And(e1, e2)</code>	like C's &&
<code>Len(Str(...))</code>	gets the length of a St
<code>Sub(Str(...), i, j)</code>	gets a substring. i and j of type Int.
<code>Lower(Str(..))</code>	reduces a string to lowercase, like Python's <code>lower()</code>
<code>Upper(Str(...))</code>	reduces a string to uppercase, like Python's <code>upper()</code>
<code>Trim(Str(...))</code>	trims whitespace from a string, like Python's <code>s.trim()</code>
<code>Replace(<string to be replace>,<replacer string>,<string>)</code>	replaces a string with another string in a string, like Python's <code>s.replace()</code>