Algorithm

Due Date: 9:20 AM, October 25

Autumn, 2012

The following problem sets are all from CLRS.

Homework 3

1. Let f(n) and g(n) be asymptotically nonnegative functions. Using the basic definition of Θ -notation, prove that $\max(f(n), g(n)) = \Theta(f(n) + g(n))$.

- 2. Guess a good asymptotic upper bound on the recurrence T(n) = 3T(n/3) + n, where T(1) = 1. Use the substitution method to verify your answer.
- 3. Answer the following questions according to the definitions of the asymptotic notations, O, Ω , and Θ .
 - (a) Asymptotically upper bound the function $f(n) = n^2 2n + 1$ by the O notation. Justify your answer by demonstrating the constants.
 - (b) Asymptotically lower bound the function $f(n) = n^2 2n + 1$ by the Ω notation. Justify your answer by demonstrating the constants.
 - (c) Show that " $f(n) = \Theta(n)$ " if and only if "f(n) = O(g(n)) and $f(n) = \Omega(g(n))$ ".
 - (d) What's wrong with the following argument?

$$\sum_{1 \le k \le n} k \cdot n = \sum_{1 \le k \le n} O(n) = n \cdot O(n) = O(n^2)$$

(In fact, we have $\sum_{1 \le k \le n} k \cdot n = n \sum_{1 \le k \le n} k = n \cdot \frac{n}{2} (n+1) = O(n^3)$.)