

Homework 6 sample solution

Due 09/13/16

September 8, 2016

1. Analyze the *worst-case* time complexity of the algorithm below. Please show all work. The $\lfloor \cdot \rfloor$ symbols represent the *floor* (“round down”) function. You may assume that this function takes $\Theta(1)$ time for any input. You may also assume it takes a constant amount of time to determine whether an integer is odd.

Note that figuring out what problem this algorithm solves is *irrelevant* to analyzing its complexity.

```
Input:  $n$ : nonnegative integer
1 Algorithm: LoopMystery
2  $sum = 0$ 
3  $t = 1$ 
4  $d = 1$ 
5  $k = n$ 
6 while  $k > 1$  do
7   for  $i = 1$  to  $k$  do
8      $t = t + d$ 
9      $sum = sum + t$ 
10  end
11  if  $k$  is odd then
12     $d = -d$ 
13  end
14   $k = \lfloor k/2 \rfloor$ 
15 end
16 return  $sum$ 
```

Answer: Lines 2–5, 8, 9, 11, 12, 14, and 16 all take $\Theta(1)$ time. Thus, the body of the for loop takes a total of $\Theta(1)$ time. Since the for loop iterates k times, it must take $\Theta(k)$ time total. This complexity dominates the cost of each iteration of the while loop, so each iteration also takes $\Theta(k)$ time. In the worst case, $k = k/2$ (not $(k - 1)/2$) in line 14, so the while loop will iterate $O(\lg n)$ times before $k = 1$. Since the time complexity of an iteration changes as k shrinks, we need to sum up the for loop iterations.

Since k is halved each iteration of the while loop, this sum is:

$$\begin{aligned} O(n) + O(n/2) + O(n/4) + \dots + O(1) &= O(n + n/2 + n/4 + \dots + 1) \\ &= O(n(1 + 1/2 + 1/4 + \dots + 1/n)) \\ &= O(n \sum_{i=1}^{\lg n} 1/2^i) \\ &= O(n(1)) \\ &= O(n) \end{aligned}$$

As the return statement in line 16 takes $\Theta(1)$ time, the total time for LoopMystery is $O(n) + \Theta(1) = O(n)$.