

1. **Input:** *data*: array of integers
Input: *t*: target value
Input: *S*: subset of *data* that sums to *t*
Output: Verifies *S* is a subset sum

```

1  Algorithm: SSD
2  Check  $S \subseteq \text{data}$            // 0(n)
3  Set sum to 0                 // 0(1)
4  for s in S                   // n iter. × 0(1) = 0(n)
5      Add s to sum             // 0(1)
6  end
7  if sum is equal to t // 0(1)
8      return true             // 0(1)
9  else
10     return false           // 0(1)
11 end

```

Since SSD runs in polynomial time $O(n + n)$, $\text{SSD} \in \text{NP}$.

2. **Input:** *data*: array of integers
Input: *n*: size of data
Input: *t*: target value
Output: *S*: subset of data that sums to *t*, or ; if no such set exists

```

1  Algorithm: SSReduction
2  if  $\sim \text{SSD}(\text{data}, t)$  then           // 0(S(n))
3      return  $\emptyset$                      // 0(1)
4  end
5  sub = 0                               // 0(1)
6  top = n                             // 0(1)
7  while sub < top do                   // n iter. × 0(S(n))
8      if  $\sim \text{SSD}(\text{data}[1..(\text{top} - 1)], t)$  then // 0(S(n))
          /* Add data[top] to subset */
9          sub = sub + 1                 // 0(1)
10         Swap data[top] and data[sub] // 0(1)
11     else
          /* Delete data[top] from array */
12         top = top - 1                 // 0(1)
13     end
14 end
15 return data[1..sub]                 // 0(1)

```

The worst-case complexity of SS is $\mathbf{O(S(n) + nS(n))}$ if the worst-case complexity of SSD is $\mathbf{O(S(n))}$.