

## Assignment 2 Posted on Feb 9, due on Feb 23

Maximum total of 35 points.

(20 points)

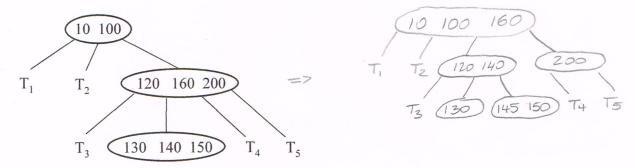
1. B-TREES

For both parts the minimum degree t = 2.

a). Consider the B-Tree below. Illustrate the operation:

B-Tree-Insert (T, 145)

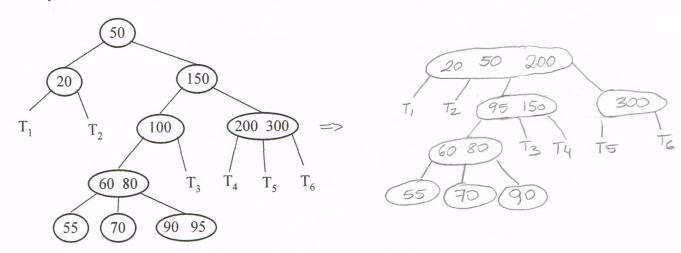
Show your work, the tree obtained after applying each step.



b). Consider the B-Tree below. Illustrate the operation:

B-Tree-Delete (T, 100)

Show your work: the rule used and the tree obtained after each step.



(continued next page)

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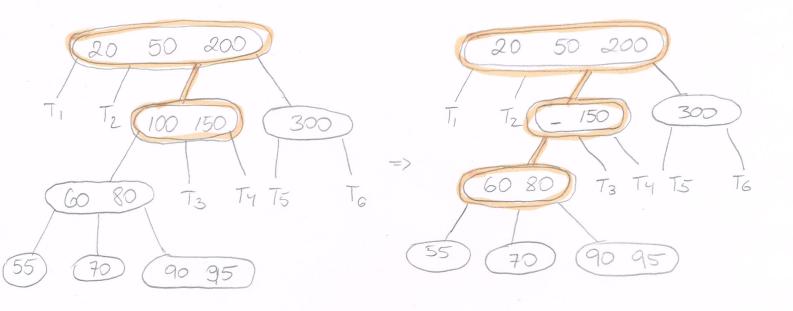
(15 points)

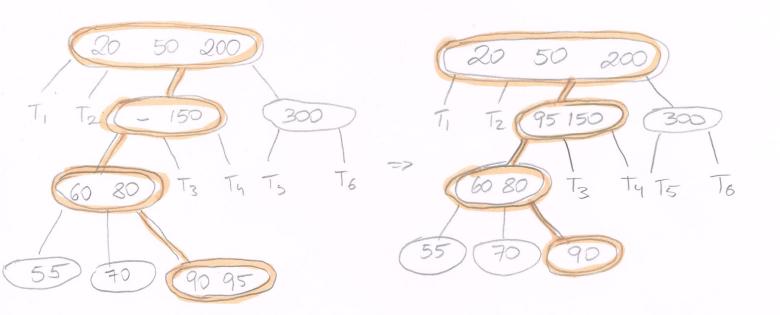
## 2. BACKTRACKING

Solve the m-Independent Set Problem (defined below) using a backtracking algorithm. Write the pseudo-code and analyze its worst-case running time. You have to use the general framework described in class. All other attempts will not be graded.

m-Independent Set Problem: Given a graph G with n nodes, where n > 2, and a value m such that 1 < m < n, find whether G has an independent set of size m. Note that G, m, and n are given as input in this problem.

(70) (90 95)





```
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  backtrack (G, m, n)
    for i from 0 to n
      for k from O to n
       { RESULTARRAY[k] = 0
       RESULTARRAY [i] =1
      for i from O to n
       If G [ C] [ S] = 0
        { ARRAYONE[j] = 1
        else
        { ARRAYONE[j]=0
      if ( rbacktrack ( G, RESULTARRAY, ARRAYONE, Index, M, n) == TRUE)
      RETURN TRUE
    RETURN FALSE
  rbacktrack (G, RESULTARRAY, ARRAYONE, Index, m,n)
 1 nextindex = index
   for j in range index to n
   F IF ARRAYONE[j] ==1
     Inextindex = 5
     lbreak
   if nextindex == n
   1 return FALSE
   for k in range o to n
     if G [nextindex][k] == 0
      ARRAYTWO = 1
      else
      { ARRAYTWO = 0
   IF (BOUND (ARRAYONE, ARRAYTWO, RESULTARRAY, nextindex, n) == m.)
   } RETURN TRUE
   ELSE
    RETURN RBACKTRACK (G, RESULTARRAY, ARRAYONE, NEXTARRAY+1, M, n)
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(2) CONTINUES

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BOUND ( ARRAYONE, ARRAYTWO, RESULTARRAY, INDEX, n)
  RESULT ARRAY [index] = 1
  for i in range from 0 to n
  F (ARRAYONE[i] == 1 AND ARRAYTWO[i] == 1)
    } ARRAY ONE = 1
    } ARRAYONE = 0
  COUNT = D
   for K in range from O to length (resultarray)
   - F If (RESULTARRAY [K] ==1)
      RETURN COUNT
```

```
BACKTRACK -> RT < O(n)
 RBACKTRACK
 1 >> 0 times
  K=2 -> In times
  K=3 -> < n(n-1) times
  k=4 \rightarrow \langle n(n-1)(n-2) \rangle times
k = n \rightarrow (n(n-1)(n-2), 2 = n! times
RT = O(n) \cdot O(n! \left( \frac{1}{n!} + \frac{1}{(n-1)!} + \frac{1}{(n-3)!} + \dots + \frac{1}{1!} \right)
      O(n) \cdot O(n! (e-1))
RT = O(nx n!)
```

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```
def backtrack(GRAPH, m, n):
    print("M parameter : " + str(m))
    for i in range(0, n):
        IARR = [0] * n
        IARR[i] = 1
        ARR1 = [0] * n
        for j in range(0, n):
            if (GRAPH[i][j] == 0):
                ARR1[j] = 1
        if(rbacktrack(GRAPH, IARR, ARR1, i+1, m, n) == True):
            print("Result : true - " + str(IARR))
            for k in range(0, n):
                if(IARR[k] == 1):
            return
    print("Result : false")
def rbacktrack(GRAPH, IARR, ARR1, index, m, n):
    nextindex = n
    for j in range(index, n):
            if(ARR1[j] == 1):
                nextindex = j
                break
    if(nextindex == n):
        return False
    ARR2 = [0] * n
    for j in range(0, n):
                ARR2[j] = 1
        return True
    else:
def bound (ARR1, ARR2, IARR, index, n):
    IARR[index] = 1
    for j in range(0, n):
            ARR1[j] = 1
        else:
            ARR1[j] = 0
    for i in range(0, n):
            if(IARR[i] == 1):
    return x
```

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