Selected Research Projects in Networking

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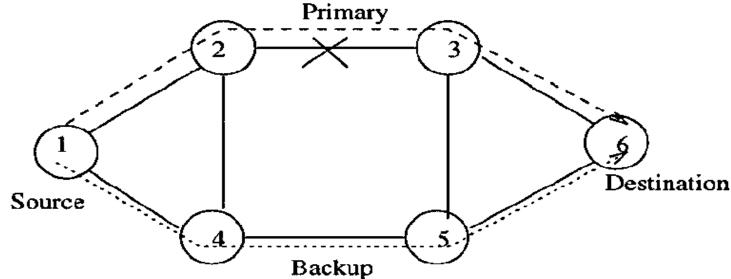
Overview of Research Areas

- Network Survivability
- Wireless Systems
- Internet (TCP/IP)
- Optical Networks



Survivability: Path Protection

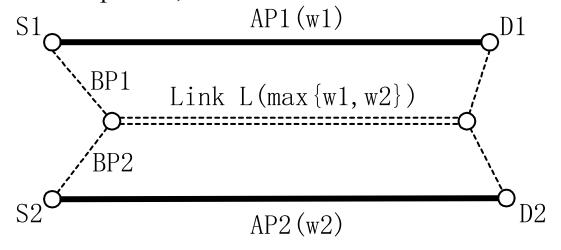
- each connection uses an *active* (or *primary*) *path* (AP) and a (disjoint) *backup path* (BP).
- allocates w units of bandwidth on AP & BP.
- protects against any single link (or node) failure (except for the *src* or *dest* node).





Shared Path Protection: Principle

• if and only if two APs are disjoint, their BPs can share backup bandwidth (BBW) on a common link L (i.e., total BBW on $L = \max\{w_1, w_2\}$ instead of $w_1 + w_2$).





Problem Statement

- Given: a network and dynamic (i.e., *online*) requests for establishment of connections
- Find: a disjoint AP and BP pair for each new request
- Optimization Objective: minimize the total bandwidth consumption (on the AP and BP)
- Constraint: cannot re-arrange existing connections



Challenges

- Using the APF heuristic (find a shortest AP first, and then a shortest (disjoint) BP) only achieves *sub-optimal* results
- So does using a shortest pair of (disjoint) paths (which can be found using a polynomial-time algorithm) due to potential bandwidth sharing among BPs
- Jointly optimize an AP/BP pair is in fact NP-hard
- Using *Integer Liner Programming* (ILP) is notoriously time consuming (even with *branch and bound* techniques)
- ILP only minimizes the *additional* total bandwidth consumption by a new request, but *not the overall* badnwidth consumption by *all the connections*.



Our Solution: APF-based heuristic with Potential Backup Bandwidth

- Similar to the APF heuristic (polynomial time)
- But, in selecting the AP first, each candidate link e will be assigned a cost $w + \beta_e(w)$ where $\beta_e(w)$ is
 - the *potential (yet-to-be-incurred) backup bandwidth*
 - proportional to w and the total bandwidth already allocated on link e to existing APs
 - obtained through rigorous statistical analysis of experimental data
- Runs *faster* than the ILP based approach and yields *a lower overall bandwidth consumption too!*



Other/Ongoing Research

- Trap Avoidance for APF based heuristics
 - May not be able to find a disjoint BP for the given AP even though a disjoint AP/BP pair does exist
- PROMISE: Protection with MultI-SEgments
 - is more bandwidth efficient and recovers from a failure faster than shared path protection
- Protect Against Dual and Multiple Failures
 - dealing with shared risk link groups (SRLGs)
 - recovery policies (which BP to use to recover the first/second failure) and their effects



Heterogeneous Wireless Technology

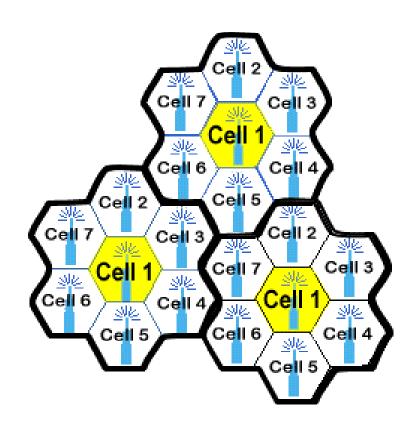
Overlay and Integration of:

- Sensors: ultra wideband /pulse radios
- Personal Area Networks (PAN): Bluetooth, IR...
- Wireless LAN: 802.11 (Ad hoc or Infrastructure)
- Cellular systems (3G and beyond): W-CDMA ...
- Satellite communications: GPS, GEOS and LEOS
- Broadband Wireless: free-space lasers for first/last mile applications



Existing Cellular Systems

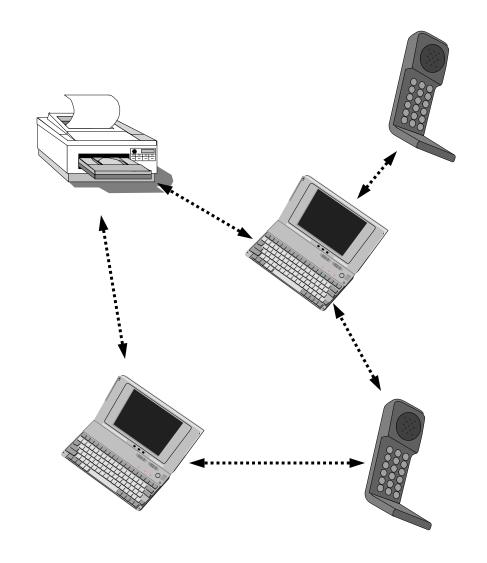
- wide coverage and good scalability
- low data rate and limited capacity (or congestion) due to interference
 - each cell has a subset of channels to facilitate their reuse in cells far away, but can't access available channels in nearby cells.





Ad hoc Networking

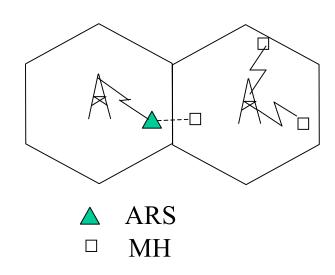
- No base station or access point (infrastructure-less)
- Easy to deploy and high data rate (Mbps and above)
- Limited coverage and scalability





Example: Integrated Cellular and Ad-hoc Relaying (iCAR)

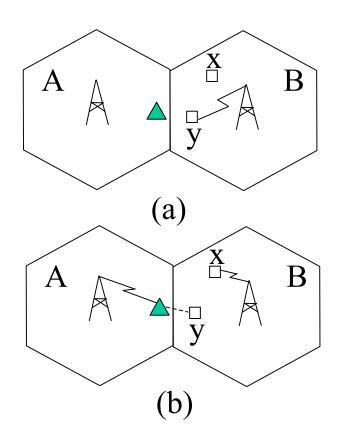
- ARS: Ad hoc Relaying Stations
- MH: Mobile Host (e.g., cell phone)
- Each ARS and MH has two interfaces (cellular and relay)
- *Primary Relaying* for blocked new or hand-off calls in a congested cell (at right) to a non-congested cell (at left)





Secondary Relaying

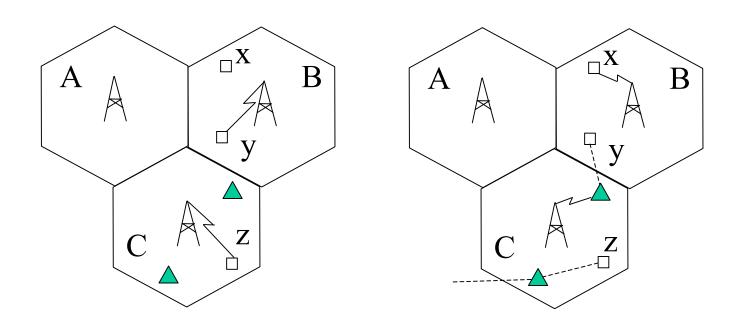
- primary relaying fails:
 - − *X* not covered by ARS
 - Reachable cell (e.g., C)
 is congested too
- Secondary Relaying: free the channel of an active MH Y by relaying its call to a neighboring cell A





Cascaded Relaying

• MH Z releases its channel to MH Y, which in turns release its channel to MH X:



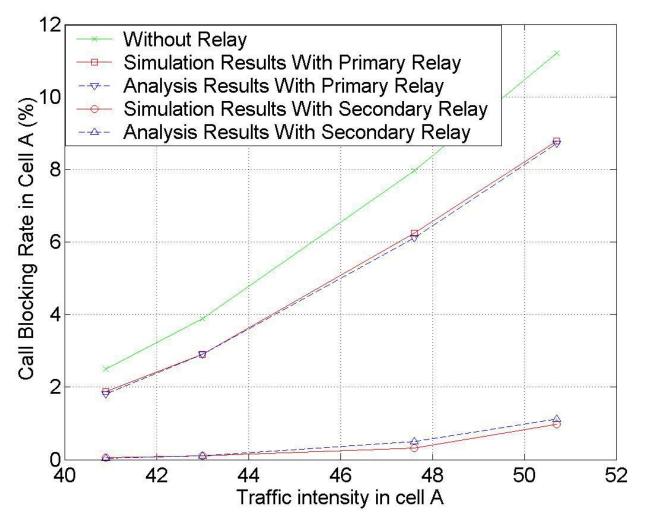


Main Results

- Theorems: iCAR does more than load-balancing
 - call blocking is minimized when the traffic load is balanced among all n cells in a conventional cellular system
 - the blocking probability in an ideal iCAR is even lower
- Formula to determine the maximal number of ARS's needed for full or just edge (boundary) coverage
- Develop strategies to place a limited # of ARS's
- Obtain the call blocking probability through analysis and simulation



Simulation and Analysis Results





Other/Ongoing Research

- Use cellular systems to help Ad hoc networks
 - wide coverage, centralized/global knowledge and authentication/signaling capability
- Integrating other technologies such as broadband wireless, sensors, Bluetooth, 802.11 etc
- Energy efficient routing
- TCP/IP over wireless medium
- Wireless security and survivability



Summary: Opportunities and Challenges

- Research on Optical Networks, Wireless Networks and Internet (including TCP/IP/G-MPLS)
- Network Architecture and Protocol Design
 - scheduling, routing, medium access control, signaling
 - Algorithm, Graph, Information & Communications Theories
- Performance Evaluation
 - Simulation (programming yourself or using existing packages)
 - Analysis: probability /stochastic process (e.g., Markov Chain)
 - Integer Linear Programming formulation (for optimization)
- Graduated 8 PhDs and dozens of MS students (employed by universities and major networking companies)
 - many of them have published papers (journals/conferences) and several have filed patent applications.
- Currently about 10 Ph.D students in LANDER open to a few good PhD students

