## **ECE/MATH 520**, Spring 2008

**Exam 2: Due Session 26** 

Name:	75 mins.; Total 50 pts
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1. (16 pts.) The purpose of this question is to derive a recursive least-squares algorithm where we *remove* (instead of add) a data point. To formulate the algorithm, suppose we are given matrices  $A_0$  and  $A_1$  such that

$$m{A}_0 = egin{bmatrix} m{A}_1 \ m{a}_1^T \end{bmatrix},$$

where  $\boldsymbol{a}_1 \in \mathbb{R}^n$ . Similarly, suppose vectors  $\boldsymbol{b}^{(0)}$  and  $\boldsymbol{b}^{(1)}$  satisfy

$$m{b}^{(0)} = egin{bmatrix} m{b}^{(1)} \\ b_1 \end{bmatrix},$$

where  $b_1 \in \mathbb{R}$ . Let  $\boldsymbol{x}^{(0)}$  be the least-squares solution associated with  $(\boldsymbol{A}_0, \boldsymbol{b}^{(0)})$ , and  $\boldsymbol{x}^{(1)}$  the least-squares solution associated with  $(\boldsymbol{A}_1, \boldsymbol{b}^{(1)})$ . Our goal is to write  $\boldsymbol{x}^{(1)}$  in terms of  $\boldsymbol{x}^{(0)}$  and the "removed" data point  $(\boldsymbol{a}_1, b_1)$ . As usual, let  $\boldsymbol{G}_0$  and  $\boldsymbol{G}_1$  be the Grammians associated with  $\boldsymbol{x}^{(0)}$  and  $\boldsymbol{x}^{(1)}$ , respectively.

- a. Write down expressions for the least-squares solutions  $x^{(0)}$  and  $x^{(1)}$  in terms of  $A_0$ ,  $b^{(0)}$ ,  $A_1$ , and  $b^{(1)}$ .
- b. Derive a formula for  $G_1$  in terms of  $G_0$  and  $a_1$ .
- c. Let  $P_0 = G_0^{-1}$  and  $P_1 = G_1^{-1}$ . Derive a formula for  $P_1$  in terms of  $P_0$  and  $a_1$ . (The formula must not contain any matrix inversions.)
- d. Derive a formula for  $\boldsymbol{A}_0^T \boldsymbol{b}^{(0)}$  in terms of  $\boldsymbol{G}_1$ ,  $\boldsymbol{x}^{(0)}$ , and  $\boldsymbol{a}_1$ .
- e. Finally, derive a formula for  $x^{(1)}$  in terms of  $x^{(0)}$ ,  $P_1$ ,  $a_1$ , and  $b_1$ . Use this and part c to write a recursive algorithm associated with successive removals of rows from  $(A_k, b^{(k)})$ .

**2.** (10 pts.) Use the penalty method to solve the following problem analytically:

minimize 
$$x_1^2 + 2x_2^2$$
  
subject to  $x_1 + x_2 = 3$ .

*Hint:* Use the penalty function  $P(x) = (x_1 + x_2 - 3)^2$ . The solution you find must be exact, not approximate.

3. (12 pts.) Consider a standard form LP problem. Suppose we start with an initial basic feasible solution  $x^{(0)}$  and we apply one iteration of the simplex algorithm to obtain  $x^{(1)}$ .

As pointed out in class, it turns out that we can express  $x^{(1)}$  in terms of  $x^{(0)}$  as

$$\mathbf{x}^{(1)} = \mathbf{x}^{(0)} + \alpha_0 \mathbf{d}^{(0)},$$

where  $\alpha_0$  minimizes  $\phi(\alpha) = f(\boldsymbol{x}^{(0)} + \alpha \boldsymbol{d}^{(0)})$  over all  $\alpha > 0$  such that  $\boldsymbol{x}^{(0)} + \alpha \boldsymbol{d}^{(0)}$  is feasible.

- a. Show that  $\boldsymbol{d}^{(0)} \in \mathcal{N}(\boldsymbol{A})$ .
- b. As usual, assume that the initial basis is the first m columns of  $\boldsymbol{A}$ , and the first iteration involves inserting  $\boldsymbol{a}_q$  into the basis, where q>m. Let the qth column of the canonical augmented matrix be  $\boldsymbol{y}_q=[y_{1q},\ldots,y_{mq}]^T$ .

Express  $d^{(0)}$  in terms of  $y_q$ .

c. Show that  $d^{(0)}$  is a descent direction if and only if  $r_q < 0$ .

**4.** (12 pts.) Suppose we are given a matrix  $A \in \mathbb{R}^{m \times n}$  and a vector  $b \in \mathbb{R}^m$  such that  $b \geq 0$ . We are interested in an algorithm that, given this A and b, is guaranteed to produce one of following two outputs: (1) If there exists x such that  $Ax \geq b$ , then the algorithm produces one such x. (2) If no such x exists, then the algorithm produces an output to declare so.

Describe in detail how to design this algorithm based on the simplex method.