

Convex optimization problems (Chap. 21)

- Consider the general problem

$$\begin{array}{ll} \text{minimize} & f(\mathbf{x}) \\ \text{subject to} & \mathbf{x} \in \Omega. \end{array}$$

- We have seen several types of FONC.
- When is a FONC sufficient for global optimality?
- Answer: In a *convex* optimization problem.

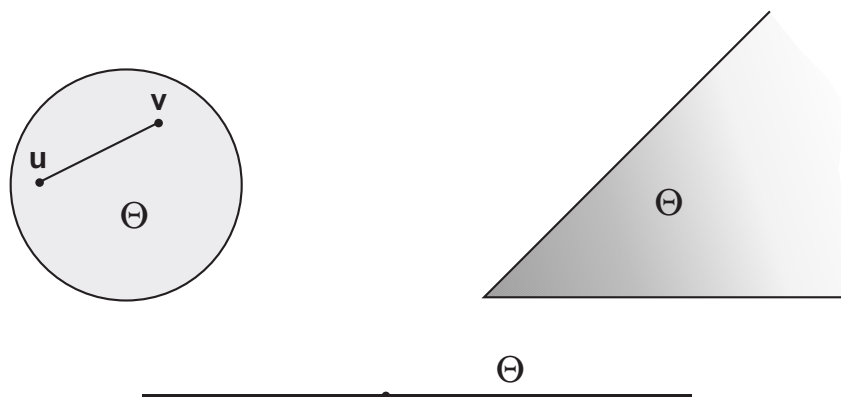
Summary of FONCs

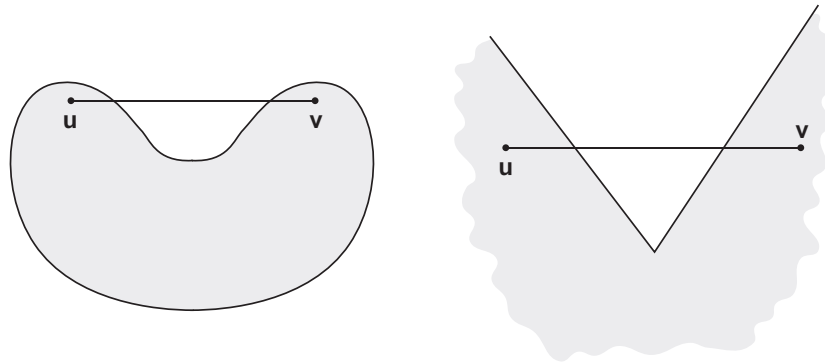
- Set constraint: $\mathbf{d}^T \nabla f(\mathbf{x}^*) \geq 0$ for all feasible directions \mathbf{d} ;
- Interior: $\nabla f(\mathbf{x}^*) = \mathbf{0}$;
- $\Omega = \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}\}$: Lagrange conditions;
- $\Omega = \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\}$: KKT conditions.

Set convexity (§4.3)

- Given: $\Omega \subset \mathbb{R}^n$.
- Definition: Ω is a convex set if, for any distinct $\mathbf{y}, \mathbf{z} \in \Omega$ and $\alpha \in (0, 1)$, we have $\alpha\mathbf{y} + (1 - \alpha)\mathbf{z} \in \Omega$.
- Convex set: the line segment joining any two points in the set lies completely inside the set.

Convex sets:



Nonconvex sets:**Examples of convex sets**

- The empty set
- A set consisting of a single point
- A line or a line segment
- A subspace
- A hyperplane
- A linear variety
- A half-space
- \mathbb{R}^n

Example: Prove that $\Omega = \{x : x \geq 0\}$ is convex.

- Let $y, z \in \Omega$, and $\alpha \in (0, 1)$.
- Want to show that $\alpha y + (1 - \alpha)z \in \Omega$.
- Consider $x = \alpha y + (1 - \alpha)z$. What does $x \in \Omega$ mean?
- To qualify as a member of Ω , each of its component must be ≥ 0 .
- Hence, we must show that each component of x is ≥ 0 .
- Each component of $x = [x_1, \dots, x_n]^T$ satisfies $x_i = \alpha y_i + (1 - \alpha)z_i$.
- Note that we have $y_i, z_i, \alpha, 1 - \alpha \geq 0$.
- Hence, $x_i \geq 0$; i.e., $x \geq 0$, which means that $x \in \Omega$.

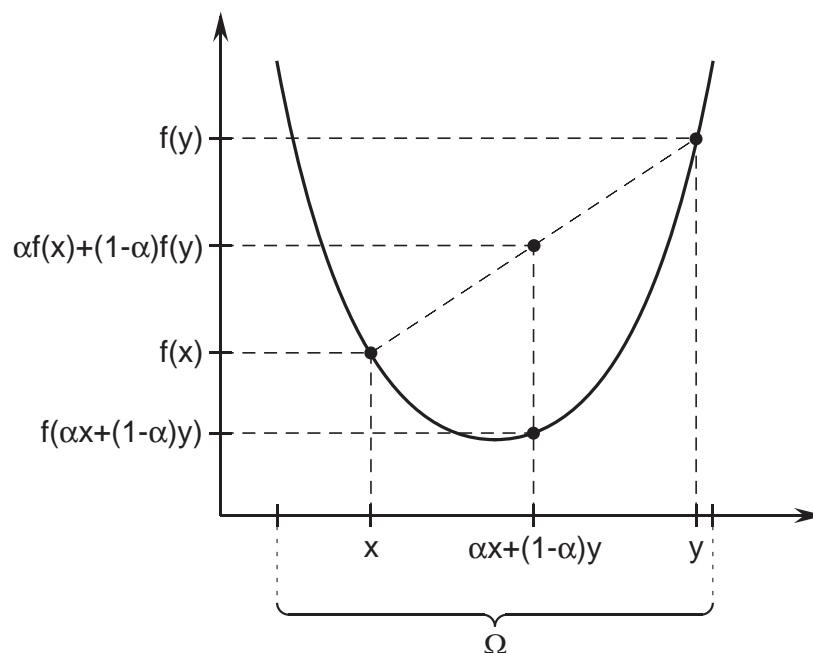
- Therefore, Ω is convex.

Exercise: Prove that $\Omega = \{\mathbf{x} : \mathbf{Ax} = \mathbf{b}\}$ is convex.

Exercise: How do we prove that a set is not convex?

Function convexity

- Given: a function $f : \Omega \rightarrow \mathbb{R}$, where Ω is convex.
- Definition: f is a convex function on Ω if, for any distinct $\mathbf{y}, \mathbf{z} \in \Omega$ and $\alpha \in (0, 1)$,
 $f(\alpha\mathbf{y} + (1 - \alpha)\mathbf{z}) \leq \alpha f(\mathbf{y}) + (1 - \alpha)f(\mathbf{z})$.
- f is strictly convex if \leq is replaced by $<$.
- Convex function: line segment joining two points on the graph lies above the graph.
- f is said to be (strictly) concave if $-f$ is (strictly) convex.



Example: (21.4)

- Consider the function $f(\mathbf{x}) = x_1 x_2$. Is f convex over $\Omega = \{\mathbf{x} : x_1 \geq 0, x_2 \geq 0\}$?
- Answer: No.
- Consider $\mathbf{y} = [2, 1]^T \in \Omega$, $\mathbf{z} = [1, 2]^T \in \Omega$, and $\alpha = 1/2$.
- We have $\mathbf{x} = \alpha\mathbf{y} + (1 - \alpha)\mathbf{z} = [3/2, 3/2]^T$. Hence, $f(\mathbf{x}) = 9/4$.

- On the other hand,

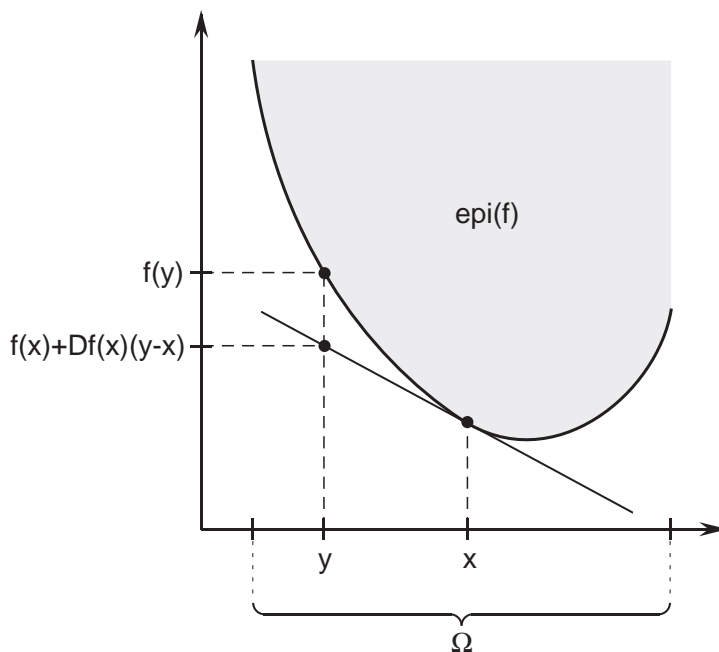
$$\alpha f(\mathbf{y}) + (1 - \alpha)f(\mathbf{z}) = 2 < f(\alpha \mathbf{y} + (1 - \alpha)\mathbf{z}).$$

Alternative way of interpreting function convexity

- Suppose $f : \Omega \rightarrow \mathbb{R}$, Ω convex and open, and $f \in \mathcal{C}^1$.
- Theorem (21.3): f is convex iff for all distinct $\mathbf{x}, \mathbf{y} \in \Omega$,

$$f(\mathbf{y}) \geq f(\mathbf{x}) + Df(\mathbf{x})(\mathbf{y} - \mathbf{x}).$$

- Interpretation: f convex means that it lies above any linear approximation of it.
- For strict convexity, replace \geq by $>$.



- Suppose $f : \Omega \rightarrow \mathbb{R}$, Ω convex and open, and $f \in \mathcal{C}^2$. Let $\mathbf{F}(\mathbf{x})$ be the Hessian of f at \mathbf{x} .
- Theorem (21.4): f is convex if and only if $\mathbf{F}(\mathbf{x}) \geq 0$ for all $\mathbf{x} \in \Omega$.
- For strict convexity, $\mathbf{F}(\mathbf{x}) > 0$ is sufficient, but not necessary (e.g., $f(x) = x^4$ is strictly convex but $f''(0) = 0$).
- If Ω is not open, but $\mathbf{F}(\mathbf{x}) \geq 0$ for all \mathbf{x} in an open set that contains Ω , then we conclude that f is convex.

Examples

- $f(x) = x^3$, $\Omega = (0, 1)$. We have $f''(x) = 6x \geq 0$ on Ω . Hence, f is convex on Ω .
- $f(x) = -x^2$, $\Omega = \mathbb{R}$. We have $f''(x) = -2 < 0$. Hence, f is strictly concave on Ω .
- For a quadratic with Hessian \mathbf{Q} , convexity on \mathbb{R}^n is equivalent to $\mathbf{Q} \geq 0$. Strict convexity is equivalent to $\mathbf{Q} > 0$.

Checking convexity for quadratics

- Proposition (21.1): Consider the quadratic function $f(\mathbf{x}) = \mathbf{x}^T \mathbf{Q} \mathbf{x}$, where $\mathbf{Q} = \mathbf{Q}^T$. Suppose Ω is a convex set. Then, the f is a convex function on Ω iff

$$(\mathbf{x} - \mathbf{y})^T \mathbf{Q} (\mathbf{x} - \mathbf{y}) \geq 0$$

for all $\mathbf{x}, \mathbf{y} \in \Omega$.

- See Example 21.5.

Convex optimization problems

- Consider

$$\begin{array}{ll} \text{minimize} & f(\mathbf{x}) \\ \text{subject to} & \mathbf{x} \in \Omega, \end{array}$$

where Ω is a convex set, and f is a convex function on Ω .

- Name: *Convex programming problem*, or *convex optimization problem*.
- Examples: LP, QP.

Theorem (21.5): In a convex programming problem, a point is a global minimizer if and only if it is a local minimizer.

Proof: The direction \Rightarrow is obvious. Hence, it remains to prove \Leftarrow :

- Suppose $\mathbf{x}^* \in \Omega$ is not a global minimizer. Hence, there is a $\mathbf{y} \in \Omega$, $\mathbf{y} \neq \mathbf{x}^*$, such that $f(\mathbf{y}) < f(\mathbf{x}^*)$.
- Draw a line between $f(\mathbf{y})$ and $f(\mathbf{x}^*)$. Every point on that line is $< f(\mathbf{x}^*)$.
- By convexity of f , the actual graph of f lies *below* the line above.
- By convexity of Ω , all points on the line segment joining \mathbf{y} and \mathbf{x}^* are in Ω .

- Moreover, all points on the line segment above (apart from the endpoint) have objective function value $< f(\mathbf{x}^*)$.
- Hence, \mathbf{x}^* cannot be a local minimizer.

Lemma (21.1): Let $g : \Omega \rightarrow \mathbb{R}$ be a convex function defined on a convex set $\Omega \subset \mathbb{R}^n$. Then, for each $c \in \mathbb{R}$, the set

$$\Gamma_c = \{\mathbf{x} \in \Omega : g(\mathbf{x}) \leq c\}$$

is a convex set.

Proof:

- Let $\mathbf{x}, \mathbf{y} \in \Gamma_c$; i.e., $g(\mathbf{x}), g(\mathbf{y}) \leq c$.
- Since g is convex, for all $\alpha \in (0, 1)$,

$$g(\alpha\mathbf{x} + (1 - \alpha)\mathbf{y}) \leq \alpha g(\mathbf{x}) + (1 - \alpha)g(\mathbf{y}) \leq c$$

- Hence, $\alpha\mathbf{x} + (1 - \alpha)\mathbf{y} \in \Gamma_c$, which implies that Γ_c is convex.

Corollary (21.1): In a convex programming problem, the set of all global minimizers is convex.

Proof:

- If a global minimizer does not exist, the result is trivial.
- If a global minimizer exists, the result follows immediately from the previous lemma by setting

$$c = \min_{\mathbf{x} \in \Omega} f(\mathbf{x}).$$

- We are now ready to prove that the FONC type conditions we have seen before are sufficient for global optimality.

Summary of FONCs

- Set constraint: $\mathbf{d}^T \nabla f(\mathbf{x}^*) \geq 0$ for all feasible directions \mathbf{d} ;
- Interior: $\nabla f(\mathbf{x}^*) = \mathbf{0}$;
- $\Omega = \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}\}$: Lagrange conditions;
- $\Omega = \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}, g(\mathbf{x}) \leq 0\}$: KKT conditions.

Set constraints

- Theorem (21.6): Consider the convex programming problem

$$\begin{aligned} & \text{minimize} && f(\mathbf{x}) \\ & \text{subject to} && \mathbf{x} \in \Omega, \end{aligned}$$

where $f \in \mathcal{C}^1$ on an open convex set that contains Ω . Suppose the point $\mathbf{x}^* \in \Omega$ satisfies

$$\mathbf{d}^T \nabla f(\mathbf{x}^*) \geq 0$$

for any feasible direction \mathbf{d} at \mathbf{x}^* . Then, \mathbf{x}^* is a global minimizer.

- Corollary (21.2): If the point \mathbf{x}^* above satisfies $\nabla f(\mathbf{x}^*) = \mathbf{0}$, then \mathbf{x}^* is a global minimizer.

Proof of Theorem:

- Consider any $\mathbf{x} \in \Omega$. We want to show that $f(\mathbf{x}) \geq f(\mathbf{x}^*)$.

- By convexity and previous theorem,

$$f(\mathbf{x}) \geq f(\mathbf{x}^*) + Df(\mathbf{x}^*)(\mathbf{x} - \mathbf{x}^*).$$

- Note that $\mathbf{d} = \mathbf{x} - \mathbf{x}^*$ is a feasible direction (because Ω is convex). Hence, by assumption,

$$Df(\mathbf{x}^*)(\mathbf{x} - \mathbf{x}^*) = \mathbf{d}^T \nabla f(\mathbf{x}^*) \geq 0.$$

- Combining the above two inequalities, we have

$$f(\mathbf{x}) \geq f(\mathbf{x}^*).$$

Equality constraints

- Let us now consider problems with equality constraint $\mathbf{h}(\mathbf{x}) = \mathbf{0}$.
- Assume that the constraint set $\Omega = \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}\}$ is convex.
- Example: $\mathbf{h}(\mathbf{x}) = \mathbf{b} - \mathbf{A}\mathbf{x}$.
- Further assume that f is a convex function, so that the problem is a convex programming problem.
- Theorem (21.7): Consider the convex programming problem

$$\begin{aligned} & \text{minimize} && f(\mathbf{x}) \\ & \text{subject to} && \mathbf{h}(\mathbf{x}) = \mathbf{0}. \end{aligned}$$

Suppose there exists a feasible point \mathbf{x}^* and a vector $\boldsymbol{\lambda}^*$ such that

$$Df(\mathbf{x}^*) + \boldsymbol{\lambda}^{*T} D\mathbf{h}(\mathbf{x}^*) = \mathbf{0}^T.$$

Then, \mathbf{x}^* is a global minimizer.

Equality and inequality constraints

- Now consider problems with both equality and inequality constraints:

$$\mathbf{h}(\mathbf{x}) = \mathbf{0}, \quad \mathbf{g}(\mathbf{x}) \leq \mathbf{0}.$$

- The constraint set is

$$\begin{aligned} \Omega &= \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}, \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\} \\ &= \{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}\} \cap \{\mathbf{x} : \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\}. \end{aligned}$$

- Note that the intersection of convex sets is convex (exercise: prove).
- Hence, Ω is convex if both the above sets are convex.
- We have already seen an example where the set $\{\mathbf{x} : \mathbf{h}(\mathbf{x}) = \mathbf{0}\}$ is convex.
- When is $\{\mathbf{x} : \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\}$ convex?
- Note that

$$\{\mathbf{x} : \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\} = \bigcap_{i=1}^p \{\mathbf{x} : g_i(\mathbf{x}) \leq 0\}.$$

- Therefore, if each g_i is convex, then by Lemma 21.1 we conclude that each $\{\mathbf{x} : g_i(\mathbf{x}) \leq 0\}$ is convex, and hence $\{\mathbf{x} : \mathbf{g}(\mathbf{x}) \leq \mathbf{0}\}$ is convex.
- Theorem (21.8): Consider the convex programming problem

$$\begin{aligned} &\text{minimize} && f(\mathbf{x}) \\ &\text{subject to} && \mathbf{h}(\mathbf{x}) = \mathbf{0} \\ &&& \mathbf{g}(\mathbf{x}) \leq \mathbf{0}. \end{aligned}$$

Suppose there exists a feasible point \mathbf{x}^* and vectors $\boldsymbol{\lambda}^*$ and $\boldsymbol{\mu}^*$ such that

- $\boldsymbol{\mu}^* \geq \mathbf{0}$;
- $Df(\mathbf{x}^*) + \boldsymbol{\lambda}^{*T} D\mathbf{h}(\mathbf{x}^*) + \boldsymbol{\mu}^{*T} D\mathbf{g}(\mathbf{x}^*) = \mathbf{0}^T$; and
- $\boldsymbol{\mu}^{*T} \mathbf{g}(\mathbf{x}^*) = 0$.

Then, \mathbf{x}^* is a global minimizer.