

EE/M 520, Spring 2007
Exam 1: Session 14

Solutions (version: March 6, 2007, 9:6)

75 mins.; Total 50 pts.

1. (15 pts.) Consider the set-constrained problem

$$\begin{array}{ll} \text{minimize} & f(\mathbf{x}) \\ \text{subject to} & \mathbf{x} \in \Omega, \end{array}$$

where $\Omega = \{[x_1, x_2]^T : x_1^2 + x_2^2 = 1\}$.

- a. Consider a point $\mathbf{x}^* \in \Omega$. Specify all feasible directions at \mathbf{x}^* .
- b. Which points in Ω satisfy the FONC for this set-constrained problem?
- c. Based on part b, is the FONC for this set-constrained problem useful for eliminating local-minimizer candidates?
- d. Suppose we use polar coordinates to parameterize points $\mathbf{x} \in \Omega$ in terms of a single parameter θ :

$$x_1 = \cos \theta \quad x_2 = \sin \theta.$$

Now use the FONC for unconstrained problems (with respect to θ) to derive a necessary condition of this sort: if $\mathbf{x}^* \in \Omega$ is a local minimizer, then $\mathbf{d}^T \nabla f(\mathbf{x}^*) = 0$ for all \mathbf{d} satisfying a “certain condition.” Specify what this “certain condition” is.

Ans.: a. There are no feasible directions at any \mathbf{x}^* .

b. All points in Ω satisfy the FONC for this set-constrained problem.

c. No, the FONC for this set-constrained problem is not useful for eliminating local-minimizer candidates.

d. Write $h(\theta) = f(g(\theta))$, where $g : \mathbb{R} \rightarrow \mathbb{R}^2$ is given by the equations relating θ to $\mathbf{x} = [x_1, x_2]^T$. Note that $Dg(\theta) = [-\sin \theta, \cos \theta]^T$. Hence, by the chain rule,

$$h'(\theta) = Df(g(\theta))Dg(\theta) = Dg(\theta)^T \nabla f(g(\theta)).$$

Notice that $Dg(\theta)$ is tangent to Ω at $\mathbf{x} = g(\theta)$. Alternatively, we could say that $Dg(\theta)$ is orthogonal to $\mathbf{x} = g(\theta)$.

Suppose $\mathbf{x}^* \in \Omega$ is a local minimizer. Write $\mathbf{x}^* = g(\theta^*)$. Then θ^* is an unconstrained minimizer of h . By the FONC for unconstrained problems, $h'(\theta^*) = 0$, which implies that $\mathbf{d}^T \nabla f(\mathbf{x}^*) = 0$ for all \mathbf{d} tangent to Ω at \mathbf{x}^* (or, alternatively, for all \mathbf{d} orthogonal to \mathbf{x}^*).

2. (12 pts.) Suppose we wish to solve the equation $\mathbf{h}(\mathbf{x}) = \mathbf{0}$, where

$$\mathbf{h}(\mathbf{x}) = \begin{bmatrix} 4 + 3x_1 + 2x_2 \\ 1 + 2x_1 + 3x_2 \end{bmatrix}.$$

Consider using an algorithm of the form $\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} - \alpha \mathbf{h}(\mathbf{x}^{(k)})$, where α is scalar constant that does not depend on k .

- Find the solution of $\mathbf{h}(\mathbf{x}) = \mathbf{0}$.
- Find the largest range of values of α such that the algorithm is globally convergent to the solution of $\mathbf{h}(\mathbf{x}) = \mathbf{0}$.
- Assuming that α is outside the range of values in part b, give an example of an initial condition $\mathbf{x}^{(0)}$ of the form $[x_1, 0]^T$ such that the algorithm is guaranteed not satisfy the descent property.

Ans.: a. We can write $\mathbf{h}(\mathbf{x}) = \mathbf{Q}\mathbf{x} - \mathbf{b}$, where $\mathbf{b} = [-4, -1]^T$ and

$$\mathbf{Q} = \begin{bmatrix} 3 & 2 \\ 2 & 3 \end{bmatrix}$$

is positive definite. Hence, the solution is

$$\mathbf{Q}^{-1}\mathbf{b} = \frac{1}{5} \begin{bmatrix} 3 & -2 \\ -2 & 3 \end{bmatrix} \begin{bmatrix} -4 \\ -1 \end{bmatrix} = \begin{bmatrix} -2 \\ 1 \end{bmatrix}.$$

- By part a, the algorithm is a fixed-step-size gradient algorithm for a problem with gradient \mathbf{h} . The eigenvalues of \mathbf{Q} are 1 and 5. Hence, the largest range of values of α such that the algorithm is globally convergent to the solution is $0 < \alpha < 2/5$.
- The eigenvectors of \mathbf{Q} corresponding to eigenvalue 5 has the form $c[1, 1]^T$, where $c \in \mathbb{R}$. Hence, to violate the descent property, we pick

$$\mathbf{x}^{(0)} = \mathbf{Q}^{-1}\mathbf{b} + c \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} -3 \\ 0 \end{bmatrix}$$

where we choose $c = -1$ so that $\mathbf{x}^{(0)}$ has the specified form.

3. (7 pts.) Consider a scalar sequence $\{x_k\}$ that converges with order of convergence p , and satisfies

$$\lim_{k \rightarrow \infty} \frac{|x_{k+1} - 2|}{|x_k - 2|^3} = 0.$$

What is the limit of $\{x_k\}$? What can you say about p ?

Ans.: The limit of $\{x_k\}$ must be 2, because it is clear from the given equation that $|x_{k+1} - 2| \rightarrow 0$. Also, we see that $|x_{k+1} - 2| = o(|x_k - 2|^3)$. Hence, we conclude that $p > 3$.

4. (16 pts.) Let $f : \mathbb{R}^n \rightarrow \mathbb{R}$ be such that $f \in \mathcal{C}^1$. Consider an optimization algorithm applied to this f , of the usual form $\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)}$, where $\alpha_k \geq 0$ is chosen according to line search. Suppose that $\mathbf{d}^{(k)} = -\mathbf{H}_k \mathbf{g}^{(k)}$, where $\mathbf{g}^{(k)} = \nabla f(\mathbf{x}^{(k)})$ and \mathbf{H}_k is symmetric.

- Show that if \mathbf{H}_k satisfies the following conditions whenever the algorithm is applied to a quadratic, then the algorithm is quasi-Newton:

- i. $\mathbf{H}_{k+1} = \mathbf{H}_k + \mathbf{U}_k$
 - ii. $\mathbf{U}_k \Delta \mathbf{g}^{(k)} = \Delta \mathbf{x}^{(k)} - \mathbf{H}_k \Delta \mathbf{g}^{(k)}$
 - iii. $\mathbf{U}_k = \mathbf{a}^{(k)} \Delta \mathbf{x}^{(k)T} + \mathbf{b}^{(k)} \Delta \mathbf{g}^{(k)T} \mathbf{H}_k$, where $\mathbf{a}^{(k)}$ and $\mathbf{b}^{(k)}$ are in \mathbb{R}^n
- b. Which (if any) among the rank-one, DFP, and BFGS algorithms satisfy the three conditions in part a (whenever the algorithm is applied to a quadratic)? (For those that do, you must specify the vectors $\mathbf{a}^{(k)}$ and $\mathbf{b}^{(k)}$.)

Ans.: a. Suppose the three conditions hold whenever applied to a quadratic. We need to show that when applied to a quadratic, for $k = 0, \dots, n-1$ and $i = 0, \dots, k$, $\mathbf{H}_{k+1} \Delta \mathbf{g}^{(i)} = \Delta \mathbf{x}^{(i)}$. For $i = k$, we have

$$\begin{aligned} \mathbf{H}_{k+1} \Delta \mathbf{g}^{(k)} &= \mathbf{H}_k \Delta \mathbf{g}^{(k)} + \mathbf{U}_k \Delta \mathbf{g}^{(k)} \quad \text{by condition i} \\ &= \mathbf{H}_k \Delta \mathbf{g}^{(k)} + \Delta \mathbf{x}^{(k)} - \mathbf{H}_k \Delta \mathbf{g}^{(k)} \quad \text{by condition ii} \\ &= \Delta \mathbf{x}^{(k)}, \end{aligned}$$

as required. For the rest of the proof ($i = 0, \dots, k-1$), we use induction on k .

For $k = 0$, there is nothing to prove (covered by the $i = k$ case). So suppose the result holds for $k-1$. To show the result for k , first fix $i \in \{0, \dots, k-1\}$. We have

$$\begin{aligned} \mathbf{H}_{k+1} \Delta \mathbf{g}^{(i)} &= \mathbf{H}_k \Delta \mathbf{g}^{(i)} + \mathbf{U}_k \Delta \mathbf{g}^{(i)} \\ &= \Delta \mathbf{x}^{(i)} + \mathbf{U}_k \Delta \mathbf{g}^{(i)} \quad \text{by the induction hypothesis} \\ &= \Delta \mathbf{x}^{(i)} + \mathbf{a}^{(k)} \Delta \mathbf{x}^{(k)T} \Delta \mathbf{g}^{(i)} + \mathbf{b}^{(k)} \Delta \mathbf{g}^{(k)T} \mathbf{H}_k \Delta \mathbf{g}^{(i)} \quad \text{by condition iii.} \end{aligned}$$

So it suffices to show that the second and third terms are both 0. For the second term,

$$\begin{aligned} \Delta \mathbf{x}^{(k)T} \Delta \mathbf{g}^{(i)} &= \Delta \mathbf{x}^{(k)T} \mathbf{Q} \Delta \mathbf{x}^{(i)} \\ &= \alpha_k \alpha_i \mathbf{d}^{(k)T} \mathbf{Q} \mathbf{d}^{(i)} \\ &= 0 \end{aligned}$$

because of the induction hypothesis, which implies \mathbf{Q} -conjugacy (where \mathbf{Q} is the Hessian of the given quadratic). Similarly, for the third term,

$$\begin{aligned} \Delta \mathbf{g}^{(k)T} \mathbf{H}_k \Delta \mathbf{g}^{(i)} &= \Delta \mathbf{g}^{(k)T} \Delta \mathbf{x}^{(i)} \quad \text{by the induction hypothesis} \\ &= \Delta \mathbf{x}^{(k)T} \mathbf{Q} \Delta \mathbf{x}^{(i)} \\ &= \alpha_k \alpha_i \mathbf{d}^{(k)T} \mathbf{Q} \mathbf{d}^{(i)} \\ &= 0, \end{aligned}$$

again because of the induction hypothesis, which implies \mathbf{Q} -conjugacy. This completes the proof.

b. All three algorithms satisfy the conditions in part a. Condition i holds, as described in class. Condition ii is straightforward to check for all three algorithms. For the rank-one and DFP algorithms, this is shown in the book. For BFGS, some simple matrix algebra establishes that it holds. Condition iii holds by appropriate definition of the vectors $\mathbf{a}^{(k)}$ and $\mathbf{b}^{(k)}$. In particular, for the rank-one algorithm,

$$\mathbf{a}^{(k)} = \frac{(\Delta \mathbf{x}^{(k)} - \mathbf{H}_k \Delta \mathbf{g}^{(k)})}{(\Delta \mathbf{x}^{(k)} - \mathbf{H}_k \Delta \mathbf{g}^{(k)})^T \Delta \mathbf{g}^{(k)}}, \quad \mathbf{b}^{(k)} = -\frac{(\Delta \mathbf{x}^{(k)} - \mathbf{H}_k \Delta \mathbf{g}^{(k)})}{(\Delta \mathbf{x}^{(k)} - \mathbf{H}_k \Delta \mathbf{g}^{(k)})^T \Delta \mathbf{g}^{(k)}}.$$

For the DFP algorithm,

$$\mathbf{a}^{(k)} = \frac{\Delta \mathbf{x}^{(k)}}{\Delta \mathbf{x}^{(k)T} \Delta \mathbf{g}^{(k)}}, \quad \mathbf{b}^{(k)} = -\frac{\mathbf{H}_k \Delta \mathbf{g}^{(k)}}{\Delta \mathbf{g}^{(k)T} \mathbf{H}_k \Delta \mathbf{g}^{(k)}}.$$

Finally, for the BFGS algorithm,

$$\mathbf{a}^{(k)} = \left(1 + \frac{\Delta \mathbf{g}^{(k)T} \mathbf{H}_k \Delta \mathbf{g}^{(k)}}{\Delta \mathbf{g}^{(k)T} \Delta \mathbf{x}^{(k)}} \right) \frac{\Delta \mathbf{x}^{(k)}}{\Delta \mathbf{x}^{(k)T} \Delta \mathbf{g}^{(k)}} - \frac{\mathbf{H}_k \Delta \mathbf{g}^{(k)}}{\Delta \mathbf{g}^{(k)T} \Delta \mathbf{x}^{(k)}}, \quad \mathbf{b}^{(k)} = \frac{\Delta \mathbf{x}^{(k)}}{\Delta \mathbf{g}^{(k)T} \Delta \mathbf{x}^{(k)}}.$$