BUMPY

SER502-Spring2019-Team11

Team 11 – bumpy

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Feature Of The Language

- Bumpy Language is developed in Prolog Completely
- We are using Top- down parsing technique
- Data structures used by the parser and interpreter: List
- Interpreter: Our interpreter is based on Reduction machine.
- It is an Imperative language and Statically typed
- It provides data types such as Integer and Boolean
- Looping construct while loop
- Decision control statement if-else statement
- inspired by C language

Operators and Constructs For BUMPY

```
• Operators: +,-,*,/,%,<,>,<=,>=, ~=, :=:, = ,and, or,

    Arithmetic Operator: +,-,*,/,%

   Addition: +
   Subtraction: -
   Multiplication: *
   Division: /
   Mod: %
• Comparison Operator: <,>,<=,>=, ~=, :=:
   Less than: <
   Greater than: >
   Less than equal to: <=
   Greater than equal to: >=
   Not equals to: ~=
   Equals to: :=:
```

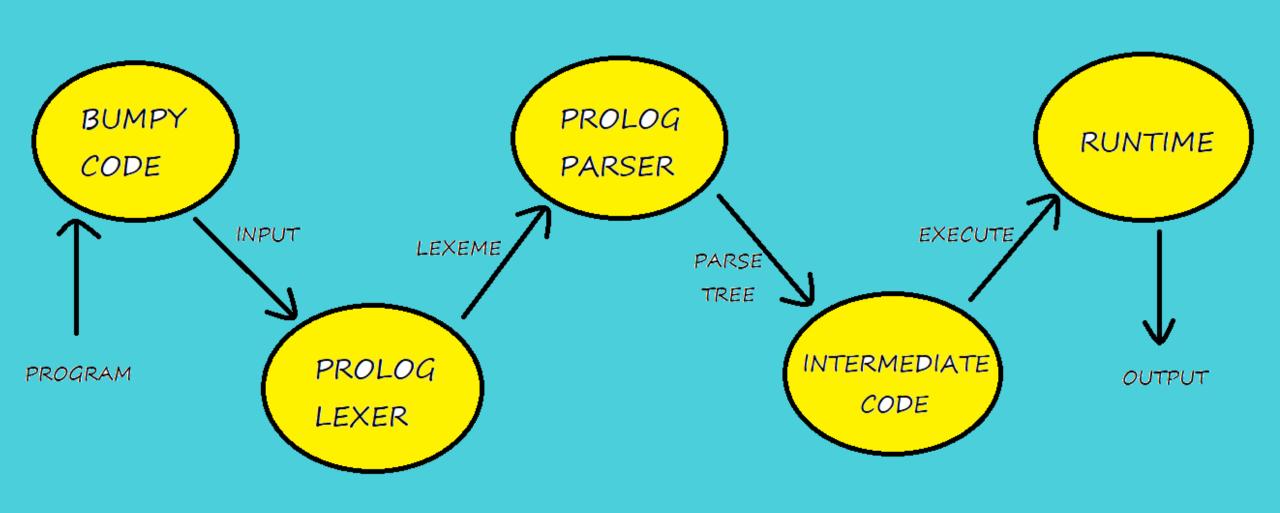
Operators and Constructs For BUMPY

- Assignment Operator: =
- Boolean Operator : and, or ,not
- Primitive types: bool, var

Bool: takes Boolean value

Var: takes integer value

- Decision Constructs: incase do otherwise endcase Incase (condition) do (process) otherwise (process) endcase
- Iterative Constructs: when repeat endrepeat When (condition) repeat (process) endrepeat



Design Flow

Lexical Analyzer and Parser

- Lexical Analyzer and Parser written in Prolog Definite Clause Grammar.
- Once the program is fed to the Lexer, it generates a list of lexemes for program by identifying keywords, eliminating spaces, tabs and new lines.
- After getting the tokens, tokens are passed through the parser which is a prolog DCG code with the grammar defined for the language using prolog predicates and it generates the parse tree for the given code.
- It Parses in a top-down fashion.

Intermediate Code

- Parse tree is itself the intermediate code as we are using Prolog for creating language.
- It gives Intermediate code as output to a separate file.
- Intermediate Code is generated using DCG as well.
- Used interpreter to evaluate the intermediate code.

Runtime/Interpreter

- Programmed in Prolog.
- The Interpreter is implemented by traversing node by node through the generated parse tree. Each node denotes an evaluation step that needs to be performed
- The environment at each node is looked up or updated via a list data structure in prolog.
- Used reduction rules to evaluate the code.

Language Grammar

- Parser → Program
- Program → Comment Block | Block
- Comment → @ Words @
- Block → start Declaration Process stop
- Words → Identifier Words | Numb Words | Identifier | Numb.
- Declaration → Datatype Identifier; Declaration | Datatype Identifier;
- Process → Assignvalue; Process | Control Process | Iterate Process |
 Print Process | ReadValue Process | Assignvalue; | Control | Iterate |
 Print | ReadValue

Language Grammar (Continued..)

- Datatype → var | bool
- Assignvalue → Identifier = Expression | Identifier is Boolexp
- Control → incase Condition do Process otherwise Process endcase
- Iterate → when Condition repeat Process endrepeat
- Print → show Expression; | show * Value *;
- ReadValue → input Identifier;
- Condition → Boolexp and Boolexp | Boolexp or Boolexp | Boolexp | Boolexp | Boolexp | Boolexp | Boolexp

Language Grammar (Continued..)

- Boolexp → Expression :=: Expression | Expression ~= Expression |
 Expression <= Expression | Expression | Expression |
 Expression | Expression | Expression :=:
 Boolexp | Expression ~= Boolexp | yes | no
- Expression → Term + Expression | Term Expression | Term
- Term → Identifier * Term | Numb * Term | Numbneg * Term | Identifier / Term | Numb / Term | Numbneg / Term | Identifier mod Term | Numbneg mod Term | Identifier | Numb | Numbneg
- Identifier → _[^a-z]alphanumeric | [^a-z]alphanumeric
- Numb → number
- Numbneg → Numb

Demonstration: Executing Prolog Files

