

# PR 1

## Database Design

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# Table of Contents

Exercise 1 . . . . .	2
Exercise 2 . . . . .	4
Exercise 3 . . . . .	4
a) Rental1 . . . . .	4
b) Rental2 . . . . .	5

## Exercise 1

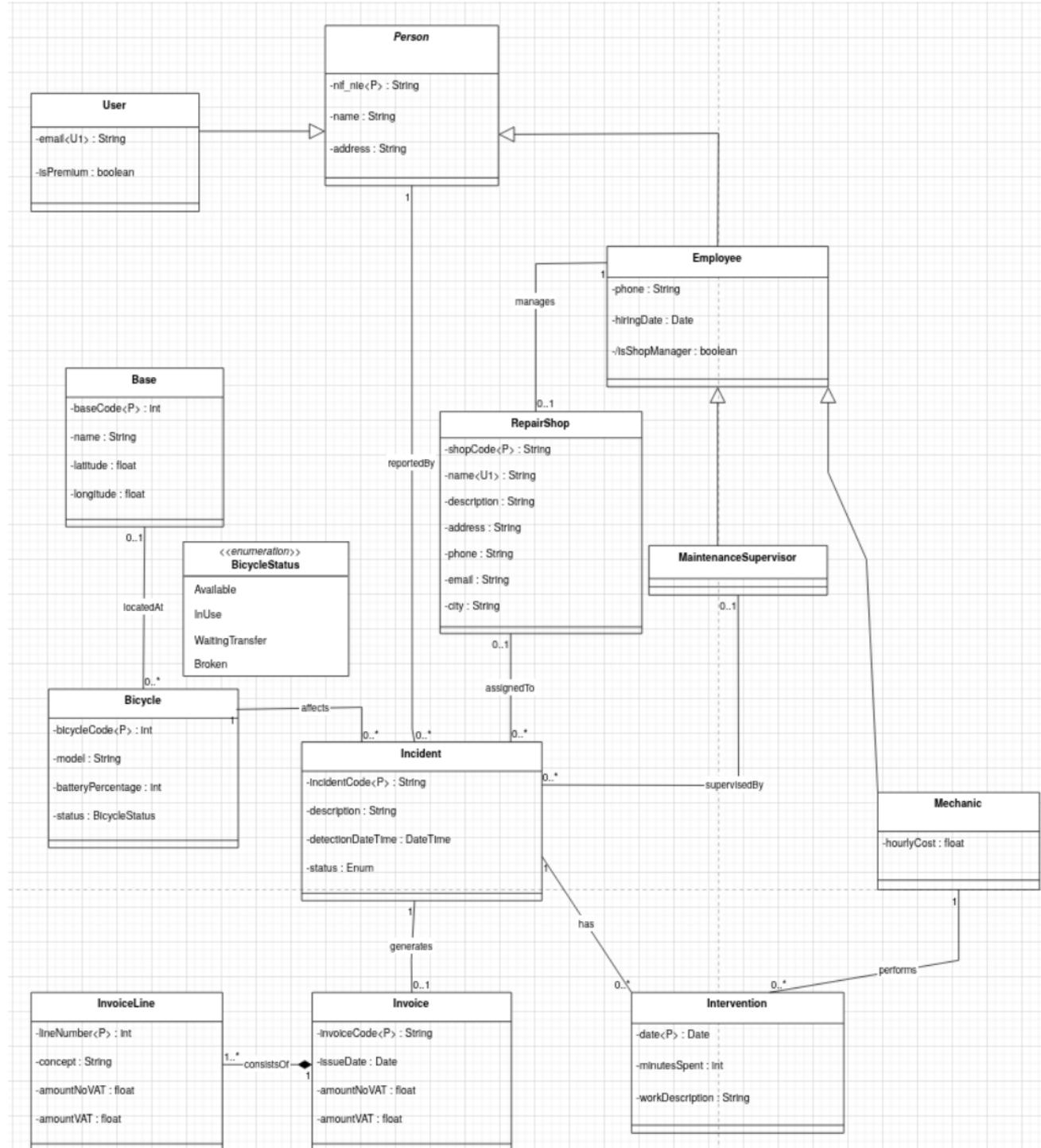


Figure 1: Conceptual Design

Constraints:

1. PK Constraints:

- **InvoiceLine**: Primary Key is composite (**InvoiceCode**, **LineNumber**).
- **Intervention**: Primary Key is composite (**IncidentCode**, **MechanicNIF**, **Date**).

**2. Data Integrity:**

- **Bicycle.batteryPercentage**: Must be integer between 0 and 100.
- **User.email**: Must be Unique.
- **RepairShop.name**: Must be Unique.

**3. Business Logic:**

- A **Bicycle** cannot be at a **Base** if its status is ‘In Use’ or ‘Broken’.
- An **Employee** cannot be both a **Mechanic** and a **MaintenanceSupervisor** (Disjoint).
- **Person** hierarchy is overlapping (Employee can be User).

## Exercise 2

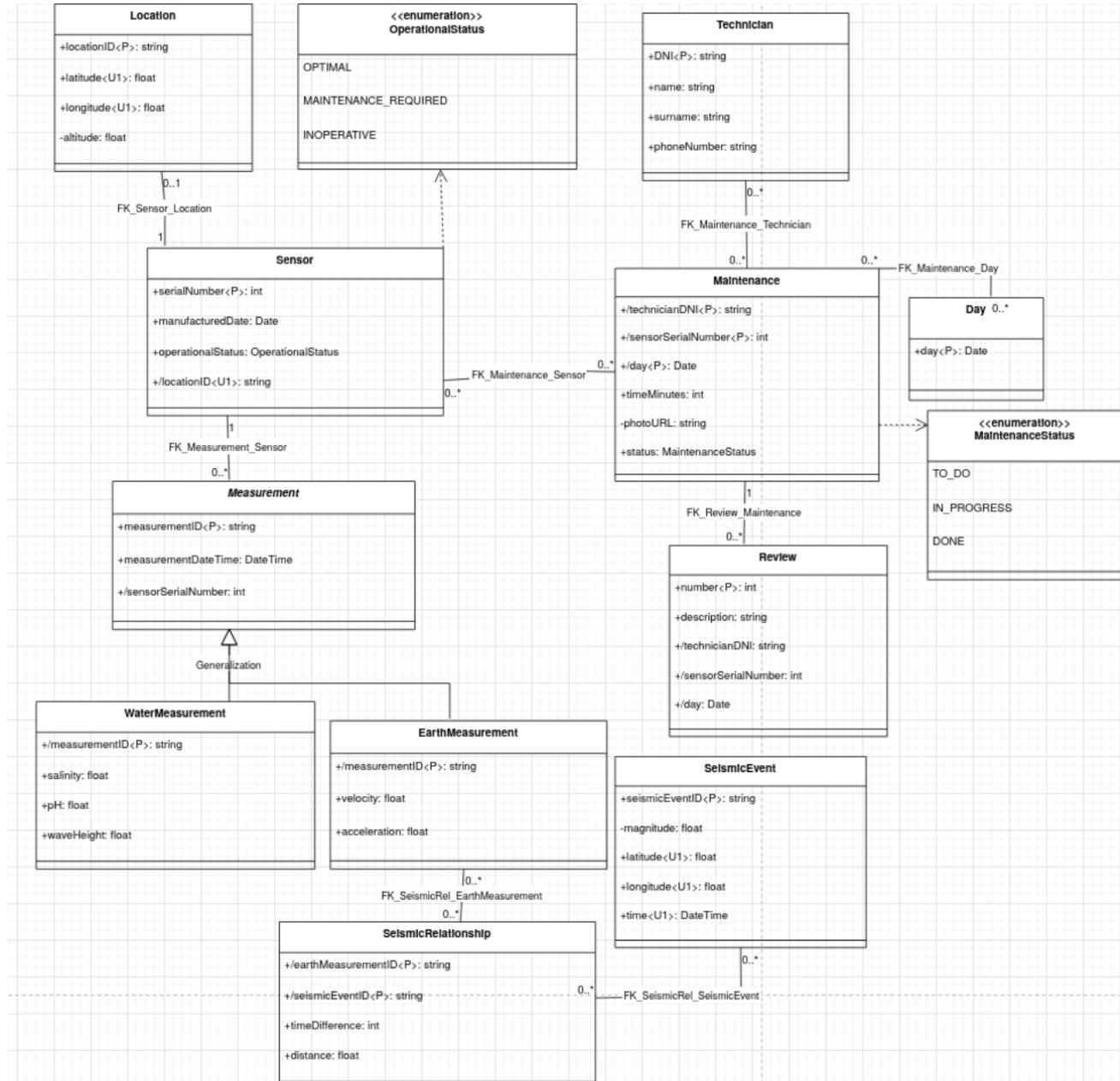


Figure 2: Relational Model

## Exercise 3

### a) Rental1

The relation **Rental1** is currently in 1NF because all the attributes are atomic (as there are no repeating groups or lists).

However, it is not in 2NF. For a relation to be in this form, non-key attributes must depend on the *entire*

primary key, not just part of it. The primary key here is `{momentStart, idUser, codeBicycle}`. Looking at the dependencies:

- The user's `name` and `phone` depend only on `idUser`, which is just a part of the primary key.
- The bicycle `model` depends only on `codeBicycle`, which is also just a subset of the key.

Because these partial dependencies exist, the relation suffers from redundancy and update anomalies.

### Process of transforming it to BCNF:

To reach BCNF, we need to decompose the relations so that every determinant is a candidate key.

### Fixing Partial Dependencies to reach 2NF

We have to separate the attributes that depend on partial keys into their own tables.

1. **User Data:** We can put `name` and `phone` in a separate relation keyed by `idUser`.
2. **Bicycle Data:** We should again put `model` and `batteryType` in a relation keyed by `codeBicycle`.
3. **Rental Data:** The original relation keeps the full primary key plus the attributes that actually depend on the specific rental instance (`momentEnd` and `price`).

After the above changes, we have the following:

- **Users** (`idUser`, `name`, `phone`)
- **Bicycles\_Temp** (`codeBicycle`, `model`, `batteryType`)
- **Rentals** (`momentStart`, `idUser`, `codeBicycle`, `momentEnd`, `price`)

### Fixing 3NF/BCNF by removing transitive dependencies

Looking at the `Bicycles_Temp` table created above, there is still an issue. The `batteryType` depends on the `model`, and the `model` depends on the `codeBicycle`.

- Dependency: `model → batteryType`

Since `model` is not a candidate key for the specific bicycle (many bikes can be the same model), this is a **transitive dependency**, which violates 3NF and BCNF. To fix this, we need to create a reference table for models>

1. Create a `Models` relation where `model` is the key and determines `batteryType`.
2. Update the `Bicycles` relation to only hold the `model` as a foreign key

### Final BCNF

After these steps, every determinant in the relations is a candidate key, as shown below:

1. **Rentals** (`momentStart`, `idUser`, `codeBicycle`, `momentEnd`, `price`)
  - *Foreign Keys:* `idUser` references `Users`, `codeBicycle` references `Bicycles`
2. **Users** (`idUser`, `name`, `phone`)
3. **Bicycles** (`codeBicycle`, `model`)
  - *Foreign Keys:* `model` references `Models`
4. **Models** (`model`, `batteryType`)

### b) Rental2

#### Current State & Justification:

The relation `Rental12` is 2NF. Unlike before, the primary key here (`numberRental`) is a single attribute (an incrementing number). 2NF violations appear when an attribute depends on only *part* of a composite primary key. Since the key has no “parts to split, it is impossible to have a partial dependency.

However the relation is not 3NF, thus also not in BCNF. This is due to it having transitive dependencies. Non-key attributes are determining other non-key attributes. For example:

- `numberRental` determines `idUser`, but `idUser` then determines `name` or `phone`. This creates a chain of “`numberRental -> idUser -> name`” which is not good practice.
- Similarly, `numberRental` determines `codeBicycle`, which determines `model`, which in its turn also determines `batteryType`.

To reach BCNF, these transitive dependencies must be removed by separating the data into distinct relations where every determinant is a candidate key.

## Transformation to BCNF

### Extracting User Info

The attributes `name` and `phone` depend on `idUser`, not on the rental transaction itself. We can remove this transitive dependency by creating a separate table for users.

- **New Relation: Users** (`idUser`, `name`, `phone`)
- **Modified Rental**: `Rental12` retains `idUser` as a foreign key but loses `name` and `phone`.

### Extracting Bicycle and Model Info

There is a nested dependency chain here: `codeBicycle -> model -> batteryType`. If we simply moved all bicycle info to one table, we would still have a 3NF violation because `batteryType` depends on `model`, not strictly on the specific `codeBicycle`. Therefore, this needs to be split into two parts:

1. **Model Data**: We create a relation to store the features that belong to a model type.
  - **New Relation: Models** (`model`, `batteryType`)
2. **Bicycle Data**: We can create a relation for the specific physical bikes, linking them to their model.
  - **New Relation: Bicycles** (`codeBicycle`, `model`)

## Final BCNF

The original `Rental12` relation is now reduced to just the transaction details, referencing the other entities via foreign keys.

1. **Rentals** (`numberRental`, `momentStart`, `idUser`, `codeBicycle`, `momentEnd`, `price`)
  - *Foreign Keys*: `idUser` references `Users`, `codeBicycle` references `Bicycles`.
2. **Users** (`idUser`, `name`, `phone`)
3. **Bicycles** (`codeBicycle`, `model`)
  - *Foreign Keys*: `model` references `Models`.
4. **Models** (`model`, `batteryType`)