

# DENTAL CLINIC DATA MANAGEMENT

**SQL PROJECT** 



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#### **Problem-Statement:**

A Dentists' Polyclinic has various dental departments such as General Dentistry, Periodontist and Endodontist. The patients can go to all these Departments.

The polyclinic has many types of doctors depending on the dental departments.

a) General dentistry: Cavities, Missing Teeth, Mobile Teeth

b) Endodontist: Root canals

c) Periodontist: Gums

The Doctor's work from Monday to Saturday, from 4:00pm-8:30pm. The appointments are given according to these timings. Salary of the doctors along with the number of years they have been working in the clinic is stored to keep a track of their performance and the number of years they have been serving, based on which the salary raise and promotions can be done.

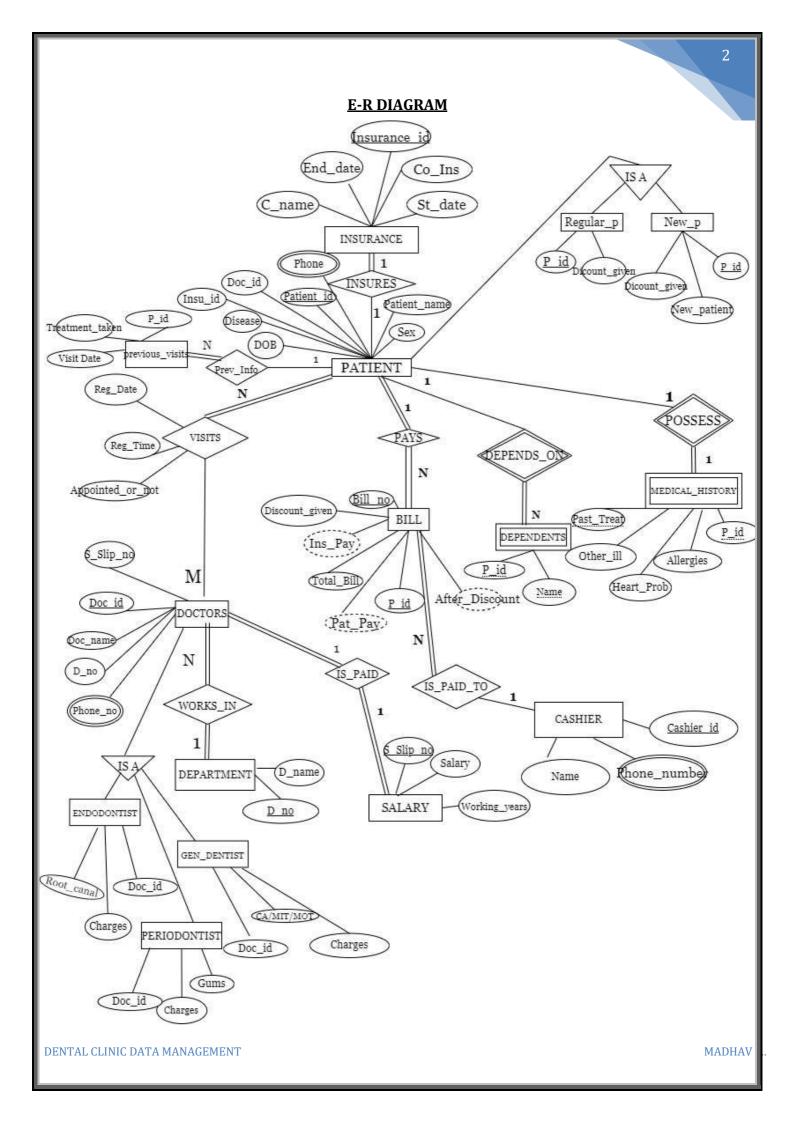
When a customer arrives in the dental clinic, we identify them based on their patient id which is unique for each patient. We ask for all the details. i.e. the name of the patient, disease, sex, phone number and we even feed their arrival time and arrival date into our system based upon which we decide whether the doctor is available or not, if the patient registers within the specified time slot (Monday-Saturday,4:00pm-8:30pm) then appointment can be made otherwise notification will be displayed asking the patient to come within the time slot.

If a customer has health insurance, which is nowadays very important as the medical costs are very high in the country and without insurance, a large portion of population can't afford the healthcare, we store the insurance ID (unique for each customer), company name, start date, end date and Co-Insurance. Co-Insurance is a percentage amount that insurance company pays for a medicinal purchase or treatment. In case the insurance of any patient expires then he needs to pay all the amount of the treatment by himself. The previous record of the patients (On what dates they have visited the clinic in the past based upon which the system automatically calculates how many times they have visited to the clinic, what all treatments did they take) are also stored in the database of the clinic .

The patients can be of two types, a regular patient and a new patient. Patients who came to the clinic 2 or more than 2 times should be considered as regular patients and they should be given a 10% discount (regular patient). A regular patient's history is accessible. If a new patient arrives then he gets the first check-up free. The patients' medical history such as heart problems, allergies, other illnesses are stored within the database so that the doctor can give a treatment accordingly. A patient can come to the polyclinic with a dependent. The information about the dependent of the patient is stored in the database as well.

Once the treatment has been completed, a bill is generated by the system. This bill is handed over to the customer by the cashier and it contains patient id, insurance information, total charges along with the treatment taken, Price after giving discount (in case of regular and new patients), Price to paid by the patient and the insurance company. The breakdown should be automatically calculated by the system and should be stored in the table ( TOTAL BILL).

Lastly the details of the cashier should also be stored in the system .



#### **SCHEMA**

#### **INSURANCE**

<u>Insurance_id</u>	Company_name	Start_date	End_date	Co_Insurnace

PRIMARY KEY: Insurance\_id

#### PATIENT1

<u>Patient_id</u>	Polyclinic_name	Patient_name	dob	Insurance_id	Sex	Disease	Doc_id
Registration_time	Registration_date						

PRIMARY KEY: Patient\_id

UNIQUE: Patient\_name

**FOREIGN KEY**: Patient1(doc\_id) references doctor\_info(doc\_id)

Patient1(insurance\_id) references insurance(insurance\_id)

#### PATIENT\_PHONE

Patient_id	Phone_number
_ · · · · · · <u>-</u> · ·	

**FOREIGN KEY**: Patient\_phone(patient\_id) references patient1(patient\_id)

#### **VISITS**

Patient id	Patient_name	registration_time	registration_date	Appointed_or_not

PRIMARY KEY: PATIENT\_ID

**FOREIGN KEY**: Visits (patient\_name) references patient1(patient\_name)

Visits (patient\_id) references patient1 (patient\_id)

#### PREVIOUS\_VISITS

Patient_id visits	prev_treatment_taken_from_this_clinic
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**FOREIGN KEY**: Previous\_visits (patient\_id) references patient1(patient\_id)

#### **NEW PATIENT**

<u>Patient id</u>	Patient_name	Insurance_id	New_patient	Discount_given

PRIMARY KEY: Patient\_id

**FOREIGN KEY:** new\_patients (patient\_id) references patient1(patient\_id);

New\_patients (patient\_name) references patient1(patient\_name)

New\_patients (insurance\_id) references patient1(insurance\_id)

#### **REGULAR PATIENT**

Patient id Patient_name	Insurance_id	Discount_given
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PRIMARY KEY: Patient\_id

**FOREIGN KEY**: regular\_patients (patient\_id) references patient1(patient\_id);

regular\_patients (patient\_name) references patient1(patient\_name)

regular\_patients (insurance\_id) references patient1(insurance\_id)

#### **DOCTOR\_INFO**

<u>Doc id</u>	Doc id	Salary_slipno	Doc_name	Dep_no	Dep_name
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PRIMARY KEY: Doc\_id

UNIQUE: Salary\_slipno

**FOREIGN KEY**: Doctor\_info (Dep\_no) references department(Dep\_no)

Doctor\_info references department(dep\_name)

Doctor\_info (salary\_slipno) references doc\_salary (salary\_slipno)

#### **DOCTOR PHONE**

Doc_id	Phone_number

**FOREIGN KEY**: DOCTOR\_PHONE (doc\_id) references doctor\_info(doc\_id)

#### **DOCTOR SALARY**

<u>Salary_slipno</u>	Salary	Number_of_years_working

PRIMARY KEY: Salary slipno

#### **DEPARTMENT**

PRIMARY KEY: dep\_no

UNIQUE: Dep\_name

#### **ENDODONTIST**

Doc_id Root_canal Charges
---------------------------

**FOREIGN KEY:** Endodontist (doc\_id) references doctor\_info(doc\_id)

#### **PERIODONTIST**

Doc id	Gums	Price
Doc_iu	dullis	11100

**FOREIGN KEY:** periodontist (doc\_id) references doctor\_info(doc\_id)

#### **GEN DENTIST**

Doc_id	cavities_OR_missing_teeth_OR_mobile_teeth	Price

**FOREIGN KEY:** gen\_dentist (doc\_id) references doctor\_info(doc\_id)

#### **TOTAL BILL**

Bill_no	Patient_id	Total_charge	After_discount	Money_insurance	Patient_pay	Cashier_id

PRIMARY KEY: bill\_no,patient\_id

**FOREIGN KEY:** TOTAL\_BILL (cashier\_id) references cashier(cashier\_id)

TOTAL\_BILL (Insurance\_id) references patient1(insurance\_id)

TOTAL\_BILL (patient\_id) references patient1(patient\_id)

TOTAL\_BILL (patient\_name) references patient1(patient\_name)

#### **DEPENDENTS**

<u>Dependent name</u>	<u>Patient id</u>	Phone_number
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PRIMARY KEY: Dependent\_name,patient\_id

**FOREIGN KEY**: Dependents(patient\_id) references patient1(patient\_id)

#### **MEDIC HISTORY**

Patient id	Past treatment	Allergies	Pain_tooth	Heart_probs	Other_illness

PRIMARY KEY: Patient\_id,past\_treatment

**FOREIGN KEY**: Medic\_history (patient\_id) references patient1(patient\_id)

#### **CASHIER**

Cashie	<u>id</u>	Name	Phone_number	Salary
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PRIMARY KEY: Cashier\_id

#### **CASHIER\_PHONE**

Cashier_id Phone_number
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**Foreign key**: Cashier\_phone(cashier\_id) references cashier(cashier\_id)

# POPULATED TABLES AND TABLE DESCRIPTION

### Insurance

	Field	Type	Null	Key	Default	Extra
•	insurance_id	int	NO	PRI	HULL	
	company_name	varchar(50)	NO	MUL	MULL	
	start_date	date	NO		HULE	
	end_date	date	NO		NULL	
	co_insurance	decimal(5,2)	YES		HULL	

	insurance_id	company_name	start_date	end_date	co_insurance
•	101	National Insurance Co.Ltd	2011-03-12	2020-04-10	55.00
	102	Go digital General Insurance	2011-09-12	2023-04-10	40.00
	103	HDFC ERGO General insurance	2010-06-01	2024-05-07	60.00
	104	HDFC ERGO General insurance	2010-06-01	2024-05-07	60.00
	105	National Insurance Co.Ltd	2008-03-09	2022-09-23	30.00
	106	National Insurance Co.Ltd	2008-03-09	2022-09-23	30.00
	NULL	NULL	HULL	NULL	NULL

### Patient1

	Field	Type	Null	Key	Default	Extra
١	patient_id	int	NO	PRI	HULL	
	polyclinic_name	varchar(20)	NO		MULL	
	patient_name	varchar(20)	NO	UNI	NULL	
	dob	date	NO		HULL	
	insurance_id	int	YES	MUL	HULE	
	sex	char(4)	NO		NULL	
	Problem_or_Disease	varchar(50)	NO		HULL	
	Dno	int	NO		NOLE	
	doc_id	int	NO	MUL	HULL	
	registration_time	time	NO		HULL	
	registration_date	date	NO		NULL	

patient_id	polydinic_name	patient_name	dob	insurance_id	sex	Problem_or_Disease	Dno	doc_id	registration_time	registration_date
1	Dental Polydinic	Mr.Smith	1967-03-25	101	M	Soft tissue Inflammation	1	100	17:00:00	2022-03-19
2	Dental Polyclinic	Mr.Andrews	1978-02-04	102	M	Gum Disease	2	300	14:00:00	2022-03-20
3	Dental Polydinic	Mrs.Rodriguez	1987-07-28	103	F	Deep Decay	1	200	17:00:00	2022-03-21
4	Dental Polydinic	Mr.Holt	1983-08-21	104	M	Cavities	3	400	21:00:00	2022-03-19
5	Dental Polydinic	Ms.Ruby	1998-01-16	105	F	Missing Teeth	3	400	17:00:00	2022-03-20
6	Dental Polydinic	Ms.Franceska	2000-03-19	106	F	Mobile Teeth	3	500	18:00:00	2022-03-19

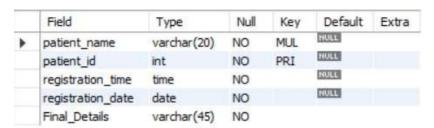
# Patient\_Phone

Field	Туре	Null	Key	Default	Extra
 patient_id	int	NO	MUL	NULL	
Phone_number	decimal(10,0)	NO		NULL	

	patient_id	Phone_number
٠	1	9821000690
	1	8999452345
	2	9811223300
	2	9786577724
	3	9013211091
	4	9210747010
	5	9900887045
	5	9900889085
	6	9601887095

(Made a table for patient\_phone as phone\_ number is a multivalued attribute in patient1 table)

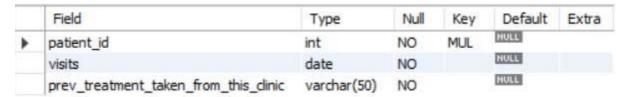
#### **Visits**



	patient_name	patient_id	registration_time	registration_date	Final_Details
•	Mr.Smith	1	17:00:00	2022-03-19	REGISTRATION CAN BE DONE
	Mr.Andrews	2	14:00:00	2022-03-20	SORRY I COME WITHIN THE SPECIFIED TIMINGS
	Mrs.Rodriguez	3	17:00:00	2022-03-21	REGISTRATION CAN BE DONE
	Mr.Holt	4	21:00:00	2022-03-19	SORRY ! COME WITHIN THE SPECIFIED TIMINGS
	Ms.Ruby	5	17:00:00	2022-03-20	SORRY! WE ARE OPEN ONLY FROM MONDAY-S
	Ms.Franceska	6	18:00:00	2022-03-19	REGISTRATION CAN BE DONE

This table is made with the help of table patient1(table made from a query) using 'create table as' statement. Depending upon the registration timings and this table informs us whether the patient can be given an appointment or not and with the help of case function we display the final details.

# Previous\_Visits



	patient_id	visits	prev_treatment_taken_from_this_dinic
١	1	2017-08-11	Root Canal
	1	2019-02-07	Root Canal
	2	2016-12-11	Gums
	2	2018-12-11	Gums
	3	2018-03-10	Cavities
	4	2018-03-23	Missing Teeth
	5	2017-10-25	Mobile Teeth

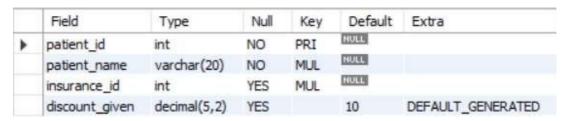
### New\_patient



	patient_id	patient_name	insurance_id	New_patient	discount_given
•	6	Ms.Franceska	106	WELCOME!YOUR FIRST CHECKUP IS FREE	100.00
	HOLL	HULL	NULL	NULL	NULL

This table is made with the help of patient1 and previous visits (i.e table made from nested queries) using 'create table as' statement where the entries in the table patient1 and previous visit table are compared and the details of the patient whose patient\_id is present in patient1 but not in previous\_visit is stored in the table new\_patient.

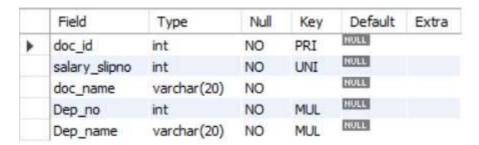
### Regular\_patients



	patient_id	patient_name	insurance_id	discount_given
١	1	Mr.Smith	101	10.00
	2	Mr.Andrews	102	10.00
	NULL	NULL	NULL	NULL

This table is made with the help of patient1 and previous visits (i.e table made from nested queries) using 'create table as' statement where the table patient1 and previous visit table are compared and the details of the patient whose patient\_id is present in previous\_visits more than 2 times is stored in the table regular\_patient.

#### **Doctors**



	doc_id	salary_slipno	doc_name	Dep_no	Dep_name
Þ	100	100	Dr. Ray	1	Endodontist
	200	101	Dr. Bing	1	Endodontist
	300	102	Dr. Stromberg	2	Periodontist
	400	103	Dr. David	3	General Dentist
	500	104	Dr. James	3	General Dentist
	HULL	HULL	NULL	HULL	NULL

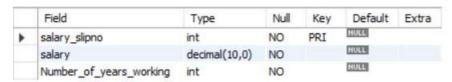
# DOCTOR\_PHONE

	Field	Туре	Null	Key	Default	Extra
٠	doc_id	int	NO	MUL	HULL	
	Phone_number	decimal(10,0)	NO		NULL	

	doc_id	Phone_number
١	100	9821054690
	100	8976452345
	200	9811223344
	200	9786574624
	300	9143211091
	400	9213447010
	500	9900887755

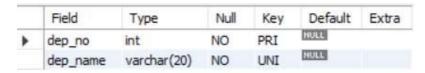
(Made a table for doctor\_phone as phone\_ number is a multivalued attribute in doctor\_info table)

# **Salary**



	salary_slipno	salary	Number_of_years_working
•	100	500000	4
	101	250200	1
	102	512200	5
	103	700000	8
	104	656666	6
٠	NULL	NULL	NULL

# **Department**



	dep_no	dep_name
•	1	Endodontist
	3	General Dentist
	2	Periodontist
	HULL	MULL

# Endodontist

	Field	Type	Null	Key	Default	Extra
Þ	doc_id	int	NO	MUL	NULL	
	root_canal	varchar(100)	NO		HULL	
	charges	int	NO		HULL	

	doc_id	root_canal	charges
١	100	Soft tissue inflammation	4000
	200	Deep decay	7000

# Periodontist

	Field	Type	Null	Key	Default	Extra
١	doc_id	int	NO	MUL	NULL	
	gums	varchar(100)	NO		NULL	
	price	int	NO		NULL	

	doc_id	gums	price
•	300	Gum Disease	6000

# Gen\_dentist

	Field	Type	Null	Key	Default	Extra
١	doc_id	int	NO	MUL	HULL	
	cavities_OR_missing_teeth_OR_mobile_teeth	varchar(20)	NO		HULL	
	PRICE	int	NO		NULL	

	doc_id	cavities_OR_missing_teeth_OR_mobile_teeth	PRICE
١	400	Cavities	2000
	400	Missing Teeth	2500
	400	Mobile Teeth	2700
	500	Cavities	3000
	500	Missing Teeth	3500
	500	Mobile Teeth	3700

#### TOTAL\_BILL



Patient_id	insurance_id	patient_name	charges	discount_given	charge_after_discount	Money_Insurance	Patient_Pay	bill_no	cashier_id
1	101	Mr.Smith	4000	400	3600	1980	1620	1	302
2	102	Mr.Andrews	6000	600	5400	2160	3240	2	301
3	103	Mrs.Rodriguez	7000	0	7000	4200	2800	3	302
4	104	Mr.Holt	2000	0	2000	1200	800	4	301
5	105	Ms.Ruby	2500	0	2500	750	1750	5	302
6	106 (333)	Ms.Franceska	3700	3700	0	0	O COLUMN	6	301 (2209

Here we made a table total\_bill with the help of patient1,endodontist,pedodontist and gen\_dentist table(table made from query) with the help of 'create table as' statement and further with the help of case function we made a column named charges and filled that particular column on the basis of patient\_id and the treatment taken, this is how we found out how much the total charge of that particular patient was and next we updated the column discount\_given,money\_insurance\_patient\_pay twith the help of 'update inner join' statement

(For discount\_given- **update** total\_bill e **INNER JOIN** regular\_patients r **ON** e.patient\_id = r.patient\_id **SET** e.discount\_given = (charges \* r.discount\_given/100.00);

(For money\_insurance- **update** total\_bill e **inner join** insurance i **on** e.insurance\_id=i.insurance\_id **SET** Money\_insurance= ((charge\_after\_discount) \* co\_insurance/100.00);)

(FOR patient\_pay- update total\_bill e inner join insurance i on e.insurance\_id=i.insurance\_id

set patient\_pay=(charge\_after\_discount)-money\_insurance)

### **Dependents**



	depen_name	phone_no	patient_id
١	Roger	9165625400	1
	Fin	9165623880	2
	Thomas	9789879765	3
	Alfie	9914323523	4
	Arthur	9678229119	5
	HULL	NULL	NULL

# **Medical History**

	Field	Type	Null	Key	Default	Extra
١	patient_id	int	NO	PRI	NULL	
	past_treatment	varchar(50)	NO	PRI	HULL	
	allergies	varchar(50)	YES		MULL	
	pain_tooth	varchar(50)	YES		NULL	
	heart_probs	varchar(50)	YES		NULL	
	other_illness	varchar(50)	YES		NULL	

	patient_id	past_treatment	allergies	pain_tooth	heart_probs	other_illness
•	1	Root Canal	Penicillin	MULL	High BP	Diabetes
	2	Root Canal	NULL	Upper Left Tooth	HULL	Rhinitis
	3	Loose Teeth	Pollen	MULL	High BP	Arthritis
	4	Decay	Pollen	Lower Left Side	HULL	HOLL
	5	Gingivitis	Lignocaine	Lower Right Side	High BP	Cardiac Problem
	HULL	RURU	NULL	HULL	NULL	RULL

# **CASHIER**

	Field	Type	Null	Key	Default	Extra
١	Name	varchar(20)	NO		NULL	
	cashier_id	int	NO	PRI	HULL	
	salary	int	NO		HULL	

	Name	cashier_id	salary
Þ	Amit	301	3000
	Rohit	302	400
	NULL	HUGH	NULL

# **CASHIER PHONE**

	Field	Туре	Null	Key	Default	Extra
١	cashier_id	int	NO	MUL	NULL	
	Phone_number	decimal(10,0)	NO		NULL	

	cashier_id	Phone_number
Þ	301	9999765642
	301	6542758545
	302	8645324455
	302	7689000678

# **QUERIES:**

```
/*Question-01 */
```

select \* from patient1 where patient\_id IN (Select patient\_id from medic\_hist where heart\_probs='High BP');

patient_id	polyclinic_name	patient_name	dob	insurance_id	sex	Problem_or_Disease	Dno	doc_id	registration_time	registration_date
1	Dental Polyclinic	Mr.Smith	1967-03-25	101	M	Soft tissue Inflammation	1	100	17:00:00	2022-03-19
3	Dental Polydinic	Mrs.Rodriguez	1987-07-28	103	F	Deep Decay	1	200	17:00:00	2022-03-21
1 3 5 HULL	Dental Polyclinic	Ms.Ruby	1998-01-16	105 HULL	F	Missing Teeth	3 MULL	400	17:00:00	2022-03-20 NULL

#### /\*Question-02 \*/

select cashier.Name,cashier\_id,salary,phone\_number from cashier,cashier\_phone where
cashier\_id=cashier\_id=cashier\_id;
select cashier.name,cashier.cashier\_id,salary,phone\_number from cashier left outer join cashier\_phone on

	name	cashier_id	salary	phone_number
Þ	Amit	301	3000	6542758545
	Amit	301	3000	9999765642
	Rohit	302	400	7689000678
	Rohit	302	400	8645324455

cashier.cashier\_id=cashier\_phone.cashier\_id;

/\*Question-03 \*/

select \* from patient1 where patient\_id NOT IN (select patient\_id from dependents);

patient_id	polydnic_name	patient_name	dob	insurance_id	sex	Problem_or_Disease	Dno	doc_id	registration_time	registration_date
6	Dental Polyclinic	Ms.Franceska	2000-03-19	106	F	Mobile Teeth	3	500	18:00:00	2022-03-19
6	DOM:	EUSE .	DOM:	THE R. LEWIS CO., LANSING, MICH.	<b>HARR</b>	CITE S	District of the last	12223	MARK TO SERVICE STREET	OH COLUMN TO SERVICE STATE OF THE PERSON NAMED IN COLUMN TO SERVICE STAT

#### /\*Question-04 \*/

```
select doctor_info.doc_id,doc_name,dep_no,dep_name,phone_number,doctor_info.salary_slipno,salary,
Number_of_years_working
from doctor_info,doctor_phone,doc_salary where doctor_info.doc_id IN
(select doc_id from doctor_phone group by doc_id having count(doc_id) >= 2 )
AND doctor_info.doc_id=doctor_phone.doc_id and
doc_salary_salary_slipno=doctor_info.salary_slipno;
```

	doc_id	doc_name	dep_no	dep_name	phone_number	salary_slipno	salary	Number_of_years_working
٠	100	Dr. Ray	1	Endodontist	9821054690	100	500000	4
	100	Dr. Ray	1	Endodontist	8976452345	100	500000	4
	200	Dr. Bing	1	Endodontist	9811223344	101	250200	1
	200	Dr. Bing	1	Endodontist	9786574624	101	250200	1

#### /\*Question-05 \*/

select \* from gen\_dentist where doc\_id IN (select doc\_id from doctor\_info where doc\_name='Dr. David');

	doc_id	cavities_OR_missing_teeth_OR_mobile_teeth	PRICE
Þ	400	Cavities	2000
	400	Missing Teeth	2500
	400	Mobile Teeth	2700

#### /\*Question-06 \*/

select endodontist.doc\_id,doc\_name,root\_canal,charges,dep\_name from endodontist,doctor\_info where
endodontist.doc\_id=doctor\_info.doc\_id;

	doc_id	doc_name	root_canal	charges	dep_name
١	100	Dr. Ray	Soft tissue inflammation	4000	Endodontist
	200	Dr. Bing	Deep decay	7000	Endodontist

#### /\*Question-07 \*/

> select \* from doctor\_info where salary\_slipno IN (select salary\_slipno from doc\_salary where salary>
 (select AVG(salary) from doc\_salary));

	doc_id	salary_slipno	doc_name	Dep_no	Dep_name
١	400	103	Dr. David	3	General Dentist
	500	104	Dr. James	3	General Dentist
	HULL	MULL	HULL	HULL	NULL

## **NORMALISATION OF TABLES:**

# 1.) <u>Insurance</u>

Insurance\_id -> company\_name,start\_date,end\_date,co\_insurance

#### **Finding Closure:**

(Insurance\_ID)+={ company\_name,start\_date,end\_date,co\_insurance, Insurance\_id}
Candidate keys={Insurance\_id}

#### Normalisation:

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 2.) Patient1

Patient\_id->patient\_id, polyclinic\_name, patient\_name,dob,insurance\_id, sex, problem\_or\_disease, Dno,Doc\_id, registration\_time, registration\_date

Patient name->patient id

Insurance id->patient id

#### Finding Closure:

(Patient\_ID)+={ Patient\_ID , polyclinic\_name Patient name, sex, problem\_or\_disease, Dno,Doc\_id,registration\_time, registration\_date}

(Patient\_name)+={ Patient\_ID, polyclinic\_name Patient name, sex, problem\_or\_disease, Dno,Doc\_id,registration\_time, registration\_date}

(Insurance\_id)<sup>+</sup>={ Patient\_ID , polyclinic\_name Patient name, sex, problem\_or\_disease, Dno,Doc\_id,registration\_time, registration\_date}

Candidate keys={ Patient\_ID, Patient\_name, Insurance\_id}

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 3.) Patient phone

Phone\_number→patient\_id

This table was already divided as I observed that the table Patient1 had multiple phone numbers associated with a single patient therefore the attribute phone\_number which is multivalue attribute was divided and kept into different table.

#### Finding closure:

(Phone\_number)+={ Phone\_number,patient\_id}

Candidate key={ Phone\_number}

#### Normalisation:

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 4.) <u>Visits:</u>

Patient\_id -> Patient name,registration\_time,registration\_date,final\_details

Registration\_time, ,registration\_date -> Final\_details

Patient\_name -> patient\_id

#### Finding Closure:

(Patient\_ID)+={ Patient\_ID ,Patient name,registration\_time, registration\_date, final\_details}

(Registration\_time, ,registration\_date)+= { registration\_time, registration\_date, final\_details }

(Patient\_name)+={ Patient\_name,patient\_id, registration\_time, registration\_date, final\_details}

Candidate key={ Patient\_ID ,Patient name }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Registration\_time, ,registration\_date being a non-prime attribute determines Final\_details which is another non-prime attribute. Hence the table is not in 3NF. To convert it into 3NF, we make an entirely new table for

Registration\_time, ,registration\_date -> Final\_details

Therefore we split table into two tables:

R<sub>1</sub> (Patient name,registration\_time,registration\_date)

R<sub>2</sub> (Registration\_time, ,registration\_date, final\_details)

Next the table  $R_1$  and  $R_2$  are further in BCNF as the determining attributes are Super keys in  $R_1$  and  $R_2$ 

# 5.) Previous Visits:

Patient\_id,visits -> prev\_treatment taken from clinic

#### **Finding Closure:**

Patient\_id, visits={prev\_treatment taken from clinic, patient\_id, visits}

Candidate key={ Patient\_id,visits }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

# 6.) New patients

```
Patient_id → patient_name,insurance_id,discount_given
```

Patient\_name → patient\_id

Insurance\_id → patient\_id

#### **Finding Closure:**

```
(Patient_ID)+={ Patient_id, patient_name,insurance_id,discount_given}
```

(Patient\_name)+={ Patient\_id, patient\_name,insurance\_id,discount\_given}

(Insurance\_id)+= { Patient\_id, patient\_name,insurance\_id,discount\_given}

Candidate key={ Patient\_ID, Patient\_name, Insurance\_id }

#### **Normalization of table:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 7.) Regular Visits

Patient\_id → patient\_name,insurance\_id,discount\_given

Patient\_name → patient\_id

Insurance\_id → patient\_id

#### Finding Closure:

(Patient\_ID)+={ Patient\_id, patient\_name,insurance\_id,discount\_given}

(Patient\_name)+={ Patient\_id, patient\_name,insurance\_id,discount\_given}

(Insurance\_id)+= { Patient\_id, patient\_name,insurance\_id,discount\_given}

Candidate key={ Patient\_ID, Patient\_name, Insurance\_id }

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

### 8.) **DOCTORS**

```
Doc_id → salary_slipno,doc_name,dep_no,dep_name
```

Salary\_slipno→doc\_id

Doc\_name→doc\_id

Dep\_no→dep\_name

Dep\_name→dep\_no

### **Finding closure:**

```
(Doc_id)+={doc_id, salary_slipno,doc_name,dep_no,dep_name}
(Salary_slipno)+={ doc_id, salary_slipno,doc_name,dep_no,dep_name}
(Doc_name)+= { doc_id, salary_slipno,doc_name,dep_no,dep_name}
(Dep_no)+= { Dep_no,dep_name}
(Dep_name)+= { Dep_no,dep_name}
```

Candidate key={ Doc\_id, Salary\_slipno, Doc\_name}

#### Normalization of table:

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Dep\_no being a non-prime attribute determines dep\_name which is another non-prime attribute. Hence the table is not in 3NF. To convert it into 3NF, we make an entirely new table for

Dep\_no→dep\_name

Therefore we split table into two tables:

R<sub>1</sub> (doc\_id, salary\_slipno,doc\_name,dep\_no)

R<sub>2</sub> (Dep\_no,dep\_name)

Next the table  $R_1$  and  $R_2$  are further in BCNF as the determining attributes are Super keys in  $R_1$  and  $R_2$ 

### 9.) **DOCTOR PHONE**

Phone\_number→doctor\_id

This table was already divided as I observed that the table Doctors had multiple phone numbers associated with a single doctor therefore the attribute phone\_number which is multivalue attribute was divided and kept into different table.

#### Finding closure:

(Phone\_number)+={ Phone\_number,doctor\_id}

Candidate key={ Phone\_number}

#### Normalisation:

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# **10.) SALARY**

Salary\_slipno→salary,no\_of\_years\_working

#### **Finding closure:**

(Salary\_slipno)+={ Salary\_slipno, salary,no\_of\_years\_working }

Candidate key={ Salary\_slipno }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

### 11.) **DEPARTMENT:**

```
Dep_no→ dep_name
```

Dep\_name→dep\_no

#### Finding closure:

```
(Dep_no)+={ Dep_no,dep_name}
```

(Dep\_name)+={ Dep\_no,dep\_name}

Candidate key={ Dep\_no,dep\_name }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 12.) Endodontist

Doc id -> rootcanal, charges

Root\_canal -> doc\_id

#### Finding closure:

```
(Doc_id) + = {root_canal,charges,doc_id}
```

(Root\_canal) + = {root,canal,doc\_id,charges}

Candidate key={ root\_canal,doc\_id }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

# 13.) Periodontist

```
Doc_id -> gums , price
Gums -> doc_id
```

#### Finding closure:

```
(Doc_id) + = {gums, price,Doc_id }
(Gums) + = {doc_id,gums,price}
Candidate key={ doc_id,gums}
```

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 14.) Gen dentist

Doc\_id ,cavities\_or\_missing teeth\_or\_mobile\_teeth -> price

#### Finding closure:

(Doc\_id ,cavities\_or\_missing teeth\_or\_mobile\_teeth )+ = { Doc\_id ,cavities\_or\_missing teeth\_or\_mobile\_teeth, price }

Candidate key={ Doc\_id ,cavities\_or\_missing teeth\_or\_mobile\_teeth }

#### Normalisation:

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

# 15.) Total Bill

Patient\_id ,bill\_no -> insurance\_id,patient\_name,charges,discount\_given,charges after discount,money\_insurance , patient\_pay,bill\_no,cashier\_id

Insurance\_id -> Patient\_id ,bill\_no

Patient name -> Patient\_id ,bill\_no

#### Finding closure:

(Patient\_id,bill\_no)+ =Patient\_id,bill\_no, insurance\_id,patient\_name, charges, discount\_given, charges after discount,money\_insurance, patient\_pay,bill\_no,cashier\_id.

(Insurance\_id) + = Patient\_id, Bill\_no, insurance\_id, patient\_name, charges, discount\_given,charges after discount,money\_insurance, patient\_pay,bill\_no,cashier\_id.

(Patient\_name) += Patient\_id, Bill\_no, insurance\_id, patient\_name, charges, discount\_given, charges after discount, money\_insurance, patient pay, bill no, cashier id.

Candidate key={ Patient\_id ,bill\_no, Insurance\_id, Patient name }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 16.) <u>Dependents</u>

Depen name , patient\_id -> phone number

Phone number -> patient\_id,

dependent name,phone number -> patient\_id

#### Finding closure:

(Depen\_name,patient\_id)+ = {Depen\_name, patient\_id, phone number}

(Phone number)+ = {phone number ,patient\_id}

(Dependent\_name, phone number) + = {Dependent\_name, phone number, patient\_id}

Candidate key={ Dependent\_name , phone number, patient\_id}

#### Normalisation:

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 17.) Medical History

Patient\_id, past\_treatment->allergies,pain\_tooth,heart\_probs,other\_illness

#### **Finding closure:**

(Patient\_id, past\_treatment)+={ allergies,pain\_tooth,heart\_probs,other\_illness,
Patient\_id, past\_treatment }

Candidate keys={ Patient\_id, past\_treatment }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

# 18.) CASHIER

Cashier\_id->Name,Salary

Name->cashier\_id

#### **Finding closure:**

(Cashier\_id)+={ Name, Salary, Cashier\_id}

(Name)+={ Name, Salary, Cashier\_id}

Candidate keys={ Name, Salary, Cashier\_id }

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.

Since the table is in 3NF and the determining attributes are Super keys, the table is also in BCNF.

### 19.) CASHIER PHONE

Phone\_number→cashier\_id

This table was already divided as I observed that the table Doctors had multiple phone numbers associated with a single doctor therefore the attribute phone\_number which is multivalue attribute was divided and kept into different table.

#### **Finding closure:**

(Phone\_number)+={ Phone\_number,cashier\_id}

Candidate key={ Phone\_number}

#### **Normalisation:**

Since all the attributes of the relation are atomic values and all the tuples have single values for the domain of their attributes, the table is in 1NF.

Since the table is in 1NF and there are no partial dependencies, it is also in 2NF.

Since the table is in 2NF and there are no transitive dependencies, it is also in 3NF.