Tommorow → 8-10 am 'Sat'

2D Arrays
& 2D ArrayLists

```
for(int i=2; i<=n; i*=i){
    cout("GFG");
}
```

$i = 2, 4, 16, 256, (256)^2, \ldots$ — n

$2^1, 2^2, 2^4, 2^8, 2^{16}, \ldots$ — 2^x

1 2 4 8 16 . . . . n

y terms

$x = 2^y \Rightarrow y = \log(x)$

$2^n = n$

$n = \log_2 n$

T.C. $= O(y)$

$= O(\log x)$

$= O(\log(\log n))$

M-1

i → 2, 4, 16, 256, $(256)^2$ . . . . n

$2^1, 2^2, 2^4, 2^8, 2^{16}$ . . . . . n

$2^{2^0}, 2^{2^1}, 2^{2^2}, 2^{2^3}, 2^{2^4}$ . . . . $2^{2^{x-1}}$

total 'x' terms hai

→ $n = 2^{2^{x-1}}$

$2^{x-1} = \log_2(n)$

$x-1 = \log_2(\log_2 n)$

$x = \log_2(\log_2 n) + 1$

T.C. = $O(x)$

= $O(\log\log n)$

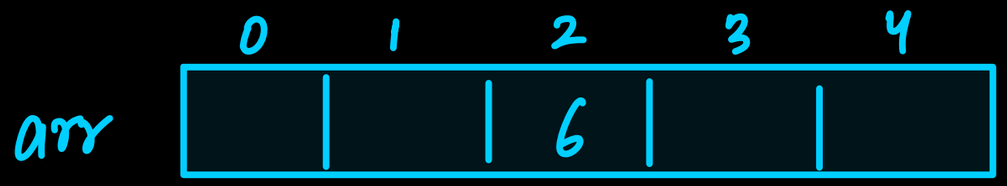
M-2

Contents

Representation of 2D Array / Matrix

```
int[] arr = new int[5];
```

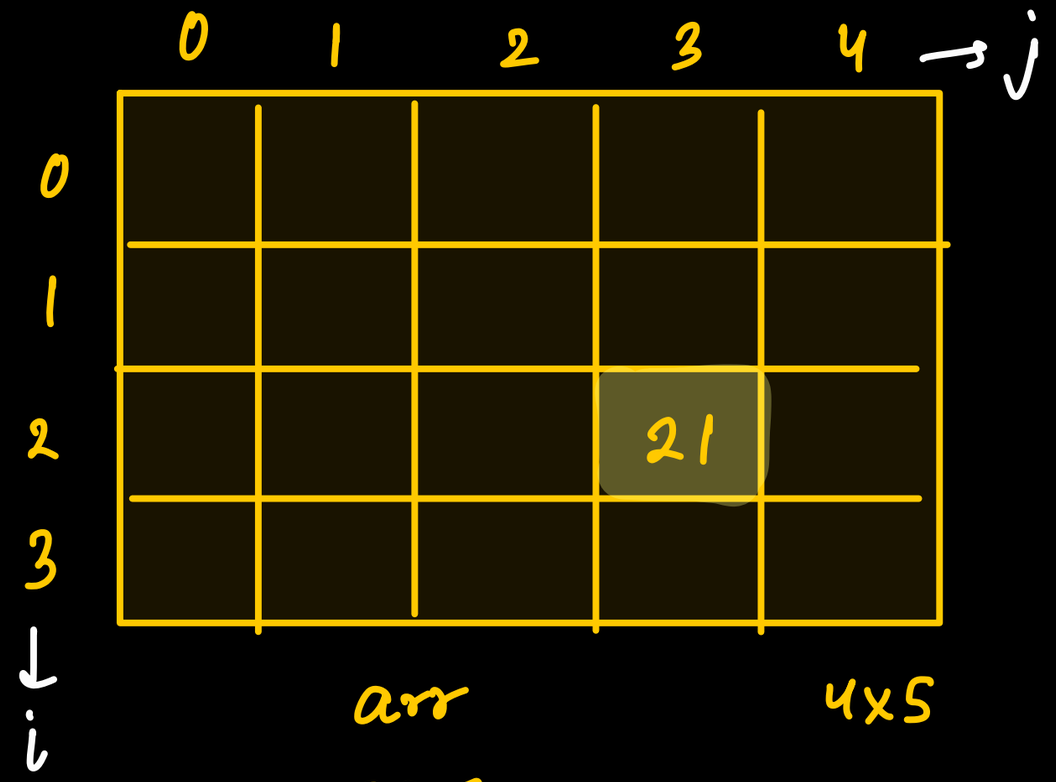
length



arr[2] = 6

```
int[][] arr = new int[4][5];
```

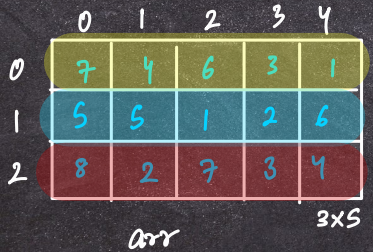
rows cols



arr[2][3] = 21

Input & Output in 2D Arrays

Array of Arrays



arr.length = 3 (no. of rows)

arr[0].length = 5 (no. of cols)

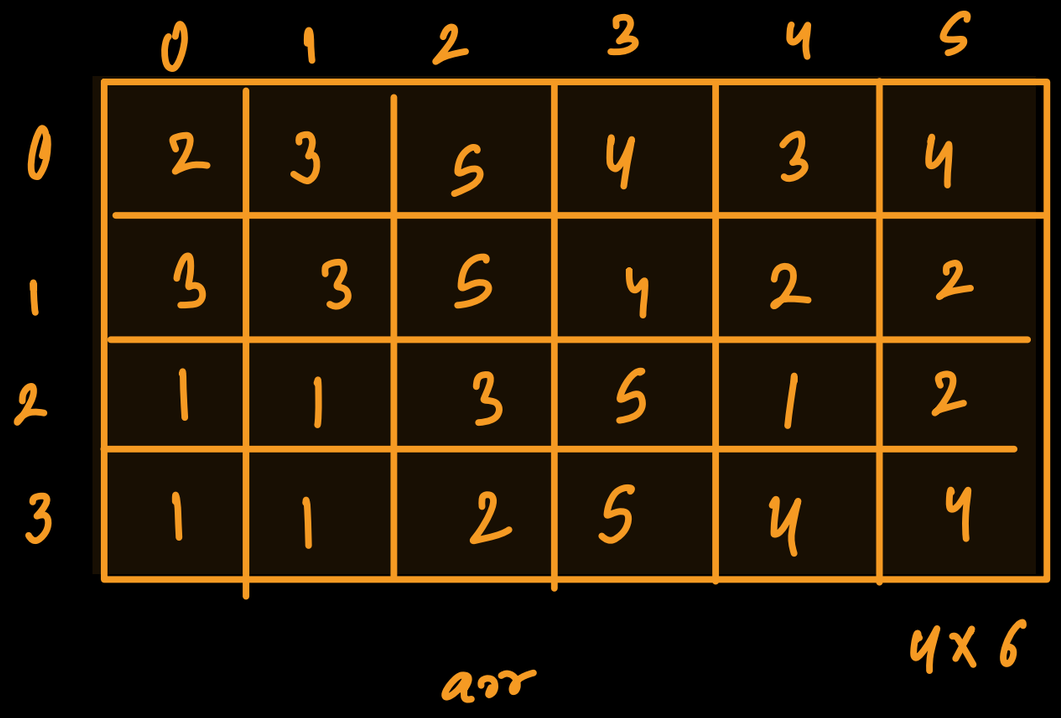
arr = { {7, 4, 6, 3, 1}, {5, 5, 1, 2, 6}, {8, 2, 7, 3, 4} }

arr[0] arr[1] arr[2]

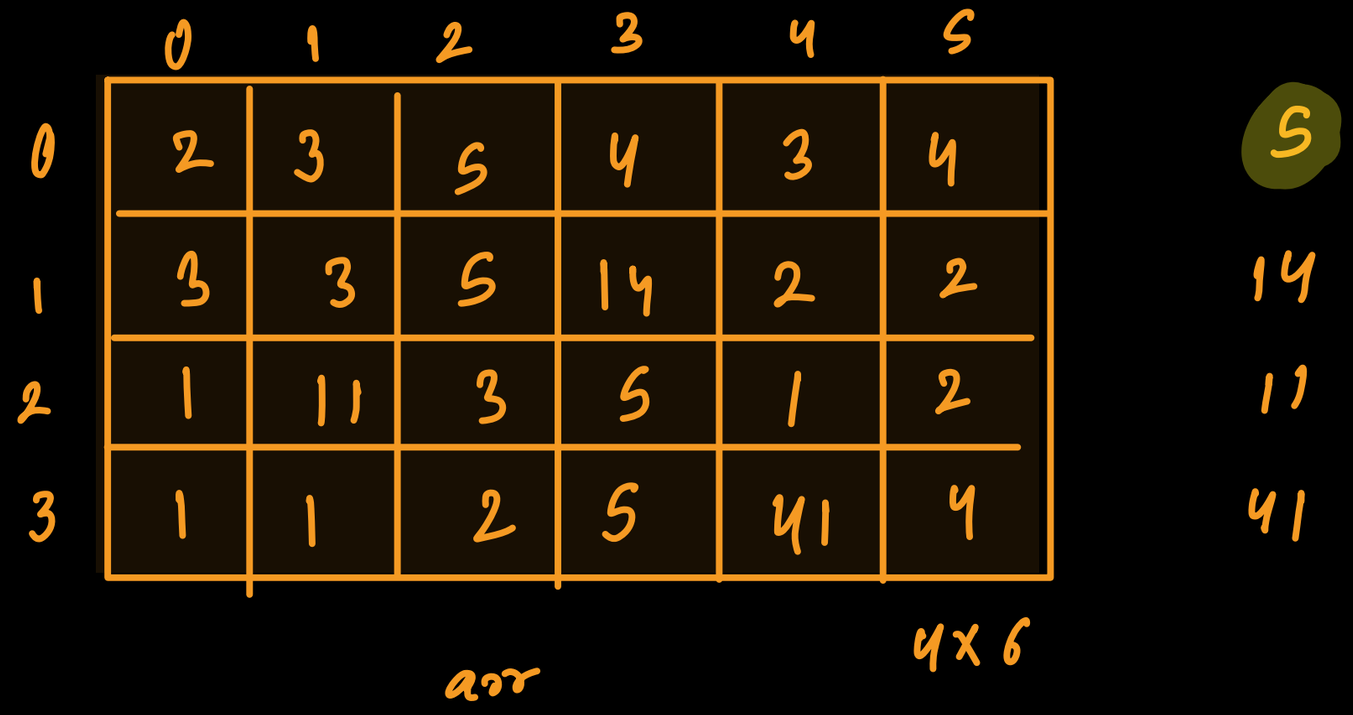
Ques: Sum of elements in given Matrix

HW: Find the maximum element in given 2D array.

Ques: Find the row with maximum sum



HW: Find the minimum element out of all the maximum elements of each row



S
14
11
41

A.S. = O(1)

Ques: Print elements of 2D Array column-wise (printing transpose)

	0	1	2	3	4	5 → j
0	2	3	5	4	3	4
1	3	3	5	14	2	2
2	1	11	3	5	1	2
3	1	1	2	5	41	4

i ↓

arr

4 x 6

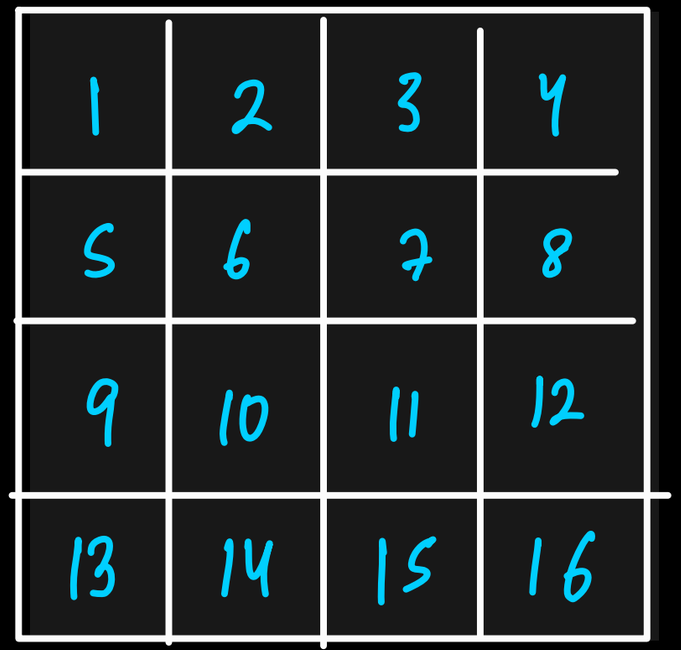
Output:

	0	1	2	3 → i
0	2	3	1	1
1	3	3	11	1
2	5	5	3	2
3	4	14	5	5
4	3	2	1	41
5	4	2	2	4

j ↓

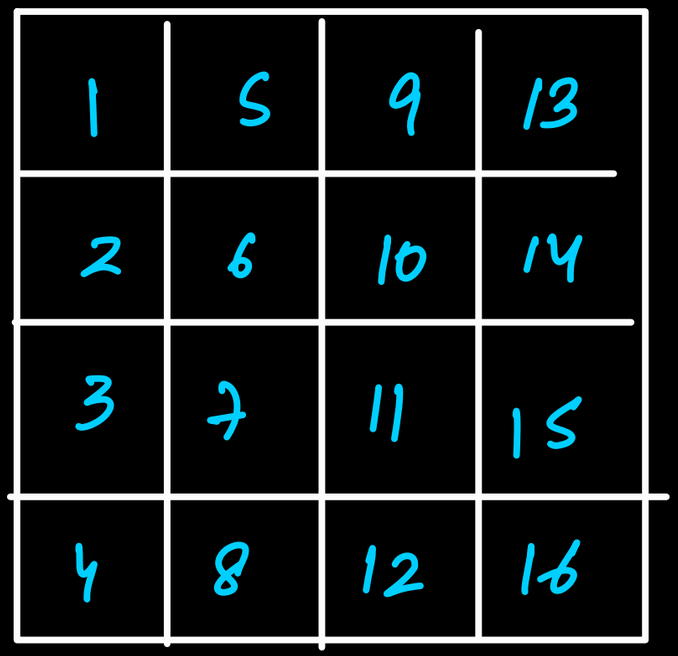
Ques: Print matrix in Snake Pattern

Ques: Transpose of Matrix



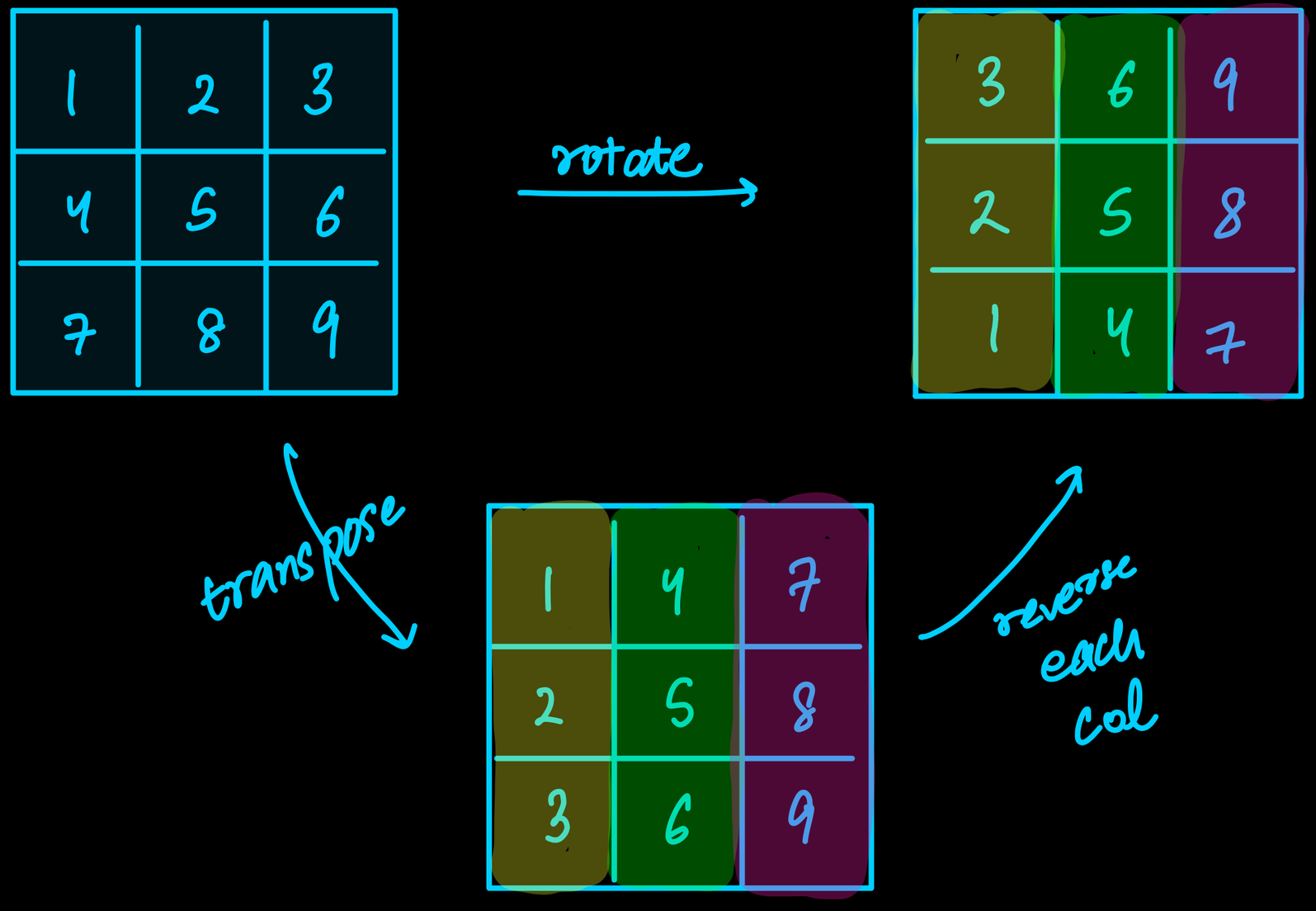
original

transformed

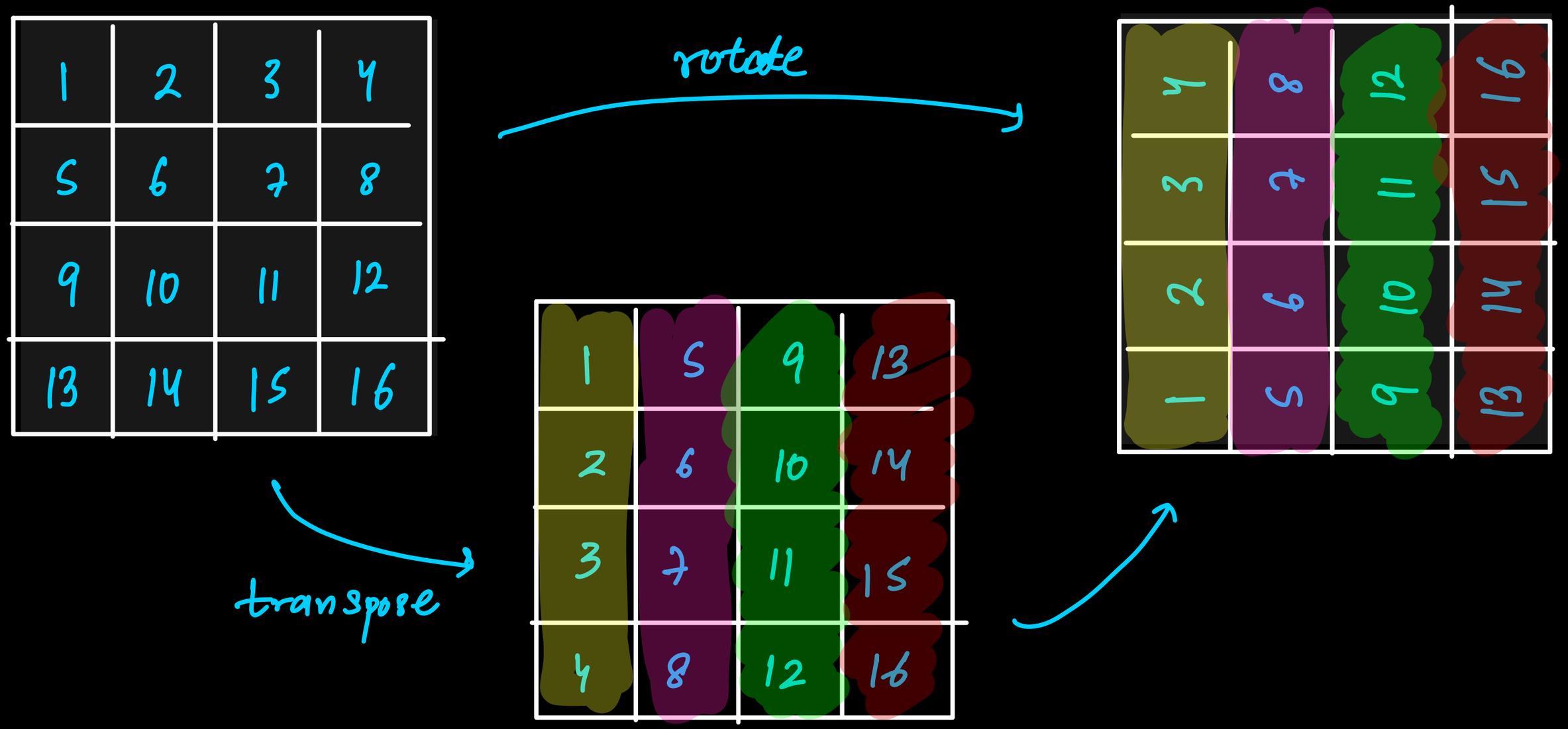


end result

Ques: Rotate by 90 degree



Ques: Rotate by 90 degree



2D ArrayLists

arr = { {1, 2, 3}, {1, 2}, {1, 2, 3, 4, 5, 6} }

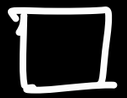
1D Arraylist : ArrayList<Integer> al = new Arraylist<>()

2D Arraylist : Arraylist < Arraylist <Integer>> x = new Arraylist<>();

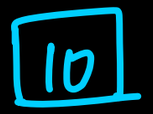
x = { }

How ArrayLists Work Internally..

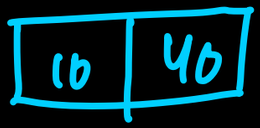
By Default → 1 capacity ka array banta hai



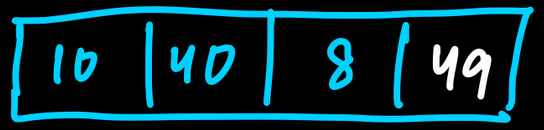
arr.add(10)



arr.add(40)

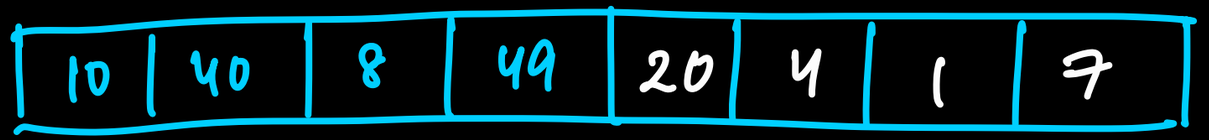


arr.add(8)



arr.add(49)

arr.add(20)

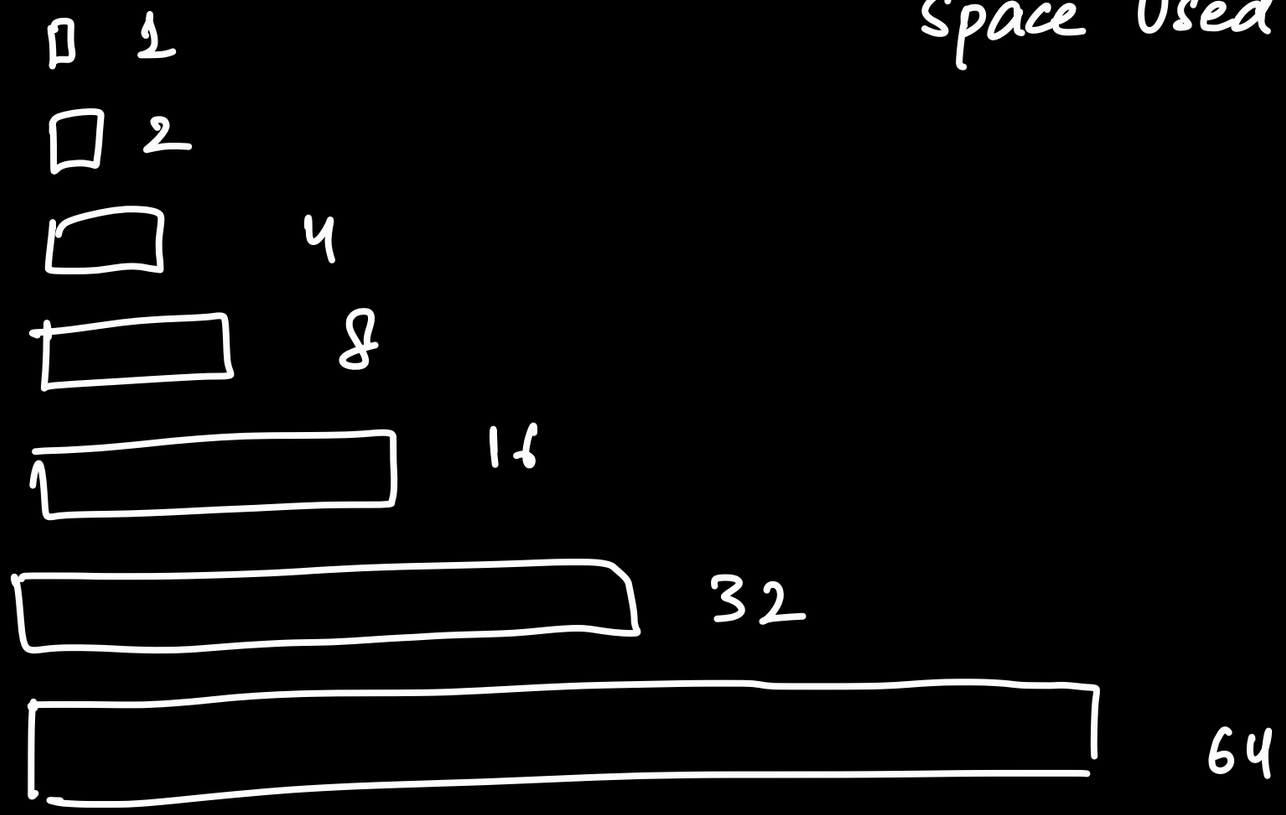


arr.add(4)

arr.add(1)

arr.add(7)

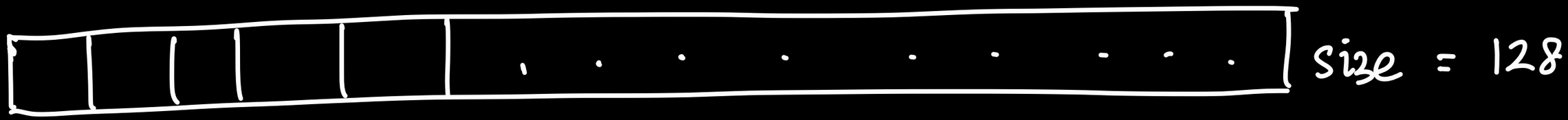
2D ArrayLists



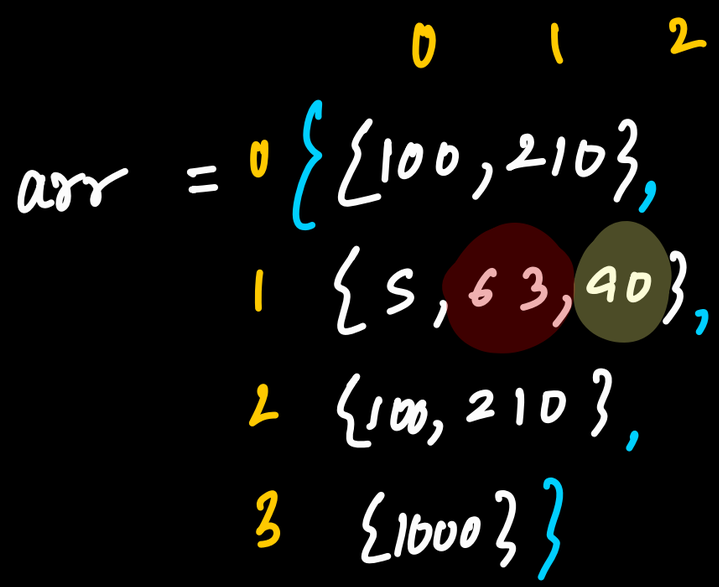
Space Used = $1+2+4+8+16 \cdots n$

$= 2n-1$

$\approx O(n)$



2D ArrayLists



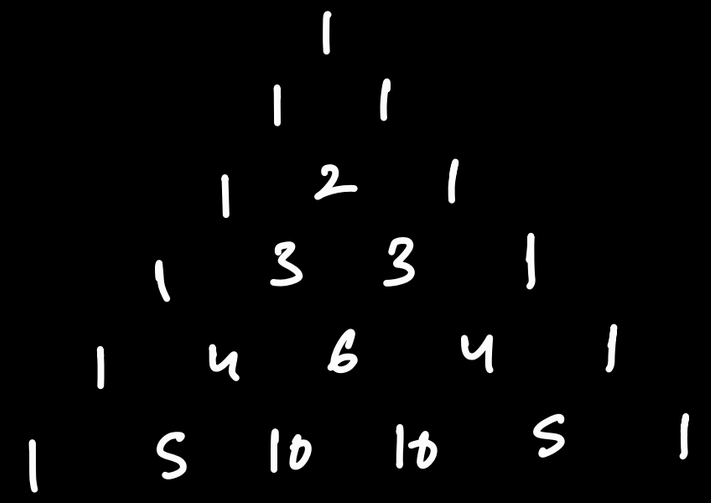
change 90 to 89

2D array me → arr[1][2] = 89

2D AL me → arr.get(1).set(2,89)

arr.get.(i).get(j) → for access

Ques: Pascal Triangle



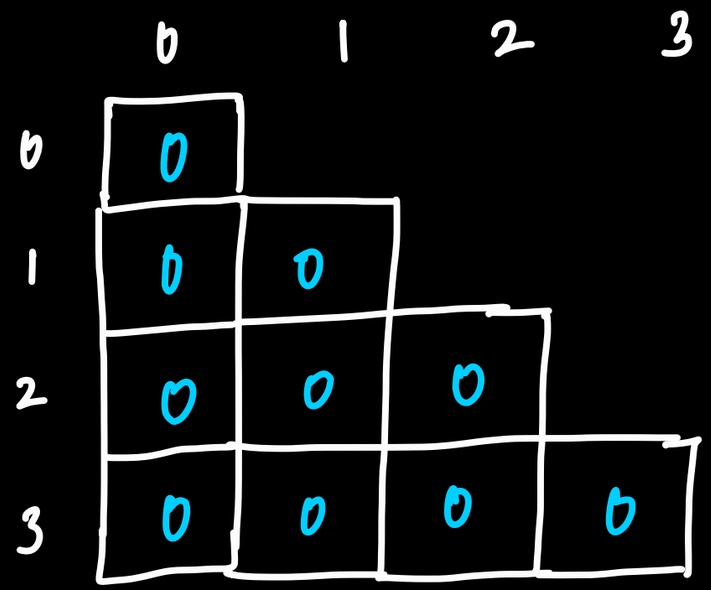
→

	0	1	2	3	4	5
0	1					
1	1	1				
2	1	2	1			
3	1	3	3	1		
4	1	4	6	4	1	
5	1	5	10	10	5	1

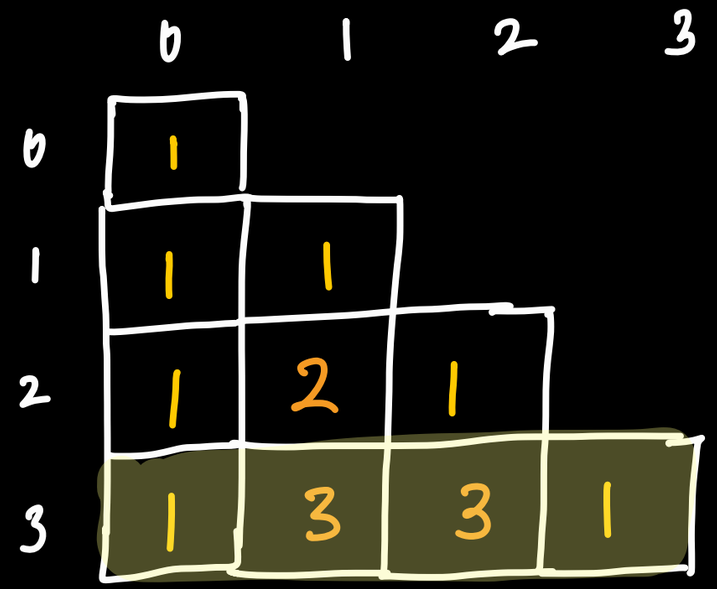
$^{i}C_{j}$

arr = { {1} , {1, 1} , {1, 2, 1} , . . . }

Ques: Pascal Triangle



Step 1



Step 2

n = 4

T.C. = $O(n^2)$

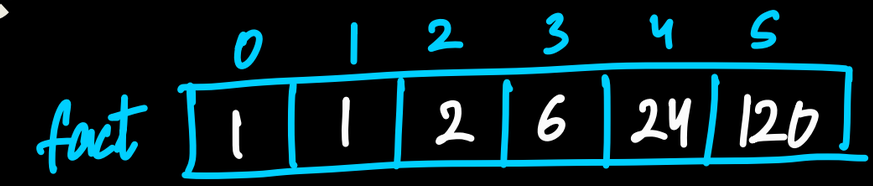
A.S. = $O(n^2)$

```
if(j==0 or j==i)
    arr[i][j] = 1
else
    arr[i][j] = arr[i-1][j] + arr[i-1][j-1]
```

Ques: Pascal Triangle

n = 6 → 1, 5, 10, 10, 5, 1

$^5C_0$ $^5C_1$ $^5C_2$ $^5C_3$ $^5C_4$ $^5C_5$




```
n--;
for(int r = 0; r <= n ; r++){
|    ans.add( nCr );
3
```

^nC_r ko fast kaise nikale

T.C. = $O(n^2)$

A.S. = $O(n)$

$$^nC_{r+1} = \frac{n-r}{r-1} \cdot {}^nC_r$$

Ques: Search in a Row-Column Sorted matrix

1	4	7	11	15
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

target = 5

Ques: Search in a Row-Column Sorted matrix

1	4	7	11	15
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

target = 18

if (arr[i][j] < target) go down

if (arr[i][j] > target) go left

Ques: Search in a Row-Column Sorted matrix

1	4	7	11	15
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

target = 20

Ques: Search in a Row-Column Sorted matrix

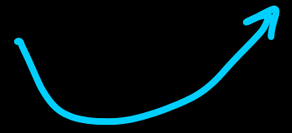
1	4	8	11	15
2	5	8	12	19
3	6	9	16	22
10	13	14	17	24
18	21	23	26	30

target = 7

while (i < m && j >= 0)

Ques: Set matrix zeroes

	0	1	2	3	4	5
0	3	6	2	1	7	3
1	8	0	-2	3	-5	-9
2	4	5	6	4	0	3
3	1	2	0	2	4	9

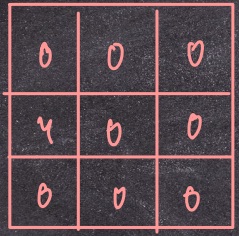


	0	1	2	3	4	5
0	3	0	0	1	0	3
1	0	0	0	0	0	0
2	0	0	0	0	0	0
3	0	0	0	0	0	0

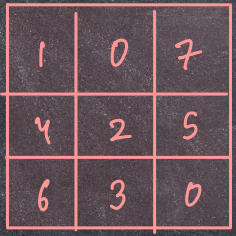
1st row, 1st col

2nd row, 4th col

3rd row, 2nd col



arr



copy

Approach 1 → make a same to same copy.
traverse in original array, make changes in duplicate array

↓

T.C. = $O(m * n * (m+n))$

A.S. = $O(m * n)$

THANKYOU

Cuties