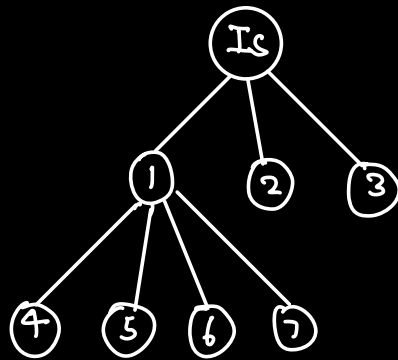


# THE SEARCH PROCESS

## TERMINOLOGY:

- Generate Children
- Expand
  - (1) Check goal
  - (2) Generate Children
- Fringe/Frontier : The unexplored nodes



Fringe: {2, 3, 4, 5, 6, 7}

Expand a fringe, say 2

↳ Check goal → not the FS

↳ Generate Children {8, 9}.

Idea: 1. Choose a node in fringe

2. Expand it

3. Repeat.

We check the goal during  
expansion not Generation!

## SEARCH STRATEGIES:

A search strategy is a criteria for deciding which node on the fringe to expand next.

## EVALUATING SEARCH STRATEGIES:

### 1. COMPLETENESS:

Does it always find a solution if one exists.

### 2. TIME COMPLEXITY:

Number of nodes generated.

### 3. SPACE COMPLEXITY:

Maximum number of nodes in memory.

### 4. OPTIMALITY:

Does it find a least-cost solution.

## PARAMETERS FOR EXPRESSING COMPLEXITY:

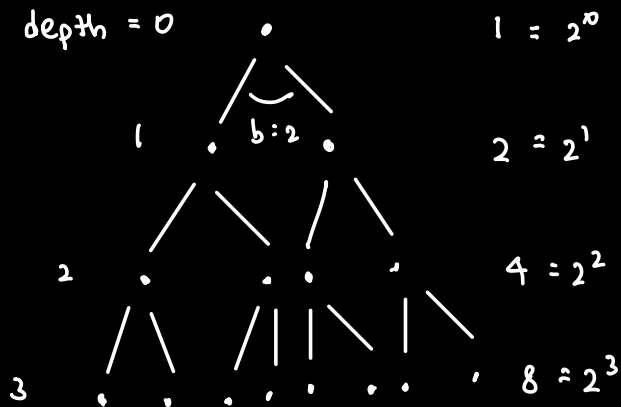
1.  $b$ : Branching factor

2.  $d$ : Depth of least-cost solution.

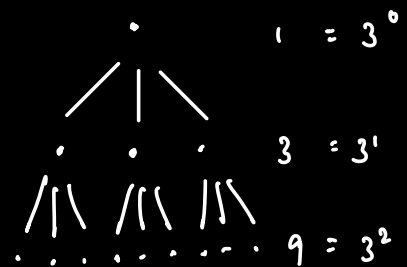
3.  $m$ : maximum depth of the search tree  
(could be infinite)

## OBSERVATION ABOUT TREES:

1. depth = 0



$b = 3$



$$N(b, d) = 1 + b + \dots + b^d$$

$$bN(b, d) = b + b^2 + \dots + b^{d+1}$$

$$(b-1)N(b, d) = b^{d+1} - 1$$

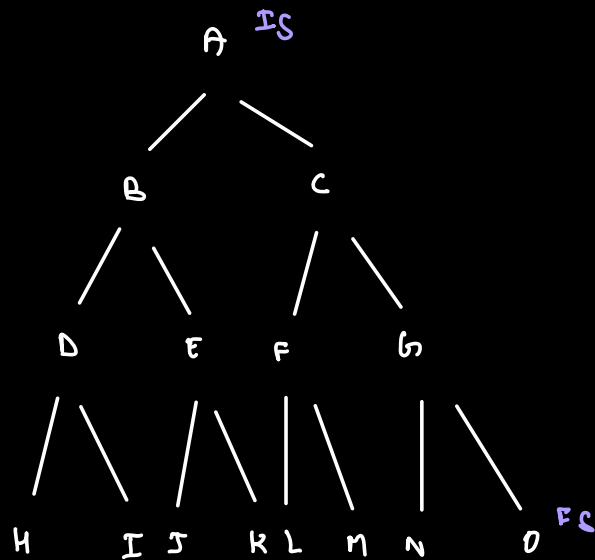
$$N(b, d) = \frac{b^{d+1} - 1}{b - 1} \approx \frac{b^{d+1}}{b - 1} = \left(\frac{b}{b-1}\right) b^d$$

$$= O(b^d) //$$

2. Last layer has more nodes than all other layers combined.

(BFS) BREADTH FIRST SEARCH:

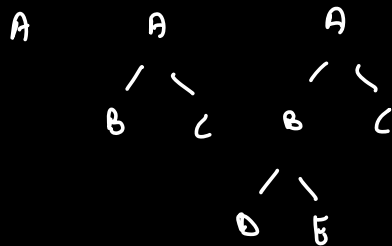
"Shallowest node first"



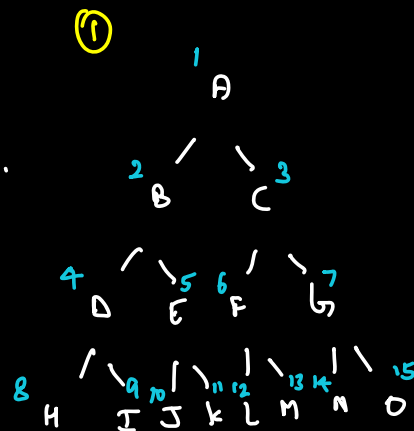
1. Number nodes according to expand

2. Implementation note

3. Evaluation Criteria

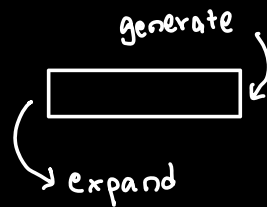


...



## ② Implementation note:

Use Queue



~~A~~

~~B~~

~~C~~

DEFB

## ③ EVALUATION CRITERIA:

1. Complete:

Expands all nodes at depth  $i$  before any node at  $i+1$ .

2. Optimality:

We expand lower levels before going to higher-levels, it is optimal.

3. Space Complexity;

$O(b^d)$ . Need full tree to check answer.

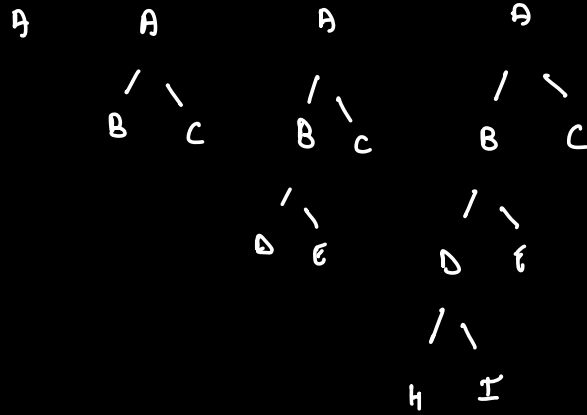
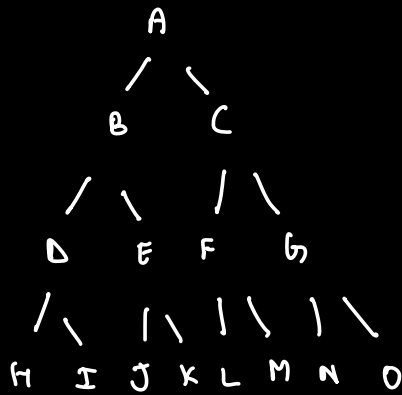
4. Time Complexity

$O(b^d)$ . Sum of nodes.

$$b^0 + b^1 + \dots + b^d.$$

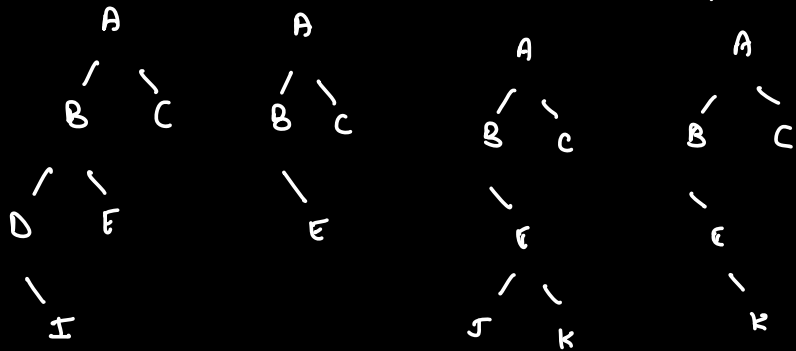
(DFS) DEPTH - FIRST SEARCH :

"deepest node first"

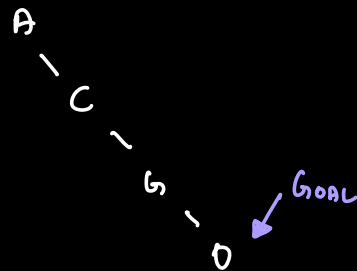


expand, not an answer

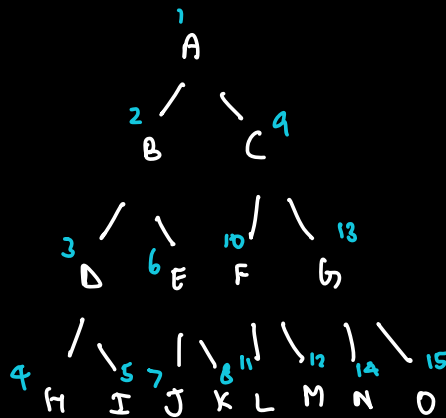
(no children)



...



1. Order of expansion:



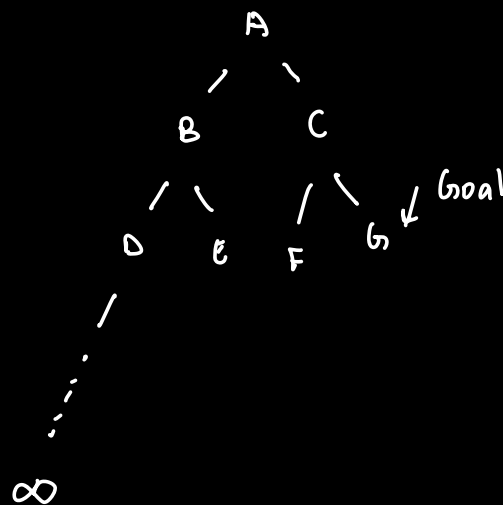
2. Implementation note:

Use a stack

3. Evaluation Metrics

1. Complete:

No, we might just infinitely expand (infinite tree).



2. OPTIMAL:

No

3. SPACE COMPLEXITY:

$$O(bm)$$

4. TIME COMPLEXITY:

$$O(b^m)$$

### LIMITED DEPTH SEARCH:

Limit depth to  $l$

1. Complete :

Yes if  $d \leq l$

2. Optimal :

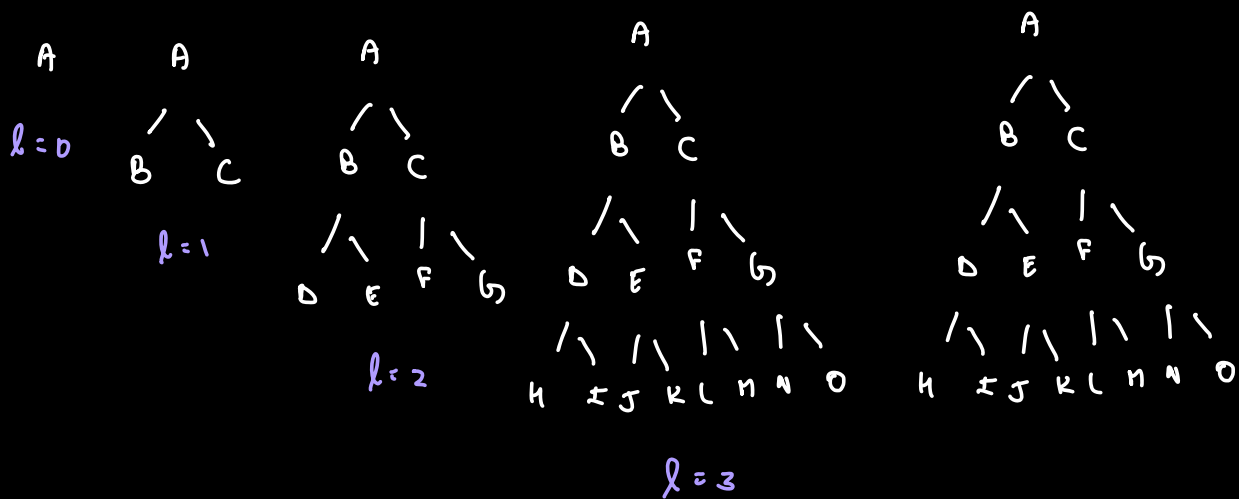
No

3. Time Complexity :  $O(b^l)$

4. Space Complexity :  $O(bl)$



## ITERATIVE DEPTH SEARCH:



1. Complete :

Yes

2. Optimal :

Yes

3. Time Complexity :  $O(b^d)$

4. Space Complexity :  $O(bd)$

TIME COMPLEXITY OF ID:

$$l=0 \quad b^0 (b/b-1)$$

$$l=1 \quad b^1 (b/b-1)$$

$$l=2 \quad b^2 (b/b-1)$$

$\vdots$

$$l = d \quad b^d (b/b-1)$$

$$Sum = (b/b-1) (b/b-1) b^d$$

$$= (b/b-1)^2 b^d$$

$$\text{For } b=10, d=5$$

$$\text{Tree: } b, d$$

$$b^d (b/b-1)$$

$$\text{ID: } b^d (b/b-1)^2$$

$$\text{Ratio: } b/b-1$$

$$b=2 \Rightarrow 2$$

$$b=3 \Rightarrow 1.5$$

$$\begin{array}{l} \rightarrow 111, 11 \\ \rightarrow 123, 456 \end{array} \left. \vphantom{\begin{array}{l} \rightarrow 111, 11 \\ \rightarrow 123, 456 \end{array}} \right\} \text{very close.}$$