

Recap:

INPUT:  $x \in \mathbb{R}^{n \times d}$

( $n > d$ )

OUTPUT:  $v$ , that  $\max_{\|v\|=1} \|xv\|$

POWER ITERATION

-  $v_0$  is a random unit vector

- for  $t=1, \dots, T$ :

$$- u = x^T \cdot (x \cdot v_{t-1})$$

$$- v_t = \frac{u}{\|u\|}$$

THEOREM:

$P_I$  will converge to a top right S.V.

After  $T = O\left(\frac{\log(d/\epsilon)}{\epsilon}\right)$  iterations,  $\|x \cdot v_T\| \geq (1-\epsilon)\sigma_1$ .

$X = U \Sigma V^T$  be its SVD

$$\sigma_1 \geq \sigma_2 \geq \sigma_3 \geq \dots \geq \sigma_d$$

$$\begin{array}{c} \sigma_1 \\ \left| \right. \\ \left| \right. \sigma_2 \\ \left| \right. \sigma_3 \\ \left| \right. \dots \\ \left| \right. \sigma_k > (1-\epsilon)\sigma_1 \\ \left| \right. \sigma_{k+1} \leq (1-\epsilon)\sigma_1 \end{array}$$

Let  $k$  be the index such that

$$\sigma_k > (1-\epsilon)\sigma_1$$

$$\sigma_{k+1} \leq (1-\epsilon)\sigma_1$$

Let  $S_k = \text{span of the first } k \text{ singular vectors}$   
 $v^{(1)}, v^{(2)}, \dots, v^{(k)}$

THEOREM:

After  $T = O\left(\frac{\log(d/\delta)}{\varepsilon}\right)$  iterations

$v_T$  is almost entirely within  $S_k$

$$\|v_T - \text{Proj}_{S_k}(v_T)\| \leq \delta$$

$\varepsilon$  such that:

$$\sigma_k > (1-\varepsilon)\sigma_1$$

$$\sigma_{k+1} \leq (1-\varepsilon)\sigma_1$$

THEOREM:

After  $T$  iterations

$$\|v_T - \text{Proj}_{S_k}(v_T)\| \leq \frac{(1-\varepsilon)^{2T}}{|\langle v_0, v^{(1)} \rangle|}$$

↳ first right singular vector

Remark:

If  $\sigma_2 \leq \sigma_1/2$

then  $\|v_T - \text{Proj}_{S_1}(v_T)\| \leq \frac{\left(\frac{1}{2}\right)^{2T}}{|\langle v_0, v^{(1)} \rangle|}$

COROLLARY:

Suppose  $\sigma_2 < (1-\epsilon)\sigma_1$ , then

$$\|v_T - \text{Proj}_{S_1}(v_T)\| \leq \frac{(1-\epsilon)^{2T}}{|\langle v_0, v^{(1)} \rangle|}$$

PROOF:

$$v_1 = \frac{x^T \cdot (x \cdot v_0)}{\|x^T \cdot (x \cdot v_0)\|}$$

$$v_2 = \frac{[x^T \cdot (x \cdot v_1)]}{\|x^T \cdot (x \cdot v_1)\|}$$

$\vdots$

$$v_t = \frac{(x^T \cdot x)^t \cdot v_0}{\|(x^T \cdot x)^t \cdot v_0\|}$$

$$X = \sum_{i=1}^d \sigma_i u^{(i)} v^{(i)T}$$

$$X^T X = \sum_{i=1}^d \sigma_i^2 \cdot v^{(i)} \cdot v^{(i)T}$$

$$(X^T X)^t = \sum_{i=1}^d \sigma_i^{2t} \cdot v^{(i)} \cdot v^{(i)T}$$

$$v_t = \frac{(x^T x)^t \cdot v_0}{\|(x^T x)^t \cdot v_0\|}$$

$v^1$	$\dots$	$v^d$	$v^{d+1}$	$\dots$
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Form an orthonormal basis

$$v^1, v^2, \dots, v^d, v^{d+1}, \dots, v^n$$

$$\rightarrow v_0 = c_1 v^1 + c_2 v^2 + \dots + c_d v^d + c_{d+1} v^{d+1} + \dots + c_n v^n$$

$$(x^T x)^t \cdot v_0 = \left( \sum_{i=1}^d \sigma_i^{2t} \cdot v^{(i)} v^{(i)T} \right) \cdot (v_0)$$

$$= \sum_{i=1}^d \sigma_i^{2t} \cdot \langle v_0, v^{(i)} \rangle \cdot v^{(i)}$$

$$\text{Proj}_{\{v^{(1)}\}}((x^T x)^t \cdot v_0) = \langle (x^T x)^t \cdot v_0, v^{(1)} \rangle$$

$$\{v^{(1)}\} = \sigma_1^{2t} \cdot \langle v_0, v^{(1)} \rangle \cdot v^{(1)}$$

$$\|(x^T x)^t \cdot v_0\|^2 = \sum_{i=1}^d \sigma_i^{4t} \cdot \langle v_0, v^{(i)} \rangle^2$$

$$\|v_t - \text{Proj}_{\{v^{(1)}\}}(v_t)\|^2 = \frac{\left\| \sum_{i=2}^d \sigma_i^{2t} \cdot \langle v_0, v^{(i)} \rangle \cdot v^{(i)} \right\|^2}{\|(x^T x)^t \cdot v_0\|^2}$$

$$= \frac{\sum_{i=2}^d \sigma_i^{4t} \cdot \langle v_0, v^{(i)} \rangle}{\sum_{i=1}^d \sigma_i^{4t} \cdot \langle v_0, v^{(i)} \rangle^2}$$

$$\text{Numerator} \leq \sum_{i=2}^d (\sigma_i (1-\epsilon))^{4t} \cdot \langle v_0, v^{(i)} \rangle^2$$

$$\text{Recall: } \sigma_2 \leq (1-\epsilon) \sigma_1$$

$$\leq \sigma_1^{4t} (1-\epsilon)^{4t} \cdot \left( \underbrace{\sum_{i=2}^d \langle v_0, v^{(i)} \rangle^2}_{\text{project a vector to}} \right)$$

project a vector to

orthonormal basis and

take their sum of

squares  $\rightarrow$  length  $\rightarrow$  unit

$$\leq \sigma_1^{4t} (1-\epsilon)^{4t}$$

$$\text{Denominator} \geq \sigma_1^{4t} \cdot \langle v_0, v^{(1)} \rangle^2$$

$$\begin{aligned} \|v_t - \text{Proj}_{\{v^{(1)}\}}(v_t)\|^2 &\leq \frac{\sigma_1^{4t} \cdot (1-\epsilon)^{4t}}{\sigma_1^{4t} \langle v_0, v^{(1)} \rangle^2} \\ &= \frac{(1-\epsilon)^{4t}}{\langle v_0, v^{(1)} \rangle^2} \end{aligned}$$

$\Rightarrow$  Corollary!

PRACTICE:

$$1. \ x_1, x_2, x_3 : x = x_1 \cdot x_2 \cdot x_3$$

All we need is ability to compute matrix-vector products for  $x$ .

Example: To compute SVD of  $x_1, x_2, x_3$

$$(x_1 \cdot (x_2 \cdot (x_3 \cdot v)))$$

LAPACK have in built support for such computations.

2. How to compute higher singular vectors?

a. (Recall): If  $x = v \Sigma v^T$

$$\sigma_1 = \|x \cdot v^{(1)}\|$$

→ first right sv

$$u^{(1)} = \frac{x \cdot v^{(1)}}{\sigma_1}$$

$$x = \sigma_1 u^{(1)} v^{(1)T} + \sigma_2 u^{(2)} v^{(2)T} + \dots + \sigma_d u^{(d)} v^{(d)T}$$

To compute 2<sup>nd</sup>:

compute 1<sup>st</sup> of  $(x - \sigma_1 u^{(1)} v^{(1)T})$

⋮

so on.

b. GENERAL PT:

→ Pick  $k$  orthonormal vectors  $y^{(1)}, y^{(2)}, \dots, y^{(k)}$ .

→ For  $t=1, \dots, T$ :

\* Let  $z^{(1)} = x^T \cdot x \cdot y_{t-1}^{(1)}$ ,  $z^{(2)} = x^T \cdot x \cdot y_{t-1}^{(2)}$ , ...

$$z^{(k)} = x^T \cdot x \cdot y_{t-1}^{(k)}$$



\* Let  $y_t^{(1)}, y_t^{(2)}, \dots, y_t^{(k)}$  be an orthonormal basis for  $\text{span}(\{z^{(1)}, z^{(2)}, \dots, z^{(k)}\})$ .

3. Let's say  $x$  had singular values

$$1, 0.99, 0.98, 0.97, 0.97, \dots, 0.97$$

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convergence will be very slow.

→ Suppose  $x = V \Sigma V^T$

\* What are singular values of  $x - \frac{1}{2}I$ :

Has singular values

$$0.5, 0.99, 0.98, 0.97, \dots, 0.97$$

\*  $x - 0.96I$  has singular values

$$0.04, 0.03, 0.02, 0.01, \dots, 0.01.$$

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multiplicative gap decreases.

4. PI gets convergence at the rate  $1/\epsilon$ .

PI with momentum gets convergence at the rate  $1/\sqrt{\epsilon}$ .

POWER ITERATION + "MOMENTUM":

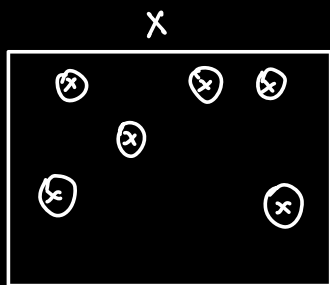
-  $v_0$  is a random unit vector

- for  $t=1, \dots, T$ :

$$- u = X^T \cdot (X \cdot v_{t-1}) - \beta \cdot v_{t-2}$$

$$- v_t = \frac{u}{\|u\|}$$

MATRIX COMPLETION / NETFLIX CHALLENGE:



→ We only see

$x_{ij}$  for  $(i,j) \in$  some set of  
observed entries

$$O \subseteq [n] \times [d]$$

INPUT: Entries of  $x$  in  $O, k$

OUTPUT: Find  $\tilde{x}$  of rank  $k$  to

$$\min \sum_{(i,j) \in O} (x_{ij} - \tilde{x}_{ij})^2$$

$$L(\tilde{x}) = \sum_{i,j \in O} (x_{ij} - \tilde{x}_{ij})^2 \rightarrow \text{convex}$$

$$\begin{array}{l} \min L(\tilde{x}) \\ \text{Rank}(\tilde{x}) \leq k \end{array}$$

→ Non-convex Optimization  
problem  
(NP-Hard)

$$\begin{array}{l} C = \{y : \text{Rank}(y) \leq k\} \\ \min L(\tilde{x}) \\ \tilde{x} \in C \end{array}$$

SINGULAR VALUE PROJECTION : (JMS 09)

→  $x_0$

→ For  $t=1, \dots, T$ :

$$y_t = x_{t-1} - \eta \nabla L(x_{t-1})$$

$$x_t = (\text{Take top-}k\text{-svo of } y_t) \rightarrow \text{Proj}_C(y_t)$$

↓

The closest  
rank  $k$  matrix.