

Lecture 2.5

Maxmin and Minmax Values coincide in
Matrix Games

Theorem: $A \in \mathbb{R}^{m \times n}$

$$\max_{x \in \Delta([m])} \min_{y \in \Delta([n])} x^T A y = \min_{y \in \Delta([n])} \max_{x \in \Delta([m])} x^T A y$$

Proof: We have already shown before,

$$\max_{x \in \Delta([m])} \min_{y \in \Delta([n])} x^T A y \leq \min_{y \in \Delta([n])} \max_{x \in \Delta([m])} x^T A y$$

To show: $\max_{x \in \Delta([m])} \min_{y \in \Delta([n])} x^A y \geq \min_{y \in \Delta([n])} \max_{x \in \Delta([m])} x^A y$. ✓

Let x^* be an optimal solution of LP1

y^* → LP2

$$\min_{y \in \Delta([n])} \max_{x \in \Delta([m])} x^A y \geq \min_{y \in \Delta([n])} x^* A y = \max_{x \in \Delta([m])} x^* A y \geq \max_{x \in \Delta([m])} \min_{y \in \Delta([n])} x^A y$$

□

Theorem: $A \in \mathbb{R}^{m \times n}$. Then the utility of both the players in any MSNE is exactly their values.

Proof: (x^*, y^*) be any MSNE.
 $v_r = \max_{x \in \Delta([m])} \min_{y \in \Delta([n])} x^* A y$, $v_c = - \min_{y \in \Delta([n])} \max_{x \in \Delta([m])} x^* A y$
From min max theorem, we know that $v_r = -v_c$
Let (x^*, y^*) be an MSNE for the matrix game A .

$$x^* A y^* \leq \max_{x \in \Delta([m])} x^* A y^* = \min_{y \in \Delta([n])} x^* A y \leq \max_{x \in \Delta([m])} \min_{y \in \Delta([n])} x^* A y = v_r$$

$$x^* A y^* \leq v_r \quad , \quad x^* A y^* \geq v_r \Rightarrow x^* A y^* = v_r$$

$$-x^* A y^* \leq -\min_{y \in \Delta([n])} x^* A y = -\max_{x \in \Delta([m])} x^* A y^* \leq -\min_{y \in \Delta([n])} \max_{x \in \Delta([m])} x^* A y = v_c$$

$$-x^* A y^* \leq v_c \quad , \quad -x^* A y^* \geq v_c \Rightarrow -x^* A y^* = -v_c.$$

	A	B
$x_A \rightarrow A$	1, -1	-1, 1
$x_B \rightarrow B$	-1, -1	1, -1

Security level of row player.

$$\begin{aligned} \max \\ x_A, x_B : \\ x_A + x_B = 1, \\ x_A, x_B \geq 0 \end{aligned}$$

$$\min \{x_A - x_B, x_B - x_A\}$$

Solution $x_A = \frac{1}{2}, x_B = \frac{1}{2}$

$(A: \frac{1}{2}, B: \frac{1}{2})$ value guaranteeing strategy for the row

player.
Similarly, $(A: \frac{1}{2}, B: \frac{1}{2})$ player.

column

$\left\{ \left(A : \frac{1}{2}, B : \frac{1}{2} \right), \left(A : \frac{1}{2}, B : \frac{1}{2} \right) \right\}$ is an MSNE.

Corollary: $A \in \mathbb{R}^{m \times n}$. $x^* \in \Delta([m])$, $y^* \in \Delta([n])$. Then,
 (x^*, y^*) is an MSNE if and only if x^* and y^* are
value guaranteeing strategies for the row and column
player respectively.

Q: Give an example of a matrix game with 10 MSNEs.

	Rock	Paper	Scissor
Rock	0, 0	-1, 1	1, -1
Paper	1, -1	0, 0	-1, 1
Scissor	-1, 1	1, -1	0, 0

$$\begin{aligned}
 & \max_{x_R, x_P, x_S} \min \left\{ x_P - x_S, x_S - x_R, x_R - x_P \right\} \\
 & \boxed{x_R + x_P + x_S = 1} \\
 & \boxed{x_R, x_P, x_S \geq 0}
 \end{aligned}$$

maximize if
 $x_P = x_S = x_R = \frac{1}{3}$

for the rock-paper-

$$\left\{ \left(R: \frac{1}{3}, P: \frac{1}{3}, S: \frac{1}{3} \right), \left(R: \frac{1}{3}, P: \frac{1}{3}, S: \frac{1}{3} \right) \right\}$$

scissor game.

