COSC 4315 - Homework 3

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Requirements

- 1. No loops may be used. Recursion must be used instead.
- $2.\,$ Must support nested functions in input, up to a depth of $3.\,$
- 3. No global variables may be used.
- 4. Code must be organized into functions.
- $5.\,$ All parameters must be passed by value. Passing by reference is disallowed.
- 6. All functions must be commented with preconditions and postconditions.
- 7. All math must operate on lists representing "infinite integers."
- 8. Program must declare invalid or malformed input.

String to Infinite Integer

Strings are read into "infinite integer" lists by starting from the end of the string and reading n characters back recursively until fewer than n characters remain (where n represents the digits per node). Each returning step of the recursion then adds its n ending characters to the resulting list of the previous step until the full integer is built. This solution naturally "aligns" digits such that any non-full nodes are shifted to the most significant digits.

Pseudocode

Algorithm 1 Building an infinite integer recursively

```
1: procedure STRToINFINT(s: string; n: int)
       if len(s) \le n then
2:
3:
          return a list composed of an integer representation of s
4:
       else
          s_i = an integer representing the last n digits of s
5:
          remove the last n digits of s
6:
          t = StrToInfInt(s, n)
7:
8:
          append s_i to the end of t
          return t
9:
       end if
11: end procedure
```

```
def StrToInfInt(s, n):
    if(len(s) <= n):
        return [int(s)]
    else:
        si = int(s[-n:])
        t = StrToInfInt(s[:-n], n)
        t.append(si)
    return t</pre>
```

Infinite Integer to String

Pseudocode

```
Algorithm 2 Returns a string representation of the integer
```

```
procedure INFINTTOSTR(a:list;i,n:int)

2: if i = len(a) then
return ""

4: else if i = 0 then
result = InfIntToStr(a, i + 1, n)

6: result = a[i] as a trimmed string + result
else

8: result = InfIntToStr(a, i + 1, n)
result = a[i] as an n-digit string + result

10: end if
return result

12: end procedure
```

```
def InfIntToStr(s, i, n):
    if i == len(s):
        return ''''
    elif i == 0:
        result = str(s[i]) + InfIntToStr(s, i + 1, n)
    else:
        result = str(s[i]).zfill(n) + InfIntToStr(s, i + 1, n)
    return result
```

Infinite Integer Addition

Two "infinite integers" are given to be added together. Starting from the least significant digits, nodes are added together (similarly to how addition by hand is performed, but using multiple digits at a time). If the sum exceeds 10^n , we will carry a 1 in the next recursive step and reduce the sum for the current step by 10^n . The last elements in both lists are removed and addition continues towards the most significant digits. Each recursive step returns the full "infinite integer" list for the addition performed so far

Pseudocode

```
Algorithm 3 Recursively adding two infinite integers together
```

```
procedure ADDINFINT(a, b : list; c, n : int)
       if a and b are empty lists then
          if c \neq 0 then
3:
              return a list containing c
          else
              return nothing
6:
          end if
          sum = c + the least significant nodes of a and b
9:
          if sum >= 10^n then
              sum = sum - 10^n
              carry = 1
12:
          else
              carry = 0
15:
          end if
          remove the least significant nodes from a and b
          result = AddInfInt(a, b, carry, n)
          append sum to result
18:
          return result
       end if
21: end procedure
```

```
\mathbf{def}\ \mathrm{AddInfInt}\left(\left.a\,,\;\;b\,,\;\;c\,,\;\;n\right)\right)
      if a == b == []:
            if c != 0:
                  return [c]
            {f else} :
                   return
      if len(a) > 0:
           ai = a[-1]
      {f else} :
            a\,i\ =\ 0
      if len(b) > 0:
           bi = b[-1]
      else:
            bi = 0
      \mathbf{sum} = \mathbf{c} + \mathbf{a}\mathbf{i} + \mathbf{b}\mathbf{i}
      if sum >= 10 ** n:
            \mathbf{sum} \ -\!\!= \ 10 \ ** \ n
            carry = 1
      {f else} :
            carry = 0
      a = a[:-1]
      b = b[:-1]
      result = AddInfInt(a, b, carry, n) or []
      \texttt{result} \; +\!\!\!= \; [\mathbf{sum}]
      return result
```

Infinite Integer Multiplication

This operation is performed similarly to multiplication by hand. The algorithm multiplies groups of digits together and sums the result. Mathematically, this can be represented by the following:

$$\sum_{i=0}^{n} \sum_{j=0}^{m} 10^{d(i+j)} b_i a_j \tag{1}$$

(1) n: total nodes in b, m: total nodes in a, d: digits per node; indices move right to left

However, using list-represented integers introduces some challenges. Instead of multiplying the nodes by $10^{d(i+j)}$, it is best to create an infinite integer with (i+j) zero-value nodes to the right. It is also important to be mindful of carry-out values; while the equation above does not need to respect carry-out, it must be ensured that the algorithm respects the specific digits per node requirement. Therefore, each step of the recursion should produce an infinite integer of the form [carry][product] (i+j)*[0] to be added to each successive step.

Pseudocode

```
Algorithm 4 Recursively multiplying two infinite integers together
```

```
procedure MULTIPLYACROSS(a: list; b, i, n: int)

if i = len(a) then

return 0 represented as an infinite integer

4: else

product_{i+1} = \text{MultiplyAcross}(a, b, i + 1, n)
product_{i} = b * a[i]
if product_{i} >= 10^{n} then

8: carry = \lfloor product_{i}/10^{n} \rfloor
product_{i} = product_{i} \mod 10^{n}
end if
```

```
append len(a) - i - 1 zero-value nodes to product_i
           prepend carry to product_i if it exists
12:
           \mathsf{return}\ product_{i+1} + product_i
       end if
   end procedure
16:
   procedure MultiplyInfInt(a, b: list; i, n: int)
       if i = len(b) then
           return 0 represented as an infinite integer
20:
       else
           result_{i+1} = MultiplyInfInt(a, b, i + 1, n)
           result_i = MultiplyAcross(a, b[i], 0, n)
           append len(b) - i - 1 zero-value nodes to the end of result_i
24:
           return result_i + result_{i+1}
       end if
   end procedure
```

```
def MultiplyAcross(a, b, i, n):
    if i = len(a):
        return [0]
    else:
        pnext = MultiplyAcross(a, b, i + 1, n)
        product = b * a[i]
        carry = 0
        if product >= 10 ** n:
            carry = math.floor(product / 10 ** n)
            product = product \ \backslash \% \ (10 \ ** \ n)
        product = [product] + [0] * (len(a) - i - 1)
        if carry > 0:
            product = [carry] + product
    return AddInfInt(product, pnext, 0, n)
def MultiplyInfInt(a, b, i, n):
    if i = len(b):
        return [0]
    else:
        resultnext = MultiplyInfInt(a, b, i + 1, n)
        result = Multiply Across(a, b[i], 0, n)
        result = result + [0] * (len(b) - i - 1)
        return AddInfInt(result, resultnext, 0, n)
```

Expression Evaluation

The goal is to take a single line of input in the form of an expression, and produce an infinite integer result. As with prior algorithms, this will be a recursive approach.

All input provided to this algorithm will be in the form [function name] followed by either an integer constant or another [function name] (...). Function calls should be recursively evaluated and reduced to infinite integer constants, which can be operated upon. Regular expressions will be used to identify sections of the input string in the form [function name]([constant], [constant]). These sub-expressions will be solved and their results re-inserted into the input expression until the entire expression is solved.

Pseudocode

```
Algorithm 5 Recursively solving an expression
    procedure SolveLine(s, n : int)
       use regular expressions to test for the desired function call form
       if expression does not exist then ▷ The input is unsolvable and thus
    invalid
           return [-1]
       else
 5:
                                                          ▷ This is the base case
           if match = s then
              \mathbf{if}s contains "add" \mathbf{then}
                  x = left + right sides of expression
                  x = left * right sides of expression
10:
              end if
              return x
```

```
else
s_1 = \text{slice } s \text{ around regex match}
15:
x = \text{SolveLine}(s_1)
\text{replace } s_1 \text{ in } s \text{ with } x
\text{return SolveLine}(s)
\text{end if}
end if
20: \text{ end procedure}
```

```
def SolveLine(s, n):
                                                    find \ = \ re.search \left( \text{`((add | multiply)} \setminus (\setminus d+, \setminus d+\setminus)) \text{'}, \cup s \right)
\cup \cup \cup \cup if \cup not \cup find:
 \cup \cup \cup \cup \cup \cup \operatorname{return} \cup [-1] 
ورين else:
\_\_\_\_\_\_ if \_ find . group (1) \_==\_\_s :
si_{\scriptscriptstyle -}=_{\scriptscriptstyle -}s . split ( '( ')
\texttt{uumResults} = \texttt{re.match} (\texttt{``([0-9]+),([0-9]+)", umResults}) = \texttt{re.match} (\texttt{``([0-9]+)", umResults}) = \texttt{r
\texttt{StrToInfInt} \ (\ numResults.\ group\ (1)\ , \texttt{\_n})\ , \texttt{\_StrToInfInt} \ (\ numResults.\ group\ (2)\ , \texttt{\_n})\ , \texttt{\_StrToInfInt} \ (\ numResults.\ group\ (2)\ , \texttt{\_n})\ , \texttt{\_Notation} \ )
 = \underbrace{\quad \text{if } si \ [0]}_{\text{constraints}} \text{ ``add'''} 
\begin{array}{ll} \text{Localize} & \text{AddInfInt} \left( \text{nums} \left[ 0 \right], \text{Lnums} \left[ 1 \right], \text{Localize} \\ \text{Localize} & \text{elif} \text{Lsi} \left[ 0 \right] \\ \text{Localize} & \text{constant} \end{array} \right), \\ \text{Localize} & \text{Localize} & \text{Localize} \\ \text{Localize} & \text{Localize} \\ \text{Localize} & \text{Localize} \\ \text{Localize} & \text{Localize} & \text{Localize} \\ \text{Localize} \\ \text{Localize} & \text{Localize} \\ \text{Localize
 \texttt{uultiplyInfInt} \ (\texttt{nums} \ [0] \ , \texttt{unums} \ [1] \ , \texttt{u} \ 0 \ , \texttt{un})
  \verb"coultreturn" \verb"cuturn" \"cuturn" \"cuturn
 ____else:
   s1 = s[find.start(1):find.end(1)]
 \verb"cutous x = SolveLine" (s1, n)
```

Processing Input

Psuedocode

```
Algorithm 6 Recursively solving all lines of input
   procedure SolveInput(s, n)
       if len(s) = 0 then
          return
       else
          result = SolveLine(s[0], n)
 6:
          if result \neq [-1] then
              output s[0] and result
          else
              output "invalid expression"
          end if
          remove s[0]
          SolveInput(s, n)
12:
       end if
   end procedure
```

```
def SolveInput(s, n):
    if len(s) == 0:
        return
    result = SolveLine(s[0], n)
    if result != [-1]:
        print(s[0], ''='', InfIntToStr(result, 0, n))
    else:
        print(''Invalid expression: '', s[0])
    SolveInput(s[1:], n)
```