

Chapter 1

Evolog Language Specification

1.1 Syntax

Every valid ASP-Core2 program is a valid Evolog program. In addition, Evolog programs may contain *action rules* and *module literals*.

Definition 1.1.1 (Action Rule). Action Rules are ASP rules that have a body as defined by the ASP-Core2 standard [CFG⁺20] and an *action head*, where an action head is of the following form:

$$h : @a[i_1, \dots, i_n] = v_r$$

where

- the head atom h is an ASP atom of form $p(t_1, \dots, t_n)$ with p and $t_1 \dots, t_n$ being a predicate symbol and a list of terms, respectively.
- the function symbol a is the name of an action function, i.e. an identifier starting with a lower-case letter
- action input terms t_1 through t_n are a list of terms
- result variable r_v is a variable.

Action result variables must not occur in the rule body.

Definition 1.1.2 (Module Literals). TBD

1.2 Semantics

1.2.1 Action Rules

Desiderata

For every Evolog Program P and answer set A , the following must be clearly defined:

- $D1$: Which actions were executed by the program?
- $D2$: For every individual action act , what led to the action being executed, i.e. of which rule body is act a consequence?

Combining $D1$ and $D2$ it follows that

- $D3$: for actions that depend on other actions, it is clearly visible in which sequence they were executed, i.e. the respective execution sequence can be unambiguously reconstructed using the answer set and program('s dependency graph).

Furthermore,

- $D4$: all state changes effected on the outside world by execution of P are reflected in each answer set (as results of actions).

Definition 1.2.1 (Expansion of action rules). Semantically, every action rule is equivalent to its *expansion*:

$file1_open(OP_RES) : @fileInputStream(PATH) = OP_RES : -file1(PATH)$

The expansion of $r1$ is:

$action_result(r1, fileInputStream, PATH, fileInputStream(PATH)) : -file1(PATH).$

$file1_open(OP_RES) : -action_result(r1, fileInputStream, PATH, OP_RES).$

Consequently, it is ensured that for each ground instance of an action rule R_a that fires, there is exactly one $action_result$ instance in every answer set. We call this atom a *witness of action act*. Requirement $D1$ is fulfilled through the existence of action witnesses. Furthermore, inspection of a program (or its dependency graph) and all action witnesses in an answer set yields the information demanded in $D2$.

Definition 1.2.2 (Applicability of action rules). In order to guarantee $D1$ and $D4$, for every (ground) action rule R_a that fires, it must hold that the corresponding *witness atom* is part of *every answer set*. Implementations may further restrict this in order to ensure static verifiability of the condition (e.g. by restricting action rules to the stratified part, i.e. common base program of a program).

This is an example, make into a proper definition

Definition 1.2.3 (Rule Identifier). Given a non-ground Evolog rule R , $id(R)$ denotes a (program-wide) unique identifier of R .

Definition 1.2.4 (Action function). An action function f_{act} maps a rule id r , a tuple S , and a list of input terms t_1, \dots, t_n to a result term t_{res} .

- Identifier r references the rule (within a program) that is the *action source* (i.e. that fires in order to trigger the action)
- State S is a ground substitution for all body variables of the action source rule, i.e. it encodes the state of the world on which the action operates.

In accordance with Definition 1.2.4, an action witness $action_result(r_1, fileInputStream, PATH, OP_RES)$ then reads as "Function $fileInputStream$, with action source r_1 , applied to input $PATH$, given world state $(PATH)$, gives result OP_RES ".

Definition 1.2.5 (Interpretations of Evolog programs). An Evolog interpretation I of a program P is a tuple (F, H) consisting of a *Frame* F and a herbrand interpretation H . The frame F defines the action functions associated with rules in P .

Definition 1.2.6 (Evolog Model). An evolog interpretation $I = (F, H)$ is a *model* of evolog program P iff H is a *stable model* of P , and all action witness atoms in H are consistent with the action function definitions in F .

Bibliography

- [CFG⁺20] Francesco Calimeri, Wolfgang Faber, Martin Gebser, Giovambattista Ianni, Roland Kaminski, Thomas Krennwallner, Nicola Leone, Marco Maratea, Francesco Ricca, and Torsten Schaub. Asp-core-2 input language format. *Theory and Practice of Logic Programming*, 20(2):294–309, 2020.