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### Danksagung

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# Acknowledgements

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### Kurzfassung

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### Abstract

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HAPTER 1

# Introduction

Intro here

### The Evolog Language

#### 2.1 Syntax

Every valid ASP-Core2 program is a valid Evolog program. In addition, Evolog programs may contain *action rules* and *module literals*.

**Definition 2.1.1** (Action Rule). Action Rules are ASP rules that have a body as defined by the ASP-Core2 standard [?] and an *action head*, where an action head is of the following form:

$$h: @a[i_1, \ldots, i_n] = v_r$$

where

- the head atom h is an ASP atom of form  $p(t_1, \ldots, t_n)$  with p and  $t_1, \ldots, t_n$  being a predicate symbol and a list of terms, respectively.
- the function symbol a is the name of an action function, i.e. an identifier starting with a lower-case letter
- action input terms  $t_1$  through  $t_n$  are a list of terms
- result variable  $r_v$  is a variable.

Action result variables must not occur in the rule body.

**Definition 2.1.2** (Module Literals). TBD

#### 2.2**Semantics**

#### 2.2.1 **Action Rules**

#### Desiderata

For every Evolog Program P and answer set A, the following must be clearly defined:

- D1: Which actions were executed by the program?
- D2: For every individual action act, what led to the action being executed, i.e. of which rule body is act a consequence?

Combining D1 and D2 it follows that

• D3: for actions that depend on other actions, it is clearly visible in which sequence they were executed, i.e. the respective execution sequence can be unambiguously reconstructed using the answer set and program('s dependency graph).

Furthermore,

• D4: all state changes effected on the outside world by execution of P are reflected in each answer set (as results of actions).

**Definition 2.2.1** (Expansion of action rules). Semantically, every action rule is equivamake into a proper lent to its expansion:

 $file1 \ open(OP \ RES) : @fileInputStream[PATH] = OP \ RES : -file1(PATH).$ 

The expansion of r1 is:

 $action\_result(r1, fileInputStream, PATH, fileInputStream(PATH)) : -file1(PATH).file1\_open$ 

Consequently, it is ensured that for each ground instance of an action rule  $R_a$  that fires, there is exactly one action result instance in every answer set. We call this atom a witness of action act. Requirement D1 is fulfilled through the existence of action witnesses. Furthermore, inspection of a program (or its dependency graph) and all action witnesses in an answer set yields the information demanded in D2.

**Definition 2.2.2** (Applicability of action rules). In order to guarantee D1 and D4, for every (ground) action rule  $R_a$  that fires, it must hold that the corresponding witness atom is part of every answer set. Implementations may further restrict this in order to ensure static verifiability of the condition (e.g. by restricting action rules to the stratified part, i.e. common base program of a program).

This is an example, definition

**Definition 2.2.3** (Rule Identifier). Given a non-ground Evolog rule R, id(R) denotes a (program-wide) unique identifier of R.

**Definition 2.2.4** (Action function). An action function  $f_{act}$  maps a rule id r, a tuple S, and a list of input terms  $t_1, \ldots, t_n$  to a result term  $t_{res}$ .

- Identifier r references the rule (within a program) that is the *action source* (i.e. that fires in order to trigger the action)
- State S is a ground susbtitution for all body variables of the action source rule, i.e. it encodes the state of the world on which the action operates.

In accordance with Definition 2.2.4, an action witness  $action\_result(r_1, fileInputStream, PATH, OP\_RES)$  then reads as "Function fileInputStream, with action source  $r_1$ , applied to input PATH, given world state (PATH), gives result  $OP_RES$ ". \_\_\_\_\_\_\_ this should be

**Definition 2.2.5** (Interpretations of Evolog programs). An Evolog interpretation I of a program P is a tuple (F, H) consisting of a Frame F and a herbrand interpretation H. The frame F defines the action functions associated with rules in P.

this should be done nicer. Actually, a function is applied to its input, given a rule body holds, that's exactly the elegance, we wanna put this into mathspeak.

# List of Figures

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# List of Algorithms

### **Bibliography**

[Tur36] Alan Mathison Turing. On computable numbers, with an application to the entscheidungsproblem. J. of Math, 58:345-363, 1936.