

Farkas established that a system of linear inequalities has a solution if and only if we cannot obtain a contradiction by taking a nonnegative linear combination of the inequalities. We state and formally prove several Farkas-like theorems in Lean 4. Furthermore, we consider the rational numbers extended with two special elements denoted by  $\perp$  and  $\top$  where  $\perp$  is below every rational number and  $\top$  is above every rational number. We define  $\perp + a = \perp = a + \perp$  of all  $a$  and we define  $\top + b = \top = b + \top$  for all  $b \neq \perp$ . For multiplication, we define  $\perp \cdot c = \perp = c \cdot \perp$  for every  $c \geq 0$  but  $\top \cdot d = \top = d \cdot \top$  only for  $d > 0$  because  $\top \cdot 0 = 0 = 0 \cdot \top$ . We extend certain Farkas-like theorems about matrices to a setting where coefficients are extended rationals and variables are nonnegative rationals.

## 1 Introduction

Rouché, Capelli, Kronecker, Fonténe, Frobenius (TODO citations) established that a system of linear equalities has a solution if and only if we cannot obtain a contradiction by taking a linear combination of the equalities. We state their theorem as follows.

**Theorem (equalityFrobenius):** Let  $I$  and  $J$  be finite types. Let  $F$  be a linearly ordered field. Let  $A$  be a matrix of type  $(I \times J) \rightarrow F$ . Let  $b$  be a vector of type  $I \rightarrow F$ . Exactly one of the following exists:

- vector  $x : J \rightarrow F$  such that  $A \cdot x = b$
- vector  $y : I \rightarrow F$  such that  $A^T \cdot y = 0$  and  $b \cdot y \neq 0$

This theorem can be given in much more general settings. In our paper, however, this is the only version we provide. This fundamental staple of linear algebra is not our main focus, but a byproduct of the other theorems we prove. The way we state the theorem exemplifies certain patterns that permeate TODO.

Farkas (TODO citation) gave a similar characterization for systems of linear equalities with nonnegative variables. We state his theorem as follows.

**Theorem (equalityFarkas):** Let  $I$  and  $J$  be finite types. Let  $F$  be a linearly ordered field. Let  $A$  be a matrix of type  $(I \times J) \rightarrow F$ . Let  $b$  be a vector of type  $I \rightarrow F$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow F$  such that  $A \cdot x = b$
- vector  $y : I \rightarrow F$  such that  $A^T \cdot y \geq 0$  and  $b \cdot y < 0$

For reasons that will be apparent later, we give another version of the last theorem.

**Theorem (equalityFarkas\_neg):** Let  $I$  and  $J$  be finite types. Let  $F$  be a linearly ordered field. Let  $A$  be a matrix of type  $(I \times J) \rightarrow F$ . Let  $b$  be a vector of type  $I \rightarrow F$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow F$  such that  $A \cdot x = b$
- vector  $y : I \rightarrow F$  such that  $(-A^T) \cdot y \leq 0$  and  $b \cdot y < 0$

The proof of equalityFarkas\_neg will come later. However, we note that equalityFarkas is an immediate corollary of equalityFarkas\_neg and that equalityFrobenius is proved by applying equalityFarkas to the matrix  $(A \mid -A)$ .

Minkowski (TODO citation) similarly established that a system of linear inequalities has a nonnegative solution if and only if we cannot obtain a contradiction by taking a nonnegative linear combination of the inequalities. We state his theorem as follows.

**Theorem (inequalityFarkas):** Let  $I$  and  $J$  be finite types. Let  $F$  be a linearly ordered field. Let  $A$  be a matrix of type  $(I \times J) \rightarrow F$ . Let  $b$  be a vector of type  $I \rightarrow F$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow F$  such that  $A \cdot x \leq b$
- nonnegative vector  $y : I \rightarrow F$  such that  $A^T \cdot y \geq 0$  and  $b \cdot y < 0$

Again, we provide one more version that will make sense later.

**Theorem (inequalityFarkas\_neg):** Let  $I$  and  $J$  be finite types. Let  $F$  be a linearly ordered field. Let  $A$  be a matrix of type  $(I \times J) \rightarrow F$ . Let  $b$  be a vector of type  $I \rightarrow F$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow F$  such that  $A \cdot x \leq b$
- nonnegative vector  $y : I \rightarrow F$  such that  $(-A^T) \cdot y \leq 0$  and  $b \cdot y < 0$

Again, inequalityFarkas is an immediate corollary of inequalityFarkas\_neg. And as for inequalityFarkas\_neg, we prove it by applying equalityFarkas\_neg to the matrix  $(1 \mid A)$  where  $1$  is the identity matrix of type  $(I \times I) \rightarrow F$ .

The next theorem generalizes equalityFarkas\_neg to structures where multiplication does not have to be commutative.

**Theorem (coordinateFarkas):** Let  $I$  be any type. Let  $J$  be a finite type. Let  $R$  be a linearly ordered division ring. Let  $A$  be an  $R$ -linear map from from  $(I \rightarrow R)$  to  $(J \rightarrow R)$ . Let  $b$  be an  $R$ -linear map from from  $(I \rightarrow R)$  to  $R$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow R$  such that, for all  $w : I \rightarrow R$ , we have  $\sum_{j:J} (A \ w)_j \cdot x_j = b \ w$
- vector  $y : I \rightarrow R$  such that  $A \ y \leq 0$  and  $b \ y > 0$

In the next generalization, we replace the “coordinate-equipped” module  $I \rightarrow R$  by a general  $R$ -module  $W$ .

**Theorem (scalarFarkas):** Let  $J$  be a finite type. Let  $R$  be a linearly ordered division ring. Let  $W$  be an  $R$ -module. Let  $A$  be an  $R$ -linear map from  $W$  to  $(J \rightarrow R)$ . Let  $b$  be an  $R$ -linear map from  $W$  to  $R$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow R$  such that, for all  $w$  in  $W$ , we have  $\sum_{j:J} (A\ w)_j \cdot x_j = b\ w$
- vector  $y : W$  such that  $A\ y \leq 0$  and  $b\ y > 0$

In the most general theorem, stated below, we replace certain occurrences of  $R$  by a linearly ordered  $R$ -module  $V$  whose ordering respects TODO. This result origins from TODO.

**Theorem (fintypeFarkasBartl):** Let  $J$  be a finite type. Let  $R$  be a linearly ordered division ring. Let  $W$  be an  $R$ -module. Let  $V$  be a linearly ordered  $R$ -module<sup>1</sup>. Let  $A$  be an  $R$ -linear map from  $W$  to  $(J \rightarrow R)$ . Let  $b$  be an  $R$ -linear map from  $W$  to  $V$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow V$  such that, for all  $w$  in  $W$ , we have  $\sum_{j:J} (A\ w)_j \cdot x_j = b\ w$
- vector  $y : W$  such that  $A\ y \leq 0$  and  $b\ y > 0$

Note that fintypeFarkasBartl subsumes scalarFarkas (as well as the other versions), since  $R$  can be viewed as a linearly ordered module over itself. We prove fintypeFarkasBartl in Section TODO, which is where the heavy lifting comes.

Until now, we have talked about known results from TODO. What follows is a new extension of the theory.

**Definition:** Let  $F$  be a linearly ordered field. We define an **extended** linearly ordered field  $F_\infty$  as  $F \cup \{\perp, \top\}$  with the following properties. Let  $p$  and  $q$  be numbers from  $F$ . We have  $\perp < p < \top$ . We define addition and scalar action on  $F_\infty$  as follows:

$+$	$\perp$	$q$	$\top$
$\perp$	$\perp$	$\perp$	$\perp$
$p$	$\perp$	$p+q$	$\top$
$\top$	$\perp$	$\top$	$\top$

$\bullet$	$\perp$	$q$	$\top$
$0$	$\perp$	$0$	$0$
$p > 0$	$\perp$	$p \cdot q$	$\top$

Informally speaking,  $\top$  represents the positive infinity,  $\perp$  represents the negative infinity, and we say that  $\perp$  is “stronger” than  $\top$  in all arithmetic operations. The surprising parts are  $\perp + \top = \perp$  and  $0 \bullet \perp = \perp$ . Because of that,  $F_\infty$  is not a field. However,  $F_\infty$  is still a densely linearly ordered abelian monoid with characteristic zero.

**Theorem (extendedFarkas):** Let  $I$  and  $J$  be finite types. Let  $F$  be a linearly ordered field. Let  $A$  be a matrix of type  $(I \times J) \rightarrow F_\infty$ . Let  $b$  be a vector of type  $I \rightarrow F_\infty$ . Assume that  $A$  does not have  $\perp$  and  $\top$  in the same row. Assume that  $A$  does not have  $\perp$  and  $\top$  in the same column. Assume that  $A$  does not have  $\top$  in any row where  $b$  has  $\top$ . Assume that  $A$  does not have  $\perp$  in any row where  $b$  has  $\perp$ . Exactly one of the following exists:

- nonnegative vector  $x : J \rightarrow F$  such that  $A \cdot x \leq b$
- nonnegative vector  $y : I \rightarrow F$  such that  $(-A^T) \cdot y \leq 0$  and  $b \cdot y < 0$

TODO everything about linear programming.

## 2 Preliminaries

We distinguish two types of vectors; implicit vectors and explicit vectors. Implicit vectors are members of a vector space; they don’t have any internal structure. Explicit vectors are functions from coordinates to values. The set of coordinates needn’t be ordered. Matrices live next to explicit vectors. They are also functions; they take a row index and a column index and they output a value at the given spot. Neither the row indices nor the column vertices are required to form an ordered set. That’s why multiplication between matrices and vectors is defined only in structures where addition forms a commutative semigroup. Consider the following example:

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} \cdot \begin{pmatrix} 7 \\ 8 \\ 9 \end{pmatrix} = \begin{pmatrix} ? \\ - \end{pmatrix}$$

We don’t know whether the value at the question mark is equal to  $(1 \cdot 7 + 2 \cdot 8) + 3 \cdot 9$  or to  $(2 \cdot 8 + 1 \cdot 7) + 3 \cdot 9$  or to any other ordering of summands. This is why commutativity of addition is necessary for the definition to be valid. On the other hand, we don’t assume any property of multiplication in the definition of multiplication between matrices and vectors; they don’t even have to be of the same type; we only require the elements of the vector to have an action on the elements of the matrix (this is not a typo – normally, we would want matrices to have an action on vectors – not in our work).

TODO.

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<sup>1</sup>We furthermore require PosSMulMono R V.

### 3 Proving the Farkas-Bartl theorem

TODO.

### 4 Inequalities and linear programming

TODO.

### 5 Proving the extended strong duality

TODO.