Abstract: Farkas established that a system of linear inequalities has a solution if and only if we cannot obtain a contradiction by taking a linear combination of the inequalities. We state and formally prove several Farkas-like theorems in Lean 4. Furthermore, we consider a linearly ordered field extended with two special elements denoted by \bot and \top where \bot is below every element and \top is above every element. We define $\bot + a = \bot = a + \bot$ of all a and we define $\top + b = \top = b + \top$ for all $b \ne \bot$. For multiplication, we define $\bot \cdot c = \bot = c \cdot \bot$ for every $c \ge 0$ but $\top \cdot d = \top = d \cdot \top$ only for d > 0 because $\top \cdot 0 = 0 = 0 \cdot \top$. We extend certain Farkas-like theorems to a setting where coefficients are from an extended linearly ordered field.

1 Introduction

Rouché, Capelli, Kronecker, Fonténe, Frobenius (TODO citations) established that a system of linear equalities has a solution if and only if we cannot obtain a contradiction by taking a linear combination of the equalities. We state their theorem as follows.

Theorem (equalityFrobenius): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F$. Let B be a vector of type $I \to F$. Exactly one of the following exists:

- vector $x: J \to F$ such that $A \cdot x = b$
- vector $y: I \to F$ such that $A^T \cdot y = 0$ and $b \cdot y \neq 0$

This theorem can be given in much more general settings. In our paper, however, this is the only version we provide. This fundamental staple of linear algebra is not our main focus, but a byproduct of the other theorems we prove. The way we state the theorem exemplifies certain patterns that permeate TODO.

Farkas (TODO citation) gave a similar characterization for systems of linear equalities with nonnegative variables. We state his theorem as follows.

Theorem (equalityFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F$. Let b be a vector of type $I \to F$. Exactly one of the following exists:

- nonnegative vector $x: J \to F$ such that $A \cdot x = b$
- vector $y: I \to F$ such that $A^T \cdot y \ge 0$ and $b \cdot y < 0$

For reasons that will be apparent later, we give another version of the last theorem.

Theorem (equalityFarkas_neg): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F$. Let B be a vector of type $I \to F$. Exactly one of the following exists:

- nonnegative vector $x: J \to F$ such that $A \cdot x = b$
- vector $y: I \to F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

The proof of equalityFarkas_neg will come later. However, we note that equalityFarkas is an immediate corollary of equalityFarkas_neg and that equalityFrobenius is proved by applying equalityFarkas to the matrix $(A \mid -A)$.

Minkowski (TODO citation) similarly established that a system of linear inequalities has a nonnegative solution if and only if we cannot obtain a contradiction by taking a nonnegative linear combination of the inequalities. We state his theorem as follows.

Theorem (inequalityFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F$. Let B be a vector of type $I \to F$. Exactly one of the following exists:

- nonnegative vector $x: J \to F$ such that $A \cdot x \leq b$
- nonnegative vector $y: I \to F$ such that $A^T \cdot y \geq 0$ and $b \cdot y < 0$

Again, we provide one more version that will make sense later.

Theorem (inequalityFarkas_neg): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F$. Let b be a vector of type $I \to F$. Exactly one of the following exists:

- nonnegative vector $x: J \to F$ such that $A \cdot x \leq b$
- nonnegative vector $y: I \to F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

Again, inequalityFarkas is an immediate corollary of inequalityFarkas_neg. And as for inequalityFarkas_neg, we prove it by applying equalityFarkas_neg to the matrix $(1 \mid A)$ where 1 is the identity matrix of type $(I \times I) \to F$.

The next theorem generalizes equalityFarkas_neg to structures where multiplication does not have to be commutative. Furthermore, we can have infinitely many equations.

Theorem (coordinateFarkas): Let I be any type. Let J be a finite type. Let R be a linearly ordered division ring. Let R be an R-linear map from from $(I \to R)$ to $(I \to R)$ to $(I \to R)$. Let R be an R-linear map from from $(I \to R)$ to R. Exactly one of the following exists:

- nonnegative vector $x: J \to R$ such that, for all $w: I \to R$, we have $\sum_{j:J} (A \ w)_j \cdot x_j = b \ w$
- vector $y: I \to R$ such that $A y \le 0$ and b y > 0

In the next generalization, we replace the partially ordered module $I \to R$ by a general R-module W.

Theorem (scalarFarkas): Let J be a finite type. Let R be a linearly ordered division ring. Let W be an R-module. Let A be an R-linear map from from W to $(J \to R)$. Let B be an B-linear map from from B to B. Exactly one of the following exists:

- nonnegative vector $x: J \to R$ such that, for all w in W, we have $\sum_{i:J} (A w)_i \cdot x_i = b w$
- vector y: W such that $A y \leq 0$ and b y > 0

In the most general theorem, stated below, we replace certain occurrences of R by a linearly ordered R-module V whose ordering respects TODO. This result origins from TODO.

Theorem (fintypeFarkasBartl): Let J be a finite type. Let R be a linearly ordered division ring. Let W be an R-module. Let V be a linearly ordered R-module¹. Let A be an R-linear map from from W to $(J \to R)$. Let B be an B-linear map from from B to B. Exactly one of the following exists:

- nonnegative (TODO terminology???) bivector $x: J \to V$ such that, for all w in W, we have $\sum_{i:J} (A w)_j \cdot x_j = b w$
- vector y: W such that $A y \leq 0$ and b y > 0

In the last branch, $Ay \le 0$ uses the partial order² on $(J \to R)$ whereäs by > 0 uses the linear order³ on V. Note that fintypeFarkasBartl subsumes scalarFarkas (as well as the other versions based on equality), since R can be viewed as a linearly ordered module over itself. We prove fintypeFarkasBartl in Section TODO, which is where the heavy lifting comes.

Until now, we have talked about known results from TODO. What follows is a new extension of the theory.

Definition: Let F be a linearly ordered field. We define an **extended** linearly ordered field F_{∞} as $F \cup \{\bot, \top\}$ with the following properties. Let p and q be numbers from F. We have $\bot . We define addition and scalar action on <math>F_{\infty}$ as follows:

+	上	q	Т
	1		1
p	1	p+q	Т
T		Т	Т

•		q	Т
0	1	0	0
p > 0	上	$p \cdot q$	Т

Informally speaking, \top represents the positive infinity, \bot represents the negative infinity, and we say that \bot is "stronger" than \top in all arithmetic operations. The surprising parts are $\bot + \top = \bot$ and $0 \bullet \bot = \bot$. Because of that, F_{∞} is not a field. However, F_{∞} is still a densely linearly ordered abelian monoid with characteristic zero.

Theorem (extendedFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F_{\infty}$. Let b be a vector of type $I \to F_{\infty}$. Assume that A does not have \bot and \top in the same row. Assume that A does not have \bot and \top in the same column. Assume that A does not have \top in any row where b has \top . Assume that A does not have \bot in any row where b has \bot . Exactly one of the following exists:

- nonnegative vector $x: J \to F$ such that $A \cdot x \le b$
- nonnegative vector $y: I \to F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

Note that extendedFarkas looks pretty much like equalityFarkas_neg and, in certain sense, generalizes it. Indeed, in Section TODO, we prove extendedFarkas using equalityFarkas_neg and some additional machinery.

Definition: Let I and J be finite types. Let R be a a partially ordered semiring. Let A be a matrix of type $(I \times J) \to R$. Let b be a vector of type $I \to R$. Let c be a vector of type $J \to R$. We say that P = (A, b, c) is a **linear program** over R. We say that its **constraints** are indexed by I. We say that its **variables** are indexed by J. We say that a nonnegative vector $x: J \to R$ is a **solution** to P if and only if $A \cdot x \le b$. We say that P **reaches** an objective value r if and only if there exists x such that x is a solution to P and $c \cdot x = r$.

TODO feasibility, dualize, weakDuality, strongDuality.

2 Preliminaries

We distinguish two types of vectors; implicit vectors and explicit vectors. Implicit vectors are members of a vector space; they don't have any internal structure. Explicit vectors are functions from coordinates to values. The set of coordinates needn't be ordered. Matrices live next to explicit vectors. They are also functions; they take a row index and a column index and they output a value at the given spot. Neither the row indices nor the column vertices are required to form an ordered set. That's why multiplication between matrices and vectors is defined only in structures where addition forms a commutative semigroup. Consider the following example:

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} \cdot \begin{pmatrix} 7 \\ 8 \\ 9 \end{pmatrix} = \begin{pmatrix} ? \\ - \end{pmatrix}$$

¹We furthermore require PosSMulMono R V, which means TODO.

²The order on $(J \to R)$ is always the coordinate-wise application of R's linear order.

 $^{^3}$ In case V has finite dimension, you can choose an arbitrary direction and project vectors from V onto it, or you can order elements of V lexicographically.

We don't know whether the value at the question mark is equal to $(1 \cdot 7 + 2 \cdot 8) + 3 \cdot 9$ or to $(2 \cdot 8 + 1 \cdot 7) + 3 \cdot 9$ or to any other ordering of summands. This is why commutativity of addition is necessary for the definition to be valid. On the other hand, we don't assume any property of multiplication in the definition of multiplication between matrices and vectors; they don't even have to be of the same type; we only require the elements of the vector to have an action on the elements of the matrix (this is not a typo – normally, we would want matrices to have an action on vectors – not in our work).

Proving the Farkas-Bartl theorem

TODO.

 $\mathbf{3}$

TODO.

4 Extended Farkas theorem

Theorem (extendedFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \to F_{\infty}$. Let b be a vector of type $I \to F_{\infty}$. Assume that A does not have \bot and \top in the same row. Assume that A does not have \bot and \top in the same column. Assume that A does not have \bot in any row where b has \bot . Assume that A does not have \bot in any row where b has \bot . Exactly one of the following exists:

- nonnegative vector $x:J\to F$ such that $A\cdot x\leq b$
- nonnegative vector $y: I \to F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

(restated)

Proof: We need to do the following steps in given order:

- 1. Delete all rows of both A and b where A has \perp or b has \top (they are tautologies).
- 2. Delete all columns of A that contain \top (they force respective variables to be zero).
- 3. If b contains \perp then the $A \cdot x \leq b$ cannot be satisfied, but y = 0 satisfies $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$. Stop here.
- 4. Assume there is no \perp in b. Use inequality Farkas_neg. In either case, extend x or y with zeros on all deleted positions.
- 5. Carefully check all four implications.

We now give a few counterexamples for situations where the preconditions are not met.

$$A = \begin{pmatrix} \bot \\ \top \end{pmatrix} \qquad b = \begin{pmatrix} -1 \\ 0 \end{pmatrix}$$

Putting x=(0) satisfies $A\cdot x\leq b$. Putting $y=\begin{pmatrix}1\\1\end{pmatrix}$ satisfies $(-A^T)\cdot y\leq 0$ and $b\cdot y<0$.

$$A = \begin{pmatrix} \bot & \top \\ 0 & -1 \end{pmatrix} \qquad b = \begin{pmatrix} 0 \\ -1 \end{pmatrix}$$

Putting $x = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$ satisfies $A \cdot x \leq b$. Putting $y = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$ satisfies $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$.

$$A = \left(\bot\right) \qquad \qquad b = \left(\bot\right)$$

Putting x=(1) satisfies $A \cdot x \leq b$. Putting y=(0) satisfies $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$.

$$A = \begin{pmatrix} \top \\ -1 \end{pmatrix} \qquad b = \begin{pmatrix} \top \\ -1 \end{pmatrix}$$

Putting x=(1) satisfies $A\cdot x\leq b$. Putting $y=\begin{pmatrix} 0\\1 \end{pmatrix}$ satisfies $(-A^T)\cdot y\leq 0$ and $b\cdot y<0$.

5 Proving the (extended) strong LP duality

TODO.

6 Related work

TODO.

7 Conclusion

TODO.