

Abstract: Farkas established that a system of linear inequalities has a solution if and only if we cannot obtain a contradiction by taking a linear combination of the inequalities. We state and formally prove several Farkas-like theorems in Lean 4. Furthermore, we consider a linearly ordered field extended with two special elements denoted by \perp and \top where \perp is below every element and \top is above every element. We define $\perp + a = \perp = a + \perp$ of all a and we define $\top + b = \top = b + \top$ for all $b \neq \perp$. For multiplication, we define $\perp \cdot c = \perp = c \cdot \perp$ for every $c \geq 0$ but $\top \cdot d = \top = d \cdot \top$ only for $d > 0$ because $\top \cdot 0 = 0 = 0 \cdot \top$. We extend certain Farkas-like theorems to a setting where coefficients are from an extended linearly ordered field.

1 Introduction

Rouché, Capelli, Kronecker, Fonténe, Frobenius (TODO citations) established that a system of linear equalities has a solution if and only if we cannot obtain a contradiction by taking a linear combination of the equalities. We state their theorem as follows.

Theorem (equalityFrobenius): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F$. Let b be a vector of type $I \rightarrow F$. Exactly one of the following exists:

- vector $x : J \rightarrow F$ such that $A \cdot x = b$
- vector $y : I \rightarrow F$ such that $A^T \cdot y = 0$ and $b \cdot y \neq 0$

This theorem can be given in much more general settings. In our paper, however, this is the only version we provide. This fundamental staple of linear algebra is not our main focus but a byproduct of the other theorems we prove. In particular, we obtain equalityFrobenius as an immediate corollary of the following theorem.

Theorem (equalityFrobenius_lt): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F$. Let b be a vector of type $I \rightarrow F$. Exactly one of the following exists:

- vector $x : J \rightarrow F$ such that $A \cdot x = b$
- vector $y : I \rightarrow F$ such that $A^T \cdot y = 0$ and $b \cdot y < 0$

The way we state the theorem exemplifies certain patterns that permeate TODO.

Farkas (TODO citation) gave a similar characterization for systems of linear equalities with nonnegative variables. We state his theorem as follows.

Theorem (equalityFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F$. Let b be a vector of type $I \rightarrow F$. Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow F$ such that $A \cdot x = b$
- vector $y : I \rightarrow F$ such that $A^T \cdot y \geq 0$ and $b \cdot y < 0$

We prove equalityFrobenius.lt by applying equalityFarkas to the matrix $(A \mid -A)$. Theorem equalityFarkas will be proved later.

Minkowski (TODO citation) similarly established that a system of linear inequalities has a nonnegative solution if and only if we cannot obtain a contradiction by taking a nonnegative linear combination of the inequalities. For reasons that will be apparent later, we give two versions of this theorem.

Theorem (inequalityFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F$. Let b be a vector of type $I \rightarrow F$. Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow F$ such that $A \cdot x \leq b$
- nonnegative vector $y : I \rightarrow F$ such that $A^T \cdot y \geq 0$ and $b \cdot y < 0$

Theorem (inequalityFarkas_neg): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F$. Let b be a vector of type $I \rightarrow F$. Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow F$ such that $A \cdot x \leq b$
- nonnegative vector $y : I \rightarrow F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

Obviously, inequalityFarkas_neg is an immediate corollary of inequalityFarkas. We prove inequalityFarkas by applying equalityFarkas to the matrix $(1 \mid A)$ where 1 is the identity matrix of type $(I \times I) \rightarrow F$.

The next theorem generalizes equalityFarkas to structures where multiplication does not have to be commutative. Furthermore, it supports infinitely many equations.

Theorem (coordinateFarkas): Let I be any type. Let J be a finite type. Let R be a linearly ordered division ring. Let A be an R -linear map from from $(I \rightarrow R)$ to $(J \rightarrow R)$. Let b be an R -linear map from from $(I \rightarrow R)$ to R . Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow R$ such that, for all $w : I \rightarrow R$, we have $\sum_{j:J} (A \ w)_j \bullet x_j = b \ w$
- vector $y : I \rightarrow R$ such that $A \ y \geq 0$ and $b \ y < 0$

In the next generalization, we replace the partially ordered module $I \rightarrow R$ by a general R -module W .

Theorem (scalarFarkas): Let J be a finite type. Let R be a linearly ordered division ring. Let W be an R -module. Let A be an R -linear map from W to $(J \rightarrow R)$. Let b be an R -linear map from W to R . Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow R$ such that, for all $w : W$, we have $\sum_{j:J} (A \ w)_j \bullet x_j = b \ w$
- vector $y : W$ such that $A \ y \geq 0$ and $b \ y < 0$

In the most general theorem, stated below, we replace certain occurrences of R by a linearly ordered R -module V whose order respects order on R . This result origins from TODO.

Theorem (fintypeFarkasBartl): Let J be a finite type. Let R be a linearly ordered division ring. Let W be an R -module. Let V be a linearly ordered R -module¹. Let A be an R -linear map from W to $(J \rightarrow R)$. Let b be an R -linear map from W to V . Exactly one of the following exists:

- nonnegative bivector $x : J \rightarrow V$ such that, for all $w : W$, we have $\sum_{j:J} (A \ w)_j \bullet x_j = b \ w$
- vector $y : W$ such that $A \ y \geq 0$ and $b \ y < 0$

In the last branch, $A \ y \geq 0$ uses the partial order² on $(J \rightarrow R)$ whereäs $b \ y < 0$ uses the linear order³ on V . Note that fintypeFarkasBartl subsumes scalarFarkas (as well as the other versions based on equality), since R can be viewed as a linearly ordered module over itself. We prove fintypeFarkasBartl in Section TODO, which is where the heavy lifting comes.

Until now, we have talked about known results. What follows is a new extension of the theory.

Definition: Let F be a linearly ordered field. We define an **extended** linearly ordered field F_∞ as $F \cup \{\perp, \top\}$ with the following properties. Let p and q be numbers from F . We have $\perp < p < \top$. We define addition and scalar action on F_∞ as follows:

+	\perp	q	\top
\perp	\perp	\perp	\perp
p	\perp	$p+q$	\top
\top	\perp	\top	\top

\bullet	\perp	q	\top
0	\perp	0	0
$p > 0$	\perp	$p \cdot q$	\top

Informally speaking, \top represents the positive infinity, \perp represents the negative infinity, and we say that \perp is “stronger” than \top in all arithmetic operations. The surprising parts are $\perp + \top = \perp$ and $0 \bullet \perp = \perp$. Because of that, F_∞ is not a field. However, F_∞ is still a densely linearly ordered abelian monoid with characteristic zero.

Theorem (extendedFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F_\infty$. Let b be a vector of type $I \rightarrow F_\infty$. Assume that A does not have \perp and \top in the same row. Assume that A does not have \perp and \top in the same column. Assume that A does not have \top in any row where b has \top . Assume that A does not have \perp in any row where b has \perp . Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow F$ such that $A \cdot x \leq b$
- nonnegative vector $y : I \rightarrow F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

Note that extendedFarkas looks pretty much like equalityFarkas_neg and, in certain sense, generalizes it. Indeed, in Section TODO, we prove extendedFarkas using equalityFarkas_neg and some additional machinery.

Definition: Let I and J be finite types. Let R be a a partially ordered semiring. Let A be a matrix of type $(I \times J) \rightarrow R$. Let b be a vector of type $I \rightarrow R$. Let c be a vector of type $J \rightarrow R$. We say that $P = (A, b, c)$ is a **linear program** over R . We say that its **constraints** are indexed by I . We say that its **variables** are indexed by J . We say that a nonnegative vector $x : J \rightarrow R$ is a **solution** to P if and only if $A \cdot x \leq b$. We say that P **reaches** an objective value r if and only if there exists x such that x is a solution to P and $c \cdot x = r$.

TODO feasibility, dualize, weakDuality, strongDuality.

2 Formalization

We distinguish two types of vectors; implicit vectors and explicit vectors. Implicit vectors are members of a vector space; they don't have any internal structure. Explicit vectors are functions from coordinates to values. The set of coordinates needn't be ordered. Matrices live next to explicit vectors. They are also functions; they take a row index and a column index and they output a value at the given spot. Neither the row indices nor the column vertices are required to form an ordered set. That's why multiplication between matrices and vectors is defined only in structures where addition forms a commutative semigroup. Consider the following example:

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} \cdot \begin{pmatrix} 7 \\ 8 \\ 9 \end{pmatrix} = \begin{pmatrix} ? \\ ? \end{pmatrix}$$

¹We furthermore require monotonicity of scalar multiplication by nonnegative elements on the left.

²The order on $(J \rightarrow R)$ is always the coordinate-wise application of R 's linear order.

³In case V has finite dimension, you can choose an arbitrary direction and project vectors from V onto it, or you can order elements of V lexicographically.

We don't know whether the value at the question mark is equal to $(1 \cdot 7 + 2 \cdot 8) + 3 \cdot 9$ or to $(2 \cdot 8 + 1 \cdot 7) + 3 \cdot 9$ or to any other ordering of summands. This is why commutativity of addition is necessary for the definition to be valid. On the other hand, we don't assume any property of multiplication in the definition of multiplication between matrices and vectors; they don't even have to be of the same type; we only require the elements of the vector to have an action on the elements of the matrix (this is not a typo – normally, we would want matrices to have an action on vectors – not in our work).

TODO.

3 Proving the Farkas-Bartl theorem

TODO.

4 Extended Farkas theorem

Theorem (extendedFarkas): Let I and J be finite types. Let F be a linearly ordered field. Let A be a matrix of type $(I \times J) \rightarrow F_\infty$. Let b be a vector of type $I \rightarrow F_\infty$. Assume that A does not have \perp and \top in the same row. Assume that A does not have \perp and \top in the same column. Assume that A does not have \top in any row where b has \top . Assume that A does not have \perp in any row where b has \perp . Exactly one of the following exists:

- nonnegative vector $x : J \rightarrow F$ such that $A \cdot x \leq b$
- nonnegative vector $y : I \rightarrow F$ such that $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$

(restated)

Proof idea: We need to do the following steps in given order:

1. Delete all rows of both A and b where A has \perp or b has \top (they are tautologies).
2. Delete all columns of A that contain \top (they force respective variables to be zero).
3. If b contains \perp , then $A \cdot x \leq b$ cannot be satisfied, but $y = 0$ satisfies $(-A^T) \cdot y \leq 0$ and $b \cdot y < 0$. Stop here.
4. Assume there is no \perp in b . Use inequalityFarkas_neg. In either case, extend x or y with zeros on all deleted positions.

4.1 Counterexamples

If A has \perp and \top in the same row, it may happen that both x and y exist:

$$A = \begin{pmatrix} \perp & \top \\ 0 & -1 \end{pmatrix} \quad b = \begin{pmatrix} 0 \\ -1 \end{pmatrix} \quad x = \begin{pmatrix} 1 \\ 1 \end{pmatrix} \quad y = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

If A has \perp and \top in the same column, it may happen that both x and y exist:

$$A = \begin{pmatrix} \perp \\ \top \end{pmatrix} \quad b = \begin{pmatrix} -1 \\ 0 \end{pmatrix} \quad x = \begin{pmatrix} 0 \end{pmatrix} \quad y = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

If A has \top in a row where b has \top , it may happen that both x and y exist:

$$A = \begin{pmatrix} \top \\ -1 \end{pmatrix} \quad b = \begin{pmatrix} \top \\ -1 \end{pmatrix} \quad x = \begin{pmatrix} 1 \end{pmatrix} \quad y = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

If A has \perp in a row where b has \perp , it may happen that both x and y exist:

$$A = \begin{pmatrix} \perp \end{pmatrix} \quad b = \begin{pmatrix} \perp \end{pmatrix} \quad x = \begin{pmatrix} 1 \end{pmatrix} \quad y = \begin{pmatrix} 0 \end{pmatrix}$$

5 Proving the (extended) strong LP duality

TODO.

6 Related work

TODO.

7 Conclusion

TODO.