Closure Properties of General Grammars Formally Verified

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Overview

- ► Lean 3
- ► Mathlib as of 2022-03-15
- ▶ 12500 lines of code
- no dependent types for data

https://github.com/madvorak/grammars/tree/publish

Symbols

Symbols

- ▶ We don't explicitly state that T and N must be finite.
- Only a finite amount of symbols will appear in rewrite rules.

Rules

```
structure grule (T : Type) (N : Type) :=
(input_L : list (symbol T N))
(input_N : N)
(input_R : list (symbol T N))
(output_string : list (symbol T N))
```

Grammars

```
structure grule (T : Type) (N : Type) :=
(input_L : list (symbol T N))
(input_N : N)
(input_R : list (symbol T N))
(output_string : list (symbol T N))
structure grammar (T : Type) :=
(nt : Type)
(initial : nt)
(rules : list (grule T nt))
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Grammar transformations

```
structure grammar (T : Type) :=
(nt : Type)
(initial : nt)
(rules : list (grule T nt))
def grammar_transforms (g : grammar T)
  (w_1 \ w_2 : list (symbol T g.nt)) :
  Prop :=
\exists r : grule T g.nt,
  r \in g.rules
  \exists u v : list (symbol T g.nt),
    w_1 = u ++ r.input_L ++
          [symbol.nonterminal r.input_N]
          ++ r.input_R ++ v
    w_2 = u ++ r.output\_string ++ v
```

Grammar derivations

```
def grammar_transforms (g : grammar T)
  (w_1 \ w_2 : list (symbol T g.nt)) :
  Prop :=
\exists r : grule T g.nt,
  r \in g.rules
  \exists u v : list (symbol T g.nt),
    w_1 = u ++ r.input_L ++
          [symbol.nonterminal r.input_N]
          ++ r.input_R ++ v
    w_2 = u ++ r.output\_string ++ v
def grammar_derives (g : grammar T) :
  list (symbol T g.nt) \rightarrow list (symbol T g.nt)
  \rightarrow Prop :=
relation.refl_trans_gen (grammar_transforms g)
```

```
def grammar_derives (g : grammar T) :
   list (symbol T g.nt) → list (symbol T g.nt)
   → Prop :=
relation.refl_trans_gen (grammar_transforms g)
```

Words generated by a grammar

```
def grammar_derives (g : grammar T) :
   list (symbol T g.nt) → list (symbol T g.nt)
   → Prop :=
relation.refl_trans_gen (grammar_transforms g)

def grammar_generates (g : grammar T)
   (w : list T) : Prop :=
grammar_derives g
   [symbol.nonterminal g.initial]
   (list.map symbol.terminal w)
```

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  (w : list T) : Prop :=
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def language (T : Type) : Type :=
set (list T)
```

Language of a grammar

```
def grammar_generates (g : grammar T)
  (w : list T) : Prop :=
grammar_derives g
  [symbol.nonterminal g.initial]
  (list.map symbol.terminal w)
def language (T : Type) : Type :=
set (list T)
def grammar_language (g : grammar T) :
  language T :=
set_of (grammar_generates g)
```

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def language (T : Type) : Type :=
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   language T :=
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Type-0 languages

```
def language (T : Type) : Type :=
set (list T)

def grammar_language (g : grammar T) :
   language T :=
set_of (grammar_generates g)

def is_TO (L : language T) : Prop :=
∃ g : grammar T, grammar_language g = L
```

Union of languages

```
def set.union (s<sub>1</sub> s<sub>2</sub> : set T) : set T := { a | a \in s<sub>1</sub> \vee a \in s<sub>2</sub> } instance : language.has_add (language T) := \langleset.union\rangle
```

Union of languages

```
def set.union (s<sub>1</sub> s<sub>2</sub> : set T) : set T := { a | a \in s<sub>1</sub> \vee a \in s<sub>2</sub> } instance : language.has_add (language T) := \langle set.union\rangle theorem T0_of_T0_u_T0 (L<sub>1</sub> L<sub>2</sub> : language T) : is_T0 L<sub>1</sub> \wedge is_T0 L<sub>2</sub> \rightarrow is_T0 (L<sub>1</sub> + L<sub>2</sub>)
```

Proof of closure of type-0 language (under union)

ightharpoonup given g_1 and $g_2 \rightsquigarrow$ construct g

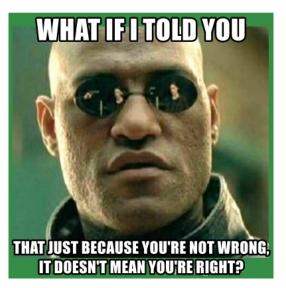
Proof of closure of type-0 language (under union)

- ightharpoonup given g_1 and $g_2 \rightsquigarrow$ construct g
- prove, for any word,
 - lacktriangle generated by g_1 or $g_2 \implies$ can be generated by g

Proof of closure of type-0 language (under union)

- ightharpoonup given g_1 and $g_2 \rightsquigarrow$ construct g
- prove, for any word,
 - ightharpoonup generated by g_1 or $g_2 \implies$ can be generated by g_1
 - ightharpoonup generated by g \Longrightarrow can be generated by g_1 or g_2

Intuitionists be like:



Wachowski, L., & Wachowski, L. (1999). The Matrix. Warner Bros. Retrieved November 11, 2021 from https://memegenerator.net/instance/47015052/matrix-morpheus-what-if-i-todeyou-that-just-because-youre-not-wrong-it-doesnt-mean-youre-right

Reversal of a language

```
def reverse_lang (L : language T) : language T := \lambda w : list T, w.reverse \in L
```

Reversal of a language

```
def reverse_lang (L : language T) : language T := \lambda \text{ w : list T, w.reverse} \in L theorem T0_of_reverse_T0 (L : language T) : is_T0 L \rightarrow is_T0 (reverse_lang L)
```

Concatenation of languages

```
def set.image2 (f : \alpha \rightarrow \beta \rightarrow \gamma) (s : set \alpha) (t : set \beta) : set \gamma := { c | \exists a b, a \in s \land b \in t \land f a b = c } instance : language.has_mul (language T) := \langle set.image2 (++)\rangle
```

Concatenation of languages

```
def set.image2 (f : \alpha \rightarrow \beta \rightarrow \gamma) (s : set \alpha) (t : set \beta) : set \gamma := { c | \exists a b, a \in s \land b \in t \land f a b = c } instance : language.has_mul (language T) := \langle set.image2 (++)\rangle theorem T0_of_T0_c_T0 (L<sub>1</sub> L<sub>2</sub> : language T) : is_T0 L<sub>1</sub> \land is_T0 L<sub>2</sub> \rightarrow is_T0 (L<sub>1</sub> * L<sub>2</sub>)
```

Kleene star of a language

```
def language.star (L : language T) : language T := { x \mid \exists S : list (list T), x = S.join \land \forall y \in S, y \in L }
```

Kleene star of a language

Related work

Chomsky hierarchy's lower levels

- Finite automata and regular expressions
 - in many languages
- Context-free grammars
 - ► (1992) Carlson et al. [Mizar]
 - ► (2007) Minamide [Isabelle]
 - (2010) Barthwal and Norrish [HOL4]
 - (2015) Firsov and Uustalu [Agda]
 - ► (2019) Ramos [Coq]

Related work

Turing-complete models

- Turing machines
 - (2001) Chen and Nakamura [Mizar]
 - ▶ (2012) Asperti and Ricciotti [Matita]
 - ► (2013) Xu et al. [Isabelle]
 - (2019) Carneiro [Lean]
 - ► (2020) Forster et al. [Coq]
 - (2023) Balbach [Isabelle]
- Random access machines
 - ► (2003) Coen [Coq]
- Lambda calculus
 - ► (2017–2021) Forster, Kunze, and their colleagues [Coq]
- Partial recursive functions
 - ► (2011) Norrish [HOL4]
 - ► (2019) Carneiro [Lean]

Future work

Chomsky hierarchy

- ► Mathlib has regular languages.
- We have context-free languages.
- We want context-sensitive languages.

Equivalence (computability theory) between TMs and grammars

- ► Turing machine → General grammar
- ▶ General grammar → Kuroda-NF grammar → Nondeterministic multi-stack machine → Turing machine