

## Security Lab – Secret Key Cryptography

### VMware

- This lab can be solved using any Java IDE. No VM is needed.

## Introduction

This lab highlights some of the properties of a secret key's work factor, and how it relates to its entropy. Recall that the work factor (in bits) is the binary logarithm of the average number of tries you need to make in order to find the correct key.

This lab simulates the effects of ransomware. This is malware that infects your computer, encrypts a few important files, sends the keys off to a command-and-control server and then deletes the original files. You will thus be left with a bunch of ciphertext files, but without the corresponding plaintext files or the key that was used to encrypt the file. You will then be contacted by the authors of the ransomware and asked to pay them money. If you do pay the money, you will sometimes get instructions how to decrypt the files. Sometimes, even the ransomware authors won't know the key, or won't care enough to actually provide you with it, even though you have paid the ransom. The bastards.

In the case of real-life ransomware, if you don't get the key somehow, your files would now be gone: this is often professionally-written software that takes great care to ensure that you cannot recover the key except from the authors. And the encryption algorithms used are normally secure so that there is no known way to break the encryption to get your files back. But perhaps the program that you are about to encounter is not so professionally written, and you can obtain the plaintext somehow?

## Exercises

### Downloading and Running the Code

**Exercise 1.** Start a shell. In a directory of your choice, type

```
git clone https://github.engineering.zhaw.ch/neut/itsec-secret-key-crypto.git
```

This should give you a directory called `itsec-secret-key-crypto`, a Java project containing our simulated ransomware. You should be able to import this project into any Java IDE; we have been using IntelliJ, but Eclipse, VS Code, and indeed any IDE should work.

**Exercise 2.** Build the project in your IDE or with maven. This should leave you either with a bunch of class files or (if you used maven to build the project) with an executable Jar file. If you want to use maven, type

```
mvn package
```

**Exercise 3.** Now create a file called "plain" in the `itsec-secret-key-crypto` directory, containing some contents of your choice. Warning: This file will be the target for the simulated ransomware, and will therefore be removed by the program, so do not use any data here that you want to keep and of which you have no backups!

Here, we create a small text file and verify that it is there:

```
$ echo "Hi, I am a quite small plaintext but not that small" > plain
$ ls -l
total 256
-rw-r--r-- 1 neut staff      80 Oct 12 15:00 itsec-secret-key-crypto.iml
-rw-r--r-- 1 neut staff 118752 Oct 15 14:19 mystery
-rw-r--r-- 1 neut staff    52 Oct 15 15:46 plain
-rw-r--r-- 1 neut staff  2156 Oct 15 10:50 pom.xml
drwxr-xr-x 4 neut staff   128 Oct 12 15:00 src/
drwxr-xr-x 9 neut staff   288 Oct 12 16:15 target/
```

**Exercise 4.** Now we run the ransomware (called Ransom) with the command-line arguments “plain” and “cipher”. For example, if you have built the program with maven, you should type:

```
java -cp target/secretkey-1.0-SNAPSHOT.jar ch/zhaw/its/lab/secretkey/Ransom -ransom
plain cipher
```

You should see something like this:

```
MUAHAHAHAHAHAHA! The original file "plain" is gone,
the key is also gone. The encryption algorithm is AES.
Now you must pay $$$ to get the files back! Resistance is futile!
```

We can verify that the file “plain” no longer exists, but that there is now a file called “cipher”:

```
$ ls -l
total 256
-rw-r--r--  1 neut  staff      80 Oct 15 15:48 cipher
-rw-r--r--  1 neut  staff      80 Oct 12 15:00 itsec-secret-key-crypto.iml
-rw-r--r--  1 neut  staff 118752 Oct 15 14:19 mystery
-rw-r--r--  1 neut  staff  2156 Oct 15 10:50 pom.xml
drwxr-xr-x  4 neut  staff   128 Oct 12 15:00 src/
drwxr-xr-x  9 neut  staff   288 Oct 12 16:15 target/
```

That file certainly looks encrypted:

```
$ cat cipher
\???.a?^w??! ?e?L1Cq}AB'M?\f??lQ?
(This output will look different on your terminal.)
```

Luckily for you, the attackers left their source code behind. Perhaps they made a mistake that enables you to recover the key and hence your original files?

## Theoretical Groundwork

**Exercise 5.** The cipher that the attackers used is simply “AES” (with no extensions to the name). Look up AES’s key length. Assuming the key was chosen uniformly at random, what is that key’s entropy? Show and explain your work.

128,192,256 bit therefore entropy =  $2^{128} \cdot 2^{192} \cdot 2^{256}$

**Exercise 6.** What is the work factor for this key length and entropy? Show and explain your work.

Work factor if all combinations are as likely is  $2^{255}$

**Exercise 7.** Make a reasonable assumption about how fast your computer can try a single key on a file of 512 bytes. You should do this by looking up your computer’s clock speed and then looking up approximate values for “cycles per byte” for AES on your CPU. Cycles per byte is a standard performance figure for cryptographic speed. Then compute how long that computer would need to exhaust the work factor. **Important note:** If you can’t find the cycles per byte for your particular computer, use the numbers for *any* modern CPU. If you’re stuck, an answer can be found on Wikipedia. Even then, absolute precision is not essential, order-of-magnitude estimates will suffice. If you still don’t know which of the numbers to choose, choose the one that means that your CPU is faster (higher cycles per second, lower cycles per byte).

```

Clock Speed intel I7-11657G: 2.8 GHz - 4.8 GHz
136 cycles per byte
256 bit per key = 32 byte
32*136 = 4352 cycles per key
4352 / 2.8 GHz * 2^255 = 8.9e70 seconds
4352 * 2^255 / 4.8e9 = 5.2.e70 = 9e61 Years

```

**Exercise 8.** As the previous calculation shows, you cannot hope to break the key by brute force alone. What assumptions would have to change so that you have a good chance of recovering the key nevertheless?

By knowing the some of the possibilities in the key, which then reduces the workforce.

**Exercise 9.** Compute the maximum entropy that the key could have so that your computer (under the assumptions from above) could break the key in 1 minute.

```

60s = entropy / key per sec
Therefore entropy can maximally be around 2^32

```

## Analysing the Code

If we look at the code, we find that the same program that so dastardly encrypted and deleted our files can also run in a decryption mode where we enter the key on the command line:

```
java -cp target/secretkey-1.0-SNAPSHOT.jar ch.zhaw.its.lab.secretkey.Ransom -pay
cipher decrypted 00112233445566778899aabbccddeeff
```

This attempts to decrypt the file `cipher` into the file `decrypted` using the key `001122...eeff`. If we could find the key, we could decrypt our file, even without paying the ransom! If we use this very key on the file `cipher`, we unfortunately get an exception:

```
Trying with key 00112233445566778899aabbccddeeff
Exception in thread "main" javax.crypto.BadPaddingException: Given final block not
properly padded. Such issues can arise if a bad key is used during decryption.
```

The program has started to decrypt the file `cipher` into `decrypted` but didn't finish the last block because it sensed that something wasn't right. This is one indication that this is not the right key. The program has left us with a partial decrypt in `decrypted`, which we can look at with, e.g., hexdump:

```
$ hexdump -C decrypted
00000000 d6 0f bb 59 d6 2a 1f 0a eb 79 93 ab d7 7e ea 69 |...Y.*...y...~.i|
00000010 31 f3 b4 db 4e d3 f1 56 ac 34 20 e5 81 a5 9b de |1...N..V.4 .....|
00000020 57 dd c9 e0 68 9c b4 16 bd 4c f1 35 c8 13 70 89 |W...h....L.5..p.|
00000030
```

This looks quite random and is definitely not our original file. This is not surprising, since the key we used for the decryption, 001122...eeff, was in all likelihood not the one that was used for encryption. When using AES, decrypting a ciphertext with the wrong key will yield random-looking plaintext. This is another indication that we haven't found the right key.

### Estimating Entropy

Recall from the slides how we defined entropy. For this part of the lab, we will slightly deviate from that definition and treat a message (e.g., a text file) as a bag of bytes. If in a file  $F$  of length  $n$  bytes, byte  $k$  ( $0 \leq k < 256$ ) appears  $n[k]$  times, we define  $f[k]$  as the relative frequency of byte  $k$ , i.e.,  $f[k] = n[k]/n$ . Then the *entropy per byte* of the file  $F$  is defined as

$$H(F) = - \sum_{k=0}^{255} f_k \log_2 f_k.$$

Should some  $f[k]$  be zero, we don't include that term in the sum (which otherwise would be infinite).

**Exercise 10.** Just as a warm-up, the Java library only has routines to compute the logarithm to base  $e$  (`java.math.log()`) and the logarithm to base 10 (`java.math.log10()`). Write a code fragment that assigns to `y` the logarithm to base 2 of `x`. Assume `x` and `y` are of type `double`.

```
public double logBase2(double x) {  
    return (java.math.log10(x)/java.math.log10(2));  
}
```

**Exercise 11.** If a file consists only of zero bytes, what is the entropy-per-byte? Show your work.

0 since the possibility is 100% that the byte is zero

**Exercise 12.** If a large file consists of uniform random bytes, what is the entropy per byte? Show your work.

$n * ((1/n) * \log_2(1/n))$  whereas  $n$  is the number of bytes?

**Exercise 13.** Use the previous two questions to find the minimum and maximum entropy per byte in a large file. Show and explain your work.

$\min 0; \max = n * ((1/n) * \log_2(1/n))$  whereas  $n$  is the number of bytes

**Exercise 14.** Write a program that takes on its command line a list of file names and that outputs, for each file, the computed entropy per byte. For example, for the file `mystery`, a file that was part of the repository, and `FileEncrypter.java`, one of the source files of the ransomware, your program should produce the following output:

```
mystery: 7.998687714859995  
src/main/java/ch/zhaw/its/lab/secretkey/FileEncrypter.java: 4.82373026044958
```

**Exercise 15.** From looking at the output of this program, which of the two files do you think contains “more random” data? Assuming that both files had been produced by the ransomware in decryption mode, which one is more likely to contain actual plaintext? Explain.

The one with the higher entropy value contain more random data  
The lower entropy value should contain plaintext

**Exercise 16.** From the preceding exercises, you should now have a good idea how to distinguish correctly decrypted text from wrongly decrypted text, provided that the original plaintext is in a natural language. Write a program that takes as input a `byte[]` and that outputs `true` if the input is likely to be a natural-language text, and `false` otherwise. Download some examples of natural-language text from the Internet and test your program. Also create some files with random data and also your program with those as well. On Unix-like systems, the special file `/dev/urandom` provides you with a convenient source of random numbers on the command line.

**Exercise 17.** Extend your program so that you use (by copy-and-paste) the `decrypt()` routine and all needed helper routines from the `FileEncrypter` class in the ransomware to do the following:

- It accepts as input a `byte[]` containing ciphertext, and a `byte[]` containing a key
- It attempts to decrypt the ciphertext
- It returns `true` if it thinks that the decrypted output is in some natural language

Now you have the basics for a program that allows you to automatically try keys until you have found one that makes the input decrypt to a natural-language plaintext.

## Cryptanalysis of the Ransomware

We have found out in the earlier sections that if the AES key is chosen randomly, we have no practical chance to recover the key by brute-force, i.e., by trying all keys until one fits. Therefore, there are only two possible ways of getting our files back: either the programmer made some security-relevant mistake while coding up the crypto in the program, or the keys aren’t as random as they should be.

We can tell you now that the implementation of the crypto is solid, there are no security-relevant mistakes to be found there. The only chance therefore is that the keys are perhaps not random enough.

As the code in `FileEncrypter` class reveals, the ransomware program uses AES in a mode known as CBC (we will discuss this in the lecture later). For this variant of AES, CBC mode splits the plaintext up into blocks of 128 bits, or 16 bytes:  $P[0]$ ,  $P[1]$ , ... But because of the way it is constructed, CBC needs an artificial zeroth block,  $P[0]$ . This block is known as the *initialization vector*, or IV. This IV is provided by the program (*FileEncrypter.java*, lines 46—50). As you can see, it is generated by a random number generator, just like the key, and is put in front of the ciphertext. So the first 16 bytes of the ciphertext should be just as random as the rest of the ciphertext file. Let's check that with hexdump:

```
$ hexdump -C mystery | head -5
00000000  98 0b ab 77 66 01 00 00  00 00 00 00 00 00 00  |...wf.....|
00000010  7e 1c 58 a0 d8 9a a5 bc  47 57 41 9b 82 e4 08 50  |~.X.....GWA...P|
00000020  ec 6e 79 8c 11 e6 db 48  6a fd 44 35 3a 16 b1 43  |.ny....Hj.D5:...C|
00000030  79 ff 69 a1 c9 88 c8 0b  52 36 09 89 7d ac 89 62  |y.i.....R6...}..b|
00000040  8f 66 12 17 d9 8a 0b 07  29 55 98 27 05 bd a9 3e  |.f.....)U.'...>|
```

Even if it only concerns the IV, which is not secret, *that* is very interesting indeed! Instead of 16 random bytes in positions 0 to 15, fully ten of those 16 bytes are zero! Something is rotten in the state of Denmark!

**Exercise 18.** Look at the code and find out how the IV is generated, exactly. How did it happen that ten out of the sixteen bytes in the IV are zero?

Since after every Byte the random number (time in ms) is shifted 8 bits to the right, after a few loops there is only zeros.

Well, now we've found out why the IV wasn't as random as perhaps it should have been. Unfortunately, this doesn't mean anything, since we only have the one mystery file, and as long as IVs aren't repeated, nonrandom IVs do not affect security (in this use case).

**Exercise 19.** Or do they? Look at the code again and find a flaw, this time a real flaw. This should give you an idea of how to extend the program you wrote above so that you can now indeed find the key, given only the ciphertext file. Explain the flaw and then write the program.

the same function is used to generate the secret key as well as for the IV, which is the current time is ms, and the time only differs a few MilliSeconds for the key and the IV so it should be possible to recreate the key by trying all times in between.

**Exercise 20.** Given your new knowledge, estimate the entropy of the key (in bits, not bits per byte). An estimate suffices, absolute precision is not required.

if we assume the time between the secret key time and IV time is 1000 ms the entropy should be around 10 bits:  $\log_2(1000)$ . As the program possibly didn't run that slow, it is most likely lower than 10 bits.

This should make it clear that the entropy of messages (or key material) depends on what we know about that message (or key material). Here, our knowledge about the key material reduced the entropy from 128 bits, which is impossible to crack, to something even a student's laptop could handle.

### Putting it All Together

**Exercise 21.** Find the key that was used to encrypt the file `mystery` and decrypt it. What is the key? What is the text? Who wrote it? In what work does it appear, and where?

## Points

In this lab, you can get up to **2 points**, by showing your instructor the program you wrote and how it hangs together, with a demonstration that it actually finds the key for the `mystery` file and correctly decrypts it.