

## HW3 程式題

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Way1: 使用迴圈遍歷

```
5   long long int count_inver(int *a, int a_size){
6       long long int inversion = 0;
7       for(int i = 0; i < a_size; ++i){
8           for(int j = i+1; j < a_size; ++j){
9               if(a[i] > a[j])
10                  inversion += 1;
11           }
12       }
13       return inversion;
14   }
```

每次都看第  $i$  個元素後面的 element 有沒有比他小，有的話就是一個 inversion，時間複雜度  $O(n^2)$

Way2: modify merge sort

```
8   ✓ long long int conut_inversion(int *a, int a_size){
9       int * temp = (int*)malloc(sizeof(int) * a_size);
10      return merge_aux(a, temp, 0, a_size-1);
11   }
```

```

13 long long int merge_inver(int *a, int *temp, int p, int q, int r){
14     int i = p; // p to q-1
15     int j = q; // q to r
16     int k = p;
17     long long int inversion = 0;
18
19     while(i <= q-1 && j <= r){
20         if(a[i] <= a[j]){ // the left one is smaller
21             temp[k++] = a[i++];
22         }
23         else{ // the right one is smaller
24             temp[k++] = a[j++];
25             inversion = inversion + (q - i); // this one must go across the whole left subarray
26         }
27     }
28
29     while(i <= q-1) // copy the left elements from left subarray to temp
30         temp[k++] = a[i++];
31
32     while (j <= r) // copy the left elements from right subarray to temp
33         temp[k++] = a[j++];
34
35     for(i = p; i <= r; i++) // copy the data from temp back to a, then a is sorted
36         a[i] = temp[i];
37
38     return inversion;
39 }

```

```

41 long long int merge_aux(int *a, int *temp, int p, int r){
42     // array a, index from p to r
43     long long int inver = 0;
44     int q;
45     if(p < r){
46         q = (p+r)/2; // middle index, separate array into two parts, left and right
47         inver += merge_aux(a, temp, p, q); // index from p to q
48         inver += merge_aux(a, temp, q+1, r); // index from q+1 to r
49         inver += merge_inver(a, temp, p, q+1, r); // merge means combine left and right subarray
50     }
51     return inver;
52 }

```

遵守 merge sort 的精神，把 array 分兩段，合併階段時如果左邊 subarray 的元素大於右邊 subarray 的元素，就更新 inversion 的數量。

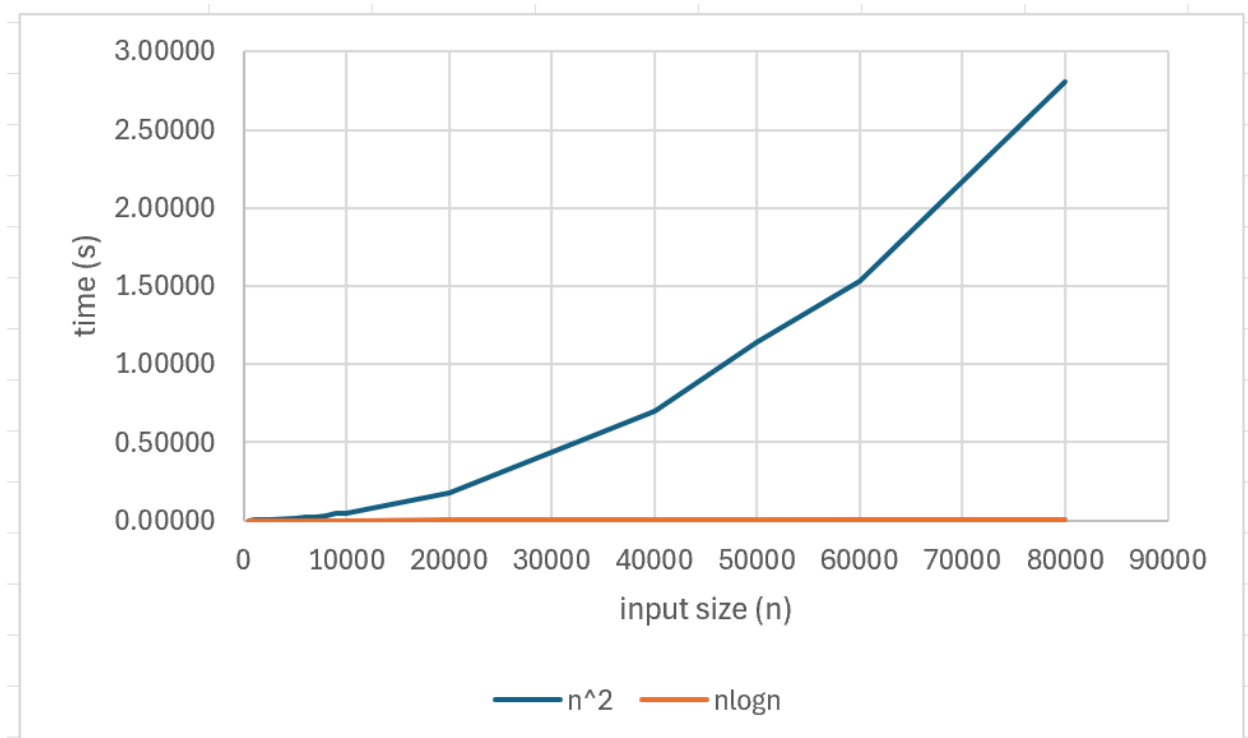
時間複雜度  $O(n \log n)$

測試環境

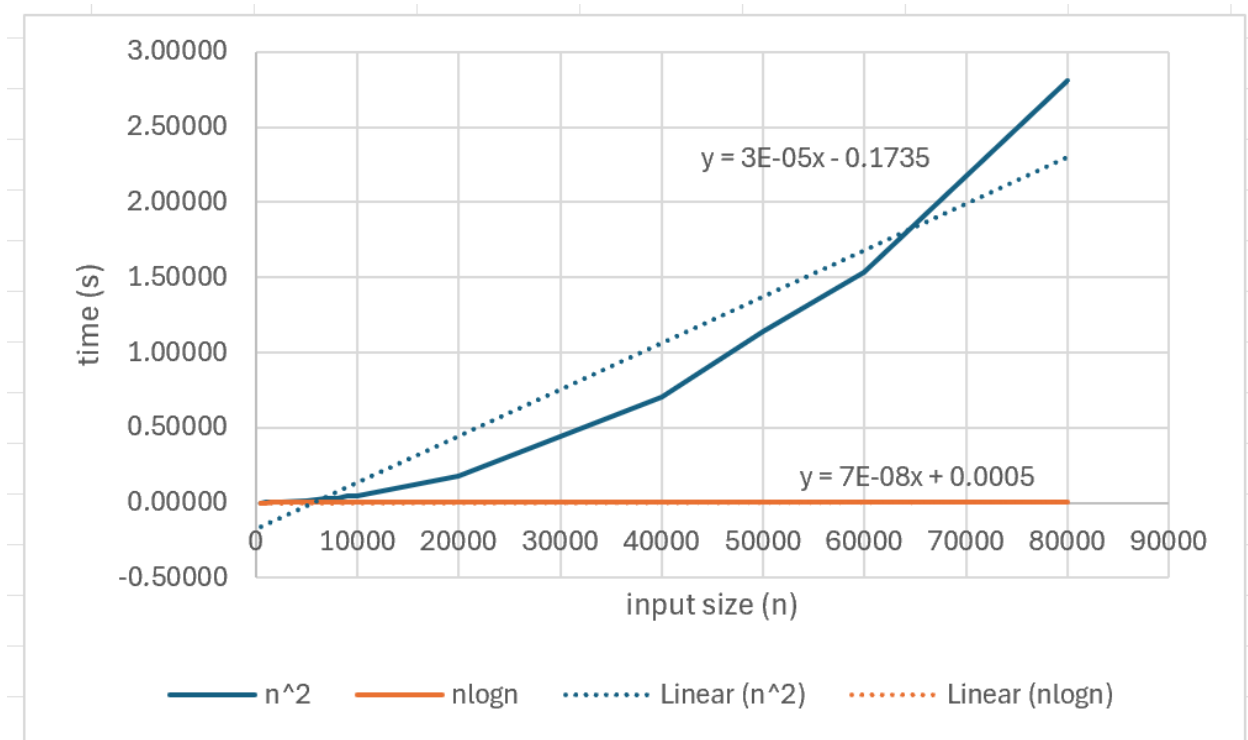
CPU: 13<sup>th</sup> Gen Intel® Core™ i7-1360P

Operating System: Windows 11

Compiler: gcc 13.2.0



一、用 excel 畫散佈圖



## 二、顯示漸進線

解聯立方程式

$$100000000y = 3000x - 17350000$$

$$100000000y = 7x + 50000$$

$$x = 5,813.564984964918$$

約在  $n=5813$  左右開始， $n^2$  的方法和  $n\log n$  的方法有顯著的時間差距