

Bit Errors

- * Reliable packet delivery.
- * Forwarding & Routing.

Reliable Delivery

1. No packets are corrupted.
2. All packets are delivered.
3. In order

Bit Error Rates.

BER ~ independent, random errors

- bursty, unpredictable.
- corrupt contiguous bits during burst.
- greater bit rate.

$10^{-6} / 10^{-7} / 10^{-8}$ BER rate

FER. $p = 10^{-6}$. $L \rightarrow \text{length}$.

$$P_{\text{Correct}} = (1-p)^L$$

$$P_{\text{FER}} = 1 - (1-p)^L. \quad \text{BER} = 10^{-6}, \quad L = 12,000 \text{ bits.}$$

$$\text{FER} = 1.19\%$$

reduce L shorter frames \rightarrow more headers / overhead.

reduce $p \rightarrow$ error control encoding

Strategy.

1. detect error / checksum.
2. respond to error
 - error correction code.
 - resend.
 - discard

Error Detection & Correction Code EDC.

PHY. Parity Bit, even number of bits

Two-Dimension Parity

$d_{1,1}$	\dots	$d_{1,n}$	P_1
$d_{2,1}$	\dots	$d_{2,n}$	P_2
$d_{3,1}$	\dots	$d_{3,n}$	P_3
$d_{4,1}$	\dots	$d_{4,n}$	P_4
q_1	\dots	q_n	r

9/25
16/25. error checking bits

Hamming (7, 4) code.

D₁D₂D₃D₄, P₁ P₂ P₃

L2.

CRC quickly filters corrupted frames

- detects single bit errors $g(x)$ two non-zero terms
- burst error $\leq g-1$. $g(x)$ start with x^{g-1} ends 1.
- Any odd number of zeros $g(x)$ even number of coefficients

L2 detects most errors before applications

L2 with CRC most errors, especially across links

L3

IP checksum.

Only check header.

check in a multiple host env.
right packet, right machine.

Transport checksum.

end-to-end at L4.

right packet
right destination

from sender
to receiver

E2E communication

Why?

Intermediate device in the Internet

introduce errors, L2, L3 headers can be replaced. / introduce errors

Less frequent, additional checksums

Space efficient: 16-bit all the way.

recalculating fast: $C = n(n + nm + m')$.

Guarantee: Single-bit error detectable.

FEC. Error correction | very expensive to
achieve low Frame Loss Rate.

Some Errors are too severe.
(bursts)

CRC32 + IP + TCP checksums.

Practical, enough FEC to put Loss Rate under 10^{-3} .

Reliable Transmission.

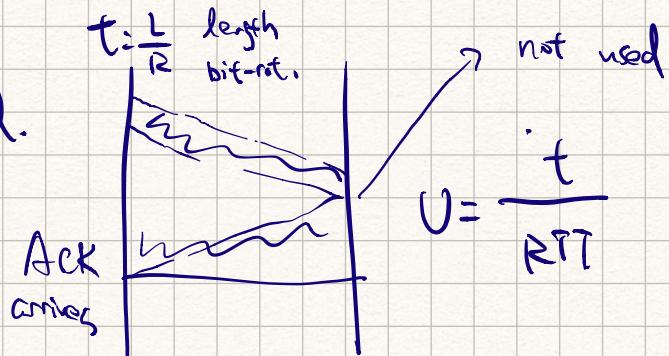
- corrupt packets
- can lose packets
- delays / reorder packets

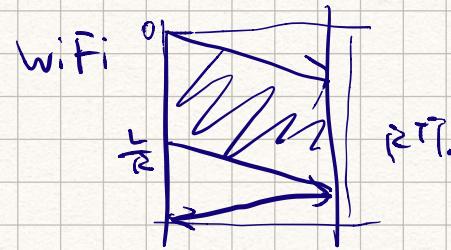
Design: Stop and wait protocol.

- send a packet
- wait until ACK or RTX expires

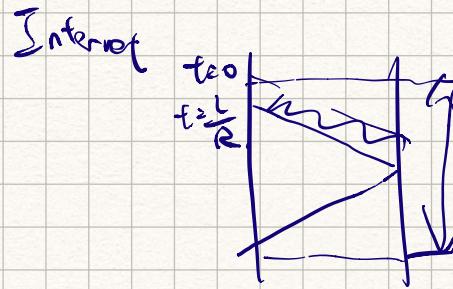
Sequence number should be large.
32-bit for TCP.

Stop-Wait Protocol.



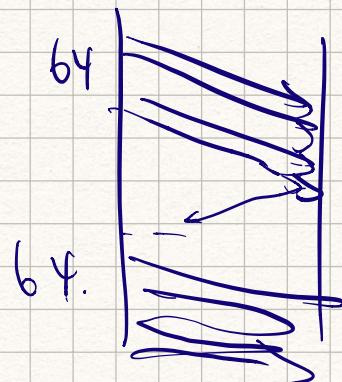
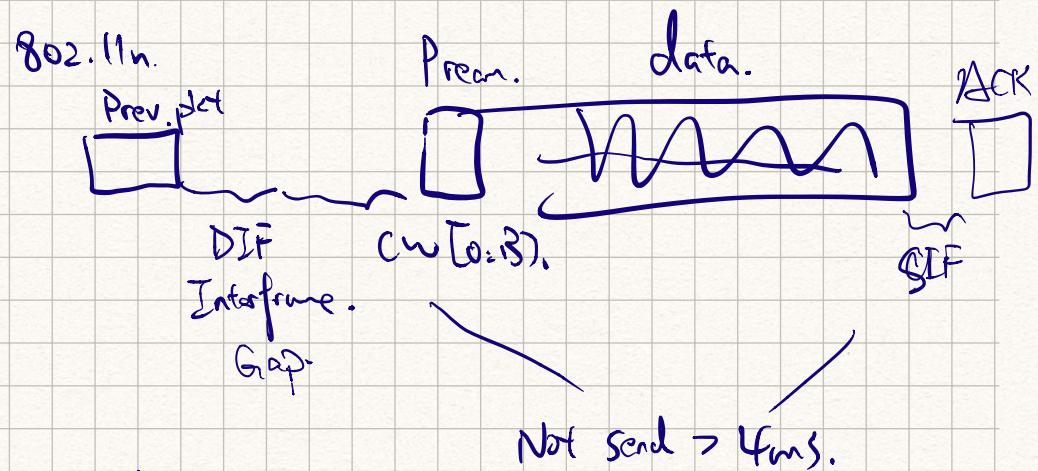


99.3 %

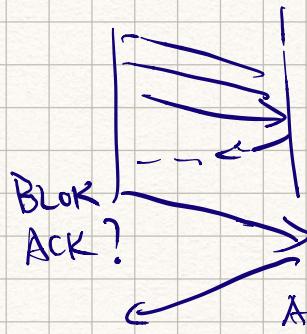


2.1 %

Aggregate Packet on WiFi



expensive to do
full resend.



BLOK
ACK?

ACK 1~64.

Reliability + Efficient across Internet.

Pipelining : Sender sends multiple packets at a time.

Go Back N.

Cumulative ACK , resend from the last ACK

- not great to discard all now packets

Pipelining with Selective ACK