## Computer Networks Lab3

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## Writing Task (30%)

How do you estimate the bandwidth and the delay of a path?

My implementation estimate the bandwidth by sampling the data size consumed each time and their time intervals, counting the bandwidth in the latest time interval and smoothing the estimated bandwidth by weighted mean. To specify, the sampling happens in <a href="lib/common/rangeclient.c:on\_body">lib/common/rangeclient.c:on\_body</a>.

As for delay, we could make use of the <a href="Ping Frame">Ping Frame</a> in HTTP2 specification.

lib/common/http2client.c has provided us with the prototype callback function handle\_ping\_frame. To handing the received ping frame, We just need to register this callback function in <a href="lib/common/htttp2client.c:setup\_steam">lib/common/htttp2client.c:setup\_steam</a>. Moreover, in order to send a ping frame periodically, we could register a timer under the event-driving loop in <a href="lib/common/http2client.c:h2o\_httpclient\_h2\_on\_connect">lib/common/http2client.c:h2o\_httpclient\_h2\_on\_connect</a>.

How do you assign jobs to the three paths?

In my implementation, the main routine of rescheduling is summarized as following:

- find one idle path
- find the path with the longest estimated remaining time(based on the bandwidth and delay we have calculated) in the other two paths
- reschedule these two path according to their bandwidth ratio. If the bandwidth of either one is unknown, just reschedule these two paths averagely.

What features (pipelining, eliminating tail byes, etc.) do you implement? And how do you implement them? You are encouraged to write down other design aspects of your implementation.

I implement http pipelining. For a http connection, when the size of remaining unreceived data is less than or equal to that of the received data in the latest frame, we could assert that this http connection is going to complete soon. At this moment, we could reschedule this path in advance, instead of rescheduling when this connection is eventually closed. Pipelining make senses especially under the condition of high latency.

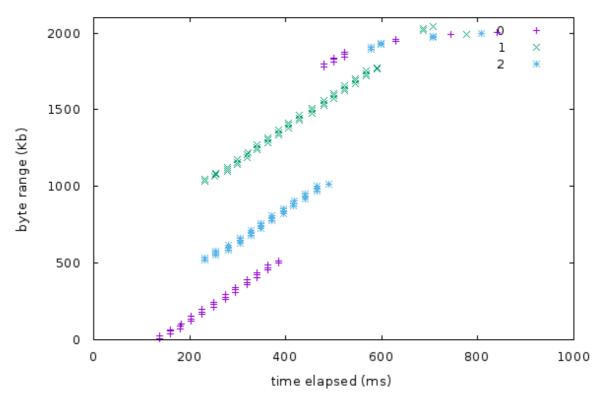
## Drawing Task (40%)

The network topology is defined the same as that in the wirteup.

```
3
0 1 10.100.1
0 2 10.100.2
0 3 10.100.3
```

1. under static network conditions

I use Mahimahi to simulate the static network conditions. Each network interface have 20ms delay.

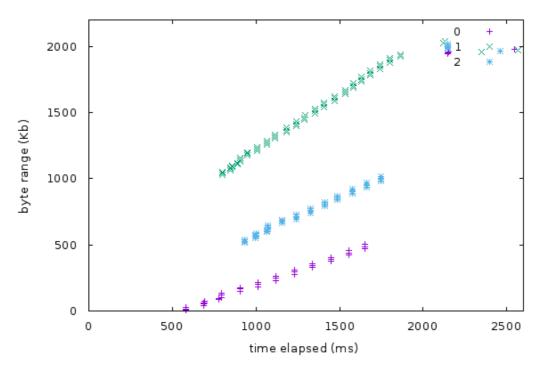


What annoys me is that some Mahimahi commands like mm-link seem to have bugs in my system platform.

```
$ mm-delay 20 mm-link --meter-all /usr/share/mahimahi/traces/Verizon-LTE-
short.up /usr/share/mahimahi/traces/Verizon-LTE-short.down
Died on std::runtime_error: xcb_get_setup: returned null pointer
Died on std::runtime_error: `downlink': process exited with failure status 1
Died on std::runtime_error: `mm-link --meter-all
/usr/share/mahimahi/traces/Verizon-LTE-short.up
/usr/share/mahimahi/traces/Verizon-LTE-short.down': process exited with
failure status 1
Died on std::runtime_error: `packetshell': process exited with failure status
```

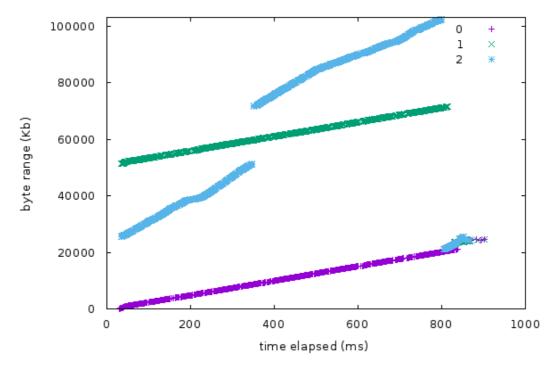
Then I tried to use to simulate changeable network conditions as following:

sudo tc qdisc add dev veth0-1 root netem delay 100ms 10ms 30%
sudo tc qdisc add dev veth0-2 root netem delay 50ms 10ms 30%
sudo tc qdisc add dev veth0-3 root netem delay 75ms 10ms 30%



3. with a high-bandwidth path and two low-bandwidth paths network configuration:

sudo tc qdisc add dev veth0-1 root tbf rate 1mbit burst 32kbit latency 50ms
sudo tc qdisc add dev veth0-2 root tbf rate 1mbit burst 32kbit latency 50ms
sudo tc qdisc add dev veth0-3 root tbf rate 10mbit burst 32kbit latency 50ms



4. with a short-delay path and two long-delay paths network configuration: sudo tc qdisc add dev veth0-1 root netem delay 10ms
sudo tc qdisc add dev veth0-2 root netem delay 3ms
sudo tc qdisc add dev veth0-3 root netem delay 10ms

