UNIVERSITY OF OSLO

COMPUTATIONAL PHYSICS

Project 2



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CHAPTER

1

INTRODUCTION

The aim of this project is to solve Schrödinger's equation numerically, using Jacobi's method, for a single electron with zero angular momentum in a harmonic oscillator potential, as well as for two electrons in a three dimensional harmonic oscillator well, both with and without Coulomb interaction. To solve Schrödinger's equation using Jacobi's method, it is first reformulated into an eigenvalue problem.

The Jacobi algorithm is implemented in c++, and the source codes developed in this project and selected results, can be found in the GitHub repository: https://github.com/birgimad/Project2_BM_MI_SC.

In order to ensure the credibility of the algorithm, various tests are run. Amongst these tests are solving the considered eigenvalue problem for a simple 2×2 case to check the correctness of the computed solution to the known values, and a comparison of the computed eigenvalues to the analytical solution.

The characteristics of the algorithm are, furthermore, examined by finding the optimal number of steps and interval that gives the best value of the lowest three eigenvalues for the single particle case. As a part of this characterization, the influence of number of steps on the number of iterations in the Jacobi method is investigated, and the consequence of changing the size of the considered interval and the number of steps both independently and dependently of each other is discussed. Furthermore, the computed Jacobi algorithm is compared to the Armadillo function for solving eigenvalue problems.

The report mainly consists of two sections. First section discusses the nature of the problem, the functionality of the algorithm and also the various tests on the algorithm. The last section is about the results, its interpretation and discussions.

2

METHOD

We present in this chapter a short discussion on the nature of the problem. The problem is to solve Schrödinger's equation numerically using the Jacobi method for an eigenvalue problem. Hence, the strategy is to reformulate Schrödinger's equation into an eigenvalue problem, which is done in Sec. 2.1. In the Jacobi method, a series of similarity transformations are made, from which a new matrix is created. The values of the elements of the new matrix are related to the angle at which the similarity matrix rotates the initial matrix, and how to choose this angle optimally is presented in Sec. 2.1.1.

The source code for the programmed Jacobi method can be found in the GitHub folder https://github.com/birgimad/Project2_BM_MI_SC, and in Sec. 2.2 the computed algorithm is tested in different ways to ensure the correctness of the code.

2.1 Nature of the problem

The aim of the first part of the project is to solve Schrödingers equations for one electron in a harmonic oscillator potential with angular momentum l=0. The radial part of the Schrödingers equation is considered which is as follows

$$\left[-\frac{\hbar^2}{2m} \frac{1}{r^2} \frac{d}{dr} r^2 \frac{d}{dr} + V(r) \right] R(r) = ER(r). \tag{2.1}$$

In order to solve this equation numerically, it is rewritten after a series of transformation and substitution as

$$-\frac{d^2}{d\rho^2}u(\rho) + \rho^2 u(\rho) = \lambda u(\rho)$$
(2.2)

Eq. (2.2) is discretized by writing the second derivative of $u(\rho)$ as

$$\frac{d^2}{d\rho^2}u(\rho) = \frac{u(\rho+h) - 2u(\rho) + u(\rho-h)}{h^2} + O(h^2)$$
(2.3)

In Eq. (2.3) h is the step length, and ρ_{max} and ρ_{min} are the maximum and minimum values of the variable ρ , respectively. For a given number of steps n, the step length is given as

$$h = \frac{\rho_{max} - \rho_{min}}{n} \tag{2.4}$$

In order to solve equation Eq. (2.2), it is transformed into a matrix eigenvalue problem

$$\mathbf{A}\mathbf{u} = \lambda \mathbf{u} \tag{2.5}$$

in which A is a tridiagonal matrix of the form

$$\mathbf{A} = \begin{pmatrix} \frac{2}{h^2} + V_1 & -\frac{1}{h^2} & 0 & 0 & \dots & 0 & 0\\ -\frac{1}{h^2} & \frac{2}{h^2} + V_2 & -\frac{1}{h^2} & 0 & \dots & 0 & 0\\ 0 & -\frac{1}{h^2} & \frac{2}{h^2} + V_3 & -\frac{1}{h^2} & 0 & \dots & 0\\ \dots & \dots & \dots & \dots & \dots & \dots\\ 0 & \dots & \dots & \dots & \dots & \dots\\ 0 & \dots & \dots & \dots & \dots & \frac{2}{h^2} + V_{n-2} & -\frac{1}{h^2}\\ 0 & \dots & \dots & \dots & \dots & -\frac{1}{h^2} & \frac{2}{h^2} + V_{n-1} \end{pmatrix}$$

$$(2.6)$$

A is obtained from Eq. (2.2), with the approximation of the derivative of $u(\rho)$ given in Eq. (2.3) when omitting all later terms, by discretizing ρ by

$$\rho_i = \rho_{min} + ih \qquad i = 0, 1, 2, \dots, n \tag{2.7}$$

This leads to the following Schrödinger equation:

$$-\frac{u(\rho_{i}+h)-2u(\rho_{i})+u(\rho_{i}-h)}{h^{2}}+\rho_{i}^{2}u(\rho_{i})=\lambda u(\rho_{i})$$
(2.8)

which can be rewritten as

$$-\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} + \rho_i^2 u_i = -\frac{u_{i+1} - 2u_i + u_{i-1}}{h^2} + V_i u_i = \lambda u_i$$
(2.9)

in which $V_i = \rho_i^2$ is the harmonic oscillator potential. When comparing this relation with the general eigenvalue problem in Eq. (2.5), it is evident that the diagonal elements of the matrix **A** is given by

$$d_i = \frac{2}{h^2} + V_i \tag{2.10}$$

while all off diagonal elements are zero apart from those neighbouring the diagonal, which are all constants with the value

$$e_i = -\frac{1}{h^2} \tag{2.11}$$

This is exactly what is given in Eq. (2.6).

2.1.1 Change in Matrix Elements after Iterations and Choice of θ

The algorithm for solving the eigenvalue problem given in Eq. (2.5) contains of multiple similarity transformations of the matrix \mathbf{A} , in which we assume a_{kl} to be the largest off-diagonal element. The matrix \mathbf{B} constructed by the similarity transformation is given by

$$\mathbf{B} = \mathbf{S}^T \mathbf{A} \mathbf{S} \tag{2.12}$$

in which S is an orthogonal transformation matrix with its non-zero matrix elements:

$$s_{kk} = s_{ll} = \cos \theta$$

$$s_{kl} = -s_{lk} = -\sin \theta$$

$$s_{ii} = 1, \quad i \neq k, i \neq l$$

After matrix multiplication with the orthogonal transformation matrix S and its transverse (as in Eq. (2.12)) the entrances of B becomes

$$b_{ii} = a_{ii}, i \neq k, i \neq l$$

$$b_{ik} = a_{ik} \cos \theta - a_{il} \sin \theta, i \neq k, i \neq l$$

$$b_{il} = a_{il} \cos \theta + a_{ik} \sin \theta, i \neq k, i \neq l$$

$$b_{kl} = a_{kl} \cos^2 \theta - 2a_{kl} \cos \theta \sin \theta + a_{ll} \sin^2 \theta$$

$$b_{ll} = a_{ll} \cos^2 \theta + 2a_{kl} \cos \theta \sin \theta + a_{kk} \sin^2 \theta$$

$$b_{kl} = (a_{kk} - a_{ll}) \cos \theta \sin \theta + a_{kl} (\cos^2 \theta - \sin^2 \theta)$$

Due to the symmetry in Eq. (2.12) with **A** being a tridiagonal symmetric matrix, $b_{lk} = b_{kl}$, $b_{ki} = b_{ik}$, and $b_{li} = b_{il}$.

Since θ can be chosen arbitrarily, we choose θ to be the angle at which b_{kl} , and hence b_{lk} , becomes zero. In this way, the largest element of **A** is eliminated, and it can be shown that this choice of θ reduces the norm of the off-diagonal elements of **A**, which ensures that the algorithm terminates towards the eigenvalues.

This yields the equation

$$0 = (a_{kk} - a_{ll})\cos\theta\sin\theta + a_{kl}(\cos^2\theta - \sin^2\theta)$$
(2.13)

By introducing $\tan \theta = \sin \theta / \cos \theta$ and the quantity

$$\tau = \frac{a_{ll} - a_{kk}}{2a_{kl}} \tag{2.14}$$

Eq. (2.13) can be rewritten as the quadratic equation in $\tan \theta$

$$\tan^2\theta + 2\tau\tan\theta - 1 = 0\tag{2.15}$$

which has the solutions

$$\tan \theta = -\tau \pm \sqrt{1 + \tau^2} \tag{2.16}$$

From the solutions for $\tan \theta$ given in Eq. (2.16), $\cos \theta$ and $\sin \theta$ can be found using the formulas

$$\cos \theta = \frac{1}{\sqrt{1 + \tan^2 \theta}}$$
 and $\sin \theta = \tan \theta \cos \theta$

If $\tau < 0$, tan θ is chosen to be

$$\tan \theta = -\tau - \sqrt{1 + \tau^2} \tag{2.17}$$

whilst if $\tau \geq 0$, $\tan \theta$ is calculated as

$$\tan \theta = -\tau + \sqrt{1 + \tau^2} \tag{2.18}$$

This choice is made to always make $\tan \theta$ the smaller of the two roots given in Eq. (2.16). Furthermore, this choice ensures that $|\tan \theta| \le 1$, yielding that $|\theta| \le \pi/4$.

This is true since $|\tau| \le 1$, because $|a_{kl}| \ge |a_{ij}|$ for all i, j, from which it follows that

$$|\tan \theta| = \left| -\tau - \sqrt{1 + \tau^2} \right| \le 1, \quad \text{for } \tau < 0$$
 (2.19)

and

$$|\tan \theta| = \left| -\tau + \sqrt{1 + \tau^2} \right| \le 1, \quad \text{for } \tau \ge 0$$
 (2.20)

since $\sqrt{1+\tau^2} \le \sqrt{2}$.

The fact that $|\theta| \le \pi/4$ ensures that $\cos \theta \ge 0$ which ultimately ensures that the difference between **A** and the new matrix **B** is minimized, since

$$||\mathbf{B} - \mathbf{A}||_F^2 = 4(1 - \cos\theta) \sum_{i=1, i \neq k, l}^n (a_{ik}^2 + a_{il}^2) + \frac{2a_{kl}^2}{\cos^2\theta}.$$
 (2.21)

When implementing the elements of **B** in the algorithm, the following source code is used:

```
// Computing the new matrix A
       for (int i = 0; i<n; i++)</pre>
       {
           if (i != row_number && i != column_number) // determining A(i,k) for new
              matrix
           {
              a_ik = A(i,row_number);
              a_il = A(i,column_number);
              A(i,row_number) = a_ik*c - a_il*s;
              A(i,column_number) = a_il*c + a_ik*s;
              A(row_number,i) = A(i,row_number);
              A(column_number,i) = A(i,column_number);
           }
       }
       A(row_number, row_number) = a_kk*pow(c,2) - 2*a_kl*c*s+a_ll*pow(s,2);
       A(column_number, column_number) = a_11*pow(c,2) + 2*a_k1*c*s+a_kk*pow(s,2);
       A(column_number, row_number) = 0.00; //By choice of theta
       A(row_number, column_number) = 0.00; //By choice of theta
```

In the above source code lines a_ij is the ij'th element of matrix \mathbf{A} , whilst A(i,j) is the ij'th element of matrix $\mathbf{B} = \mathbf{S}^T \mathbf{A} \mathbf{S}$. row_number and column_number are the row number and column number of the element of \mathbf{A} with the maximum absolute value. As seen in the source code, the elements $b_{kl} = b_{lk}$ with the largest absolute value are forced to zero by the choice of θ to save computational time.

2.1.2 Off-diagonal norm

Similarity transformations is used to reduce the off-diagonal norm. The norm found by Eq. (2.22), is wanted as small as possible and smaller than a given test value ε . Ideally the norm should get to zero, but that is difficult because when the elements gets small there can be problems with round-off errors.

2.2. Tests University of Oslo

The value ε is therefore set so that it gives the smallest values possible without problems round-off errors, typically set around 10^{-8} .

$$off(\mathbf{A}) = \sqrt{\sum_{i=1}^{n} \sum_{j=1, j \neq i}^{n} a_{ij}^{2}}$$
 (2.22)

This norm is compared after each transformation so that the new matrix \mathbf{B} is as close to a diagonal matrix as possible. But this is a very time consuming approach and requires many calculations. So instead Eq. (2.23) is used as it is less time consuming.

$$max(a_{kl}^2) > \varepsilon \tag{2.23}$$

This is possible because if the biggest element squared is smaller than ε then all other values will be equally small or smaller. Which means that they are so small that the possibility for round-off errors is big. As it is not possible to get ensure that the process further gives correct values the transformation loop should stop.

2.2 Tests

Several tests are run to ensure that the algorithm runs correctly. The two test presented in this section is a test of the computed $find_max$ function to find the off diagonal element with the largest absolute value of a symmetric matrix and a comparison of the computed solution for a simple 2×2 case. In Chap. 3 the results gained for the one particle case using the computed Jacobi algorithm is furthermore compared to the known analytical solution.

2.2.1 Test of the *find_max* function

To find the maximum absolute value of the elements of the matrix for which we want to solve the eigenvalue problem using the Jacobi method, the following c++ source code is used.

```
}
}
return;
```

The programmed function $find_max$ finds the entrance with the maximal absolute value amongst the entrances above the diagonal. The initial guess of the maximum absolute value of the off diagonal elements is set to a_{12} (notice that the first row/column of the matrix in the code is 0, whilst it is 1 in the text). The two for loops then run through all the elements above the diagonal, and if the absolute value of that element is greater than the absolute value of the until then computed maximal value, the new value max is set equal to the value of that entrance. Since the matrix $\bf A$ for this project is symmetric, it is not necessary to run through the elements below the diagonal.

To check that the *find_max* function runs as expected, a random matrix **A**, with the maximum absolute value above the diagonal being $a_{25} = 6$, is considered.

$$\mathbf{A} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 2 & 3 & 4 & 5 & 6 \\ 0 & -1 & -2 & -3 & -4 \\ 0 & -3 & -6 & -9 & 0 \\ -1 & 0 & 1 & 2 & 3 \end{pmatrix}$$
 (2.24)

When running the function for the matrix **A** in Eq. (2.24) and an initial guess that the greatest absolute value can be found as the element $a_{12} = 2$, the function outputs the maximum value:

```
max value = 6

row number = 2

column number = 5
```

which is exactly what was expected when considering the investigated matrix. When running the function for $-\mathbf{A}$ the element with the greatest absolute value is once again found to be a_{25} . In this case the element has a value of $a_{25} = -6$, which was to be expected.

2.2.2 Testing against simple 2×2 case

To check that the Jacobi function runs correctly, consider the case with matrix dimensionality 2×2 case with $\rho_{min} = 0$ and $\rho_{max} = 6$, yielding a step length of

$$h = \frac{\rho_{max} - \rho_{min}}{n+1} = \frac{6-0}{3} = 2$$

With the potential described by ρ^2 , this case gives that the matrix **A**, for which the eigenvalue problem is solved, takes the form

$$\mathbf{A} = \begin{pmatrix} \frac{2}{2^2} + 2^2 & -\frac{1}{2^2} \\ -\frac{1}{2^2} & \frac{2}{2^2} + 4^2 \end{pmatrix} = \begin{pmatrix} 4.5 & -0.25 \\ -0.25 & 16.5 \end{pmatrix}$$

2.2. Tests University of Oslo

The entrance with the greatest absolute value of the off diagonal element is the $a_{12} = a_{21}$, which means that τ introduced in Eq. (2.14)

$$\tau = -\frac{16.5 - 4.5}{2 \cdot 0.25} = -24$$

and hence

$$\tan \theta = 24 - \sqrt{1 + 24^2} \approx -0.0208$$

which yields that

$$\cos \theta = \frac{1}{\sqrt{1 + 48^2}} \approx 0.9998$$
 and $\sin \theta = -0.0208 \cdot 0.9998 \approx -0.0208$ (2.25)

From the values for $\cos \theta$ and $\sin \theta$, the diagonal elements of the constructed matrix **B** after one similarity transformation described in Eq. (2.12) take the form

$$b_{11} \approx 4.5 \cdot 0.9998^2 - 2 \cdot 0.25 \cdot 0.0208 \cdot 0.9998 + 16.5 \cdot 0.0208^2 \approx 4.495$$

 $b_{22} \approx 16.5 \cdot 0.9998^2 + 2 \cdot 0.25 \cdot 0.0208 \cdot 0.9998 + 4.5 \cdot 0.0208^2 \approx 16.51$

giving

$$\mathbf{B} \approx \left(\begin{array}{cc} 4.495 & 0 \\ 0 & 16.51 \end{array} \right)$$

Which means that the first and second eigenvalues are 4.495 and 16.51, respectively. When running the computed Jacobi function for this 2×2 example, this is exactly what is gained. (see GitHub: https://github.com/birgimad/Project2_BM_MI_SC/blob/master/Results/Simple2x2case.txt)

3

RESULTS AND DISCUSSION

When running the computed Jacobi method for both the single particle in a harmonic oscillator case and the two particles in a harmonic oscillator potential with varying oscillation frequencies both with and without a repulsive coulomb potential, it is found that the computed eigenvalues vary when changing ρ_{max} and n, which was to be expected since the step length h introduced in Eq. (2.4) depends on both of these factors.

Sec. 3.1 and 3.2 are dedicated to dealing with the single particle case, whilst Sec. 3.3 the probability distributions for the two particles case are plotted for different oscillation frequencies after finding acceptable values for ρ_{max} and n.

Some selected results from running the code can be found in the GitHub folder https://github.com/birgimad/Project2_BM_MI_SC/tree/master/Results.

3.1 Dependence of ρ_{max} and n on Eigenvalue

For a single electron moving in a three-dimensional harmonic oscillator potential, the analytical solution for the first three eigenvalues to the rewritten Schrdinger's equation stated in Eq. (2.2) are $\lambda_0 = 3$, $\lambda_1 = 7$, and $\lambda_2 = 11$, for l = 0.

In the computed code for solving the eigenvalue problem numerically using the Jacobi method, the two parameters ρ_{max} and n can be modified to give a more or less accurate numerical solution to the problem. If not considering the computational time, limit in memory, and round-off errors, the obvious preferable choice of ρ_{max} and n would be to make both numbers infinite. This is, however, not a realistic possibility, and this sections is therefore dedicated to find and discuss on the optimal choices for ρ_{max} and n to obtain acceptable values for the eigenvalues of Eq. (2.5) with the matrix **A** defined in Eq. (2.6).

Changing ρ_{max} causes the interval $[\rho_{min}, \rho_{max}]$, in which the wave function is considered, to change. Since the wave function goes to zero as the distance goes to infinity, it is acceptable to neglect the contribution from some ρ_{max} . It is an advantage to decrease ρ_{max} and hence making the interval smaller, in the sense that then a smaller n is needed to create a sufficient step length and ultimately a good enough "resolution". However, if this ρ_{max} is too close to ρ_{min} the neglected part of the eigenfunction can actually not be neglected, if an acceptable result is wished for.

In the figure below, the dependence of different integer valued ρ_{max} on the three first eigenvalues gained by computed Jacobi method is plotted for n = 100.

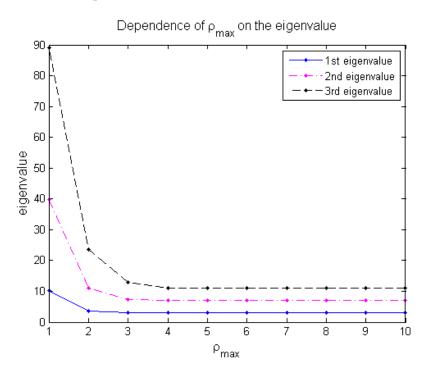


Figure 3.1. The computed eigenvalue for the ground state and the first two excited states of the one electron in a harmonic oscillator potential case plotted as functions of ρ_{max} for $\rho_{min}=0$. For small values of ρ_{max} (that is $\rho_{max} \leq 3$ for this case) the eigenvalue is greatly varying with ρ_{max} , but for greater values of ρ_{max} the eigenvalue seems to stabilize as a function of ρ_{max} towards the analytical results for the eigenvalues of this problem.

From the figure it is evident that if $\rho_{max} < 3$ the eigenvalues are varying dramatically. This happens due to neglect of strongly contributing parts of the eigenfunction. When ρ_{max} is so large that the greatest contributive parts of the eigenfunction is in the interval from ρ_{min} to ρ_{max} , the variation of the eigenvalue decreases strongly. However, when increasing ρ_{max} for the same value of n the step length is increased, as well, which once again influences the precision of the computed eigenvalue.

The ρ_{max} that, with n = 100, gives the most accurate result for all three of the first eigenvalues is $\rho_{max} = 5$. Since this ρ_{max} gives the most accurate result for a relatively small n, this is chosen as the optimal ρ_{max} in this and the following sections for this specific problem.

With this ρ_{max} , we wish to find the integer value of n that gives the first three eigenvalues with four leading digits. This optimal n is found by steady increment of n, as seen in Fig. 3.2.

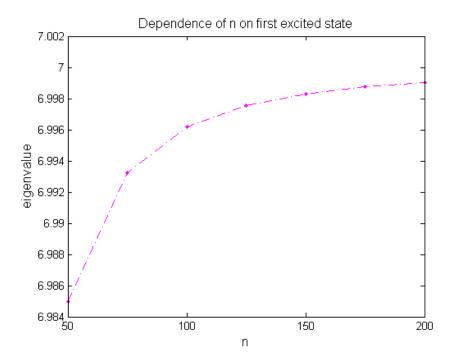


Figure 3.2. Plot of the second eigenvalue of the one electron in a simple harmonic oscillator case. The second eigenvalue is plotted as a function of the number of steps n for $\rho_{min} = 0$ and $\rho_{max} = 5$. The analytical solution for the first excited state of this case is 7, and from the graph it is evident that the numerical solution gained from the computed Jacobi method approaches the analytical solution, as n is increased. One can, however, suspect that a further increment of n can lead to greater round-off errors, and hence a worse result.

The eigenvalue of the first exited state is seen to be asymptotic to the analytical solution $\lambda_2 = 7$, and at a matrix size of n = 200, the eigenvalue of the first exited state has the numerical solution 6.99904. This yields an accuracy up to four leading digits, which is also found to be the case for the ground state and the third eigenstate. Hence, the optimal ρ_{max} and n is 5 and 200, respectively.

3.2 Computation Time Compared to Alternative Algorithm

When computing the time needed for the Jacobi rotation algorithm computed in this project to solve the eigenvalue problem with a matrix $\bf A$ of dimensionality 200×200 , $\rho_{max} = 5$, and $\varepsilon = 10^{-8}$ the elapse time if found to be 17 sec. This is a much greater value than the elapsed time for solving the same eigenvalue problem using the Armadillo function eig_sym , since the Armadillo function has a displayed computational time of 0 sec (the computational time is of course not 0 sec, but the precision of the displayed time is so low that the number is displayed as 0). It is hence clear that the computed Jacobi algorithm consumes more time than the precomputed Armadillo function and as we increases the size of the matrix the elapsed time to get the solution also increases. This makes Jacobi rotation algorithm less efficient.

The following figures show the relation between the number of steps n and the number of iterations of the while loop in the Jacobi algorithm computed for this project with the same ρ_{max} and ε as described in the section above.

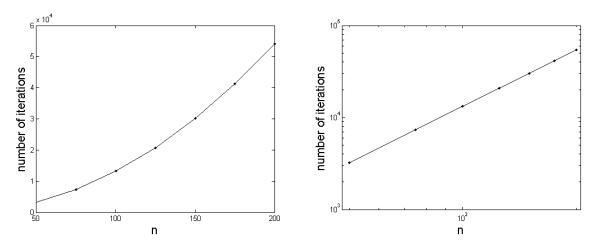


Figure 3.3. The number of iterations of the while loop in the computed Jacobi method as a function of the number of steps n plotted in both a normal graph and as a log-log plot. From the log-log plot it is evident that the number of iterations in the while loop is proportional to n^2 , since the slope of the log-log plot is found to equal 2.

In the log-log plot, the relation is shown as a straight line with a slope of approximately 2, which is calculated directly from the data for n=100 and n=200. This yield that the number of iterations increases as the number of steps squared (that is, n^2), and hence an increment of n has a great impact on the number of iteration. E.g. with n=100 the number of iterations for the while loop in the Jacobi method is 13,200, whilst when n is doubled to 200, the function runs through the while loop 54,071 times before the requirement of $max(a_{kl}^2) \le \varepsilon$ in Eq. (2.23) is fulfilled.

The slowness of the Jacobi algorithm has the effect that it cannot be run for too large matrices, and hence there is a significant limit for how small the step length h can be made, and ultimately a limit to the precision of the algorithm.

It is, however, evident that the Jacobi method implemented in this project can be improved for solving the specific eigenvalue problem described by Eq. (2.5) and (2.6) by taking into account that the matrix $\bf A$ is tridiagonal and has constant values in the entrances adjacent to the diagonal. However, this improvement of the algorithm is not in the scope of this project.

3.3 New potential and repulsive Coulomb interaction

Until now the considered potential for one particle has been described by ρ^2 . When introducing another particle the potential is changed so that the new potential is given by Eq. (3.1), where ω is an oscillator frequency and $\frac{1}{\rho}$ is a repulsive Coulomb potential.

$$V_{\rho} = \omega^2 \rho^2 + \frac{1}{\rho} \tag{3.1}$$

First considering the case without a repulsive Coulomb interaction. The aim is finding a value of ρ_{max} so that the eigenvalues are considered stable.

For the eigenvalue to be stable the change should be small. This is achieved when $\rho_{max}=10$, since then the change is about 0.02 between the surrounding steps (when the step length of ρ_{max} is 1). Furthermore, the change between 10 and 10.1 is about 0.002 so small changes in ρ_{max} doesn't give big changes in the eigenvalue, and the eigenvalue can be considered stable. In the table, the relation between ρ_{max} and n is kept constant so that the step length n is constant. Running for different values of ρ_{max} with the ratio gives Tab. 3.1, which shows that when ρ_{max} is 10 the results are stable for the 1st eigenvalue, and hence this value of ρ_{max} is used to determine the eigenfunctions of the two particle case.

The eigenvectors are used for determining the wave function. In Fig. 3.4 and 3.5 they are plotted squared to show the probability density af a function of ρ , which gives the most probable distance between the

the particles, since ρ is related to the distance r by $\rho = (\frac{1}{\alpha})r$ where $\alpha = (\frac{\hbar^2}{mk})^{1/4}$ is the Bohr radius. This comes from the step between Eq. (2.1) and (2.2) in Sec. 2.1.

Table 3.1. Table for finding best ρ_{max} for the 1st eigen-

ρ_{max}	N	1st eigenvalue
1	10	9.80273
2	20	2.46292
3	30	1.09594
4	40	0.617001
:	÷	:
9	90	0.124119
10	100	0.101504
10.1	101	0.0996154
11	110	0.0849622

Plotting the distance as a function of ρ for different values of the frequency ω_r without a repulsive coulomb potential for only the ground state, Fig. 3.4 shows that for a higher frequency ω_r the most probable distance between the two particles are bigger and the spread in area which the particle can be in is also more stretched out.

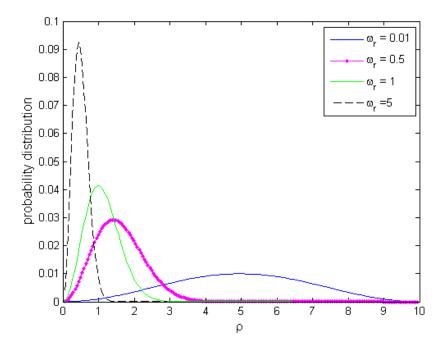


Figure 3.4. The probability density as a function of distance ρ without repulsive Coulomb interaction. Where $\rho_{min} = 0$, $\rho_{max} = 10$ and n = 200

When adding a repulsive Coulomb potential $\frac{1}{\rho}$ to the potential $V(\rho)$ in Fig. 3.5 the probability density is pushed to a longer distance, although it might be a bit difficult to see in Fig. 3.5 and Fig. 3.4. Physically

it makes sense that the repulsive Coulomb interaction should increase the distance between the electrons as it gives a greater contribution when the electrons are close. The choice of ρ_{max} is good because most of the probability distribution is inside the distance ρ_{max} . Also if there is no repulsive Coulomb interaction and the frequency $\omega = 1$ then the system acts as in the one particle case, and the choice of $\rho_{max} = 5$ as discussed for in Sec. 3.1 includes most of the particle distributions.

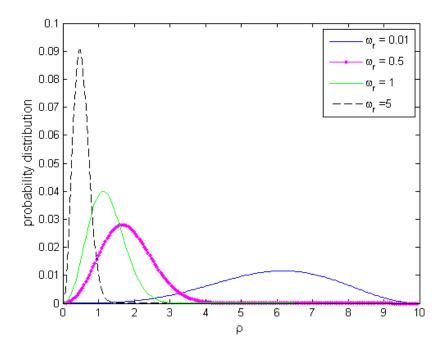


Figure 3.5. The probability density as a function of distance ρ with a repulsive Coulomb interaction. Where $\rho_{min} = 0$, $\rho_{max} = 10$ and n = 200

CHAPTER



CONCLUSION

An algorithm was developed to implement Jacobis method to solve Schrödinger's equation for electrons in a three dimensional harmonic oscillator potential well with and without a repulsive coloumb interaction. Optimal results for the various tests on the algorithm ensured its validity. However the algorithm seems to be less efficient when compared to the Armadillo function in terms of computational speed. A study is made on how the number of points, number of similarity transformations, and the dimensionality of matrix should be chosen so as to get better accuracy for the first three eigenvalues.

Using the computed algorithm for the Jacobi method a accuracy of up to 4 leading digits for the first three eigenvalues compared to the analytical solution is found when running the code for an interval of 5 and 200 steps. This, however, yields 54,000 iterations in the while loop of the Jacobi method.

As a final result the probability distributions of two particles in a harmonic oscillator potential with and without repulsive coloumb interaction is plotted for different oscillator frequencies, and it is found that the most probable distance between the two particles decrease with increasing frequency, and is increased when the repulsive interaction is added compared to without this interaction.