

# Theory of Computation: CS-202

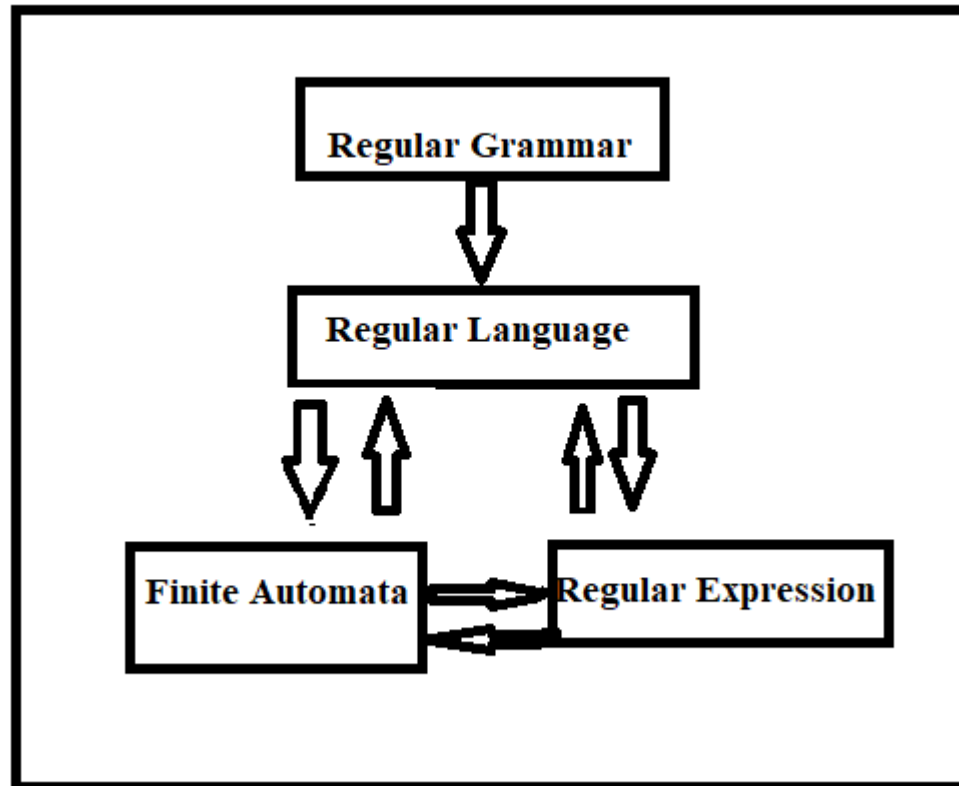
## Regular Expression

# Outline

- Regular Expressions

- Regular Expression to Finite Automata

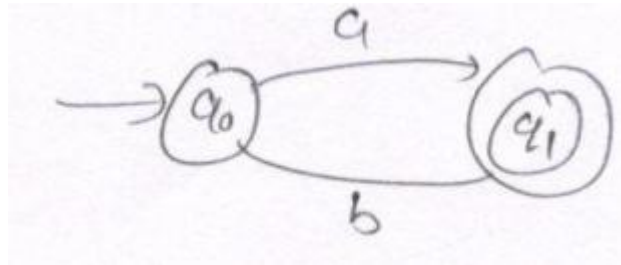
- Finite Automata to Regular Expression



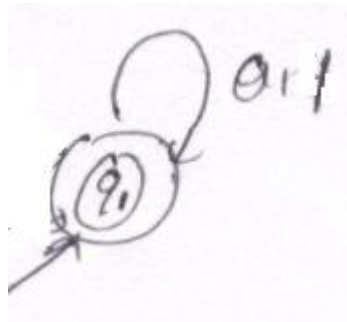
# Conversion from R.E. to finite Automata

## Convert the following R.E. to finite Automata

1. R. E.  $r = (a+b)$

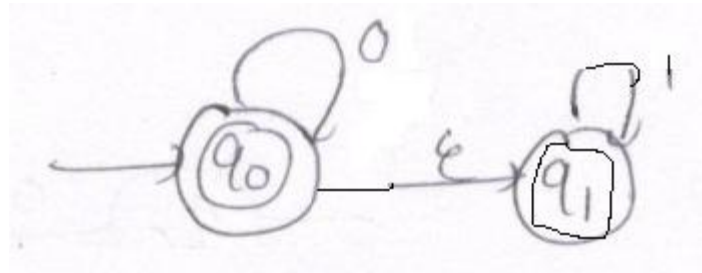


2. R. E.  $r = (0+1)^*$



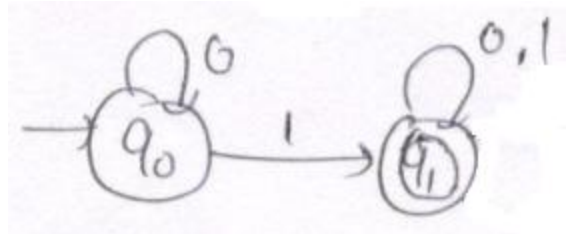
3. R. E.

$r = 0^*1^*$



4. R. E.

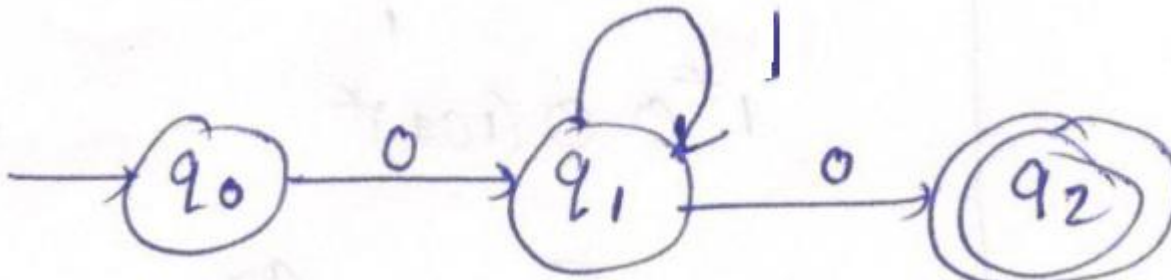
$r = 0^*1(0+1)^*$



# Conversion from Finite Automata to R.E.

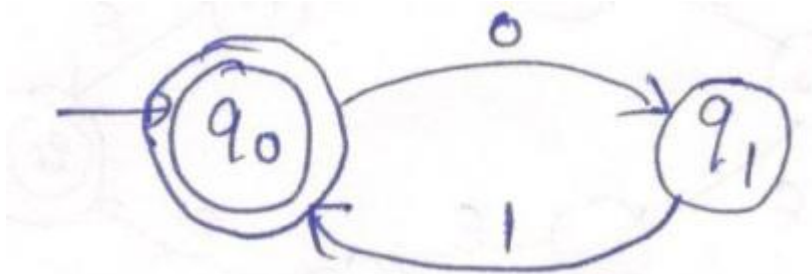
Convert the following FA to R.E.

1.



R. E.  $r = 01^*0$

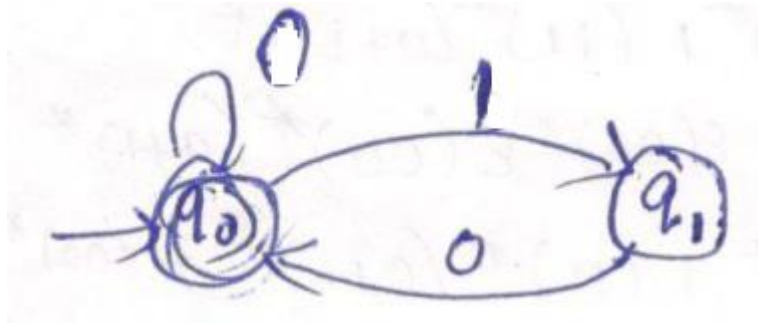
2.



R. E.  $r = (01)^*$

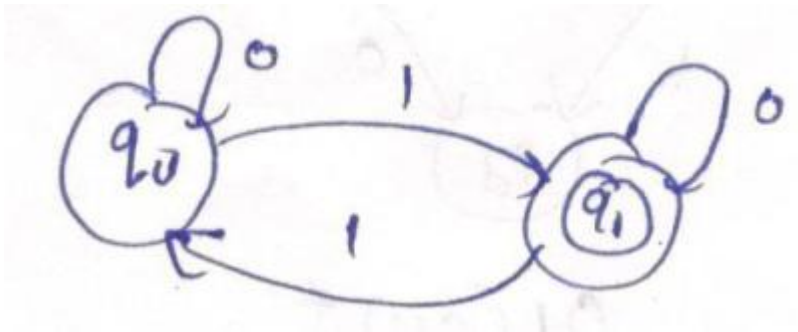


3.



R. E.  $r = (0^* + (10)^*)^*$  or  $((0+10)^*)$

4.



R. E.  $r = 0^*1(0+10^*1)^*$

## Identities for Regular Expression

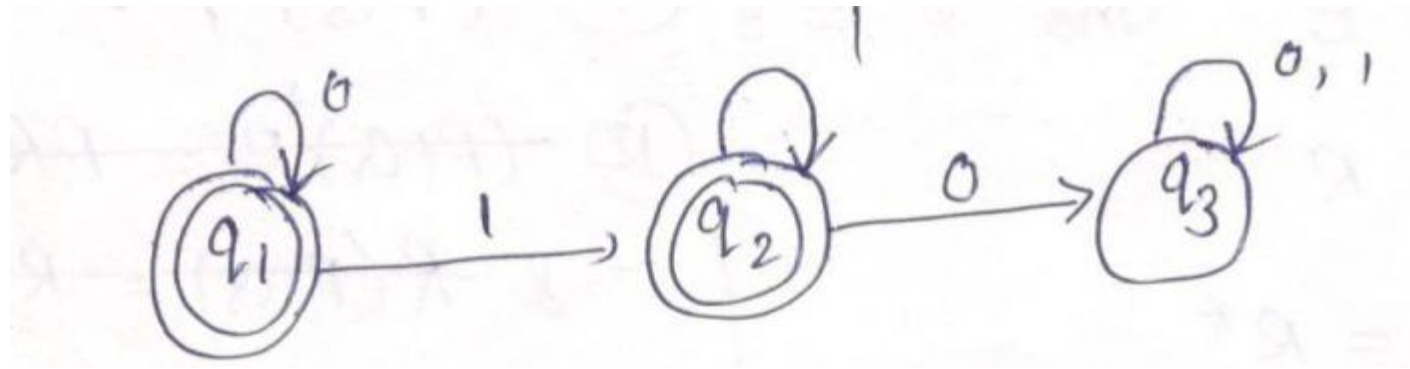
If P, Q and R are Regular Expression, then

1.  $\Phi + R = R$  the identity for union
2.  $\epsilon R = R$   $\epsilon = R$  the identity for concatenation
3.  $\Phi R + R \Phi = \Phi$  the annihilator for concatenation
4.  $\epsilon^* = \epsilon$  and  $\Phi^* = \epsilon$
5.  $R + R = R$
6.  $R^* R^* = R^*$
7.  $RR^* = R^*R$
8.  $(R^*)^* = R^*$
9.  $\epsilon + RR^* = \epsilon + R^*R = R^*$
10.  $(PQ)^*P = P(QP)^*$
11.  $(P+Q)^* = (P^*Q^*) = (P^*+Q^*)^*$
12.  $(P+Q)R = PR + QR$  &  $R(P+Q) = RP + RQ$

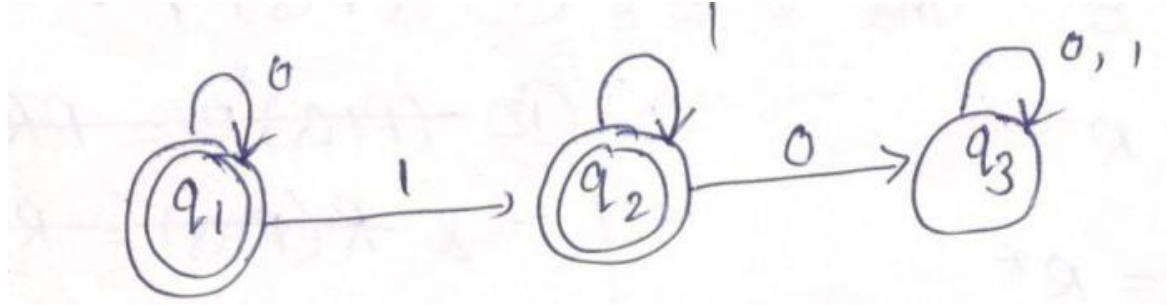
## Arden's Theorem:

Let  $P$  and  $Q$  be two R.E. over  $\Sigma$ . If  $P$  does not contain  $\varepsilon$ , then the equation  $R=Q+RP$  has a unique solution  $R=QP^*$

Construct R.E. corresponding to the state diagram



Construct R.E. corresponding to the state diagram



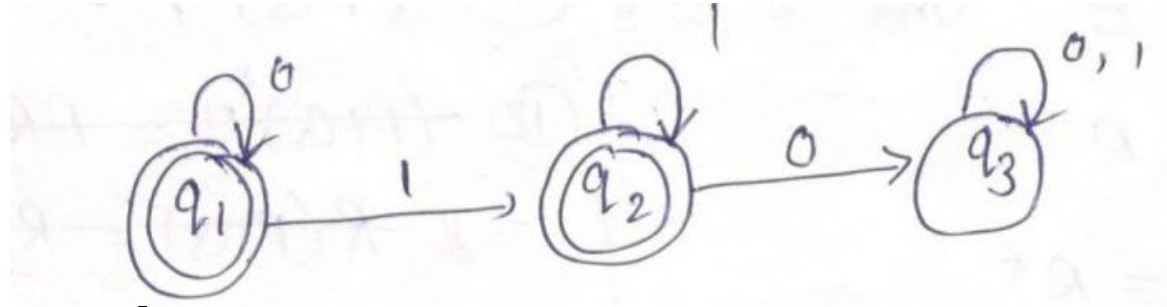
The three equation for  $q_1$ ,  $q_2$  and  $q_3$  can be written as:

$$q_1 = q_1 0 + \epsilon \dots\dots\dots (1)$$

$$q_2 = q_1 1 + q_2 1 \dots\dots\dots (2)$$

$$q_3 = q_2 0 + q_3 0 + q_3 1 \dots\dots\dots (3)$$

## Construct R.E. corresponding to the state diagram



$$R = Q + RP$$

$$R = QP^*$$

From eq. 1

$$q_1 = q_1 0 + \epsilon \dots\dots\dots(1)$$

$$q_1 = \epsilon 0^* \quad \text{using Arden theorem}$$

From eq. 2

$$q_2 = q_1 1 + q_2 1 \dots\dots\dots(2)$$

$$q_2 = \epsilon 0^* 1 + q_2 1 = 0^* 1 + q_2 1$$

$$\Rightarrow q_2 = (0^* 1) 1^*$$

now from eq. 3  $q_3 = q_2 0 + q_3 0 + q_3 1 \dots\dots\dots(3)$

$$q_3 = (0^* 1) 1^* 0 + q_3 0 + q_3 1 = (0^* 1) 1^* 0 + (0 + 1) q_3$$

## Construct R.E. corresponding to the state diagram

$$R=Q+RP$$

$$R=QP^*$$

As  $q_1$  and  $q_2$  are the final states, so we need not to go for state  $q_3$

So, the require regular expression

$$R=q_1+q_2$$

$$R=\epsilon 0^* + (0^* 1)1^*$$

$$R=0^* + (0^* 1)1^* \quad \text{by identity 3}$$

$$R=0^* + (0^* 1)1^*$$

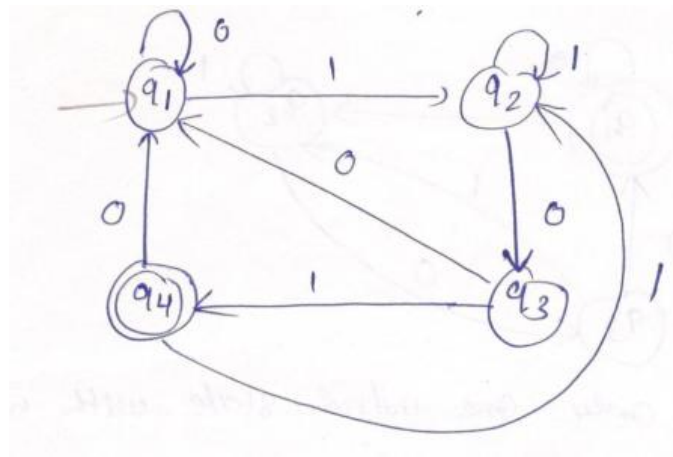
$$R=0^*(\epsilon + 11^*)$$

$$R=0^*(1^*) \quad \text{by identity 9}$$

$$R=0^*1^*$$

# Practice Problems

1. Convert the R.E  $(a+bc^*d)^*$  to F.A.
2. Convert F.A to R.E using Arden's theorem.





## Suggested readings

1. An introduction to FORMAL LANGUAGES and AUTOMATA by PETER LINZ.
2. Introduction to Automata Theory, Languages, And Computation by JOHN E. HOPCROFT, RAJEEV MOTWANI, JEFFREY D. ULLMAN
3. Theory of computer science: automata, languages and computation **by** K.L.P MISHRA

**Thank you**