



**Birzeit University**

**Faculty of Engineering and Technology**

**Department of Electrical and Computer Engineering**

**First Semester – 2023/2024**

**ENCS4320 - Applied Cryptography**

**Homework # 2 - Due Saturday, January 27, 2024**

<b>Section A</b>	<b>Symmetric Key Ciphers (PRFs, PRPs, Block Ciphers), Group Theory, and Number Theory</b>
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**Question 1 (12 points):**

Consider a **Feistel cipher** with **four rounds**. Then the plaintext is denoted as  $P = (L_0, R_0)$  and the corresponding ciphertext is  $C = (L_4, R_4)$ . What is the **simplest form** of the ciphertext  $C$ , in terms of  $L_0$ ,  $R_0$ , and the subkeys, for each of the following round functions? (You should clearly show steps about how you get the answer)

- A)  $F(R_{i-1}, K_i) = 0$
- B)  $F(R_{i-1}, K_i) = R_{i-1}$
- C)  $F(R_{i-1}, K_i) = R_{i-1} \oplus K_i$

**Question 2 (5 points):**

Within a single round, the Data Encryption Standard (**DES**) employs both confusion and diffusion.

- A) What is the difference between confusion and diffusion in cryptography?
- B) Give one source of confusion within a DES round.
- C) Give one source of diffusion within a DES round.

**Question 3 (8 points):**

Compute  $(345^{28567} \times 23^{567} + 1078) \bmod 29$  given that 29 is a prime.

**Question 4 (9 points):**

Using the Extended Euclidean algorithm,

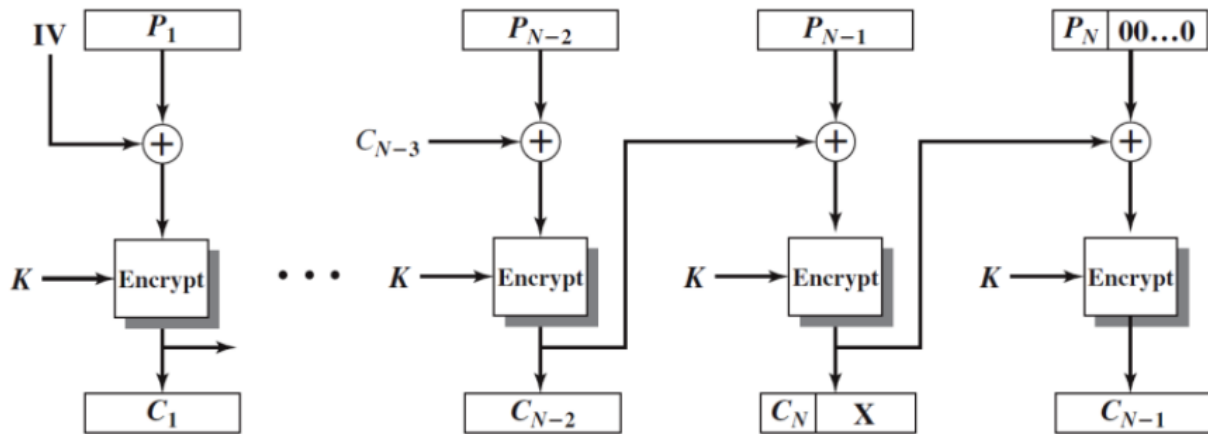
- A) Find the *greatest common divisor* of **19** and **999**, that is, **gcd(999, 19)**. Show your work clearly step by step.
- B) Express the **gcd(999, 19)** as a *linear combination* of **999** and **19**.
- C) Compute the *multiplicative inverse* of **19 mod 999**, which is a number between **0** and **998**.
- D) Compute the *multiplicative inverse* of **999 mod 19**, which is a number between **0** and **18**.

<b>Section B</b>	<b>Block Cipher Modes, Integrity, and Authenticated Encryption</b>
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**Question 5 (14 points):**

Bob wishes to encrypt some plaintext and stores the resulting ciphertext on his hard drive. Specifically, he wants the ciphertext to be the *same length* as the original plaintext. For this purpose, he employed the **ciphertext stealing (CTS)** mode, the implementation of which is shown in the figure below. Initially, the plaintext is divided into independent blocks of length  $S$  bits, giving the plaintext blocks  $P_1, P_2, \dots, P_N$ . Assume that the last block of plaintext (i.e.,  $P_N$ ) is  $L$  bits long, where  $L < S$ . The encryption sequence is as follows:

1. Encrypt the first  $(N - 2)$  blocks using the traditional cipher block chaining (**CBC**) technique.
2. XOR  $P_{N-1}$  with the previous ciphertext block  $C_{N-2}$  to create  $Y_{N-1}$ .
3. Encrypt  $Y_{N-1}$  to create  $E_{N-1}$ .
4. Select the first  $L$  bits of  $E_{N-1}$  to create  $C_N$ .
5. Pad  $P_N$  with  $(S - L)$  zeros at the end and exclusive-OR with  $E_{N-1}$  to create  $Y_N$ .
6. Encrypt  $Y_N$  to create  $C_{N-1}$ .



- A) Describe how to decrypt the ciphertext ( $C_1, \dots, C_{N-1}, C_N$ ), that is, show the decryption sequence.
- B) If a single bit error occurs in the storage of ciphertext  $C_i$ , which plaintext blocks, if any, will be correctly restored by the decryption algorithm? Explain your answer.

## Section C

### Hash Functions

#### Question 6 (6 points):

Suppose that  $H(m)$  is a secure hash function that generates a 12-bit output.

- A) How many collisions would you expect to find if you hash 1024 randomly selected messages?
- B) What is the expected number of hashes that must be computed to find 25 collisions? That is, what is the expected number of hashes that must be computed to find pairs  $(x_i, y_i)$ ,  $x_i \neq y_i$ , with  $H(x_i) = H(y_i)$ , for  $i = 1, 2, \dots, 25$ ?

#### Question 7 (6 points):

Consider the following hash function  $H(m)$ , which receives as an input a message in the form of a sequence of decimal numbers,  $m = (D_1, D_2, \dots, D_l)$ .

- A)** If  $H(m)$  is defined as  $(\sum_{i=1}^l D_i) \bmod n$ , for some predefined large value  $n$ . Does this hash function satisfy the pre-image resistance (one-way property) requirement? Explain your answer.
- B)** If  $H(m)$  is defined as  $(\sum_{i=1}^l D_i^2) \bmod n$ , for some predefined large value  $n$ . Does this hash function satisfy the collision resistance requirement? Explain your answer.
- C)** Calculate the hash function of part (B) for  $m = (189, 632, 900, 722, 349)$  and  $n = 989$ .

***GOOD LUCK***