

# COP 5536 Project Report

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## Programming Language - C++

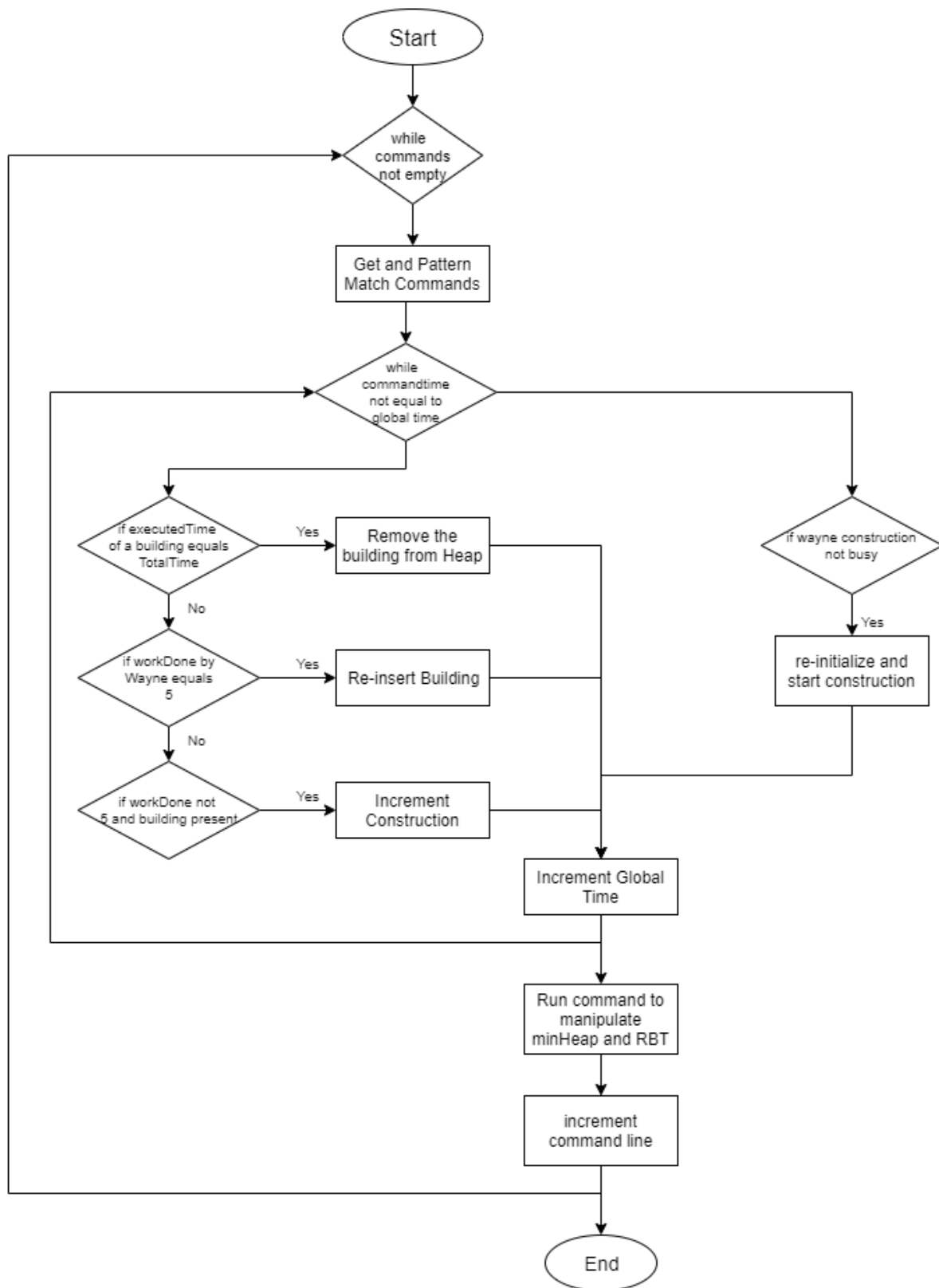
### Files Created for the Project

1. mainfile.cpp
2. minheap.cpp
3. minheap.h
4. redblacktree.cpp
5. redblacktree.h

**mainfile.cpp** – Contains the main flow of the program

- *int main(int argc, char \*argv[])*
  - Both the Min Heap and Red-Black Tree are initialized at the beginning
  - Input File is read and the commands are stored in a vector
  - The commands are iterated over a while loop
  - The core logic of the program works here
- *void runInputLine(string cmd, string parseArgs, minheap\* mh, redblacktree\* rbt, ofstream& opF)*
  - Command (Insert/Display), parseArgs(values), Minheap object, RBT object and output file instance are sent as arguments into this function.
  - The command is recognized using regex pattern matching and the Insert and Display function runs accordingly.
  - If the command is “Insert”, a new Red-black Node and min Heap node are created. The reference to red-black node is store in min heap array.

## Flowchart of Business Lo



## Min Heap [minheap.cpp and minheap.h]

**Class Node:** Used for declaring variables and constructor of Min Heap

Variables

- *cBuildingnum* – Building Number stored
- *cExecutedtime* – Defines the time for which work is done on the building
- *cTotaltime* – Total time required for the completion
- *cRBTnode* – Reference of the corresponding Red-black node

### Function Definitions

- *bNode(long int buildingNum, long int executedTime, long int totalTime, redblacknode\* rbn);*
  - Constructor of the class which is called whenever a new node is created.
  - All the variable values of min heap node are initialized in the constructor

**Class MinHeap:** Contains all the functions of Min Heap data structure

- *Void insertNode(bNode\* value)*
  - Insert a new Node in the MinHeap Array
  - Bubbleup function called which conserves the min heap property of the structure
- *Void deleteMin()*
  - Deletes the minimum node from the Min Heap
  - Bubbledown function called which conserves the min heap property of the structure
- *bNode\* getMinFromHeap()*
  - Gets the minimum node from the Min Heap
- *Void bubbleDown(int index)*
  - When the min node is deleted, we replace it with the last element.
  - If replaced element is greater than any of its child node, swap the element with its smallest child(Min-Heap) or with its greatest child(Max-Heap).
  - Do above step recursively.
- *Void bubbleUp()*
  - When a new element is inserted, if inserted element is smaller than its parent, swap the element with its parent.
  - Do the above step recursively.

## Red-Black Tree [redblacktree.cpp and redblacktree.h]

**Class redblacknode:** Used for declaring variables and constructor of Red-Black Tree

### Variables

- *cBuildingnum* - Building Number stored
- *cExecutedtime* - Defines the time for which work is done on the building
- *cTotaltime* - Total time required for the completion
- *cLeft* – Store the reference to left child
- *cRight* – Store the reference to right child
- *cParent* – Store the reference to parent
- *cColor*– Store the color of the node. (0 – black node, 1 – red node)

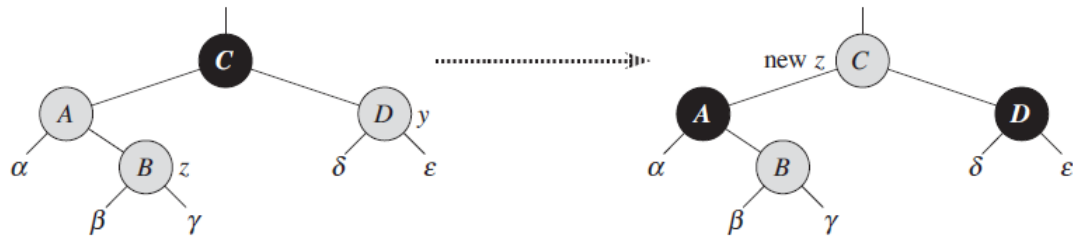
### Function Definitions

- *redblacknode(long int buildingNum, long int executedTime, long int totalTime, redblacknode\* rbn);*
  - Constructor of the class which is called whenever a new node is created.
  - All the variable values of RBT node are initialized in the constructor

**Class RedBlackTree:** Contains all the functions of Min Heap data structure

- *void insertNode(redblacknode\* node)*
  - Function for inserting a new node in RBT
  - The insertion is the same as that of Simple Binary Search Tree
  - After the node has been inserted, we fix the Insert by calling the reinsert function which fixes the RBT and conserves it's property.
- *void redblacktree::reInsert(redblacknode\* node)*
  - This function is used to fix the imbalance in the RBT caused by the insertion of new node
  - There are 3 cases that are handled by this function

➤ Case1:

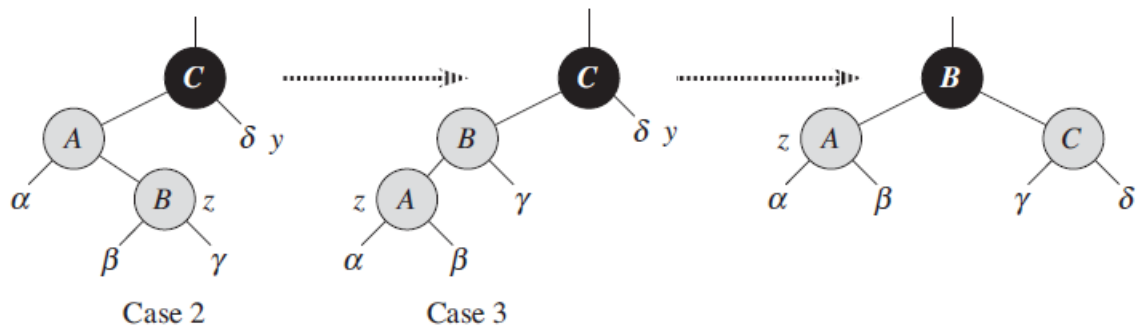


Picture Reference: CLRS

Explanation:

- Here z and its parent z->p are both red.
- Each of the subtrees has a black root, and each has the same black edges.
- The code for case 1 changes the colors of some nodes, preserving property 5. All downward paths from a node to a leaf have the same number of blacks.
- The while loop continues with node z's grandparent z->p->p as the new z.

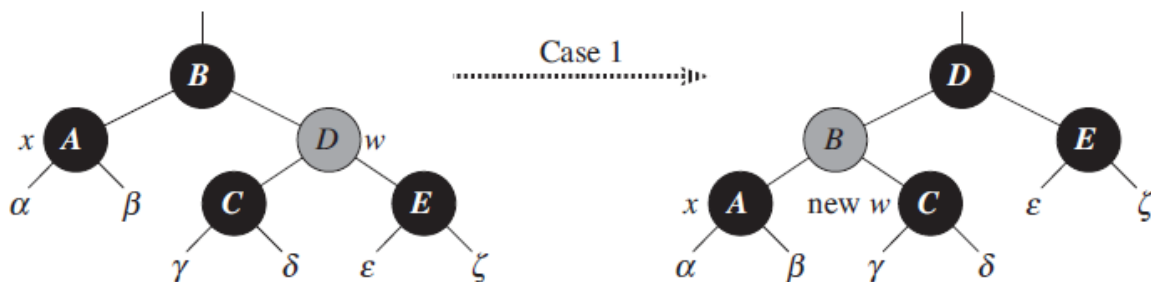
➤ Case2 and Case3:



Explanation:

- In cases 2 and 3, the color of z's uncle y is black.
- In case 2, node z is a right child of its parent. We immediately use a left rotation to fix this issue.
- In case 3, we execute some color changes and a right rotation, which preserve property 5, and then, since we no longer have two red nodes in a row, we are done.

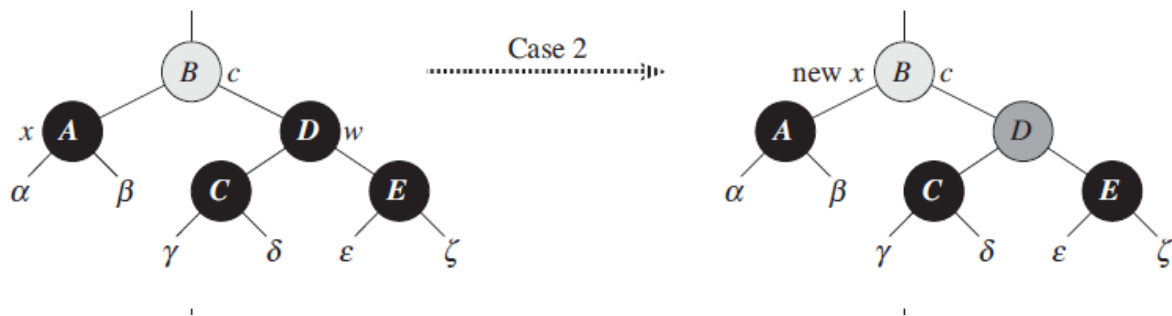
- *redblacknode\* findMin(redblacknode\* node)*
  - This function returns the minimum value node from the Red-Black Tree
- *void redblacktree::lRotate(redblacknode\* node)*
  - Function performs left rotation in the RBT
- *void redblacktree::rRotate(redblacknode\* node)*
  - Function performs right rotation in the RBT
- *Void deleteNode(redblacknode\* node)*
  - A standard BST delete is performed at first
  - Following cases are handled:
    1. When we want to delete a node z which has fewer than two children, then is removed from the tree
    2. When a node z has 2 children
  - After deletion, redelete() function is called to fix the imbalances caused by the deletion.
- *Void reDelete(redblacknode\* node)*
  - Following Cases are handled in this function
    1. x's sibling w is red



Reference: CLRS

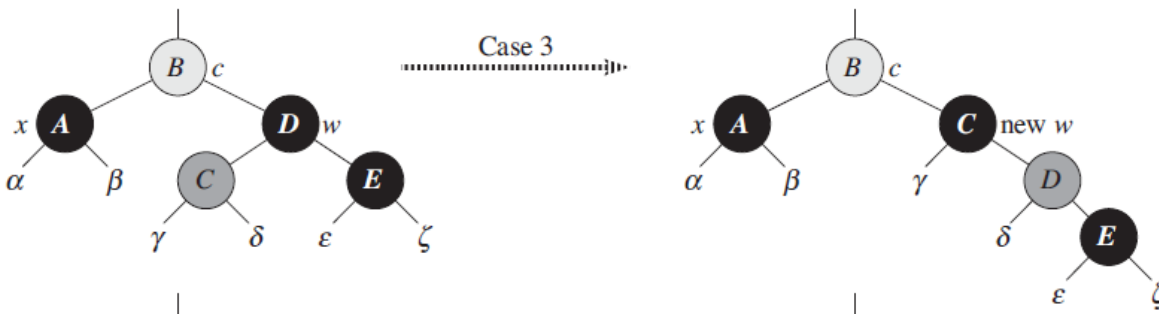
Here a Left rotation is performed in the redelete() function

2.  $x$ 's sibling  $w$  is black, and both of  $w$ 's children are black



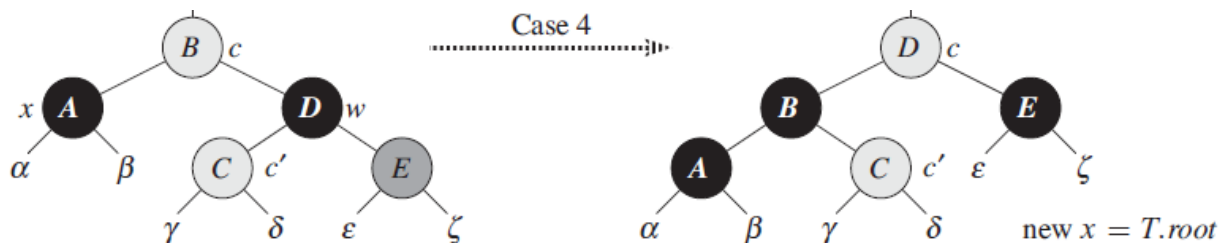
- Take one black off from  $x$  (making  $x$  singly black) and off  $w$  (making  $w$  red).
- Move that black to  $x$ .parent.
- Do the next iteration with  $x$ .parent as the new  $x$ .

3.  $x$ 's sibling  $w$  is black,  $w$ 's left child is red, and  $w$ 's right child is black



- Make  $w$  red and  $w$ 's left child black.
- Then right rotate on  $w$ .
- The new sibling  $w$  of  $x$  is black with a red right child, and we go immediately into case 4.

4.  $x$ 's sibling  $w$  is black, and  $w$ 's right child is red



- Make  $w$  be  $x$ .parent's color ( $c$ ).

- Make x.parent black and w's right child black.
- Then left rotate on x.parent.

All the cases are recursively fixed in a while loop till the RBT properties are not conserved.

- *Void rbtRelink(redblacknode\* n1, redblacknode\* n2)*
  - This function replaces one subtree as a child of its parent with another subtree.
- *Redblacknode\* searchTree(long int bnum)*
  - This function takes a building number as input and searches it in the red-black tree
  - It calls searchTreeHelper() which is a recursive function
- *Void searchTreeInRange(long int bnum1, long int bnum2, vector<redblacknode\*>\* vect)*
  - This function takes two building numbers as input
  - It called searchTreeInRangeHelper() which is a recursive function that helps in finding the building numbers between the given inputs.

## References:

- Introduction to Algorithms [CLRS]