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# Affine Objects

Master thesis

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Voor hooyo.

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# 1 Introduction

## 1.1 Motivation

Quasi-coherent sheaves and affine schemes have a special relationship. One can find out that a sheaf of modules on any scheme is quasi-coherent by looking at the sections at the affine open subsets. One can find out that a quasi-compact scheme is affine by proving all higher cohomology of any quasi-coherent sheaf vanishes [7, Tag 01XF].

## 1.2 Aim

This thesis will investigate and formalize this special relationship in a definition that only depends on the category of quasi-coherent sheaves on a scheme. This generalizes affine schemes to any ringed site, which gives us a notion of ‘affine’ object, which we will call **caffine**, shorthand for categorically affine. We will investigate how fruitful this definition is in other contexts.

## 1.3 Outline

In chapter 2 we will build up all the necessary theory that is well-known, but not assumed. We will start with some general categorical results and definitions. Then we will introduce the generalisation of topology for categories, the category of sites and the central notion of a sheaf. The induced topology on an over/slice category will get special attention because we will use it extensively.

In chapter 3 we will introduce restrictive morphisms and show how this notion interacts with affine schemes. It will turn out any morphism between affine schemes is restrictive.

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In Chapter 4 we will define affine objects. Then we will show that the affine objects in the category of schemes are exactly the affine schemes. Lastly we will look at some examples of categories that lack any affine objects.

In Chapter 5 I will expand on possible avenues of further research.

### 1.4 Assumptions

The intended audience are my peers, so some mathematical maturity is assumed. I assume familiar with the usual undergraduate curriculum in topology and algebra. Furthermore basic categorical notions such as limits and adjunctions are assumed. See the first 3 chapters of [6].

## 2 Preliminaries

This chapter introduces all the basic notions that are needed but not assumed to be known to the reader. We will start with a discussion of some purely categorical notions like over categories and presheaves. Secondly, we will introduce a notion of a topology on a category and look at some constructions that are relevant for us. Then we will introduce modules on ringed sites. Lastly, the notion of a scheme is introduced.

### 2.1 Basic Category Theory

Some categorical notions like presheaves and over categories will be introduced in this section. See [1] and [4].

**Definition 2.1.1** (Presheaf category). Let  $C$  be a category. We define the category of presheaves on  $C$  as the category of contravariant functors to the category of sets  $\text{Set}$ . We will denote it by  $\hat{C}$ .

Define the functor  $h : C \rightarrow \hat{C}$  as follows for  $a \in C$ .

$$a \mapsto \text{Hom}(-, a),$$

for  $f : a' \rightarrow a$

$$f \mapsto f \circ -.$$

This functor is fully faithful by the Yoneda lemma.

**Notation 2.1.2.** Let  $I, C$  be categories. Let  $L : I \rightarrow C$  be a functor. The limit over this functor will be denoted by  $\lim_{i \in I} L(i)$ . The colimit will be denoted by  $\text{colim}_{i \in I} L(i)$ .

**Definition 2.1.3** (Sections functor). For any  $a \in C$  define the functor

$$\Gamma(a; -) : \hat{C} \rightarrow \text{Set}$$

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by

$$\mathfrak{F} \rightarrow \mathfrak{F}(\mathfrak{a}).$$

For any presheaf  $\mathfrak{G}$ , we define

$$\begin{aligned} \Gamma(\mathfrak{G}; -) : \hat{C} &\rightarrow \text{Set}, \\ \mathfrak{F} &\rightarrow \text{Hom}(\mathfrak{G}, \mathfrak{F}). \end{aligned}$$

We will mostly use this for the terminal presheaf, which will allow us to compute the global sections.

**Definition 2.1.4** (Over/Under categories). Let  $C$  and  $C'$  be categories. Let  $F : C \rightarrow C'$  and  $z \in C'$ . Define the category  $C_z$  and  $C^z$  as

$$\begin{aligned} \text{Obj}(C_{z,F}) &:= \{(\mathfrak{a}, w) \mid \mathfrak{a} \in C, w : F(\mathfrak{a}) \rightarrow z\}, \\ \text{Hom}((\mathfrak{a}, w), (\mathfrak{b}, v)) &:= \{f : \mathfrak{a} \rightarrow \mathfrak{b} \mid v \circ F(f) = w\}, \end{aligned}$$

and

$$\begin{aligned} \text{Obj}(C^{z,F}) &:= \{(\mathfrak{a}, w) \mid \mathfrak{a} \in C, w : z \rightarrow F(\mathfrak{a})\}, \\ \text{Hom}((\mathfrak{a}, w), (\mathfrak{b}, v)) &:= \{f : \mathfrak{a} \rightarrow \mathfrak{b} \mid F(f) \circ w = v\}. \end{aligned}$$

We get faithful functors  $C_z \rightarrow C : (\mathfrak{a}, w) \rightarrow \mathfrak{a}$  and  $C^z \rightarrow C : (\mathfrak{a}, w) \rightarrow \mathfrak{a}$ . We will denote them both by  $u$ .

**Definition 2.1.5** (direct image). Let  $f : C \rightarrow D$ . Define the direct image functor  $f_* : \hat{D} \rightarrow \hat{C}$  as

$$f_* = - \circ f.$$

When  $u$  is the forgetful functor  $C_{\mathfrak{a}} \rightarrow C$  then we denote  $u_* \mathfrak{F}$  sometimes by  $\mathfrak{F}|_{\mathfrak{a}}$ .

**Definition 2.1.6** (inverse image of presheaves). Let  $C, D$  be categories. Let  $f : C \rightarrow D$  be a functor. Define the inverse image functor  $f^{-1} : \hat{C} \rightarrow \hat{D}$  as follows. Let  $\mathfrak{F} \in \hat{C}$ . For any  $d \in D$

$$f^{-1}(\mathfrak{F})(d) = \text{colim}_{C^{d,f}} \mathfrak{F}.$$

See [7, Tag 00VC].



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**Lemma 2.1.7.** *The functor  $f_*$  is left adjoint to  $f^{-1}$ .*

*Proof.* See [7, Tag 00VE]. ■

**Corollary 2.1.8.** *The functor  $f_*$  commutes with arbitrary colimits.*

### 2.2 Topology

In this section we will define a notion of a topology on a category and look at the related notions of sheaves, sites and restriction of sites.

See [4] for more details.

#### 2.2.1 Basic

**Definition 2.2.1** (Sieve). Let  $C$  be a category and  $a \in C$ . Define the maximal sieve  $\max(a)$  on  $a$  to be the collection of all morphisms to  $a$ . In formula,

$$\max(a) = \{f \in C \mid \text{Codom}(f) = a\}.$$

A sieve  $S$  on  $a$  is a subcollection of  $\max(a)$  such that  $f \circ g \in S$  for any  $f \in S$  and any  $g$ .

**Definition 2.2.2** (Sieve category). Let  $C$  be a category and  $a \in C$ . The sieve category  $\text{Sieves}(a)$  consists of all the sieves on  $a$  as objects and inclusions of sieves as morphisms.

**Definition 2.2.3** (Pullback of sieve). Let  $C$  be a category and  $a, b \in C$ . Let  $S$  be a sieve on  $a$ . Let  $f : b \rightarrow a$ . The sieve  $f^*S$  on  $b$  is given by  $f^*S = \{g \in \max b : fg \in S\}$  for any  $c \in C$ . To show that this is actually a sieve on  $b$ , let  $k : c \rightarrow c'$  and  $h \in f^*S$ . Hence  $fh \in S$  and so  $fhk \in S$ . Conclude that  $hk \in f^*S$ . This defines a functor  $f^* : \text{Sieves}(a) \rightarrow \text{Sieves}(b)$ .

**Definition 2.2.4** (Grothendieck Topology). A Grothendieck topology  $\mathcal{T}$  is a family  $\mathcal{T}(a)$  of ‘covering’ sieves for every  $a \in C$  with the following conditions:

1.  $\max(a) \in \mathcal{T}(a)$
2.  $f^*R \in \mathcal{T}(a')$  if  $R \in \mathcal{T}(a)$  for all  $f : a' \rightarrow a$

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3. if  $f^*R \in \mathcal{T}(a')$  for all  $f \in S$  with  $S \in \mathcal{T}(a)$  then  $R \in \mathcal{T}(a)$

**Definition 2.2.5** (Basis). Let  $C$  be a category with pullbacks. A Grothendieck pretopology, or basis,  $\mathcal{B}$  is a collection  $\mathcal{B}(a)$  of ‘covering’ families  $\{f_i : a_i \rightarrow a\}$  of morphisms for every  $a \in C$  with the following conditions.

1. every isomorphism is a covering singleton family,
2. (Stability) The pullback of a covering family is a covering family. If  $\{f_i : a_i \rightarrow a\}$  is covering and  $g : b \rightarrow a$ , then  $\{f'_i : a_i \times_a b \rightarrow b\}$  is covering.
3. (Transitivity) If  $\{f_i : a_i \rightarrow a\}$  is a covering family and  $\{f_{ij} : a_{ij} \rightarrow a_i\}$  for every  $i$ , then  $\{f_{ij} : a_{ij} \rightarrow a\}$  is a covering family.

Generating a topology from a basis: take any sieve containing a covering family to be a covering sieve.

*Remark 2.2.6.* Sometimes you have a set of families  $S$  you would like to generate the topology. So one can take the smallest basis and the topology containing these families. This will be called ‘the topology generated by  $S$ ’.

### 2.2.2 Sites

**Definition 2.2.7** (Site). A site  $(C, \mathcal{T})$  is a category  $C$  with a Grothendieck topology  $\mathcal{T}$ .

**Definition 2.2.8** (continuous functor). Let  $G : C \rightarrow D$  be a functor between sites. Let  $c \in C$  and let  $R$  be a covering sieve on  $c$ . The functor  $G$  is said to preserve covers if  $G(R)$  generates a covering sieve. It is enough to check that  $G$  sends covering families to covering families, if the topology of the sites is defined by a basis. If  $G$  also preserves pullbacks then we call it continuous. See [7, Tag 00WU].

**Definition 2.2.9** (cocontinuous functor). The functor  $G$  is said to lift covers or be cocontinuous if for every  $R \in \text{Cov}_G(c)$  there is some  $S \in \text{Cov}_c$  such that  $G(S) \subset R$ . See [7, Tag 00XJ].

**Definition 2.2.10.** The category Sites has as objects sites and the morphisms are continuous functors. When no confusion can arise then we will use  $C$  to denote the whole site  $(C, \mathcal{T})$ .

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*Remark 2.2.11.* In most resources a morphism of sites is defined to preserve all finite limits. To get the results that we want we only need preservation of pullbacks and we need the forgetful functor  $C_a \rightarrow C$  to be a morphism of sites, hence this slightly weaker notion than usual.

**Definition 2.2.12** (Ringed sites). A ringed site  $(C, \mathcal{T}, \mathfrak{O})$  is the site  $(C, \mathcal{T})$  with the sheaf of rings  $\mathfrak{O}$  on it. A morphism of ringed sites is a morphism pair  $(f, f^\#) : (C, \mathcal{T}, \mathfrak{O}) \rightarrow (D, \mathcal{S}, \mathfrak{U})$  with  $f : (C, \mathcal{T}) \rightarrow (D, \mathcal{S})$  a morphism of sites and  $f^\# : f_*\mathfrak{U} \rightarrow \mathfrak{O}$  a morphism of sheaf of rings. We will denote the pair  $(f, f^\#)$  by just  $f$  when no confusion can arise.

**Example 2.2.13** (small site). Let  $X$  be a topological space. Let the category  $\text{Open}(X)$  consist of the opens of  $X$  as objects and inclusion as the morphisms. Define a basis on this site consisting of the families  $\{U_i \rightarrow U\}$  such that  $\bigcup_i U_i = U$ . A continuous map  $f : X \rightarrow Y$  induces a morphism of sites  $\text{Open}(Y) \rightarrow \text{Open}(X)$  by sending  $U \mapsto f^{-1}(U)$ . In this way you can embed the category of topological spaces into the category of sites.

If  $X$  is a ringed space we can turn the site  $\text{Open}(X)$  into a ringed site by setting the sheaf of rings of  $X$  to be the sheaf of rings of  $\text{Open}(X)$ .

### 2.2.3 Sheaves

We will introduce the very important notion of a sheaf. See [4] for a more detailed treatment.

**Definition 2.2.14** (Matching family). Let  $C$  be a category. Let  $\mathfrak{F}$  be a presheaf on  $C$ . Let  $a \in C$  be an object. Let  $R$  be a sieve on  $a$ . A family  $\{x_i\}_{i \in R}$  with  $x_i \in \Gamma(\text{Dom}(i); \mathfrak{F})$  indexed by a sieve  $R$  and such that  $x_{g \circ i} = \mathfrak{F}(g)(x_i)$  for any  $g : b \rightarrow \text{Dom}(i)$  and  $b \in C$  is called a ‘matching family’.

**Definition 2.2.15** (Amalgamation). An amalgamation of a matching family  $\{x_i\}_R$  is an element  $x \in \Gamma(a; \mathfrak{F})$  such that  $\mathfrak{F}(i)(x) = x_i$ .

**Definition 2.2.16** (Sheaf). Let  $(C, \mathcal{T})$  be a site. Let  $\mathfrak{F} \in \hat{C}$ .

A presheaf that admits a unique amalgamation for every matching family is called a sheaf. The category  $\text{Shv}(C)$  is the full subcategory in  $\hat{C}$  of all sheaves.

**Definition 2.2.17** (Plus construction). Let  $(C, \mathcal{T})$  be a site. Let  $a, a' \in C$  and  $f : a \rightarrow a'$ . Let  $\mathfrak{F} \in \hat{C}$ . Define the functor  $(-)^+ : \hat{C} \rightarrow \hat{C}$  as follows.

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For all  $\mathfrak{a} \in \mathcal{C}$ ,

$$F^+(\mathfrak{a}) = \frac{\{(\mathbf{R}, \varphi) \mid \mathbf{R} \in \mathcal{T}(\mathfrak{a}), \varphi \in \Gamma(\mathbf{R}; \mathfrak{F})\}}{\sim},$$

for all morphisms  $f \in \mathcal{C}$ ,

$$F^+(f)([(\mathbf{R}, \varphi)]) = [(f^*\mathbf{R}, \varphi h(f))].$$

The equivalence relation is defined as  $(\mathbf{R}, \varphi) \sim (\mathbf{S}, \phi)$  if  $\varphi = \phi$  on some covering sieve  $Q \subset \mathbf{R} \cap \mathbf{S}$ .

Let  $L : \mathfrak{F} \rightarrow \mathfrak{F}'$ . Then

$$(L^+)_\mathfrak{a}([(R, \varphi)]) = [(R, L \circ \varphi)]$$

This functor comes with a natural transformation  $\omega : \text{Id} \rightarrow (-)^+$  defined by

$$\omega_{\mathfrak{F}, \mathfrak{a}}(\mathbf{x}) = [(\max(\mathfrak{a}), \mathbf{y})]$$

where

$$\mathbf{y}(\mathfrak{i}) = \mathfrak{F}(\mathfrak{i})(\mathbf{x}).$$

**Lemma 2.2.18.** *Let  $\mathfrak{F}$  be a presheaf,  $\mathfrak{G}$  a sheaf and  $g : \mathfrak{F} \rightarrow \mathfrak{G}$  a morphism in  $\hat{\mathcal{C}}$ . Then  $g$  factors through  $\omega_{\mathfrak{F}}$  via a unique  $g'$ .*

*Proof.* See [5, p. 24]. ■

**Lemma 2.2.19.** *For every presheaf  $\mathfrak{F}$ ,  $F^+$  is separated. If  $\mathfrak{F}$  is separated, then  $F^+$  is a sheaf.*

*Proof.* See [5, p. 24]. ■

**Definition 2.2.20.** Define  $\text{sh} = (-)^+ \circ (-)^+$ . This functor is left adjoint to the inclusion  $\hat{\mathcal{C}} \rightarrow \text{Shv}(\mathcal{C})$  with unit

$$\omega_{\mathfrak{F}}^2 = \omega_{\mathfrak{F}^+} \circ \omega_{\mathfrak{F}}.$$

**Lemma 2.2.21.** *Let  $f : \mathcal{C} \rightarrow \mathcal{D}$  be a continuous functor. Let  $\mathfrak{D}$  be a sheaf on  $\mathcal{D}$ . Then  $f_*\mathfrak{D}$  is a sheaf.*

*Proof.* See [7, Tag 00WW]. ■

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### 2.2.4 Relative topology

We will look at what the induced topology on a over category looks like and what this implies for restriction of sheaves. See [7, Tag 03A4] for a more detailed treatment.

*Remark 2.2.22.* Let  $C$  be a category and  $a, b \in C$ . Let  $f : b \rightarrow a \in C_a$ . The map  $\max(f) \rightarrow \max(b)$  sending a morphism to  $f$  to its underlying morphism in  $C$  is a bijection. Moreover composition in  $C$  and  $C_a$  are the same, so this identification respects pullback of sieves. This observation yields us that  $\text{Sieves}(b) = \text{Sieves}(f)$ .

Whenever  $R$  is a sieve on  $b$ , we will denote the corresponding sieve on  $f$  by  $R_f$ .

**Definition 2.2.23** (Relative topology). Let  $(C, \mathcal{T})$  be a site. Let  $a \in C$ .

Define the induced topology  $\mathcal{T}_a$  on  $C_a$  by, for each  $f \in C_a$

$$\mathcal{T}_a(f) = \mathcal{T}(\text{Dom}(f)).$$

The identification from Remark 2.2.22 implies that  $\mathcal{T}_a$  is a Grothendieck topology.

**Definition 2.2.24** (Oversite). Let  $Y = (C, \mathcal{T})$  be a site. Let  $a \in C$ . Define the site  $Y_a$  to be the category  $C_a$  with the induced topology  $\mathcal{T}_a$ . We will denote it by just  $C_a$ .

**Lemma 2.2.25.** *The functor  $u : C_a \rightarrow C$  is a morphism of sites that is cover lifting.*

*Proof.* Clearly  $u$  preserves pullbacks. By Remark 2.2.22 it is also immediate that  $u$  preserves covers and lifts covers. ■

**Definition 2.2.26** (Over ringed site). Let  $Y = (C, \mathcal{T}, \mathfrak{O})$  be a ringed site. Let  $a \in C$ . Define the ringed site  $Y_a$  to be the site  $C_a$  with the structure sheaf  $u_*\mathfrak{O}$  where  $u : C_a \rightarrow C$  the forgetful functor. This makes sense because  $u$  is continuous, see Lemma 2.2.21.

Whenever  $u$  is used as a morphism between ringed sites it stands for the morphism of ringed sites  $(u, \text{Id})$  as defined in Definition 2.2.12.

**Definition 2.2.27.** Let  $(C, \mathcal{T}, \mathfrak{O})$  be a ringed site. Let  $a, b \in C$  and  $f : a \rightarrow b$ .

Consider the induced functor  $C_a \xrightarrow{F} C_b$ . Let  $c$  be the pullback of some pair  $d \rightarrow g \leftarrow e$ . Then  $F(c)$  is the pullback of the pair  $F(d) \rightarrow F(g) \leftarrow F(e)$ , so  $F$  preserves pullbacks.

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Let  $c \rightarrow b \in C_b$ . Let  $R$  be a covering sieve on  $c$ . By Lemma 2.5.9 and the fact that  $u \circ F = u$ , we get that  $F(R)$  generates a covering sieve, so  $F$  is continuous. With the same reasoning one can deduce that  $F$  lifts

Let  $\mathfrak{D}$  be the structure sheaf of  $C_a$  and  $\mathfrak{U}$  be the structure sheaf of  $C_b$ . Note that  $\Gamma(c \xrightarrow{g} b \xrightarrow{f} a; \mathfrak{D}) = \Gamma(c; \mathfrak{D}) = \Gamma(g; \mathfrak{U})$ . Define the morphism of ringed sites  $(F, \text{Id})$ , which we will denote by  $F$  or just  $f$ .

### 2.3 Modules

Presheaf modules and sheaf modules on a ringed site will be introduced in this section. Next we will introduce two main functors  $\lambda$  and  $\Lambda$ . The functors  $\lambda$  and  $\Lambda$  introduced here will be used extensively. See [7, Tag 03A4] for more detail.

**Definition 2.3.1** (Presheaf modules). Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site.

A presheaf module on this ringed site is a presheaf of sets  $\mathfrak{F}$  on  $C$  together with a map of presheaves

$$\mathfrak{D} \times \mathfrak{F} \rightarrow \mathfrak{F}$$

such that for every object  $a \in C$  the map  $\Gamma(a; \mathfrak{D}) \times \Gamma(a; \mathfrak{F}) \rightarrow \Gamma(a; \mathfrak{F})$  defines a  $\Gamma(a; \mathfrak{D})$ -module structure on  $\Gamma(a; \mathfrak{F})$ .

A morphism

$$\mathfrak{F} \rightarrow \mathfrak{G}$$

is a morphism of presheaf modules if

$$\begin{array}{ccc} \mathfrak{D} \times \mathfrak{F} & \longrightarrow & \mathfrak{F} \\ \downarrow & & \downarrow \\ \mathfrak{D} \times \mathfrak{G} & \longrightarrow & \mathfrak{G} \end{array}$$

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commutes. The category of presheaf modules on  $C$  will be denoted  $\text{PMod}(\mathfrak{D})$ . The full subcategory of sheafs in  $\text{PMod}(\mathfrak{D})$  will be denoted  $\text{Mod}(\mathfrak{D})$ .

**Definition 2.3.2.** Let  $\mathfrak{F}$  be a sheaf of modules on  $(C, \mathcal{T}, \mathfrak{D})$ . It is called quasi-coherent if the following holds. For any object  $a \in C$  there exists a covering sieve  $S$  such that for any map  $f : b \rightarrow a$  in  $S$  there exists a presentation

$$\mathfrak{D}|_b^{\oplus I} \rightarrow \mathfrak{D}|_b^{\oplus J} \rightarrow \mathfrak{F}|_b \rightarrow 0$$

It is enough to have presentation for a generating set of  $S$ .

Quasi-coherent modules form a full subcategory of the category of sheafs of modules over  $(C, \mathcal{T}, \mathfrak{D})$  which are denoted by  $\text{Qcoh}(\mathfrak{D})$ .

**Lemma 2.3.3.** *Let  $f : C \rightarrow D$  be a morphism of ringed sites. Let  $\mathfrak{F}$  be a quasi-coherent sheaf of modules on  $D$ . Then  $f_*\mathfrak{F}$  is quasi-coherent.*

*Proof.* ■

**Definition 2.3.4.** Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $R = \Gamma(1; \mathfrak{D})$ . Let  $M, N$  be an  $R$ -module.

Define

$$\lambda : R\text{-Mod} \rightarrow \text{PMod}(\mathfrak{D})$$

by for all  $a \in C$ ,

$$\lambda(M)(a) = M \otimes_R \Gamma(a; \mathfrak{D}),$$

for all  $f : b \rightarrow a \in C$ ,

$$\lambda(M)(f) = \text{Id} \otimes \mathfrak{D}(f),$$

for all  $g : M \rightarrow N \in R\text{-Mod}$ ,

$$\lambda(g) = (a : g \otimes \text{Id}).$$

**Lemma 2.3.5.** *Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $R = \Gamma(1; \mathfrak{D})$ . The functor  $\lambda$  is left adjoint to*

$$\Gamma(1; -) : \text{PMod}(\mathfrak{D}) \rightarrow R\text{-Mod}.$$

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*Proof.* Let  $M, N$  be  $R$ -modules. Let  $\mathfrak{F}, \mathfrak{G} \in \text{PMod}(\mathfrak{D})$  be presheaf modules.

Let  $\varphi : \lambda(M) \rightarrow \mathfrak{G}$  be a morphism of presheaf modules. Let  $\phi : M \rightarrow \Gamma(1; \mathfrak{G})$  be a morphism of  $R$ -modules.

Define

$$\begin{aligned} \alpha : \text{Hom}(\lambda(M), \mathfrak{G}) &\rightarrow \text{Hom}(M, \Gamma(1; \mathfrak{G})) \\ \varphi &\mapsto \varphi_1 \end{aligned}$$

where  $\varphi_1$  is the component of  $\varphi$  on the global sections.

Define

$$\beta : \text{Hom}(M, \Gamma(1; \mathfrak{G})) \rightarrow \text{Hom}(\lambda(M), \mathfrak{G})$$

by, for each  $a \in C$

$$\beta(\phi)_a = \phi \otimes_R \Gamma(a; \mathfrak{D}).$$

We will show that  $\beta$  and  $\alpha$  are mutually inverse.

Let  $\varphi \in \text{Hom}(\lambda(M), \mathfrak{G})$  and  $a \in C$ . Let  $d = \beta(\alpha(\varphi))$ . Let  $m \otimes g \in M \otimes_R \Gamma(a; \mathfrak{D})$ . Let  $p : \lambda(M)(1) \rightarrow \lambda(M)(a)$  be the projection map. Let  $q : \mathfrak{G}(1) \rightarrow \mathfrak{G}(a)$  be the projection map. Then  $d_a(m \otimes g) = \varphi_1(m) \otimes g$  and

$$\begin{aligned} \varphi_a(m \otimes g) &= g\varphi_a(m \otimes 1) \text{ by linearity} \\ &= g\varphi_a(p(m)) \\ &= gq(\varphi_1(m)) \text{ by naturality of } \varphi \\ &= g(\varphi_1(m) \otimes 1) \\ &= \varphi_1(m) \otimes g. \end{aligned}$$

Hence  $d = \varphi$ . In words, the natural transformations from presheaves of the form  $\lambda(M)$  are uniquely determined by their global sections component.

Let  $d = \alpha(\beta(\phi))$ . Let  $m \in M$ . Then  $d(m) = (\phi \otimes_R R)(m) = \phi(m)$ . Hence  $d = \phi$ , which makes  $\alpha$  and  $\beta$  mutual inverses. Naturality is straightforward to check and is omitted. ■



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**Definition 2.3.6.** Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Define

$$\Lambda : \mathbf{R}\text{-Mod} \rightarrow \mathbf{Mod}(\mathfrak{D})$$

by  $\text{sh} \circ \lambda$ .

It follows by that we have the adjunction  $\Lambda \dashv \Gamma(1; -)$ .

This functor is the generalisation of [7, Tag 01BH] to general sites, so  $\Lambda$  on  $\text{Open}(X)$  for some ringed space  $X$  will coincide with the construction defined in the Stacks Project.

### 2.4 Restricting and module functors

**Lemma 2.4.1.** *Let  $f : (D, \mathcal{S}, \mathfrak{U}) \rightarrow (C, \mathcal{T}, \mathfrak{D})$  be a morphism of ringed site. Suppose  $f^\# : f_* \mathfrak{D} \rightarrow \mathfrak{U}$  be an isomorphism. We get a natural isomorphism  $f_* \circ \lambda \Rightarrow \lambda \circ (- \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(1; \mathfrak{U}))$ .*

*Proof.* Define the natural transformation  $t : \lambda \circ (- \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(1; \mathfrak{U})) \Rightarrow \lambda \circ f_*$ , by for each  $\Gamma(1; \mathfrak{D})$ -module  $M$  and for each  $a \in D$ ,

$$\begin{aligned} t_{M,a} : M \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(1; \mathfrak{U}) \otimes_{\Gamma(1; \mathfrak{U})} \Gamma(a; \mathfrak{U}) &\rightarrow M \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(a; \mathfrak{D}), \\ m \otimes r \otimes s &\mapsto m \otimes r f^{\#, -1}(s). \end{aligned}$$

Every component  $t_{M,f}$  is an isomorphism by basic commutative algebra and the fact that  $f^\#$  is an isomorphism ■

**Lemma 2.4.2.** *Let  $F : (C, \mathcal{T}) \rightarrow (D, \mathcal{S})$  be a morphism of sites that lift covers. Then  $F_* \circ (-)^+ = (-)^+ \circ F_*$  and hence also  $F_* \circ \text{sh} = \text{sh} \circ F_*$ .*

*Proof.* Let  $c \in C$ . Let  $R \in \text{Cov}_c$ . Let  $\mathfrak{G}$  be a presheaf on  $D$ . Let  $s = \{s_i \mid i \in R\}$  be a matching family in  $F_* \mathfrak{G}$  on  $c$  indexed by  $R$ . So  $s_i \in \Gamma(F(\text{Dom}(i)); \mathfrak{G})$ . By the cover preserving assumption we know  $F(R)$  generates a covering sieve  $S$  on  $F(c)$ . Define the natural transformation  $T : F_* \circ (-)^+ \rightarrow (-)^+ \circ F_*$  by

$$s \mapsto \{t_j \mid j \in S\}$$

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where  $t_j = F(h)s_i$  for some factorisation  $j = F(i)h$ . Such a factorisation always exists since  $S$  is generated by  $F(R)$ . This assignment is well-defined since  $F$  preserves pullbacks, see proof of [3, Lemma 2.3.3].

We will prove that  $T$  is injective. Let  $s' = \{s'_i \mid i \in R'\}$ . Assume  $T(s) = T(s')$  then  $\{t_j \mid j \in S\} = \{t'_{j'} \mid j' \in S'\}$ , hence there exists covering sieve  $Q \subset S \cap S'$  such that  $t_k = t'_k$  for every  $k \in Q$ . Since  $F$  lifts covers, there exists covering sieve  $P \subset R \cap R'$  on  $c$  such that  $F(P)$  generates a sieve  $P' \subset Q$ . Hence  $t_k = F(h)s_k = F(h)s'_k = t'_k$  with  $k \in P'$  and  $k = F(k)h$ . Take  $h = \text{Id}$  to conclude that  $s = s'$  on  $P$ , hence they are in the same equivalence class.

Now comes surjectivity. Let  $t = \{t_j \mid j \in S\} \in \dots$  on  $F(c)$ . We get a covering sieve  $R$  such that  $F(R) \subset S$ . Consider this matching family now on the covering sieve generated by  $F(R)$ , which amounts to taking a different representing element from the equivalence class. So now  $F(R)$  generates  $S$ . For any  $j \in S$  we get a factorisation  $j = F(i)h$  for some  $i \in R$ . We want to have  $s_i = t_{F(i)}$ . This produces a matching family because restriction in  $F_*\mathfrak{G}$  is exactly the same as in  $\mathfrak{G}$ . Note that now  $T(s) = t$ . Hence we got surjectivity. ■

**Corollary 2.4.3.** *We have  $F_*\omega = \omega$  and hence also  $F_*\omega^2 = \omega^2$*

**Lemma 2.4.4** ( $\Lambda$  commutes with restriction). *Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $f : (C, \mathcal{T}) \rightarrow (D, \mathcal{S})$  be a morphism of sites that lift covers. Let  $f^\#$  be an isomorphism.*

*We have a natural isomorphism*

$$f_* \circ \Lambda \Rightarrow \Lambda \circ (- \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(1; \mathfrak{U})).$$

*Proof.* Follows from Lemmas 2.4.1 and 2.4.2. ■

**Lemma 2.4.5.** *Let  $f : (D, \mathcal{S}, \mathfrak{U}) \rightarrow (C, \mathcal{T}, \mathfrak{D})$  be a morphism of ringed sites that is cover lifting. Suppose  $f^\#$  be an isomorphism. Let  $\mathfrak{F}$  be a quasi-coherent module on  $C$ . Let  $M = \Gamma(1; \mathfrak{F})$ . Consider  $\epsilon : \Lambda(M) \rightarrow \mathfrak{F}$ . Then  $f_*\epsilon : \Lambda(M \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(1; \mathfrak{U})) \rightarrow f_*\mathfrak{F}$  is the counit of the adjunction  $\Lambda \dashv \Gamma(1; -)$  on  $D$ .*

*Proof.* Let  $a \in D$ . We will show that  $f_*\epsilon$  corresponds to the same morphism  $\lambda(M \otimes_{\Gamma(1; \mathfrak{D})}) \rightarrow f_*\mathfrak{F}$  as the counit. By the universal property of the sheaffication, this implies that  $f_*\epsilon$  is the counit.

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We have  $(f_*e)_a(\omega^2(m \otimes r)) = rm$ , which is the same as for the counit. ■

**Example 2.4.6.** Let  $(C, \mathcal{T}, \mathfrak{O})$  be a ringed site. Let  $a \in C$ . Consider the morphism of ringed sites  $u : C_a \rightarrow C$ . Note that  $u^\# = \text{Id}$  and by Lemma 2.2.25 it lifts covers. So all the results from this section apply. This is the most important case.

### 2.5 Schemes

We will recap some definitions in scheme theory that we use later. See [8, 2] for thorough treatments of scheme theory.

**Definition 2.5.1** (Spectrum of a ring). Let  $R$  be a ring. The spectrum  $\text{Spec } R$  of  $R$  is the locally ringed space defined as follows. The underlying set is the set of prime ideals of  $R$ . The (Zariski) topology is generated by the basis of distinguished opens  $D(f) = \{\mathfrak{p} \subset R \mid f \notin \mathfrak{p}\}$ . The sheaf of rings is given on this basis by

$$D(f) \mapsto R_f.$$

A distinguished open  $D(f)$  of  $\text{Spec}(R)$  viewed as locally ringed space is isomorphic to  $\text{Spec}(R_f)$ , where the inclusion  $\text{Spec}(R_f) \rightarrow \text{Spec}(R)$  corresponds to the canonical map  $R \rightarrow R_f$ .

The functor

$$\text{Spec} : \text{Rng} \rightarrow \text{LRSpaces}$$

is left adjoint to

$$\Gamma(1; -) : \text{LRSpaces} \rightarrow \text{Rng},$$

see [7, Tag 01I1].

**Definition 2.5.2** (Scheme). We call the locally ringed space  $\text{Spec}(R)$  an affine scheme.

A scheme  $S$  is a locally ringed space that admits a covering of affine schemes. A morphism of schemes is a morphism of locally ringed spaces. The category of schemes we will denote by  $\text{Sch}$ .

**Definition 2.5.3** (Sheaf of algebras). A sheaf of algebras  $\mathfrak{F}$  on a ringed site  $(C, \mathcal{T}, \mathfrak{O})$  is a sheaf of rings that comes with a (structure) morphism of sheaves of rings  $\mathfrak{O} \rightarrow \mathfrak{F}$ .

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**Definition 2.5.4** (Relative spec). Let  $X$  be a scheme. Let  $\mathfrak{S}$  be a sheaf of algebras on  $X$  that is quasi-coherent as a sheaf of modules.

Define the relative spectrum of  $\mathfrak{S}$  over  $X$  to be the scheme

$$\mathrm{Rspec} \mathfrak{S} \rightarrow X$$

that you get by glueing the spectra  $\mathrm{Spec}(\Gamma(V; \mathfrak{S})) \rightarrow V \subset X$  for every affine open  $V$ . See [7, Tag 01LW].

**Definition 2.5.5** (Tilde functor). Let  $\mathrm{Spec}(R)$  be an affine scheme. Let  $M$  be a  $R$ -module. Define  $\widetilde{M}$  to be the unique sheaf on  $\mathrm{Spec}(R)$  with

$$\widetilde{M} : D(f) \mapsto M_f.$$

See [7, Tag 01HR].

*Remark 2.5.6.* This functor (with this notation) is commonly used in algebraic geometry texts. By [7, Tag 01I7] and uniqueness of left adjoints  $\wedge$  and  $\widetilde{\phantom{x}}$  are canonically isomorphic on spaces, hence on small Zariski sites. We will only use  $\wedge$  in subsequent sections.

**Definition 2.5.7.** A morphism  $f : X \rightarrow Y$  of schemes is called affine if for any affine open  $U \subset Y$  the open subscheme  $f^{-1}(U)$  is affine.

**Definition 2.5.8.** Let  $X$  be a scheme. Let  $f \in \Gamma(X; \mathfrak{D})$ . Define  $X_f$  to be the pullback of  $D(f) \rightarrow \mathrm{Spec}(\Gamma(X; \mathfrak{D})) \leftarrow X$ .

**Lemma 2.5.9.** *Let  $X$  be a scheme. If  $X_{f_i}$  is affine for all  $i$  and  $(f_i)_{i \in I} = (1)$  then  $X$  is affine.*

*Proof.* Let  $\bigcup_i D(f_i) = \mathrm{Spec}(\Gamma(X; \mathfrak{D}))$ . Consider the canonical  $\epsilon : X \rightarrow \mathrm{Spec}(\Gamma(X; \mathfrak{D}))$ . By definition  $\epsilon^{-1}(D(f_i)) = X_{f_i}$ . By assumption  $X_{f_i}$  is affine. By [7, Tag 01S8] the morphism  $\epsilon$  is affine, hence  $X$  is affine. ■

## 2.6 Small Zariski site

In this section we will work on the small Zariski site corresponding to a scheme. We will prove that the unit  $\eta$  and  $\epsilon$  of the adjunction  $\wedge \dashv \Gamma(1; -)$  are isomorphisms in this case.

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**Lemma 2.6.1.** *Let  $R$  be a ring. Let  $M$  be a  $R$ -module. Consider  $\lambda(M), \Lambda(M)$  as sheaf modules on  $\text{Open}(\text{Spec}(R))$ . Then*

$$\omega_{\lambda(M), \text{Spec}(R)}^2 : \Gamma(\text{Spec}(R); \lambda(M)) \rightarrow \Gamma(\text{Spec}(R); \Lambda(M))$$

*is an isomorphism.*

*Proof.* Let  $\mathfrak{p} \subset R$  be a prime ideal. As stated in Definition 2.3.6, we may use results from [7, Tag 01BH] in this setting. We will use that  $\Lambda(M)_{\mathfrak{p}} = M \otimes_R R_{\mathfrak{p}}$ . By naturality, localized at  $\mathfrak{p}$ , the map  $\omega^2$  sends  $m$  to  $m \otimes 1 \in M \otimes_R R_{\mathfrak{p}}$ , hence is the inverse of the multiplication map which is an isomorphism. Hence globally  $\omega^2$  is an isomorphism. ■

**Corollary 2.6.2.** *Let  $X$  be a scheme. Consider  $\lambda(M), \Lambda(M)$  as sheaf modules on  $\text{Open}(X)$ . Let  $\text{Spec}(R)$  be an open subset of  $X$ . Let  $M$  be a  $\Gamma(X; \mathfrak{D})$ -module. Then  $\omega_{\lambda(M), \text{Spec}(R)}^2 : \Gamma(\text{Spec}(R); \lambda(M)) \rightarrow \Gamma(\text{Spec}(R); \Lambda(M))$  is an isomorphism.*

*Remark 2.6.3.* Let us be in the situation of Definition 2.2.15. Note that  $\eta_M$  is the composition of the unit of  $\lambda \dashv \Gamma(1; -)$  and  $\text{sh} \dashv i$ . The first unit is the identity and the second is  $\omega_{\lambda(M), \text{Spec}(R)}^2$ .

**Corollary 2.6.4.** *Consider the adjunction  $\Lambda(-) \dashv \Gamma(1; -)$  on  $\text{Open}(\text{Spec}(R))$ . The unit  $\eta_M : M \rightarrow \Gamma(\text{Spec}(R); \Lambda(M))$  is an isomorphism.*

**Lemma 2.6.5.** *Let  $\mathfrak{F}$  be a sheaf of modules on scheme  $X$ .  $\mathfrak{F}$  is quasi-coherent on  $X$  if and only if for any open  $\text{Spec}(R) \subset X$  the sheaf  $\mathfrak{F}|_{\text{Spec}(R)}$  is isomorphic to  $\Lambda(M)$  with  $M = \Gamma(\text{Spec}(R); \mathfrak{F})$ .*

*Proof.*  $\Rightarrow$ : By assumption we get local presentations indexed by a covering. Let  $\bigcup_{i \in I} U_i = X$  be this covering. Assume without loss of generality that it is an affine open covering. Let  $U_i = \text{Spec}(R_i)$ . Let  $\mathfrak{D}_{U_i}^{\oplus K} \rightarrow \mathfrak{D}_{U_i}^{\oplus J} \rightarrow \mathfrak{F}|_{U_i} \rightarrow 0$  be one of the given presentations. Taking global sections gives us an exact sequence

$$R_i^{\oplus K} \rightarrow R_i^{\oplus J} \rightarrow \Gamma(U_i; \mathfrak{F}) \rightarrow 0.$$

Tensoring it with the localisation  $R_{i,f}$  for any  $f \in R_i$  yields

$$R_{i,f}^{\oplus K} \rightarrow R_{i,f}^{\oplus J} \rightarrow \Gamma(U_i; \mathfrak{F}) \otimes R_{i,f} \rightarrow 0.$$

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Taking sections at  $D(f)$  from the sheaf sequence yields

$$R_{i,f}^{\oplus K} \rightarrow R_{i,f}^{\oplus J} \rightarrow \Gamma(D(f); \mathfrak{F}) \rightarrow 0.$$

Hence  $\mathfrak{F}|_{U_i}$  is the unique sheaf with  $D(f) \mapsto \Gamma(U_i; \mathfrak{F})_f$ , which we defined to be  $\Lambda(\Gamma(U_i; \mathfrak{F}))$ . By the affine communication lemma, this property holds for any affine and not just for the affines in this covering.

$\Leftarrow$ : Let  $M = \Gamma(X; \mathfrak{F})$ . Take a presentation  $R^{\oplus I} \rightarrow R^{\oplus J} \rightarrow M$  and apply  $\Lambda(-)$ . Then note that  $\Lambda(-)$  commutes with arbitrary colimits since it is a left adjoint, see lemma ?. We have  $\Lambda(R) = \mathfrak{O}_{\text{Spec}(R)}$  so we get a presentation  $\mathfrak{O}_{\text{Spec}(R)}^{\oplus I} \rightarrow \mathfrak{O}_{\text{Spec}(R)}^{\oplus J} \rightarrow \mathfrak{F}$  on every affine open subset  $\text{Spec}(R) \subset X$ , hence  $\mathfrak{F}$  is quasi-coherent.  $\blacksquare$

**Corollary 2.6.6.** *Consider the adjunction  $\Lambda(-) \dashv \Gamma(1; -)$ . For quasi-coherent sheaf  $\mathfrak{F}$  the counit  $\epsilon_{\mathfrak{F}} : \Lambda(M) \rightarrow \mathfrak{F}$  is an isomorphism.*

## 2.7 Big Zariski Site

In this section we will introduce the big Zariski ringed site and look at how quasi-coherence and  $\Lambda$  behave on this site. We will prove that the category of quasi-coherent sheaves on an affine scheme  $\text{Spec}(R) \in \text{Sch}$  is equivalent to the category of  $R$ -modules.

**Definition 2.7.1** (Big Zariski site). Define the big Zariski site to be  $(\text{Sch}, \mathcal{T}, \mathfrak{O})$  with the following components. The underlying category is  $\text{Sch}$ . The topology  $\mathcal{T}$  is generated by the basis consisting of the covering families  $\{X_i \xrightarrow{f_i} X\}$  where  $f_i$  is an open immersion and  $\bigcup_i f_i(X_i) = X$ . The sheaf of rings  $\mathfrak{O}$  sends  $(U, \mathfrak{O}) \rightarrow (X, \mathfrak{O})$  to  $\Gamma(U; \mathfrak{O})$ .

We will mostly be interested in the site  $\text{Sch}_X$  for a scheme  $X$ .

**Definition 2.7.2** (k). Let  $X$  be a scheme. Define the functor  $k : \text{Open}(X) \rightarrow \text{Sch}_X$  by  $U \mapsto ((U, \mathfrak{O}_U), i)$  where  $i : U \rightarrow X$  is the inclusion of the open subscheme into  $X$ .

We will show that it preserves limits and covers. The terminal  $X \in \text{Open}(X)$  is sent to the terminal  $X \rightarrow X$ . Let  $U \rightarrow V$  and  $W \rightarrow V$  be two morphism in  $\text{Open}(X)$ . We have  $k(U \cap W) = U \cap W \rightarrow X$  which is the pullback of  $k(U) \rightarrow k(V)$  and  $k(W) \rightarrow k(V)$ .

Let  $S = \{D(f_i) \rightarrow \text{Spec}(R)\}$  be one of the generating family in  $\text{Open}(X)$ . Note that  $k(D(f_i))$  is isomorphic to the object  $\text{Spec}(R_{f_i}) \rightarrow \text{Spec}(R)$ . Hence  $k(S)$  generates a covering sieve on  $\text{Sch}$  and hence on  $\text{Sch}_X$ . We established that  $k$  is a morphism of sites.

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Consider  $(\text{Open}(X), \mathcal{T}, \mathfrak{O})$  and  $(\text{Sch}_X, \mathcal{S}, \mathfrak{U})$ . Note that  $k_*\mathfrak{U} = \mathfrak{O}$  by construction. Define the morphism of ringed sites  $(k, \text{Id})$ . We will denote this by  $k$ .

**Lemma 2.7.3.** *Let  $X$  be a scheme. The morphism of sites  $k : \text{Open}(X) \rightarrow \text{Sch}_X$  is cover lifting.*

*Proof.* Let  $\mathfrak{U} \in \text{Open}(X)$ . Let  $W = \{\mathfrak{U}_i \xrightarrow{f_i} k(\mathfrak{U})\}$  be a covering family on  $k(\mathfrak{U})$ . Then  $V = \{f_i(\mathfrak{U}_i) \rightarrow \mathfrak{U}\}$  is a covering family on  $\mathfrak{U}$  such that  $k(V)$  generates the same sieve as  $W$ . Hence  $k$  is cover lifting.  $\blacksquare$

**Lemma 2.7.4.** *Let  $X$  be a scheme. Consider  $k : \text{Open}(X) \rightarrow \text{Sch}_X$ . Then  $k_*\Lambda = \Lambda$*

*Proof.* By definition  $k^\#$  is an isomorphism. By Lemma 2.7.3 we get that  $k$  is cover lifting. Then we can apply Lemma 2.4.4. Let  $\mathfrak{U}$  be the structure sheaf of  $\text{Sch}_X$ . Let  $\mathfrak{O}$  be the structure sheaf of  $\text{Open}(X)$ . Note that the structure morphism  $\Gamma(1; k_*\mathfrak{U}) \rightarrow \Gamma(1; \mathfrak{O})$  is an isomorphism, so Lemma 2.4.4 implies this result.  $\blacksquare$

**Lemma 2.7.5.** *Let  $X = \text{Spec}(R)$  be a scheme. Let  $\mathfrak{F}$  be a quasi-coherent sheaf on  $\text{Sch}_X$ . Let  $M = \Gamma(X; \mathfrak{F})$ . Then  $\epsilon_{\mathfrak{F}} : \Lambda(M) \rightarrow \mathfrak{F}$  is an isomorphism.*

*Proof.* By Lemma 2.3.3,  $k_*$  sends quasi-coherent sheafs to quasi-coherent sheafs. so  $k_*\mathfrak{F} = \Lambda(M)$ , where  $k : \text{Open}(X) \rightarrow \text{Sch}_X$ . Hence  $\Gamma(D(f); \mathfrak{F}) = \Gamma(D(f); \Lambda M) = M_f$  for any  $f \in R$ .

Let  $D(f_i)|_{\mathfrak{O}} \xrightarrow{\alpha_i} D(f_i)|_{\mathfrak{O}}^{\oplus K} \rightarrow D(f_i)|_{\mathfrak{F}} \rightarrow 0$  be a presentation with  $(f_i) = (1)$ . Note that the presheaf cokernel  $\text{Coker}(\alpha_i)$  of the sheaf morphism is  $\lambda(\text{Coker}(\alpha_{i, D(f_i)}))$  where  $\alpha_{i, D(f_i)}$  is the component at  $D(f_i)$  of  $\alpha_i$ . So  $D(f_i)|_{\mathfrak{F}} = \Lambda(M_f)$  since  $\omega^2$  is iso for affines.

Hence

$$\begin{aligned} \Gamma(D(f) \times_X \text{Spec}(S); \mathfrak{F}) &= \Gamma(D(f) \times_X \text{Spec}(S); \Lambda \mathfrak{F}) \\ &= \Gamma(\text{Spec}(S_f); \Lambda M_f) \\ &= M_f \otimes S_f \\ &= M \otimes S_f. \end{aligned}$$

By the sheaf property it follows that  $\Gamma(\text{Spec}(S); \mathfrak{F}) = \Gamma(X; \mathfrak{F}) \otimes S$ , hence the counit  $\Lambda(M) \rightarrow \mathfrak{F}$  is an isomorphism.  $\blacksquare$

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**Lemma 2.7.6.** *Let  $X = \operatorname{Spec}(R)$  be a scheme. Let  $M$  be a  $R$ -module. We work on the ringed site  $\operatorname{Sch}_X$ . The unit  $\eta : M \rightarrow \Gamma(X; \Lambda(M))$  is an isomorphism*

*Proof.* Let  $\Lambda_a(M)$  be  $\Lambda(M)$  over  $\operatorname{Open}(X)$ . By Lemma 2.7.4, we have  $\Lambda_a(M) = k_* \Lambda(M)$ . Hence  $\Gamma(X; \Lambda(M)) = \Gamma(X; \Lambda_a(M))$ . Moreover  $\Gamma(X; \Lambda(M)) = M$  by Corollary 2.6.4. ■



### 3 Restrictive morphisms

This section will introduce the notion of a restrictive morphism. We will see some examples, non-examples and results in the category of schemes and see that this notion is closely related to affinenes. We will always work over small sites in this section. For some of the examples and results see the chapter on quasi-coherent modules in [8].

**Definition 3.0.1** (Restrictive morphism). Let  $(C, \mathcal{T}, \mathfrak{D})$ . A morphism  $f : a \rightarrow b \in C$  is called restrictive if for every quasi-coherent module  $\mathfrak{G}$  on  $C_b$  the morphism

$$\widehat{f} : \Gamma(b; \mathfrak{G}) \otimes_{\Gamma(b; \mathfrak{D})} \Gamma(a; \mathfrak{D}) \rightarrow \Gamma(a; \mathfrak{G}) \quad (3.1)$$

is an isomorphism.

*Remark 3.0.2.* Let us be in the setting of Definition 3.0.1. Assume  $\mathfrak{G} = \Lambda(M)$  where  $M = \Gamma(1; \mathfrak{G})$ . Consider  $\omega_{\mathfrak{G}, a}^2 : \lambda(M)(a) \rightarrow \Lambda(M)(a)$ . This map is equal to  $\widehat{f}$ .

**Example 3.0.3.** In  $\text{Sch}_{\text{Spec}(A)}$  the morphism  $\text{Spec}(A_f) \rightarrow \text{Spec}(A)$  is restrictive. Let  $\mathfrak{G}$  be a quasi-coherent sheaf on  $\text{Sch}_{\text{Spec}(A)}$ . This implies that  $\mathfrak{G} = \Lambda(\Gamma(\text{Spec}(A); \mathfrak{G}))$ , see ?. The morphism

$$\Gamma(\text{Spec}(A); \mathfrak{G}) \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(\text{Spec}(A_f); \mathfrak{D}) \rightarrow \Gamma(\text{Spec}(A_f); \mathfrak{G}) = \Gamma(A; \mathfrak{G})_f,$$

$$m \otimes r \rightarrow rm$$

is an isomorphism by basic commutative algebra.

**Example 3.0.4.** Let  $R$  be a ring. Consider the open immersion  $U = \text{Spec}(R[x, y]) \setminus \{(x, y)\} \rightarrow \text{Spec}(R[x, y])$  and the quasi-coherent sheaf  $\mathfrak{G} = \Lambda(\frac{R[x, y]}{xy})$ . The global sections of this sheaf are  $\frac{R[x, y]}{xy}$ , as shown in Lemma 2.6.1.

Define  $U_1 = D(x) \rightarrow U$  and  $U_2 = D(y) \rightarrow U$ . Note that these cover  $U$  together. We have  $\Gamma(U_1; \mathfrak{G}) = 0$  and  $\Gamma(U_2; \mathfrak{G}) = 0$ , since  $\frac{R[x, y]}{xy} \cdot x = 0$  and  $\frac{R[x, y]}{xy} \cdot y = 0$ . Hence since  $\mathfrak{G}$  is a sheaf, we get  $\Gamma(U; \mathfrak{G}) = 0$ .

### 3 Restrictive morphisms

The sections over  $U$  of  $\Lambda(R[x, y])$  are (also)  $R[x, y]$ . See [8, p. 4.4.1]. We conclude that  $\Gamma(1; \mathfrak{G}) \otimes_{\Gamma(1; \mathfrak{D})} \Gamma(U; \mathfrak{D}) \rightarrow \Gamma(U; \mathfrak{G})$  is not an isomorphism.

**Lemma 3.0.5.** *Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $f : b \rightarrow a \in C$  and  $g : c \rightarrow b \in C$  be morphisms.*

1. *If  $fg$  and  $f$  are restrictive, then  $g$  is.*
2. *if  $f$  and  $g$  are restrictive, then  $fg$ .*

*If  $\Gamma(b; \mathfrak{D}) \rightarrow \Gamma(c; \mathfrak{D})$  is faithfully flat then  $fg$  and  $g$  restrictive implies  $f$  is restrictive.*

*Proof.* Consider the following diagram

$$\begin{array}{ccc} \Gamma(a; \mathfrak{G}) \otimes_{\Gamma(a; \mathfrak{D})} \Gamma(c; \mathfrak{D}) & \xrightarrow{\widehat{f} \otimes_{\Gamma(b; \mathfrak{D})} \Gamma(c; \mathfrak{D})} & \Gamma(b; \mathfrak{G}) \otimes_{\Gamma(b; \mathfrak{D})} \Gamma(c; \mathfrak{D}) \\ \parallel & & \downarrow \widehat{g} \\ \Gamma(a; \mathfrak{G}) \otimes_{\Gamma(a; \mathfrak{D})} \Gamma(c; \mathfrak{D}) & \xrightarrow{\widehat{fg}} & \Gamma(c; \mathfrak{G}), \end{array}$$

where  $\widehat{\phantom{x}}$  is as in lemma ?.

This diagram commutes: going either direction sends  $g \otimes r$  to  $rg$ . The results follows from commutativity. ■

**Lemma 3.0.6** (coproduct). *Let  $X_1, X_2, Y$  be a schemes.  $X_1 \xrightarrow{f_1} Y$  and  $X_2 \xrightarrow{f_2} Y$  are restrictive morphisms if and only if the corresponding morphism  $X_1 \sqcup X_2 \xrightarrow{f_1 \sqcup f_2} Y$  is restrictive.*

*Proof.*  $\Rightarrow$ :

Note that  $\Gamma(X_1 \sqcup X_2; -) = \Gamma(X_1; -) \times \Gamma(X_2; -)$  by the sheaf property. We will show that

$$\widehat{f_1 \sqcup f_2} : \Gamma(Y; \mathfrak{G}) \otimes_{\Gamma(Y; \mathfrak{D})} \Gamma(X_1 \sqcup X_2; \mathfrak{D}) \rightarrow \Gamma(X_1 \sqcup X_2; \mathfrak{G})$$

is an isomorphism. Tensor commutes over products, so this becomes

### 3 Restrictive morphisms

$$\widehat{f}_1 \times \widehat{f}_2 : (\Gamma(Y; \mathfrak{G}) \otimes_{\Gamma(Y; \mathfrak{D})} \Gamma(X_1; \mathfrak{D})) \times (\Gamma(Y; \mathfrak{G}) \otimes_{\Gamma(Y; \mathfrak{D})} \Gamma(X_2; \mathfrak{D})) \rightarrow \Gamma(X_1; \mathfrak{G}) \times \Gamma(X_2; \mathfrak{G}).$$

By assumption  $\widehat{f}_1$  and  $\widehat{f}_2$  are isos, so their product is.

$\Leftarrow$ : By Lemma 3.0.7, the canonical morphism  $X_i \rightarrow X_1 \sqcup X_2$  is restrictive. By assumption  $X_1 \sqcup X_2 \rightarrow Y$  is restrictive. Composing these morphisms yields  $f_i$ , by Lemma 3.0.5 this yields a restrictive morphism. ■

**Lemma 3.0.7.** *Let  $X, Y$  be a schemes. The canonical morphism  $X \rightarrow X \sqcup Y$  is restrictive.*

*Proof.* Let  $\mathfrak{G}$  be a quasi-coherent sheaf on  $X \sqcup Y$ . By the sheaf property  $\Gamma(X \sqcup Y; \mathfrak{G}) = \Gamma(X; \mathfrak{G}) \times \Gamma(Y; \mathfrak{G})$ . The same holds also for  $\mathfrak{D}$ .

We are considering the morphism

$$\Gamma(X \sqcup Y; \mathfrak{G}) \otimes_{\Gamma(X \sqcup Y; \mathfrak{D})} \Gamma(X; \mathfrak{D}) \rightarrow \Gamma(X; \mathfrak{G}).$$

By the previous remark about disjoint unions and the sheaf property and some basic commutative algebra one sees that this becomes

$$(\Gamma(X; \mathfrak{G}) \otimes_{\Gamma(X \sqcup Y; \mathfrak{D})} \Gamma(X; \mathfrak{D})) \times (\Gamma(Y; \mathfrak{G}) \otimes_{\Gamma(X \sqcup Y; \mathfrak{D})} \Gamma(X; \mathfrak{D})) \rightarrow \Gamma(X; \mathfrak{G}).$$

Since  $\Gamma(Y; \mathfrak{G}) \otimes_{\Gamma(X \sqcup Y; \mathfrak{D})} \Gamma(X; \mathfrak{D}) = 0$ , we are left with

$$\begin{aligned} & (\Gamma(X; \mathfrak{G}) \otimes_{\Gamma(X \sqcup Y; \mathfrak{D})} \Gamma(X; \mathfrak{D})) \rightarrow \Gamma(X; \mathfrak{G}) \\ & g \otimes r \rightarrow rg. \end{aligned}$$

Note that  $\Gamma(X; \mathfrak{G})$  already is an  $\Gamma(X; \mathfrak{D})$ -module and conclude that hence this morphism is an isomorphism. ■

**Lemma 3.0.8** (Restrictive to affines). *If  $f : X \rightarrow \text{Spec}(\mathbb{R})$  is a restrictive open immersion, then  $X$  is affine.*

### 3 Restrictive morphisms

*Proof.* Since  $X$  is an open in  $\text{Spec}(R)$ , we get a distinguished covering  $\bigcup_i D(f_i) = X$  with  $f_i \in R$  and  $i \in I$ . We will prove that  $(f_i|_X)_{i \in I} = (1)$  in  $S = \Gamma(X; \mathfrak{D})$ .

Then we invoke the result in Lemma 2.5.9 that states the following. For a scheme  $Y$  let  $Y_f$  be the support of  $f \in \Gamma(Y; \mathfrak{D})$  as in lemma ?. if  $Y_{g_j}$  are affine and  $(g_j)_{j \in J} = (1)$  then  $Y$  is affine.

Note that  $D(f_i) = X_{f_i|_X}$ . Consider  $M = \frac{R}{(f_i)}$  as an  $R$ -module and look at  $\Lambda(M)$ . By restrictiveness we get  $M \otimes_R S = \Lambda(M)(S)$  and by  $M \otimes_R R_{f_i} = \Lambda(M)(D(f_i)) = M_{f_i} = 0$ . Hence  $\Lambda(M)(S) = 0$  by the sheaf axiom. This implies that  $(f_i|_X)_{i \in I} = (1)$  in  $S$ . So  $X$  is affine. ■

**Lemma 3.0.9.** *Any morphism  $\text{Spec}(S) \xrightarrow{f} \text{Spec}(R) \in \text{Sch}_{\text{Spec}(R)}$  between affine schemes is restrictive.*

*Proof.* Let  $\mathfrak{G}$  be a quasi-coherent module on  $\text{Sch}_{\text{Spec}(R)}$ . Set  $M = \Gamma(\text{Spec}(R); \mathfrak{G})$ . We want to prove that

$$\hat{f} : M \otimes_R S \rightarrow \Gamma(\text{Spec}(S); \mathfrak{G})$$

is an isomorphism.

By Lemma 2.7.5, we get  $\mathfrak{G} = \Lambda(M)$ . As said in Remark 3.0.2, in this case  $\hat{f} = \omega_{\text{Spec}(S)}^2$ . By Definition 2.2.15, we know that  $\omega^2$  is an isomorphism at affine schemes. ■

**Example 3.0.10** (Affine non-restrictive map). One might expect(or want) that any property of all maps between affine schemes also hold for affine maps between any schemes. This is not the case for restrictiveness, so it is not local on the target.

Consider the canonical inclusions  $\mathbb{A}^1 \rightarrow \mathbb{P}^1$  and the shifted quasi-coherent module  $\mathfrak{D}(-1)$ . This module is locally free of degree 1, this is often called an invertible module.

The global sections of the module  $\mathfrak{D}(-1)$  are the elements of degree  $-1$  in the global sections of  $\mathfrak{D}$ . There are no such elements, hence the global sections are the zero module.

On  $\mathbb{A}^1$  all invertible modules are isomorphic to the structure sheaf. See [8, p. 14.2.8].

Similary any inclusion  $\text{Spec}(\kappa(p)) \rightarrow \mathbb{P}^1$  of a point is not restrictive which can be shown with the same argument.

### *3 Restrictive morphisms*

This is a (more opaque) way of saying that on projective space not every quasi-coherent sheaf is generated by global sections.

## 4 Caffine objects

In this chapter we will define the generalisation of an affine scheme to general ringed sites. This will be called ‘caffine’. We will prove that caffine and affine are equivalent for schemes. Lastly we will see some examples and uses.

### 4.1 Caffine objects

In this section we will introduce caffine objects, see some examples and non-examples and prove some properties. Most of these results will be generalisations of their counterpart for affine schemes.

**Definition 4.1.1** (Caffine object). Let  $(C, \mathcal{T}, \mathcal{O})$  be a ringed site. Let  $\alpha \in C$  be an object. We call  $\alpha$  caffine if the unit  $\eta$  and co-unit  $\epsilon$  of the adjunction  $\Lambda \dashv \Gamma(\alpha; -)$  on  $C_\alpha$  are natural isomorphisms for any  $\Gamma(\alpha; \mathcal{O})$ -module  $M$  and any quasi-coherent  $\mathfrak{F}$  on  $C_\alpha$ .

In other words, that  $\Lambda : \Gamma(\alpha; \mathcal{O})\text{-Mod} \rightarrow \text{Qcoh}(\mathcal{O})$  is an equivalence of categories with  $\Gamma(\alpha; -)$  as pseudo-inverse and the unit and co-unit as witnessing natural isomorphisms.

**Example 4.1.2** (Examples of caffine objects). The main example to keep in mind is  $\text{Spec}(\mathbb{R}) \in \text{Sch}$ . See Lemmas 2.7.5 and 2.7.6 for proofs that the unit and co-unit are isomorphisms.

**Example 4.1.3.** The scheme  $\mathbb{P}^1$  is not caffine in  $\text{Sch}$ . The counit at the quasi-coherent sheaf of modules  $\mathcal{O}(-1)$  has signature  $\epsilon : \Lambda(0) \rightarrow \mathcal{O}(-1)$ . Since  $\mathcal{O}(-1)$  is locally free of degree 1, this cannot be an isomorphism.

**Definition 4.1.4** (caffine cover). Let  $(C, \mathcal{T}, \mathcal{O})$  be a ringed site. A family of maps  $\{\alpha_i \rightarrow \alpha\}$  is called a caffine covering of  $\alpha$  if every  $\alpha_i$  is caffine and the family is a covering family.

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**Definition 4.1.5.** We say that a ringed site  $(C, \mathcal{T}, \mathfrak{D})$  has enough affines if any object admits a affine covering.

**Lemma 4.1.6.** *Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $a \in C$ . Let  $\{b_i \rightarrow a\}$  be a affine covering on  $a$ . Assume every map  $b_i \xrightarrow{i} a$  is restrictive. Then the counit  $\epsilon$  of the adjunction  $\Lambda_a(-) \dashv \Gamma(a; -)$  is a natural isomorphism.*

*Proof.* Let  $\mathfrak{F}$  be a quasi-coherent sheaf module. Set  $M = \Gamma(a; \mathfrak{F})$ . Set  $M_i = \Gamma(b_i; \mathfrak{F})$ . Set  $\beta_i = i_* \epsilon_{\mathfrak{F}, a}$ . By Definition 2.2.27 the induced functor of  $i$  satisfies all requirements of Lemma 2.4.5, so we have  $i_* \epsilon_{\mathfrak{F}, a} \cong \epsilon_{i_* \mathfrak{F}, b_i}$ . Since  $b_i$  is affine  $\beta_i$  is an isomorphism. Hence  $\epsilon$  is an isomorphism between sheafs, since it is an isomorphism restricted along every morphism from a covering. ■

**Lemma 4.1.7.** *Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $a \in C$ . Let  $M$  be a  $\Gamma(a; \mathfrak{D})$ -module. The component*

$$\omega_{\lambda(M), a}^2 : \lambda(M)(a) \rightarrow \Lambda(M)(a)$$

*at  $Id_a$  of the sheafification morphism*

$$\omega_{\lambda(M)}^2 : \lambda(M) \rightarrow \Lambda(M)$$

*is equal to the unit of  $\Lambda \dashv \Gamma(1; -)$  in  $C_a$ .*

*Proof.* Consider the following maps, which you get by repeatedly calling on an adjunction.

$$Id : \Lambda(M) \rightarrow \Lambda(M)$$

$$\omega_{\lambda(M)}^2 : \lambda(M) \rightarrow \Lambda(M) \text{ use sheafification adjunction}$$

$$\omega_{\lambda(M), a}^2 M \rightarrow \Gamma(a; \Lambda(M)) \text{ take sections at } a$$

We took the adjunct of  $Id$  with respect to the sheafification adjunction as the first step. Then we took the adjunct of the result wrt the  $\lambda$  adjunction. Hence we get the adjunct of  $Id$  wrt the  $\Lambda$  adjunction. so the last map is actually the unit of the  $\Lambda$  adjunction. ■

**Corollary 4.1.8.** *Let  $(C, \mathcal{T}, \mathfrak{D})$  be a ringed site. Let  $a \in C$  be affine. Then  $\omega_{\lambda(M), a}^2$  is an isomorphism for any  $\mathfrak{D}(a)$ -module  $M$ .*

## 4 Affine objects

**Theorem 4.1.9** (Morphism between affines is restrictive). *Let  $Y = (C, \mathcal{T}, \mathfrak{D})$ . Let  $f : b \rightarrow a \in C$  be a morphism between affine objects, then  $f$  is restrictive.*

*Proof.* Let  $\mathfrak{F}$  be a quasi-coherent module on  $Y_a$ . Let  $M = \Gamma(a; \mathfrak{F})$ . Since  $a$  is affine, we have  $\mathfrak{F} = \Lambda(M)$ .

We have to show that

$$\widehat{f} : \Gamma(a; \mathfrak{F}) \otimes_{\Gamma(a; \mathfrak{D})} \Gamma(b; \mathfrak{D}) \rightarrow \Gamma(b; \mathfrak{F})$$

is an isomorphism, where  $\widehat{f}$  as defined in Definition 3.0.1.

Note that  $\widehat{f}$  is the component at  $b$  of the natural transformation  $\omega_{\lambda(\Gamma(1; \mathfrak{F}))}^2$ , see Remark 3.0.2. Since  $b$  is affine, this component is an isomorphism by Corollary 4.1.8. ■

## 4.2 Affine schemes

Let  $X \in \text{Sch}$ .

Set  $R = \Gamma(X; \mathfrak{D})$ . We will prove that the canonical morphism

$$\epsilon_X : X \rightarrow \text{Spec}(R)$$

is an isomorphism if  $X$  is a affine object. This morphism sends  $x \in X$  to  $\ker x$  as defined later.

We will start this section with some general results that work for any scheme.

**Definition 4.2.1** (Support of module). Let  $\mathfrak{M}$  be a  $\mathfrak{D}$ -module. The support  $\text{Supp}(\mathfrak{M})$  of this module is the subspace

$$\{x \in X \mid \mathfrak{M}_x = 0\} \subset X.$$

**Definition 4.2.2** (Set cut out by ideal). Define the set cut out by an ideal  $I$  of  $R$  to be

$$V_X(I) = \text{Supp}(\Lambda_X(\frac{R}{I})).$$

**Definition 4.2.3.** For a global section  $a \in R$  define

$$D_X(a) = \{x \in X \mid a \notin \ker(x)\}.$$



#### 4 Affine objects

The following lemma will be the tool we use to prove that  $X$  is affine if it is affine.

**Lemma 4.2.4.** *The following are equivalent:*

1.  $X$  is affine
2.  $\epsilon_X$  is an isomorphism
3.  $\epsilon_X$  is a homeomorphism

*Proof.*  $1 \Rightarrow 2$ : Assume that  $X = \text{Spec } R$ . Then  $\ker \mathfrak{p} = \mathfrak{p}$  for  $\mathfrak{p} \in X$ , so the topological part of  $\epsilon_X$  is an homeomorphism. The sheaf part of  $\epsilon_X$  sends the global section  $1 \in \mathcal{O}(X)$  to  $1 \in R$ , since the sheaf morphism is a  $\mathcal{O}_{\text{Spec}(R)}$ -algebra map, it has to be the identity.

$2 \Rightarrow 3$ : By definition.

$3 \Rightarrow 1$ : Let  $\text{Spec } A_i = U_i \subset X$  be open affines and suppose  $\bigcup_i U_i = X$ . Assume it is a finite affine cover, which can be done since  $X$  is quasi-compact. Using our base, we get a cover of  $U_i = \bigcup_j D_X(a_{ij})$  with  $a_{ij}$  global sections. Observe that  $D_X(a_{ij}) \subset U_i$ , hence  $D_{U_i}(a_{ij}|_{U_i}) = D_X(a_{ij})$  which makes them affine. Continuing like this, we get a finite cover of affines  $D_X(a_{ij})$  of  $X$ . Since

$$\epsilon_X(X) = \epsilon_X\left(\bigcup_{ij} D_X(a_{ij})\right) = \bigcup_{ij} D_{\text{Spec}(R)}(a_{ij}) = \text{Spec } R,$$

we have  $(a_{ij}) = (1)$ . Now both requirements of Lemma 2.5.9 are satisfied, hence  $X$  is affine. ■

**Lemma 4.2.5 (Stalks).** *Let  $M$  be a  $\Gamma(X; \mathcal{O})$ -module. Let  $x \in X$ . Then  $\Lambda(M)_x = M_{\mathcal{O}_x}$ .*

*Proof.* See [7, Tag 01BH]. ■

**Lemma 4.2.6.** *For  $x \in X$  TFAE:*

1.  $x \in V_X(I)$
2.  $I_{\mathcal{O}_x} \neq \mathcal{O}_x$
3.  $I \subset \ker(x)$ .

*Proof.*  $1 \Rightarrow 2$ :

#### 4 Affine objects

Assume  $x \in V_X(I)$ . We have  $\Lambda_X(\frac{R}{I}) = \frac{\Lambda_X(R)}{\Lambda_X(I)}$  and taking stalks is exact so  $\frac{\Lambda_X(R)}{\Lambda_X(I)}_x = \frac{\Lambda_X(R)_x}{\Lambda_X(I)_x}$ . Lastly by Lemma 4.2.5 we get  $\Lambda_X(\frac{\mathfrak{D}(X)}{I})_x = \frac{\mathfrak{D}_x}{I\mathfrak{D}_x}$ . Since  $x \in V_X(I)$  we know  $\frac{\mathfrak{D}_x}{I\mathfrak{D}_x} \neq 0$ , hence  $I\mathfrak{D}_x \neq \mathfrak{D}_x$ .

2  $\Rightarrow$  3:

Assume  $I\mathfrak{D}_x \neq \mathfrak{D}_x$ . Then  $I\mathfrak{D}_x$  is proper hence contained in the unique maximal ideal of the local ring  $\mathfrak{D}_x$ , therefore  $I \mapsto 0$  in  $k(x)$  or equivalently  $I \subset \ker(x)$ .

3  $\Rightarrow$  1:

Assume  $I \subset \ker(x)$ . Then  $I$  maps into  $\mathfrak{m}_x$ , hence  $I\mathfrak{D}_x \subset \mathfrak{m}_x$ . Therefore

$$\frac{\mathfrak{D}_x}{\Lambda_X(I)_x} = \frac{\mathfrak{D}_x}{I\mathfrak{D}_x} \neq 0.$$

■

**Corollary 4.2.7.** *If  $y \in I$  then  $D_X(y) \cap V_X(I) = \emptyset$ .*

*Proof.* Assume  $y \in I$ . Let  $z \in V_X(I)$ , then  $y \in \ker(z)$  by Lemma 4.2.6. This implies  $z \notin D_X(y)$ . ■

**Corollary 4.2.8.**  $V_X(I) \cup V_X(J) = V_X(IJ)$

*Proof.* Let  $z \in V_X(I) \cup V_X(J)$ . Then  $I \subset \ker(z)$  and  $J \subset \ker(z)$  by Lemma 4.2.6, hence  $IJ \subset \ker(z)$ . Apply Lemma 4.2.6 again to get  $z \in V_X(IJ)$ . Let  $z \in V_X(IJ)$ . Then  $IJ \subset \ker(z)$  by Lemma 4.2.6. The ideal  $\ker(z)$  is prime, so  $I \subset \ker(z)$  or  $J \subset \ker(z)$ . Invoke Lemma 4.2.6 again to get  $z \in V_X(I) \cup V_X(J)$ . ■

**Lemma 4.2.9.** *Let  $I \subset \Gamma(X; \mathfrak{D})$  be an ideal. The set  $V_X(I)$  is closed.*

*Proof.* Consider the exact sequence

$$\mathfrak{D}(X) \rightarrow \frac{\mathfrak{D}(X)}{I} \rightarrow 0.$$

The functor  $\Lambda$  is a left adjoint hence right exact so

$$\mathfrak{D} \xrightarrow{f} \Lambda\left(\frac{\mathfrak{D}(X)}{I}\right) \rightarrow 0$$

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is exact. Hence the sequence

$$\mathfrak{D}_x \xrightarrow{f_x} \Lambda\left(\frac{\mathfrak{D}(X)}{I}\right)_x \rightarrow 0$$

is exact. So we get that  $f_x(1_x)$  generates  $\Lambda\left(\frac{\mathfrak{D}(X)}{I}\right)_x$ .

Note that  $f_x(1_x) = f(1)_x$  by definition of  $f_x$ , hence  $f(1)_x$  is a generating element. Hence  $\Lambda\left(\frac{\mathfrak{D}(X)}{I}\right)_x \neq 0$  if and only if  $f(1)_x \neq 0$ .

This implies  $V_X(I) = \text{Supp}(f(1))$  which makes  $V_X(I)$  closed as the support of a global section. ■

From now on assume that  $X$  is a affine object in  $\text{Sch}$ .

**Lemma 4.2.10.** *Every closed set  $W \subset X$  can be written as  $V_X(I)$  for some ideal  $I \subset \Gamma(x; \mathfrak{D})$ .*

*Proof.* Let  $\mathfrak{I}$  be some ideal sheaf inducing a closed subscheme structure on  $W$ . This is by definition a quasi-coherent module. Let  $I = \Gamma(x; \mathfrak{I})$ . Since  $X$  is affine, we get  $\mathfrak{I} = \Lambda(I)$ . Let  $\mathfrak{D}_W$  be the structure sheaf of this closed subscheme. Consider the closed immersion  $W \xrightarrow{i} X$ . By construction  $\mathfrak{I} \rightarrow \mathfrak{D} \rightarrow i_*\mathfrak{D}_W \rightarrow 0$  is exact, hence  $i_*\mathfrak{D}_W = \frac{\mathfrak{D}}{\mathfrak{I}}$ . Hence  $V_X(I) = \text{Supp}(i_*\mathfrak{D}_W) = W$ . ■

**Lemma 4.2.11.** *The sets  $D_X(a)$  form a basis for the topology of  $X$ , with  $a \in R$ .*

*Proof.* Let  $x \in D_X(a)$ . Then  $a \notin \ker x$ , hence  $\frac{1}{a}$  exists in  $\mathfrak{D}_x$ . We get an open neighbourhood  $V$  of  $x$  such that  $\frac{1}{a} \in \Gamma(V; \mathfrak{D})$ . Let  $y \in V$ , then clearly  $a \notin \ker y$ . We have shown that  $x \in V \subset D_X(a)$ , hence  $D_X(a)$  is open.

Let  $U \subset X$  be any open. Let  $x \in U$ . By Lemma 4.2.10 we get  $I$  such that  $V_X(I) = U^c$ . It follows that  $x \notin V_X(I)$  and  $I \not\subset \ker(x)$  by Lemma 4.2.6. So we get a  $g \in I$  with  $g \notin \ker(x)$ . We get  $x \in D_X(g)$  and  $D_X(g) \subset U$ , because  $D_X(g) \cap V_X(I) = \emptyset$  by Corollary 4.2.7.

Let  $a, b \in R$ . Note that  $D_X(ab) = D_X(a) \cap D_X(b)$  since  $\ker(x)$  is a prime ideal. So the opens  $D_X(a)$  form a basis. ■

**Lemma 4.2.12.** *The map  $e_X : X \rightarrow \text{Spec}(R)$  is surjective.*

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*Proof.* Let  $p \in \text{Spec } R$  be a point in the target of  $\epsilon_X$ . Then  $\Lambda_X(\kappa(p))$  is a quasi-coherent sheaf of modules. In fact  $\kappa(p) \otimes_{\Gamma(X; \mathcal{O})} \mathcal{O}(U)$  is a  $\mathcal{O}(U)$  algebra for any open  $U \subset X$ , hence  $\Lambda_X(\kappa(p))$  is a quasi-coherent sheaf of algebras. Hence we can compute the relative spec  $\text{Rspec}(\Lambda_X(\kappa(p))) \xrightarrow{h} X$ .

Since  $X$  is affine we have  $\Lambda_X(\kappa(p))(X) = \kappa(p)$ . So

$$\text{Rspec}(\Lambda_X(\kappa(p))) \xrightarrow{\epsilon_X \circ h} \text{Spec}(R)$$

factors as

$$\text{Rspec}(\Lambda_X(\kappa(p))) \rightarrow \text{Spec}(\kappa(p)) \rightarrow \text{Spec}(R).$$

This gives us a commutative square

$$\begin{array}{ccc} X & \xrightarrow{\epsilon_X} & \text{Spec } R \\ \uparrow & & \uparrow \\ \text{Rspec}(\Lambda_X(\kappa(p))) & \longrightarrow & \text{Spec } \kappa(p) \end{array}$$

Since  $\Lambda_X(\kappa(p))$  is not the zero sheaf hence the structure sheaf of  $\text{Rspec}(\Lambda_X(\kappa(p)))$  is non-zero. This implies that  $\text{Rspec}(\Lambda_X(\kappa(p)))$  is not the empty scheme. Therefore the point  $p$  is in the image of  $\text{Rspec}(\Lambda_X(\kappa(p))) \rightarrow \text{Spec}(\kappa(p)) \rightarrow \text{Spec}(R)$ , hence also in the image of  $\epsilon_X$ . ■

**Lemma 4.2.13.** *Let  $p \subset R$  be a prime ideal. The closed set  $V_X(p)$  is irreducible. This implies that  $\epsilon_X$  is injective.*

*Proof.* Let  $z \in X$ . Assume that  $\epsilon_X(z) = p$ . So  $\ker z = p$ . Let  $y \in V_X(p)$ . Then  $\ker(z) \subset \ker(y)$ , hence if  $y \in D_X(a)$  then  $z \in D_X(a)$ . Therefore  $z$  is contained in any open subset of  $V_X(p)$ , hence it is irreducible. Note that  $z$  is its unique generic point since the biggest closed subset of  $V_X(p)$  not containing  $z$  is the empty set. This shows that  $\epsilon_X$  is injective. ■

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**Lemma 4.2.14.** *The counit  $\epsilon_X$  is open, hence a homeomorphism.*

*Proof.* Let  $a \in R$ . Let  $x \in X$ . Let  $p = \epsilon_X(x)$ . Since  $\epsilon_X$  is a morphism of locally ringed spaces, if  $a \notin \ker x$  then  $a \notin \ker p = p$ . Note that  $\ker -$  is defined for arbitrary schemes so also works on  $\text{Spec}(R)$ . Hence  $\epsilon_X(D_X(a)) = \{\epsilon_X(x) \mid a \notin \ker(x)\} = \epsilon_X(X) \cap D_{\text{Spec } R}(a) = D_{\text{Spec } R}(a)$ . Our map  $\epsilon_X$  is continuous and open, so a homeomorphism. ■

**Proposition 4.2.15.** *The affine scheme  $X$  is affine.*

*Proof.* By Lemmas 4.2.12 to 4.2.14 it follows that  $\epsilon$  is an homeomorphism and so by Lemma 4.2.4 it follows that  $X$  is affine and is  $\epsilon$  an isomorphism. ■

### 4.3 Enough affines

For schemes a quasi-coherent module often is defined by a module that over affine schemes is of the form  $\Lambda(M)$  for some module  $M$ . We will look how often this generalizes, which we will do through the notion of ‘enough affines’. We will see some examples where it does not work as in Sch and some results about when it does.

**Definition 4.3.1.** We say that a ringed site  $(C, \mathcal{T}, \mathcal{O})$  has enough affines if any object admits a affine covering.

**Lemma 4.3.2.** *Any ringed site  $(C, \mathcal{T}, \mathcal{O})$  that has a finite poset as underlying category, has enough affines.*

*Proof.* Let  $x_0 \in C$ . If  $x_0$  is covered by the maximal sieve only or the maximal sieve and the empty sieve, it is affine and we are done. Assume otherwise. Let  $S = \{y_i \rightarrow x_0\}$  be a non-maximal, non-empty cover of  $x_0$ . Then  $S$  does not contain isomorphisms.

We can associate to any non-maximal non-empty covering sieve  $S$  of an element  $x_0$ , the set of all NA-chains  $x_0 \leftarrow x_1 \leftarrow \dots \leftarrow x_n$ . An NA-chain, associated to  $R$ , is a chain of maps ending in  $x_0$  such that  $x_i \leftarrow x_{i+1}$  is contained in a non-maximal, non-empty cover of  $x_i$ , where  $x_0 \leftarrow x_1$  is contained in  $R$ .

By finiteness of  $C$ , any chain of maps is bounded by the size of  $C$  or contains a cycle. If a chain contains a cycle, it contains isomorphisms. By construction, no isomorphism can be present in a NA-chain. Therefore the length of any NA-chain is bounded by  $\|C\|$ .

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Let  $H$  be a NA-chain associated to  $S$  of maximal length  $m$ . Then the last map  $\dots \leftarrow h \leftarrow g$  in  $H$  has an affine object  $g$  as domain, because  $H$  cannot be increased and so  $g$  has no non-maximal, non-empty coverings which makes it affine. Also the non-maximal, non-empty covering of  $h$  where this map appears must be an affine covering by applying the same reasoning to the other objects occurring in it. Hence all objects occurring at the  $(m-1)$ th place in any NA-chain admits an affine cover. Let  $i \leq m-1$ . Assume all elements at the  $(i-1)$ th place admit an affine cover. Let  $b$  be a object occurring at the  $(i-1)$ th place in a chain. It is either affine or all objects in any non-maximal, non-empty cover occur at the  $i$ th place in some chain hence admit an affine cover. Therefore any non maximal, non empty cover on  $b$  can be refined to an affine cover. This provides us with an affine cover of  $b$ . By reversed induction,  $x_0$  admits a affine cover. ■

**Example 4.3.3.** The category  $C$  is  $\mathbb{Z} \times \mathbb{Z}$  with the usual ordering. An element  $(i, j)$  is only non-trivially covered by  $\{(i, j-1) \rightarrow (i, j), (i-1, j) \rightarrow (i, j)\}$ . Let  $k$  be any field. Let  $R = k[x_{ij} | i, j \in \mathbb{Z}]$ . Define the structure sheaf as  $\mathfrak{D}(i, j) = R[x_{kl}^{-1} | i \leq k \text{ \& } j \leq l]$ .

Fix  $(a, b) \in C$ . Consider the over category  $C_{(a,b)}$  at this point. Let  $(i, j) \rightarrow (a, b)$  be an object of  $C_{(a,b)}$ . So  $i \leq a$  and  $j \leq b$ . Define the presheaf of modules  $F(i, j) = \mathfrak{D}(i, j) / (x_{a-1,b} x_{a,b-1})$  on  $X$ . We have  $a > i$  or  $b > j$  or  $(i = a \text{ and } j = b)$ . If  $a > i$  or  $b > j$ , then  $x_{a-1,b}$  or  $x_{a,b-1}$  is invertible in  $\mathfrak{D}(i, j)$ , hence  $F(i, j) = 0$  in both cases. This presheaf is zero everywhere except at  $(a, b)$ , because both  $x_{a-1,b}$  and  $x_{a,b-1}$  are not invertible in  $\mathfrak{D}(a, b)$  hence sheafifies to the zero sheaf. In other words:  $\bigwedge_{(x_{a-1,b} x_{a,b-1})} \mathfrak{D}(a, b) = 0$ , Hence  $(a, b)$  is not affine, which shows that  $C$  has no affine objects.

**Example 4.3.4.** Let  $C$  be as in Example 4.3.3. Consider  $G = \mathfrak{D}(i, j)[y_{k,l} | k \leq i \text{ \& } l \leq j]$ . Let  $\bigoplus_{k \in I} \mathfrak{D} \xrightarrow{\alpha} G$  be any sheaf map. Let  $\alpha(e_k)$  be the image of the generators  $e_k \in \bigoplus_{k \in I} \mathfrak{D}$  in the global sections. The section  $y_{1,1} \in G(1, 1)$  cannot be written as a finite sum  $\sum_k \lambda_k \alpha(e_k)$  for scalars  $\lambda_k \in \mathfrak{D}(i, j)$  for any  $(i, j)$ . This shows that  $\alpha$  is not surjective hence  $G$  is not quasi-coherent (locally presentable).

**Example 4.3.5.** This is the example in [7, Tag 01BL].

Let  $L = (\mathbb{R}, \mathfrak{D}_R)$  be the real line with the euclidean topology and the sheaf of continuous real valued functions as structure sheaf. Let

$$X = \bigcup_{i=0}^{\infty} L_i$$

with  $[i, x] \sim [j, y]$  if and only if  $i = j$  and  $x = y$  or  $y = x = 0$ . The real lines are glued to each other at zero. Define the open  $U_n \subset X$  as  $U_n \cap L_i = (-\frac{1}{n}, \frac{1}{n})$ . These opens form

#### 4 Affine objects

a basis of neighbourhoods of 0. Let  $f : \mathbb{R} \rightarrow \mathbb{R}$  be any continuous function such that  $f(x) = 0$  if  $x \in (-1, 1)$  and  $f(x) = 1$  if  $x \in (-\infty, -2) \cup (2, \infty)$ . Let  $f_n(x) = f(nx)$ .

Define the sheaf map

$$\bigoplus_i \mathfrak{D}_R \xrightarrow{\alpha} \bigoplus_{ij} \mathfrak{D}_R,$$

$$e_i \mapsto \sum_j f_i 1_{L_j} e_{ij}.$$

To prove that this is well-defined, we need to show that the sum  $\sum_j f_i 1_{L_j} e_{ij}$  is locally finite for every  $i$ . Let  $[k, y] \in X$ . If  $y \neq 0$ , then

$$W_{[k,y]} = \{[k, z] \in X \mid z \in (y - \delta, y + \delta)\} \subset L_k$$

is open in  $X$  and  $\alpha_{W_{[k,y]}}(e_i) = f_i e_{ik}$  for any  $\delta < |y|/2$ . If  $y = 0$ , then  $\alpha_{U_n}(e_i) = 0$  if  $n > i$  because  $f_i$  is zero on  $U_n$ . Hence we found a cover on which our sum is locally finite, which makes  $\alpha$  well-defined.

The adjunction  $\Lambda \dashv \Gamma(X; -)$  implies that  $\Lambda$  commutes with arbitrary colimits. Moreover

$$\mathfrak{D}_X \cong \Lambda(\Gamma(X; \mathfrak{D})),$$

so

$$\bigoplus_i \mathfrak{D}_R \cong \Lambda(\bigoplus_i \Gamma(X; \mathfrak{D})).$$

This shows that  $\alpha$  is a morphism between objects that are of the form  $\Lambda(M)$ .

Let  $\beta : \bigoplus_i \Gamma(U; \mathfrak{D}) \rightarrow \bigoplus_{ij} \Gamma(U; \mathfrak{D})$  for some open  $U$ . Then  $\Lambda(\beta)(e_i) = \sum_{j \in J_i} a_{ij} e_{ij}$  where  $J_i$  is finite for every  $i$ .

Assume that  $\alpha = \Lambda(\beta)$  over some neighbourhood  $U$  of 0. Then there exists a  $m$  such that  $U_m \subset U$ . Let  $k > 2m$ . Then  $f_k \neq 0$  on  $U_m$ , hence  $f_k 1_{L_j} \neq 0$  on  $U_m$  for every  $j$  and so no coefficients vanish of  $\alpha_{U_m}(e_k) = \sum_j f_k 1_{L_j} e_{kj}$ . This contradicts  $\alpha = \Lambda(\beta)$ . Hence  $\alpha$  does not come from a module map locally.

*Remark 4.3.6.* This is an example of a ringed space where  $\Lambda$  is not an equivalence locally, since it shows that  $\Lambda$  is not full. We will categorify this example next.

#### 4 Caffeine objects

**Definition 4.3.7** (Neighbourhood site). Let  $X$  be a topological space and  $\text{Open}(X)$  the corresponding category of open subsets. Let  $y \in X$ . Define the category  $N(y)$  to be the full subcategory of  $\text{Open}(X)$  of all neighbourhoods  $U$  of  $y$ . A family  $\{U_i \rightarrow U\}$  is covering on  $U$  in  $N(y)$  if it is covering in  $\text{Open}(X)$ . Note that  $i : N(y) \rightarrow \text{Open}(X)$  is a morphism of sites.

Assume that  $X$  is a ringed space coming with a structure sheaf  $\mathcal{O}$ . Define  $\mathcal{O}_{X,y} = i_* \mathcal{O}$ .

**Example 4.3.8.** Let  $X, f_j$  and  $U_n$  be as in Example 4.3.5. Note that  $\bigoplus_i \mathcal{O}_{X,y}$  is quasi-coherent.

Set  $y = 0 \in X$ . Define the sheaf map

$$\bigoplus_i \mathcal{O}_{X,y} \xrightarrow{\alpha} \bigoplus_{ij} \mathcal{O}_{X,y},$$

$$e_i \mapsto \sum_j f_i 1_{L_j} e_{ij}.$$

Fix  $i$ . We will prove that  $\alpha_X(e_i)$  is a well-defined global section. Choose  $m > i$ . Let  $V_k = L_k \cup U_m$  and note that  $\bigcup_k V_k = X$ . By construction  $f_i$  is zero on  $U_m$ , hence  $f_i 1_{L_j}$  is zero on  $V_k$  if  $k \neq j$  and so  $\sum_j f_i 1_{L_j} e_{ij} = f_i 1_{L_k} e_{ik}$  on  $V_k$ . This shows that  $\alpha_X(e_i)$  is a well-defined section on any element of the cover  $\{V_k \rightarrow X\}$  and this family is matching since the sections are functions and the 'restriction' maps are actual restriction.

Let  $V \in N(y)$ . Assume there exists  $\beta : \bigoplus_i \Gamma(V, \mathcal{O}_{X,y}) \rightarrow \bigoplus_{ij} \Gamma(V, \mathcal{O}_{X,y})$  such that  $\Lambda(\beta) = \alpha_V$ . Then  $\alpha_V(e_i) = \sum_j f_i 1_{L_j} e_{ij}$  is not just locally finite over some cover that may depend on  $i$ , but actually finite globally on  $V$  for all  $i$ . So almost all  $f_i 1_{L_j}$  are zero on  $V$ . Note that  $y \in V$ , so  $U_d \subset V$  for some  $d$ . Let  $i > 2d$ , then  $f_i \neq 0$  on  $(-\frac{1}{d}, \frac{1}{d})$  and so  $f_i 1_{L_j} \neq 0$  on  $U_d$  for any  $j$ . Hence  $\alpha_V(e_i) = \sum_j f_i 1_{L_j} e_{ij}$  is not a finite sum for  $i > 2d$ . This contradicts our assumption, so  $\alpha$  does not come

The restriction of any quasi-coherent sheaf is quasi-coherent. Observe that  $\alpha$ , and its restrictions, is a morphism between quasi-coherent sheafs but does not come from a map of modules by the previous contradiction. Therefore  $\Lambda$  is not full for any  $V$  and so no object  $V$  is caffeine in  $N(y)$ .



## 5 Further research

All the loose ends and open questions are collected in this section. Let  $C$  be a ringed site. Let  $\alpha \in C$  be a affine object. Let  $\mathfrak{F}$  be a quasi-coherent sheaf on  $C_\alpha$ .

### Cohomology

The vanishing theorem of Serre tells us that all higher cohomology of quasi-coherent sheaves are zero.[7, Tag 01XB]

**Question 5.0.1.** *Does  $H^i(\alpha, \mathfrak{F}) = 0$  for all  $i > 0$  hold?*

### Restrictiveness

Recall the definition of a restrictive morphism from Definition 3.0.1.

**Question 5.0.2.** *Is the pullback of a restrictive morphism restrictive?*

### Determining if object is affine

It would be nice to have methods to see whether an object is affine. We used some counting arguments in the finite case in section 3.3, which might be extendable for objects that are quasi-compact in some sense categorical sense.

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