

Affine Objects

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1 Preliminaries

1.1 Topology

Definition 1 (Sieve). Let C be a category and $a \in C$. A sieve S on a is a subpresheaf of $h(a)$. Explicitly, for each $c \in C$, $S(c)$ is a subset of $\text{Hom}(c, a)$ such that $fg \in S(\text{Dom}(g))$ for all $f \in S(c)$ and for all $g \in h(c)$.

The maximal sieve on a , which is $h(a)$, will be denoted by $\max(a)$.

Definition 2 (Sieve category). Let C be a category and $a \in C$. The sieve category $\text{Sieves}(a)$ is the subobject poset of the presheaf $h(a)$.

Definition 3 (Pullback of sieve). Let C be a category and $a, b \in C$. Let S be a sieve on a . Let $f : b \rightarrow a$.

For any $c \in C$ the sieve f^*S on b is given by $f^*S(c) = \{g \in \text{Hom}(c, b) : fg \in S(c)\}$.

To show that this is actually a subpresheaf of $h(b)$, let $k : c \rightarrow c'$ and $h \in f^*S(c')$. Hence $fh \in S(c')$ and so $fhk \in S(c)$. Conclude that $hk \in f^*S(c)$.

This defines a functor $f^* : \text{Sieves}(a) \rightarrow \text{Sieves}(b)$.

Definition 4 (Grothendieck Topology). A Grothendieck topology \mathcal{T} is a family $\mathcal{T}(a)$ of 'covering' sieves for every $a \in C$ with the following conditions:

- $\max(a) \in \mathcal{T}(a)$
- $f^*R \in \mathcal{T}(a')$ if $R \in \mathcal{T}(a)$ for all $f : a' \rightarrow a$
- if $f^*R \in \mathcal{T}(a')$ for all $f \in S$ with $S \in \mathcal{T}(a)$ then $R \in \mathcal{T}(a)$

Definition 5 (Basis). Let C be a category with pullbacks. A Grothendieck pretopology \mathcal{B} is a collection $\mathcal{B}(a)$ of families $\{f_i : a_i \rightarrow a\}$ of 'covering' morphisms for every $a \in C$ with the following conditions.

Are these later called Axioms 1-3? If so, they should be labeled.

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- every isomorphism is a covering singleton family,
- (Stability) The pullback of a covering family is a covering family. If $\{f_i : a_i \rightarrow a\}$ is covering and $g : b \rightarrow a$, then $\{f'_i : a_i \times_a b \rightarrow b\}$ is covering.
- (Transitivity) If $\{f_i : a_i \rightarrow a\}$ is a covering family and $\{f_{ij} : a_{ij} \rightarrow a_i\}$ for every i , then $\{f_{ij} : a_{ij} \rightarrow a\}$ is a covering family.

Generating a topology from a basis: take any sieve containing a covering family to be a covering sieve.

1.1.1 Sheaves

Does this mean τ is a Grothendieck topology, or a pretopology?

Definition 6 (Sheaves). Let (C, \mathcal{T}) be a site. Let $\mathcal{F} \in \hat{C}$.

A presheaf that admits a unique amalgamation for every matching family is called a sheaf. The category $\text{Shv}(C)$ is the full subcategory in \hat{C} of all sheaves. Let i be the inclusion functor $\text{Shv}(C) \rightarrow \hat{C}$.

In other words, we call \mathcal{F} a sheaf if the map

$$\mathcal{F}(a) \rightarrow \mathcal{F}(R)$$

$$a; x \mapsto \{\mathcal{F}(i)(x)\}_{i \in R}$$

is an isomorphism...

for each $a \in \mathcal{C}$ and $R \in \mathcal{T}(a)$?

namespace collision.

try an align* environment with <to and &\mapsto.

Definition 7 (Plus construction). Let (C, \mathcal{T}) be a site. Let $a, a' \in C$ and $f : a \rightarrow a'$. Let $\mathcal{F} \in \hat{C}$. Define the functor $(-)^+ : \hat{C} \rightarrow \hat{C}$ as follows

On objects:

$$(\mathcal{F})^+(a) = \frac{\{(R, \varphi) \mid R \in \mathcal{T}(a), \varphi \in \Gamma(R; \mathcal{F})\}}{\sim}$$

which notation do you use?

Again, these are too soon. Better:

"... for all $a \in \mathcal{C}$ "

$$(\mathcal{F})^+(f)([(R, \varphi)]) = [(f^*R, \varphi h(f))].$$

"... for all morphisms f in \mathcal{C} "

The equivalence relation is defined as:

$$(R, \varphi) \sim (S, \phi)$$

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← also a covering sieve?

if $\varphi = \phi$ on some $Q \subset R \cap S$

Let $L: \mathcal{F} \rightarrow \mathcal{F}'$. Then

$$((L)^+)_a([(R, \varphi)]) = [(R, L \circ \varphi)]$$

This functor comes with a natural transformation $\omega: \text{Id} \rightarrow (-)^+$ defined by

$$\omega_{\mathcal{F}, a}(x) = [(\max(a), y)]$$

$$y(i) = \mathcal{F}(i)(x).$$

Lemma 8. Let \mathcal{F} be a presheaf, \mathcal{G} a sheaf and $g: \mathcal{F} \rightarrow \mathcal{G}$ a morphism in $\hat{\mathcal{C}}$. Then g factors through $\omega_{\mathcal{F}}$ via a unique g' .

Lemma 9. For every presheaf \mathcal{F} , $(\mathcal{F})^+$ is separated.

Lemma 10. If \mathcal{F} is separated, then $(\mathcal{F})^+$ is a sheaf.

/ reference proofs, e.g. M-M (sheaves in geometry + logic)

Definition 11. Define $\text{sh} = (-)^+ \circ (-)^+$.

Lemma 12 (Sheafification adjunction). Let $\mathcal{Y} = (C, \mathcal{T})$ be a site. The functor sh is left adjoint to the inclusion $\hat{\mathcal{Y}} \rightarrow \text{Shv}(C)$ with unit

$$\omega_{\mathcal{F}}^2 = \omega_{(\mathcal{F})^+} \circ \omega_{\mathcal{F}}$$

should not be italicized.
Are you using \text or \mathrm?

1.1.2 Relative topology

Definition 13 (Relative topology). Let (C, \mathcal{T}) be a site. Let $a \in C$.

Set $\mathcal{T}_a(f) = \{R^f : R \in \mathcal{T}(b)\}$. Define the induced topology \mathcal{T}_a on C_a by, for each $f \in C_a$

what is this?

$$\mathcal{T}_a(f) = Q^f(\mathcal{T}(\text{Dom}(f))).$$

what is this?

Lemma 14. \mathcal{T}_a defines a Grothendieck topology

Proof. Axiom 1: Q^f is an equivalence of posets. So the terminal object is sent to the terminal object. Hence $\max(f) \in \mathcal{T}_a(f)$.

Axiom 2 & 3 are consequences of: Q^f is an equivalence and Q^f commutes with sieve pullback. ■

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Definition 15 (Oversite). Let $Y = (C, \mathcal{T})$ be a site. Let $a \in C$. Define the site Y_a to be the category C_a with the induced topology \mathcal{T}_a .

Definition 16 (Natural transformation s). Let (C, \mathcal{T}) be a site. Let $a, b \in C$ and $f : b \rightarrow a$.

Let $\{\chi_i\}$ be a compatible family indexed by a sieve R on b . The same set $\{\chi_i\}$ is a compatible family on f indexed by $Q^f(R)$. Define the natural transformation

$$s : u \circ (-)^+ \rightarrow (u \circ -)^+$$

"by"

$$s_{\mathcal{F}} : u \circ (\mathcal{F})^+ \rightarrow (u \circ \mathcal{F})^+$$

$$s_{\mathcal{F},f}([\chi_i]_{i \in R}) = [\chi_i]_{i \in Q^f(R)}.$$

Lemma 17 (Restriction commutes with plus). *The transformation $s_{\mathcal{F}}$ is an isomorphism of functors.*

Proof. We use the variables from definition ?. The morphism $s_{\mathcal{F},f}$ has an inverse, again sending the set to itself and applying Q^f on the indexing sieve. These are mutual inverses, so $s_{\mathcal{F},f}$ is an isomorphism for each f . ■

Lemma 18 (s and ω commute). *Let \mathcal{F} be a presheaf on (C, \mathcal{T}) . Let $f : b \rightarrow a \in C$. Then $\omega_{u \circ \mathcal{F}} = s_{\mathcal{F}} \circ (u \circ \omega_{\mathcal{F}})$.*

what is u ? Maybe better expressed diagrammatically.

Proof. For any section $\chi \in \Gamma(b; \mathcal{F})$. Let $\chi_i = \mathcal{F}(i)(\chi)$ for any morphism $i \in C$. Note that $Q^f(\max(f)) = \max(b)$. This implies that the compatible family $\{\chi_i\}$ indexed by the maximal sieve on f is sent by s to the same set $\{\chi_i\}$ indexed by the maximal sieve on b . In diagram form, that

$$\begin{array}{ccc} u \circ \mathcal{F} & & \\ \downarrow u \circ \omega_{\mathcal{F}} & \searrow \omega_{u \circ \mathcal{F}} & \\ u \circ (\mathcal{F})^+ & \xrightarrow{s} & (u \circ \mathcal{F})^+ \end{array}$$

commutes. ■

are these defined anywhere?

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Corollary 19. Let $(C, \mathcal{T}, \mathcal{D})$ be a ringed site. Let $b \in C$. Let M be a $\Gamma(1; \mathcal{D})$ -module. The morphism $\omega_{\lambda(M), 1, b}^2: \lambda(M)(b) \rightarrow \Lambda(M)(b)$ is isomorphic to

$$\omega_{\lambda(M \otimes \mathcal{D}(b)), b, Id_b}^2: \lambda(\lambda(M \otimes \mathcal{D}(b))) \rightarrow \Lambda(\lambda(M \otimes \mathcal{D}(b))).$$

Definition 20. Let P be the plus functor and $U = u \circ -$. Define

$$s^2: U \circ P \circ P \rightarrow P \circ P \circ U$$

as

$$s_{\mathcal{F}}^2 = P(s_{\mathcal{F}}) \circ s_{P(\mathcal{F})}.$$

where do these functors go?

Lemma 21. Let \mathcal{F} be a presheaf on (C, \mathcal{T}) . Let $f: b \rightarrow a \in C$. Let \mathcal{F} be a presheaf on C . Then $\omega_{u \circ \mathcal{F}}^2 = s_{\mathcal{F}}^2 \circ (u \circ \omega_{\mathcal{F}}^2)$.

Proof. Let $a \in C$. We have

$$\omega_{u \circ \mathcal{F}}^2 = \omega_{(u \circ \mathcal{F})+} \circ \omega_{u \circ \mathcal{F}},$$

$$s_{\mathcal{F}}^2 = (s_{\mathcal{F}})^+ \circ s_{(\mathcal{F})+}.$$

are you using P or $+$?
I would just stick with $+$ everywhere.

$$u \circ \omega_{\mathcal{F}}^2 = u \circ (\omega_{(\mathcal{F})+} \circ \omega_{\mathcal{F}}).$$

Let $x \in \Gamma(b; u \circ \mathcal{F})$.

Then

$$\omega_{u \circ \mathcal{F}, f}^2(x) = ((x_i)_{i \in \max(\text{Dom}(j))})_{j \in \max(f)}$$

with $x_i = \mathcal{F}i(x)$. We also have

$$u \circ \omega_{\mathcal{F}, f}^2(x) = ((x_i)_{i \in \max(\text{Dom}(j))})_{j \in \max(b)},$$

where $\text{Dom}(j) \in C_a$. Apply $s_{(\mathcal{F})+}$ on this to get

$$((x_i)_{i \in \max(\text{Dom}(j))})_{j \in \max(f)},$$

where $\text{Dom}(j) \in C$. Lastly, apply $(s_{\mathcal{F}})^+$ to get

$$((x_i)_{i \in \max(\text{Dom}(j))})_{j \in \max(f)}$$

where $\text{Dom}(j) \in C_a$. The statement of the lemma is established. ■

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Corollary 22. Let $Y = (C, \mathcal{T})$. Let $a, b \in C$. Let $f : b \rightarrow a$. Sheafifying and restricting commute via the iso

$$s^2 : sh_b \circ *|_b \rightarrow *|_b \circ sh_a.$$

Lemma 23 (λ commutes with restriction). Let $(C, \mathcal{T}, \mathcal{D})$ be a ringed site. Let $f : b \rightarrow a \in C$. The functors $u \circ \lambda$ and $\lambda \circ (- \otimes \Gamma(a; \mathcal{D}))$ are the same functor.

where is f here?

Lemma 24 (λ counit commute with restriction). Let C be a category. Let $a \in C$. Let ϵ be the counit of the adjunction $\lambda \dashv \Gamma(1; -)$ on C . Let ϵ_a be the counit of the adjunction $\lambda_a \dashv \Gamma(1; -)$ on C_a . Let \mathcal{F} be a presheaf module on C . By lemma ?, $\lambda(\Gamma(1; \mathcal{F}))|_a = \lambda(\Gamma(1; \mathcal{F}) \otimes \Gamma(a; \mathcal{D}))$. We have $u \circ (\epsilon) = \epsilon_a$ on this presheaf module.

Proof. Let $f : b \rightarrow a \in C_a$. Let $m \otimes r \in \lambda(\Gamma(1; \mathcal{F}) \otimes \Gamma(a; \mathcal{D}))(f) = \Gamma(1; \mathcal{F}) \otimes \Gamma(b; \mathcal{D})$. Then both maps do $m \otimes r \mapsto rm$. ■

Lemma 25 (λ commutes with restriction). Let $(C, \mathcal{T}, \mathcal{D})$ be a ringed site. Let $f : b \rightarrow a \in C$. The functors $u \circ \Lambda$ and $\Lambda \circ (- \otimes \Gamma(a; \mathcal{D}))$ are isomorphic functors.

again

Lemma 26 (Λ counit commute with restriction). Let C be a category. Let $a \in C$. Let ϵ be the counit of the adjunction $\Lambda \dashv \Gamma(1; -)$ on C . Let ϵ_a be the counit of the adjunction $\Lambda_a \dashv \Gamma(1; -)$ on C_a . Let \mathcal{F} be a presheaf module on C . By lemma ?, $\Lambda(\Gamma(1; \mathcal{F}))|_a = \Lambda(\Gamma(1; \mathcal{F}) \otimes \Gamma(a; \mathcal{D}))$. We have $u \circ (\epsilon) = \epsilon_a$ on this presheaf module.

Proof.

Lemma 27. Let (C, \mathcal{T}) be a site. Let $a \in C$. Consider the counit ϵ of $\Gamma(1; -) \dashv \Lambda(-)$ on C and the counit ϵ_a of the same adjunction on C_a . Then $\epsilon|_a = \epsilon_a$.

are these different lemmas?

Proof. Let \mathcal{F} be a sheaf module on C . Let $f : b \rightarrow a \in C$. We want to show that $\epsilon_{\mathcal{F}, a, f} : \Lambda(\Gamma(1; \mathcal{F})) \rightarrow \mathcal{F}$ is the same map as $\epsilon_{\mathcal{F}}|_{af} : \Lambda(\Gamma(1; \mathcal{F})) \rightarrow \mathcal{F}$

$$\alpha : \lambda(\Gamma(1; \mathcal{F})) \rightarrow \mathcal{F},$$

$$m \otimes r \mapsto mr$$

composed with $\omega_{\lambda(\Gamma(1; \mathcal{F}))}^2$. ■

2 Affine objects

2.1 Affine objects

Definition 28 (Affine object). Let $Y = (C, \mathcal{T}, \mathcal{O})$ be a ringed site. Let $a \in C$ be an object. We call a affine if the unit η and co-unit ϵ of the adjunction $\Gamma(1; -) \dashv \Lambda(-)$ on Y_a are natural isomorphisms.

Example 29 (Examples of affine objects). The main example to keep in mind is $\text{Spec}(R) \in \text{Sch}$.

Definition 30 (affine cover). Let $(X, \mathcal{T}, \mathcal{O})$ be a ringed site. A family of maps $\{a_i \rightarrow a\}$ is called a affine covering of a if every a_i is affine and the family is a covering family.

Definition 31. We say that a ringed site $(C, \mathcal{O}, \mathcal{T})$ has enough affines if any object admits a affine covering.

Lemma 32. Let $(C, \mathcal{T}, \mathcal{O})$ be a ringed site. Let $a \in C$. Let $\{b_i \rightarrow a\}$ be a affine covering on a . Assume every map $b_i \rightarrow a$ is restrictive. Then the counit ϵ of the adjunction $\Gamma(1; -) \dashv \Lambda(-)$ on Y_a is a natural isomorphism.

Proof. Let \mathcal{F} be a quasi-coherent sheaf module. Set $M = \Gamma(a; \mathcal{F})$. Set $M_i = \Gamma(b_i; \mathcal{F})$. Set $\beta_i = \epsilon_{\mathcal{F}, a}|_{b_i}$. By lemma ? $\beta_i \cong \epsilon_{b_i}$, hence β_i is an isomorphism. ■

Lemma 33. Let $(C, \mathcal{T}, \mathcal{O})$ be a ringed site. Let $a \in C$. Let M be a $\Gamma(a; \mathcal{O})$ -module. The component

$$\omega_{\lambda(M), a}^2 : \lambda(M) \rightarrow \Lambda(M)$$

at Id_a of the sheafification morphism

$$\omega_{\Lambda(M)}^2 : \lambda(M) \rightarrow \Lambda(M)$$

these have the same domain and codomain?

which is the left adjoint? see p.8.

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is equal to the unit of $\Lambda \dashv \Gamma(1; -)$ in C_a .

Proof. Consider the following maps, which you get by repeatedly calling on an adjunction.

$$\text{Id} : \Lambda(M) \rightarrow \Lambda(M)$$

$$\omega_{\Lambda(M)}^2 : \lambda(M) \rightarrow \Lambda(M) \text{ use sheafification adjunction, see lemma ..}$$

$$\omega_{\lambda(M), a}^2 M \rightarrow \Gamma(a; \Lambda(\mathcal{M})) \text{ take sections at } a$$

We took the adjunct of Id with respect to the sheafification adjunction as the first step. Then we took the adjunct of the result wrt the λ adjunction. Hence we get the adjunct of Id wrt the Λ adjunction. so the last map is actually the unit of the Λ adjunction. ■

Corollary 34. Let $(C, \mathcal{T}, \mathcal{D})$ be a ringed site. Let $a \in C$ be affine. Then $\omega_{\lambda(M), a}^2$ is an isomorphism for any $\mathcal{D}(a)$ -module M .

Theorem 35 (Morphism between affines is restrictive). Let $Y = (C, \mathcal{T}, \mathcal{D})$. Let $f : b \rightarrow a \in C$ be a morphism between affine objects, then f is restrictive.

Proof. Let \mathcal{F} be a quasi-coherent module on Y_a . Let $M = \Gamma(a; \mathcal{F})$. Since a is affine, we have $\mathcal{F} = \Lambda(M)$.

We have to show that the adjunct, along the extension of scalars adjunction, of $\mathcal{F}(f)$

$$\Gamma(a; \mathcal{F}) \otimes_{\Gamma(a; \mathcal{D})} \Gamma(b; \mathcal{D}) \rightarrow \Gamma(b; \mathcal{F})$$

is an isomorphism.

This adjunct is the component at b of the natural transformation $\omega_{\lambda(\Gamma(1; \mathcal{F})), a}^2$. Since b is affine, this component is an isomorphism. ■

