Theoretical Computer Science Cheat Sheet		
Definitions		Series
f(n) = O(g(n))	iff $\exists$ positive $c, n_0$ such that $0 \le f(n) \le cg(n) \ \forall n \ge n_0$ .	$egin{aligned} egin{aligned} egin{aligned\\ egin{aligned} egi$
$f\left(n\right) {=} \Omega \ \left(g(n)\right)$	$\begin{array}{l} \text{iff } \exists \ positive \ c, n_0 \ such \ that } \\ f\left(n\right) \geq cg(n) \geq 0 \ \forall n \geq n_0. \end{array}$	In general:
$f(n)=\Theta(g(n))$	$\begin{array}{c} \text{iff } f\left(n\right) = O(g(n)) \ \text{ and } \\ f\left(n\right) = \Omega \ (g(n)). \end{array}$	$\begin{bmatrix} X^{h} & i^{m} = \frac{1}{m+1} (n+1)^{m+1} - 1 - \frac{X^{h}}{i=1} (i+1)^{m+1} - i^{m+1} - (m+1)i^{m} \end{cases}$
f(n) = o(g(n))	$\text{iff } \lim_{n\to\infty} \ f\left(n\right)/g(n) {=} 0 \ .$	$\mathbf{X}^{-1}\mathbf{i}^{\mathbf{m}} = \frac{1}{\mathbf{m}+1} \frac{\mathbf{X}^{\mathbf{n}-?} \mathbf{m}+1}{\mathbf{k}} \mathbf{B}_{\mathbf{k}} \mathbf{n}^{\mathbf{m}+1-\mathbf{k}}.$
$\lim_{n \to \infty} a_n = a$	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	Geometric series:
sup S	$ \begin{array}{c} \text{least } b \in R \text{ such that } b \geq s, \\ \forall s \in S. \end{array} $	$egin{aligned} egin{aligned} egin{aligned\\ egin{aligned} egi$
inf S	greatest $b \in R$ such that $b \le s$ , $\forall s \in S$ .	
$\liminf_{n\to\infty}a_n$	$\lim_{n\to\infty}\inf\{a_i\mid i\geq n, i\in N\}.$	Harmonic series: $\begin{array}{cccccccccccccccccccccccccccccccccccc$
$\limsup_{\substack{n \to \infty \\ \frac{n}{2}?}} a_n$	$\lim_{n\to\infty}\sup\{a_i\mid i\geq n, i\in N\}.$	$H_n = rac{X^n}{i} rac{1}{i}, \qquad rac{X^n}{i H_i} = rac{n(n+1)}{2} H_n - rac{n(n-1)}{4}. \ X^n \qquad \qquad rac{X^n}{X^n} rac{?}{i} rac{}{i} rac{?}{i} rac{?}{i} rac{?}{i} rac{?}{i} rac{?}{i} rac$
k	Combinations: Size k subsets of a size n set.	$\begin{array}{cccccccccccccccccccccccccccccccccccc$
? ? n k	Stirling numbers (1st kind): Arrangements of an n element set into k cycles.	1. $\frac{n}{k} = \frac{n!}{(n-k)!k!}$ , 2. $\frac{X^n}{k} = 2^n$ , 3. $\frac{n}{k} = \frac{n}{n-k}$ ,
? ? k	Stirling numbers (2nd kind): Partitions of an n element set into k non-empty sets.	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
? ? k	1st order Eulerian numbers: Permutations $\pi_1 \pi_2 \pi_n$ on $\{1, 2,, n\}$ with k ascents.	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
<b>??                                   </b>	2nd order Eulerian numbers	<b>10.</b> $ = (-1)^k $ $ = 1 $ $ = 1 $
$C_n$	$\begin{array}{c} \textbf{Catalan Numbers: Binary} \\ \textbf{trees with } n+1 \ \textbf{vertices.} \end{array}$	<b>12.</b> $n = 2^{n-1} - 1$ , <b>13.</b> $n = k + n - 1 + n - 1$ ,
$\begin{array}{ c c c c c c c c c c c c c c c c c c c$		
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<b>31.</b> $=$ $(-1)^{n-k-m}k!$ , <b>32.</b> $=1$ , $0$ $=1$		
$\begin{bmatrix} 34. \\ 1 \end{bmatrix} = (k+1) \\ 1 \end{bmatrix} + (2n-1-k) \\ 1 \end{bmatrix} $		
36. $\frac{x}{x-n} = \frac{y}{x}$	$egin{array}{cccccccccccccccccccccccccccccccccccc$	37. $\frac{?}{m+1} = \frac{X}{k} \cdot \frac{?}{n} \cdot \frac{?}{k} \cdot \frac{X^{k} = 0}{k} \cdot \frac{R}{N} \cdot \frac{Z^{n}}{N}$

Trees

Every tree with n

vertices has n-1

ity: If the depths of the leaves of

a binary tree are

and equality holds only if every in-

ternal node has 2

inequal-

edges.

Kraft

 $\overset{d_1,\dots,d}{\overset{_n}{X}^n}:$ 

## Theoretical Computer Science Cheat Sheet Identities Cont. $n^{\frac{n-k}{2}} = n!$ $39. \quad x = \frac{1}{2} \frac{$ 38. $= \frac{1}{k} \cdot \frac{1}{k} \cdot \frac{m+1}{m+1} \cdot \frac{m+1}$ **45.** (n - m)!48.

## Recurrences

Master method:

$$T(n) = aT(n/b) + f(n), a \ge 1, b > 1$$

If  $\exists$ ? > 0 such that f (n) = O(n<sup>log<sub>b</sub> a-?</sup>) then

$$T(n) = \Theta (n^{\log_b a}).$$

If 
$$f(n) = \Theta(n^{\log_b a})$$
 then  $T(n) = \Theta(n^{\log_b a} \log_2 n)$ .

If  $\exists$ ?> 0 such that  $f(n) = \Omega$  ( $n^{\log_b a + ?}$ ), and  $\exists c < 1 \text{ such that af } (n/b) \le cf(n)$ for large n, then

$$T(n)=\Theta (f(n)).$$

Substitution (example): Consider the following recurrence

$$T_{i+1} = 2^{2^i} \cdot T_i^2, T_{i-1} = 2$$
.

Note that T<sub>i</sub> is always a power of two. Let  $t_i = \log_2 T_i$ . Then we have

$$t_{i+1} = 2^{i} + 2 t_{i}, t_{i-1} = 1$$
.

Let  $u_i = t_i/2^i$ . Dividing both sides of the previous equation by  $2^{i+1}$  we get

$$rac{t_{i+1}}{2^{i+1}} = rac{2^i}{2^{i+1}} + rac{t_i}{2^i}.$$

Substituting we find

$$u_{i+1} = \frac{1}{2} + u_i, u$$
  $_1 = \frac{1}{2},$ 

which is simply  $u_i = i/2$ . So we find that  $T_i$  has the closed form  $T_i = 2^{i2^{i-1}}$ . Summing factors (example): Consider the following recurrence

$$T(n)=3 T(n/2) + n, T(1) = 1.$$

Rewrite so that all terms involving T are on the left side

$$T(n) - 3T(n/2) = n.$$

Now expand the recurrence, and choose a factor which makes the left side "telescope"

$$egin{aligned} & rac{1}{2} \mathrm{T(n)} - 3 \mathrm{T(n/2)} = n \ & rac{2}{3} \mathrm{T(n/2)} - 3 \mathrm{T(n/4)} = n/2 \ & \vdots & \vdots & \vdots \ & rac{3 \log_2 n - 1}{2} \mathrm{T(2)} - 3 \mathrm{T(1)} = 2 \end{aligned}$$

Let  $m = log_2 n$ . Summing the left side we get  $T(n) - 3^{m}T(1) = T(n) - 3^{m} =$  $T(n) - n^k$  where  $k = \log_2 3 \approx 1.58496$ . Summing the right side we get

$$\sum_{i=0}^{n} \frac{1}{2^{i}} 3^{i} = n \sum_{i=0}^{n} \frac{1}{2^{i}} \frac{3}{2}^{i}.$$

$$\begin{split} \text{Let } c &= \frac{3}{2}. \text{ Then we have} \\ n &\overset{{}^{n}\!\!X^{\!\!-1}}{\overset{{}^{\circ}}{=}} n \overset{?}{=} \frac{c^{m}-1}{c-1} \overset{?}{=} \\ &= 2 \, n (c^{\log_{2}n}-1) \\ &= 2 \, n (c^{(k-1)\log_{c}n}-1) \\ &= 2 \, n^{k}-2n. \end{split}$$

and so T(n)=3  $n^k-2n$ . Full history recurrences can often be changed to limited history ones (example): Consider

$$T_i = 1 + T_j, T_0 = 1.$$

Note that

$$T_{i+1} = 1 + X^i$$
 $T_{j}$ .

Subtracting we find

And so 
$$T_{i+1} = 2 T_i = 2^{i+1}$$
.

Generating functions:

1. Multiply both sides of the equation by x<sup>i</sup>.

sons.

- 2. Sum both sides over all i for which the equation is valid.
- 3. Choose a generating function G(x). Usually  $G(x) = \int_{i=0}^{\infty} x^{i} g_{i}$ .
- 3. Rewrite the equation in terms of the generating function G(x).
- 4. Solve for G(x).
- 5. The coefficient of  $x^i$  in G(x) is  $g_i$ . Example:

$$g_{i+1} = 2 g_i + 1, g_0 = 0.$$

$$\begin{array}{c} \text{Multiply and sum:} \\ X \\ g_{i+1}x^i = 2g_ix^i + X \\ i \geq 0 \\ i \geq 0 \\ D \end{array} x^i.$$

$$\begin{array}{ll} & \text{i} \geq 0 & \text{i} \geq 0 \\ \text{We choose } G(x) = & P \\ \text{in terms of } G(x) : & \\ & \frac{G(x) - g_0}{x} = 2 \; G(x) + \sum_{i \geq 0}^{X^i} x^i. \end{array}$$

$$\frac{G(x)}{x} = 2 G(x) + \frac{1}{1-x}$$

Solve for 
$$G(x)$$
: 
$$G(x) = \frac{x}{(1-x)(1-2x)}.$$

Expand this using partial fractions:

So 
$$g_i = 2^i - 1$$
.