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 Games
   vector<int> *pvec[MAX];
                           vector<int> G[MAX];
                          9 int sz[MAX], color[MAX], color_counter[MAX];
 Geo
                          pair<ll,int> Info[MAX];
   pair<ll,int>dfs(int u,int p=-1,bool keep=false)
   Half Plane Intersection . . . . . . . . . . . . . . . . .
                             int i,j,k,child,hchild=-1;
for(i=0; i<G[u].size(); i++)</pre>
   Line Segment Intersection . . . . . . . . . . . . .
                          10
   Minimum Perimeter Triangle . . . . . . . . . . .
                          10
                               if(G[u][i]==p) continue;
                               if (hchild==-1||sz[hchild]<sz[G[u][i]])</pre>
   hchild=G[u][i];
   Pair of Intersecting Segments . . . . . . . . . .
                          11
   11
                             for(i=0; i<G[u].size(); i++)</pre>
   12
                               if(G[u][i]==p||G[u][i]==hchild) continue;
                          12
   dfs(G[u][i],u,false);
   12
                             if(hchild!=-1)
 13
                               Info[u] = dfs(hchild, u, true);
                          13
                               pvec[u] =pvec[hchild];
 5.13 radial sweep - anikda unga bunga . . . . . . .
```

```
else
   pvec[u] = new vector < int > ();
pvec[u]->push_back(u);
color_counter[color[u]]++;
if(color_counter[color[u]]>Info[u].second)
   Info[u].second=color_counter[color[u]];
   Info[u].first=color[u];
else if(color_counter[color[u]] == Info[u].second)
   Info[u].first=Info[u].first+color[u];
for(i=0; i<G[u].size(); i++)</pre>
   if(G[u][i]==p||G[u][i]==hchild) continue;
   child=G[u][i];
   for(j=0; j<(*pvec[child]).size(); j++)</pre>
       k=(*pvec[child])[j];
       pvec[u]->push_back(k)
       color_counter[color[k]]++;
       if(color_counter[color[k]]>Info[u].second)
           Info[u].second=color_counter[color[k]];
           Info[u].first=color[k];
 else if(color_counter[color[k]] == Info[u].second)
           Info[u].first=Info[u].first+color[k];
if(!keep)
   for(j=0; j<(*pvec[u]).size(); j++)</pre>
       k=(*pvec[u])[j];
       color_counter[color[k]]--;
return Info[u];
  Divide and Conquer Optimization
```

```
vector<long long> dp_before(n), dp_cur(n);
long long C(int i, int j);
// compute dp_cur[1], ... dp_cur[r] (inclusive)
void compute(int 1, int r, int opt1, int optr) {
    if (1 > r)
         return:
    int mid = (1 + r) >> 1;
    pair<long long, int> best = {LLONG_MAX, -1};
    for (int k = optl; k <= min(mid, optr); k++) {
best = min(best, {(k ? dp_before[k - 1] : 0) + C(k, mid)}</pre>
    dp_cur[mid] = best.first;
    int opt = best.second;
    compute(l, mid - 1, optl, opt);
    compute(mid + 1, r, opt, optr);
    int solve() {
    for (int i = 0; i < n; i++)
         dp_before[i] = C(0, i);
    for (int i = 1; i < m; i++) {
         compute(0, n - 1, 0, n - 1);
         dp_before = dp_cur;
```

```
}
return dp_before[n - 1];
}
```

1.3 Li Chao Tree

```
class LiChaoTree{
   bool minimize;
   int lines;
   struct Node{
       complex<11> line;
       Node *children[2]
       Node(complex<ll> in= {0,100000000000000000})
           children[0]=0;
           children[1]=0;
       }
   } *root;
   11 dot(complex<11> a, complex<11> b){
       return (conj(a) * b).real();
   ll f(complex<ll> a, ll x){
       return dot(a, \{x, 1\});
   void clear(Node* &node){
       if(node->children[0]){
           clear(node->children[0]);
       if(node->children[1]){
           clear(node->children[1]);
       delete node;
   void add_line(complex<ll> nw, Node* &node,ll 1,ll r){
       if(node==0){
           node=new Node(nw);
           return:
       11 m = (1 + r) / 2;
       bool lef = (f(nw, 1) < f(node->line, 1) &&
minimize) | | ((!minimize) \& f(nw, 1) > f(node->line, 1));
       bool mid = (f(nw, m) < f(node->line, m) &&
minimize) | | ((!minimize) &&f(nw, m) > f(node->line, m));
       if(mid) swap(node->line, nw);
       if(r - 1 == 1) return;
       else if(lef != mid)
         add_line(nw, node->children[0], 1, m);
         add_line(nw, node->children[1], m, r);
 ll get(ll x, Node* &node, ll l, ll r){
   1Im = (1 + r) / 2;
   if(r - 1 == 1)
     return f(node->line, x);
   else if (x < m) {
     if(node->children[0]==0) return f(node->line, x);
     if(minimize) return min(f(node->line, x)
                      get(x, node->children[0], 1, m));
     else return max(f(node->line, x),
                  get(x, node->children[0], 1, m));
   else{
       if(node->children[1]==0) return f(node->line, x);
       if(minimize) return min(f(node->line, x),
                    get(x, node->children[1], m, r));
       else return max(f(node->line, x),
                    get(x, node->children[1], m, r));
```

```
public:
LiChaoTree(ll l=-1e9-1,ll r=1e9+1,bool mn=false){
    L=1; R=r; root=0; minimize=mn; lines=0;
}
void AddLine(pair<ll,ll> ln){
    add_line({ln.first,ln.second},root,L,R);
    lines++;
}
int number_of_lines(){
    return lines;
}
ll getOptimumValue(ll x){
    return get(x,root,L,R);
}
~LiChaoTree(){
    if(root!=0) clear(root);
}
};
```

11.4 zero matrix

```
int zero matrix(vector<vector<int>> a) {
    int n = a.size();
    int m = a[0].size();
    int ans = 0;
    vector\langle int \rangle d(m, -1), d1(m), d2(m);
    stack<int> st;
    for (int i = 0; i < n; ++i) {</pre>
       for (int j = 0; j < m; ++j) {
   if (a[i][j] == 1)</pre>
                  d[j] = i;
        for (int j = 0; j < m; ++j) {
            while (!st.empty() && d[st.top()] <= d[j])</pre>
                st.pop();
            d1[j] = st.empty() ? -1 : st.top();
            st.push(j);
        while (!st.empty())
            st.pop();
        for (int j = m - 1; j \ge 0; --j) {
            while (!st.empty() && d[st.top()] <= d[j])</pre>
                st.pop();
            d2[j] = st.empty() ? m : st.top();
            st.push(j);
        while (!st.empty())
            st.pop();
        for (int j = 0; j < m; ++j)
      ans = \max(ans, (i - d[j]) * (d2[j] - d1[j] - 1));
    return ans;
```

2 DS

2.1 BIT 2D

```
const int mx = 1002,my = 1002;
long long bit[4] [mx] [my];
void update( int x, int y, int val, int i ) {
   int y1;
   while( x<=mx ) {
     y1=y;
     while( y1<=my)
        bit[i][x][y1] += val, y1 += (y1&-y1);
     x += (x&-x);
}</pre>
```

```
long long query( int x, int y, int i ) {
 long long ans=0; int y1;
 while (x>0)
  y1 = y;
   while( y1>0 )
    ans += bit[i][x][y1], y1 -= (y1&-y1);
  x = (x\&-x);
 return ans;
// add value k from (x1,y1) to (x2,y2) inclusive
void add( int x1, int y1, int x2, int y2, int k) {
update(x1,y1,k,0);
 update(x1,y2+1,-k,0);
 update(x2+1,y1,-k,0);
 update(x2+1,y2+1,k,0);
 update(x1,y1,k*(1-y1),1);
 update(x1, y2+1, k*y2, 1);
 update(x2+1,y1,k*(y1-1),1);
 update(x2+1,y2+1,-y2*k,1);
 update(x1,y1,k*(1-x1),2);
 update(x1,y2+1,k*(x1-1),2);
 update(x2+1,y1,k*x2,2);
 update(x2+1,y2+1,-x2*k,2)
 update(x1,y1,(x1-1)*(y1-1)*k,3);
 update(x1, y2+1, -y2*(x1-1)*k, 3);
 update(x2+1,y1,-x2*(y1-1)*k,3);
 update(x2+1,y2+1,x2*y2*k,3);
// get value from (x1,y1) to (x2,y2) inclusive
long long get( int x1, int y1, int x2, int y2 ) {
LL v1=query(x2,y2,0)*x2*y2 +
      query(x2, y2, 1)*x2 +
      query(x2,y2,2)*y2 +
      query(x2,y2,3);
 LL v2=query(x2,y1-1,0)*x2*(y1-1) +
      query(x2,y1-1,1)*x2 +
      query(x2,y1-1,3) +
      query(x2,y1-1,2)*(y1-1);
 LL v3=query(x1-1,y2,0)*(x1-1)*y2 +
      query(x1-1,y2,2)*y2+
      query(x1-1,y2,1)*(x1-1) +
      query(x1-1,y2,3);
 LL v4=query(x1-1,y1-1,0)*(x1-1)*(y1-1) +
      query(x1-1,y1-1,1)*(x1-1) +
      query(x1-1,y1-1,2)*(y1-1) +
      query(x1-1,y1-1,3);
 LL ans=v1-v2-v3+v4;
 return ans;
```

2.2 CD - hellbent

```
vector <int> g[N]; int n, child[N], done[N];
void dfs_size(int u, int par) {
  child[u] = 1;
  for (int v: g[u]) {
    if (done[v] or v == par) continue;
    dfs_size(v, u); child[u] += child[v];
  }
}
int dfs_find_centroid(int u, int par, int sz) {
  for (int v: g[u]) {
    if (!done[v] and v != par and child[v] > sz) {
      return dfs_find_centroid(v,u,sz);
    }
  }
  return u;
```

```
}
void solve (int u) {/**problem specific things */}
void dfs_decompose(int u) {
    dfs_size(u, -1);
    int centroid=dfs_find_centroid(u,-1,child[u]/2);
    solve(centroid);
    done[centroid] = 1;
    for (int v : g[centroid]) {
        if (!done[v]) dfs_decompose(v);
    }
}
```

2.3 Hld - cpalgo

```
vector<int> parent, depth, heavy, head, pos;
int dfs(int v, vector<vector<int>> const& adj) {
   int size = 1:
   int max_c_size = 0;
   for (int c : adj[v]) {
       if (c != parent[v]) {
           parent[c] = v, depth[c] = depth[v] + 1;
           int c_size = dfs(c, adj);
           size += c size:
           if (c_size > max_c_size)
              max_c_size = c_size, heavy[v] = c;}}
   return size:}
decompose(int v,int h,vector<vector<int>> const& adj){
   head[v] = h, pos[v] = cur_pos++;
   if (\text{heavy}[v] \stackrel{!}{!} = -1)
       decompose(heavy[v], h, adj);
   for (int c : adj[v]) {
       if (c != parent[v] && c != heavy[v])
           decompose(c, c, adj);}
void init(vector<vector<int>> const& adj) {
   int n = adj.size();
   parent = vector<int>(n);
   depth = vector<int>(n);
   heavy = vector\langle int \rangle (n, -1);
   head = vector<int>(n);
   pos = vector<int>(n);
   cur_pos = 0;
   dfs(0, adj);
   decompose(0, 0, adj);
int query(int a, int b) {
  int res = 0;
  for (; head[a] != head[b]; b = parent[head[b]]) {
     if (depth[head[a]] > depth[head[b]])
         swap(a, b);
int cur_heavy_path_max =
segment_tree_query(pos[head[b]], pos[b]);
     res = max(res, cur_heavy_path_max);
  if (depth[a] > depth[b])
     swap(a, b);
  int last_heavy_path_max =
  segment_tree_query(pos[a], pos[b]);
 res = max(res, last_heavy_path_max);
 return res;
```

2.4 Implicit Treap

```
#include<bits/stdc++.h>
#include<math.h>
#include<vector>
#include<stdlib.h>
using namespace std;
```

```
#define MAX 200005
#define MOD 998244353
template <class T>
class implicit_treap
   struct item
       int prior, cnt;
       T value;
       bool rev;
       item *1,*r;
       item(T v)
           value=v;
           rev=false;
           1=NULL;
           r=NULL:
           cnt=1;
           prior=rand();
       }
   } *root,*node;
   int cnt (item * it)
       return it ? it->cnt : 0:
   void upd_cnt (item * it)
           it->cnt = cnt(it->1) + cnt(it->r) + 1;
   void push (item * it)
       if (it && it->rev)
           it->rev = false;
           swap (it->1, it->r);
           if (it->1) it->1->rev ^= true;
           if (it->r) it->r->rev ^= true:
   }
   void merge (item * & t, item * 1, item * r)
       push (1);
       push (r):
       if (!1 || !r)
           t = 1 ? 1 : r;
       else if (l->prior > r->prior)
           merge (1->r, 1->r, r), t = 1;
           merge (r->1, 1, r->1), t = r;
       upd_cnt (t);
   }
   void split (item * t, item * & 1, item * & r, int key,
       if (!t)
           return void( 1 = r = 0 );
       push (t):
       int cur_key = add + cnt(t->1);
       if (key <= cur_key)</pre>
           split (t->1, 1, t->1, key, add), r = t;
           split (t\rightarrow r, t\rightarrow r, r, key, add + 1 + cnt(t\rightarrow l))
1 = t;
       upd_cnt (t);
   void insert(item * &t.item * element.int kev)
```

```
item *1.*r:
       split(t,l,r,key);
       merge(1,1,element);
       merge(t,1,r);
       1=NŬLL;
       r=NULL;
   T elementAt(item * &t,int key)
       push(t);
       T ans;
       if(cnt(t->1)==key) ans=t->value;
       else if(cnt(t->1)>key) ans=elementAt(t->1,key);
       else ans=elementAt(t->r,key-1-cnt(t->l));
       return ans;
   void erase (item * & t, int key)
       push(t);
       if(!t) return;
       if (key == cnt(t->1))
          merge (t, t->1, t->r);
       else if(key<cnt(t->1))
          erase(t->1,key);
          erase(t->r, key-cnt(t->1)-1);
       upd_cnt(t);
   void reverse (item * &t, int 1, int r)
       item *t1, *t2, *t3;
       split (t, t1, t2, 1);
       split (t2, t2, t3, r-l+1);
       t2->rev ^= true;
       merge (t, t1, t2);
       merge (t, t, t3);
   void cyclic_shift(item * &t,int L,int R)
       if(L==R) return;
       item *1,*r,*m;
       split(t,t,l,L);
       split(1,1,m,R-L+1);
       split(1,1,r,R-L);
       merge(t,t,r);
       merge(t,t,1);
       merge(t,t,m);
       1=NŬLL;
      r=NULL:
       m=NULL;
   void output (item * t,vector<T> &arr)
       if (!t) return;
int add = push (t);
output (t->1,arr);
       arr.push_back(t->value);
       output (t->r,arr);
public:
   implicit_treap()
       root=NULL;
   void insert(T value,int position)
       node=new item(value);
       insert(root, node, position);
```

```
void erase(int position)
       erase(root, position);
   void reverse(int 1,int r)
       reverse(root,1,r);
     elementAt(int position)
       return elementAt(root,position);
   void cyclic_shift(int L,int R)
       cyclic_shift(root,L,R);
   int size()
       return cnt(root);
   void output(vector<T> &arr)
       output(root,arr);
};
```

Mo Algorithm

```
void remove(int idx);
void add(int idx);
int get_answer();
// TODO: extract the current answer of the data structure
int block_size;
struct Query {
  int 1, r,k, idx;
  bool operator<(Query other) const{</pre>
   if(l/block_size!=other.l/block_size)return(l<other.l);</pre>
   return (1/block_size&1)? (r<other.r) : (r>other.r);
vector<int> mo_s_algorithm(vector<Query> queries) {
   vector<int> answers(queries.size());
   sort(queries.begin(), queries.end());
    // TODO: initialize data structure
    int cur_1 = 0;
   int cur_r = -1;
   // invariant: data structure will always
   //reflect the range [cur_1, cur_r]
   for (Query q : queries) {
       while (cur_1 > q.1) {
           cur_1--;
           add(cur_1);
       while (cur_r < q.r) {</pre>
           cur_r++;
           add(cur_r);
       while (cur_1 < q.1) {
           remove(cur_1);
           cur_l++;
       while (cur_r > q.r) {
           remove(cur_r);
           cur_r--;
       answers[q.idx] = get_answer();
   return answers:
```

12.6 Treap

```
#include<bits/stdc++.h>
#include<math.h>
#include<vector>
#include<stdlib.h>
using namespace std;
#define MAX 400005
#define MOD 998244353
#define INF 2000000000
template <class T>
class treap
    struct item
       int prior, cnt;
       T key;
       item *1,*r;
       item(T v)
           kev=v
           1=ŇULĹ:
           r=NULL;
           cnt=1;
           prior=rand();
    } *root,*node;
    int cnt (item * it)
       return it ? it->cnt : 0;
    void upd_cnt (item * it)
       if (it)
           it->cnt = cnt(it->1) + cnt(it->r) + 1;
    void split (item * t, T key, item * & 1, item * & r)
       if (!t)
           1 = r = NULL:
       else if (key < t->key)
           split (t->1, key, 1, t->1), r = t;
           split (t\rightarrow r, key, t\rightarrow r, r), l = t;
       upd_cnt(t);
    }
    void insert (item * & t, item * it)
       if (!t)
           t = it;
       else if (it->prior > t->prior)
           split (t, it->kev, it->l, it->r), t = it;
           insert (it->key < t->key ? t->l : t->r, it);
       upd_cnt(t);
    void merge (item * & t, item * 1, item * r)
       if (!1 || !r)
           t = 1 ? 1 : r;
       else if (l->prior > r->prior)
           merge (1->r, 1->r, r), t = 1;
           merge (r->1, 1, r->1), t = r;
       upd_cnt(t);
    void erase (item * & t, T key)
```

```
if (t->kev == kev)
          merge (t, t->1, t->r);
          erase (key < t->key ? t->1 : t->r, key);
       upd_cnt(t);
   T elementAt(item * &t,int key)
       T ans:
       if(cnt(t->1)==key) ans=t->key;
       else if(cnt(t->1)>key) ans=elementAt(t->1,key);
       else ans=elementAt(t->r,key-1-cnt(t->l));
       upd cnt(t):
       return ans;
   item * unite (item * 1, item * r)
       if (!1 || !r) return 1 ? 1 : r;
       if (l->prior < r->prior) swap (l, r);
       item * lt, * rt;
       split (r, 1->key, lt, rt);
       1->1 = unite (1->1, 1t);
       1->r = unite (1->r, rt);
       upd_cnt(1);
       upd_cnt(r);
       return 1:
   void heapify (item * t)
       if (!t) return;
       item * max = t
       if (t->l != NULL && t->l->prior > max->prior)
          max = t->1:
       if (t->r != NULL && t->r->prior > max->prior)
          max = t->r;
       if (max != t)
          swap (t->prior, max->prior);
          heapify (max);
   }
   item * build (T * a, int n)
       if (n == 0) return NULL;
       int mid = n / 2;
       item * t = new item (a[mid], rand ());
       t->1 = build (a, mid);
       t->r = build (a + mid + 1, n - mid - 1);
       heapify (t);
       return t:
   void output (item * t,vector<T> &arr)
       if (!t) return:
       output (t->1,arr);
       arr.push_back(t->key);
       output (t->r,arr);
public:
   treap()
       root=NULL:
   treap(T *a, int n)
       build(a,n);
   void insert(T value)
```

```
node=new item(value);
       insert(root, node);
   void erase(T value)
       erase(root, value);
     elementAt(int position)
       return elementAt(root, position);
    int size()
       return cnt(root);
   void output(vector<T> &arr)
       output(root,arr);
   int range_query(T 1,T r) //(1,r]
       item *previous,*next,*current;
       split(root,1,previous,current);
       split(current,r,current,next);
       int ans=cnt(current);
       merge(root, previous, current);
       merge(root,root,next);
       previous=NÚLL;
       current=NULL;
       next=NULL;
       return ans;
};
```

Wavelet Tree

```
struct wavelet_tree{
 int lo , hi ;
wavelet_tree *l=0,*r=0;
  vi b,c;//c for prefix sum
//nos are in range [x,y] , array indices [from,to)
 wavelet_tree(int *from, int *to, int x, int y){
   lo = x, hi = y;
   if( from >= to) return;
   if( hi == lo ) {
     b.reserve(to-from+1);
     b.pb(0);
     c.reserve(to-from+1);
     c.pb(0);
     for(auto it = from; it != to; it++){
       b.pb(b.back() + 1);
       c.pb(c.back()+*it);
     return :
   int mid = (lo+hi)/2;
   auto f = [mid](int x){
     return x <= mid;</pre>
   b.reserve(to-from+1);
   b.pb(0);
   c.reserve(to-from+1);
   c.pb(0);
   for(auto it = from; it != to; it++){
     b.pb(b.back() + f(*it));
     c.pb(c.back() + *it);
   auto pivot = stable_partition(from, to, f);
   1 = new wavelet_tree(from, pivot, lo, mid);
```

```
r = new wavelet_tree(pivot, to, mid+1, hi);
 }// swap a[i] with a[i+1] , if a[i]!=a[i+1]
 void swapadjacent(int i){
   if(lo == hi) return
   b[i] = b[i-1] + b[i+1] - b[i];

c[i] = c[i-1] + c[i+1] - c[i];
   if(b[i+1]-b[i] == b[i] - b[i-1]){
     if(b[i]-b[i-1])
           return this->l->swapadjacent(b[i]);
     else return this->r->swapadjacent(i-b[i]);
//count of nos in [1, r] Less than or equal to k
 int LTE(int 1, int r, int k){
   if(1 > r \text{ or } k < 10) \text{ return } 0;
   if(hi <= k) return r - l + 1;</pre>
   int lb = b[1-1], rb = b[r];
   return this->l->LTE(lb+1,rb,k)+
           this->r->LTE(1-1b,r-rb,k);
//count of nos in [l, r] equal to k
 int count(int 1, int r, int k){
   if(l > r or k < lo or k > hi) return 0;
   if(lo == hi) return r - l + 1;
   int lb = b[1-1], rb = b[r], mid = (lo+hi)/2;
   if(k <= mid) return this->l->count(lb+1,rb,k);
   return this->r->count(l-lb, r-rb, k);
//sum of nos in [l ,r] less than or equal to k
 int sumk(int 1, int r, int k){
   if(1 > r \text{ or } k < 10) \text{ return } 0;
   if (hi \leq k) return c[r]-c[l-1];
   int lb = b[1-1], rb = b[r];
   return this->l->sumk(lb+1,rb,k)+
           this->r->sumk(l-lb,r-rb,k);
  ~wavelet_tree(){
   if(1) delete 1;
   if(r) delete r;
      ordered set
```

```
#include <ext/pb_ds/assoc_container.hpp> // Common file
using namespace __gnu_pbds;
//find_by_order(k) > returns iterator to the kth largest
// element counting from 0 , order_of_key(val) --> the
//number of items that are strictly smaller than our item
typedef tree< int, null_type, less<int>, rb_tree_tag,
tree_order_statistics_node_update> ordered_set;
```

sparse table 2d

```
int st[K][K][N][N]; int lg[N];
void pre() {
 lg[1] = 0;
 for (int i=2; i<N; i++) lg[i] = lg[i/2]+1;
int query(int 11, int r1, int 12, int r2) {
 int xx = lg[12-l1+1], yy = lg[r2-r1+1];
 return max(max(st[xx][yy][11][r1],
            st[xx][yy][12-(1<(xx)+1][r1]),
        \max(\text{st}[xx][yy][11][r2-(1<<yy)+1],
        st[xx][yy][12-(1<<xx)+1][r2-(1<<yy)+1]);
void build() {
 for (int x=0; x<K; x++) {
   for (int y=0; y<K; y++) {</pre>
     for (int i=1; i<=n; i++) {
```

```
for (int j=1; j<=m; j++) {
            if (i+(1<< x)-1>n | j+(1<< y)-1>m)
              continue;
            if (!x&&!y) st[0][0][i][j]=flag[i][j];
else_if_(x>0) st[x][y][i][j] =
\max(\text{st}[x-1][y][i][j], \text{st}[x-1][y][i+(1<<(x-1))][j]);
            else if (\tilde{y}>0) st[x][\tilde{y}][i][j] =
\max(\text{st}[x][y-1][i][j], \text{st}[x][y-1][i][j+(1<<(y-1))]);
  }
```

Flow

Dinic's Algorithm

```
struct FlowEdge {
   int v, u;
   long long cap, flow = 0;
   FlowEdge(int v, int u, long long cap) :
     v(v), u(u), cap(cap) {}
struct Dinic {
 const long long flow_inf = 1e18;
 vector<FlowEdge> edges;
 vector<vector<int>> adj;
 int n, m = 0;
 int s, t;
 vector<int> level, ptr;
 queue<int> q;
 Dinic(int n, int s, int t) : n(n), s(s), t(t) {
     adj.resize(n);
     level.resize(n);
     ptr.resize(n);
 void add_edge(int v, int u, long long cap) {
     edges.emplace_back(v, u, cap);
     edges.emplace_back(u, v, 0);
     adj[v].push_back(m);
     adj[u].push_back(m + 1);
     m += 2;
 bool bfs() {
     while (!q.empty()) {
         int v = q.front();
         q.pop();
         for (int id : adj[v]) {
            if (edges[id].cap - edges[id].flow < 1)</pre>
            if (level[edges[id].u] != -1)
            level[edges[id].u] = level[v] + 1;
            q.push(edges[id].u);
     return level[t] != -1;
 long long dfs(int v, long long pushed) {
   if (pushed == 0)
       return 0;
   if (v == t)
       return pushed;
   for(int&cid=ptr[v];cid<(int)adj[v].size();cid++){</pre>
     int id = adj[v][cid];
     int u = edges[id].u;
     if (level[v] + 1 != level[u]
           || edges[id].cap - edges[id].flow < 1)</pre>
```

```
long long tr = dfs(u.
         min(pushed, edges[id].cap - edges[id].flow));
     if (tr = 0)
         continue:
     edges[id].flow += tr;
     edges[id ^ 1].flow -= tr;
   return 0;
 long long flow() {
     long long f = 0;
     while (true) {
         fill(level.begin(), level.end(), -1);
         level[s] = 0;
         q.push(s);
         if (!bfs())
             break;
         fill(ptr.begin(), ptr.end(), 0);
         while (long long pushed = dfs(s, flow_inf)) {
            f += pushed;
     return f;
};
```

3.2 Edmond's Blossom Algorithm

```
GETS:
V->number of vertices
E->number of edges
pair of vertices as edges (vertices are 1..V)
output of edmonds() is the maximum matching
match[i] is matched pair of i
(-1 if there isn't a matched pair)
#include <bits/stdc++.h>
using namespace std;
const int M=500;
struct struct_edge
   struct_edge* n;
typedef struct_edge* edge;
struct_edge pool[M*M*2];
edge top=pool,adj[M];
int V,E,match[M],qh,qt,q[M],father[M],base[M];
bool inq[M],inb[M],ed[M][M];
void add_edge(int u,int v)
   top->v=v,top->n=adj[u],adj[u]=top++;
   top->v=u, top->n=adj[v], adj[v]=top++;
int LCA(int root,int u,int v)
   static bool inp[M];
   memset(inp,0,sizeof(inp));
   while(1)
       inp[u=base[u]]=true;
       if (u==root) break;
       u=father[match[u]];
   while(1)
```

```
if (inp[v=base[v]]) return v;
       else v=father[match[v]];
void mark_blossom(int lca,int u)
   while (base[u]!=lca)
       int v=match[u];
       inb[base[u]]=inb[base[v]]=true;
       u=father[v];
       if (base[u]!=lca) father[u]=v;
void blossom_contraction(int s,int u,int v)
   int lca=LCA(s,u,v);
   memset(inb,0,sizeof(inb));
   mark_blossom(lca,u);
   mark_blossom(lca,v);
   if (base[u]!=lca)
       father[u]=v;
   if (base[v]!=lca)
       father[v]=u;
   for (int u=0; u<V; u++)</pre>
       if (inb[base[u]])
           base[u]=lca:
           if (!inq[u])
              inq[q[++qt]=u]=true;
int find_augmenting_path(int s)
 memset(inq,0,sizeof(inq));
 memset(father,-1,sizeof(father));
 for (int i=0; i<V; i++) base[i]=i;</pre>
  inq[q[qh=qt=0]=s]=true;
  while (qh<=qt)</pre>
   int u=q[qh++];
   for (edge e=adj[u]; e; e=e->n)
     int v=e->v:
     if (base[u]!=base[v]&&match[u]!=v)
       if((v==s)||(match[v]!=-1 && father[match[v]]!=-1))
           blossom_contraction(s,u,v);
       else if (father[v]==-1)
         father[v]=u;
         if (match[v] == -1)
            return v;
         else if (!inq[match[v]])
             inq[q[++qt]=match[v]]=true;
   }
 return -1;
int augment_path(int s,int t)
   int u=t,v,w;
   while (u!=-1)
       v=father[u];
       w=match[v];
       match[v]=u;
       match[u]=v;
```

3.3 Hungarian Algorithm

```
class HungarianAlgorithm{
 int N,inf,n,max_match;
 int *1x,*1y,*xy,*yx,*slack,*slackx,*prev;
 int **cost:
 bool *S,*T;
 void init_labels(){
     for(int x=0;x<n;x++) lx[x]=0;</pre>
     for(int y=0; y< n; y++) ly[y]=0;
     for (int x = 0; x < n; x++)
         for (int y = 0; y < n; y++)
             lx[x] = max(lx[x], cost[x][y]);
 void update labels(){
     int x, y, delta = inf; //init delta as infinity
     for (y = 0; y < n; y++)//calculate delta using slack</pre>
         if (!T[y])
             delta = min(delta, slack[y]);
     for (x = 0; x < n; x++) //update X labels
         if (S[x]) lx[x] = delta;
     for (y = 0; y < n; y++) //update Y labels
   if (T[y]) ly[y] += delta;</pre>
     for (y = 0; y < n; y++) //update slack array
    if (!T[y])</pre>
             slack[y] -= delta;
 void add_to_tree(int x, int prevx)
//x - current vertex,prevx - vertex from X before x in
// the alternating path,
//so we add edges (prevx, xy[x]), (xy[x], x)
   S[x] = true; //add x to S
   prev[x] = prevx; //we need this when augmenting
   for (int y = 0; y < n; y++) //update slacks, because
     //we add new vertex to S
       if (lx[x] + ly[y] - cost[x][y] < slack[y])
           slack[y] = lx[x] + ly[y] - cost[x][y];
           slackx[y] = x;
 void augment(){ //main function of the algorithm
   if (max_match == n) return; //check wether matching
```

```
// is already perfect
   int x, y, root; //just counters and root vertex
   int q[N], wr = 0, rd = 0; //q - queue for bfs, wr,
   //rd - write and read
//pos in queue
   //memset(S, false, sizeof(S)); //init set S
   for(int i=0;i<n;i++) S[i]=false;</pre>
   //memset(T, false, sizeof(T)); //init set T
   for(int i=0;i<n;i++) T[i]=false;</pre>
   //memset(prev, -1, sizeof(prev)); //init set prev
   //- for the alternating tree
   for(int i=0;i<n;i++) prev[i]=-1;</pre>
   for (x = 0; x < n; x++){ //finding root of the tree
       if (xy[x] == -1)
           q[wr++] = root = x;
          prev[x] = -2;
           S[x] = true;
           break;
   for (y = 0; y < n; y++){//initializing slack array}
       slack[y] = lx[root] + ly[y] - cost[root][y];
       slackx[y] = root;
   while (true){ //main cycle
       while (rd < wr) //building tree with bfs cycle
           x = q[rd++]; //current vertex from X part
           for (y = 0; y < n; y++) //iterate through
            //all edges in equality graph
              if (cost[x][y] == lx[x]+ly[y] && !T[y])
                  if (yx[y] == -1) break; //an
          //exposed vertex in Y found, so
//augmenting path exists!
                  T[y] = true; //else just add y to T,
                  q[wr++] = yx[y]; //add vertex yx[y]
         //, which is matched with y, to the queue
                  add_to_tree(yx[y], x);
         //add edges (x,y) and (y,yx[y]) to the tree
           if (y < n) break; //augmenting path found!
       if (y < n) break; //augmenting path found!</pre>
       update_labels(); //augmenting path not found,
       //so improve labeling
       wr = rd = 0;
       for (y = 0; y < n; y++)
//in this cycle we add edges that were added to the
//equality graph as a result of improving the labeling,
// we add edge (slackx[y], y) to the tree if and only if
// !T[y] && slack[y] == 0, also with this edge we add
// another one (y, yx[y]) or augment the matching, if y
// was exposed
           if (!T[y] && slack[y] == 0){
              if (yx[y] == -1){
//exposed vertex in Y found - augmenting path exists!
                  x = slackx[y]; break;
              else{
                  T[y] = true; //else just add y to T,
                  if (!S[yx[y]]){
                      q[wr++] = yx[y];
//add vertex yx[y], which is matched with y, to the queue
```

```
add_to_tree(yx[y], slackx[y]);
     //and add edges (x,y) and (y,yx[y]) to the tree
           }
       if (y < n) break; //augmenting path found!
   if (y < n){ //we found augmenting path!</pre>
       max_match++; //increment matching
//in this cycle we inverse edges along augmenting path
       for (int cx = x, cy = y, ty; cx != -2;
cx = prev[cx], cy = ty){
           ty = xy[cx];
yx[cy] = cx;
           xy[cx] = cy;
       augment(); //recall function,
       // go to step 1 of the algorithm
 }//end of augment() function
 HungarianAlgorithm(int vv,int inf=1000000000)
   N=vv;
   n=N;
   max_match=0;
   this->inf=inf;
   lx=new int[N];
   ly=new int[N];//labels of X and Y parts
   xy=new int[N];//xy[x] - vertex that is matched with x,
   yx=new int[N];//yx[y] - vertex that is matched with y
   slack=new int[N];//as in the algorithm description
   slackx=new int[N];//slackx[y] such a vertex, that
   //1(\operatorname{slackx}[y]) + 1(y) - w(\operatorname{slackx}[y], y) = \operatorname{slack}[y]
   prev=new int[N];
    //array for memorizing alternating paths
   S=new bool[N];
T=new bool[N];//sets S and T in algorithm
   cost=new int*[N];//cost matrix
   for(int i=0; i<N; i++){</pre>
       cost[i]=new int[N];
  ~HungarianAlgorithm(){
   delete []lx;
delete []ly;
   delete []xy;
   delete [] yx;
   delete []slack;
   delete []slackx;
   delete []prev;
   delete []S;
   delete []T;
   int i;
   for(i=0; i<N; i++)</pre>
       delete [](cost[i]);
   delete []cost;
 void setCost(int i,int j,int c){
     cost[i][j]=c;
 int* matching(bool first=true){
   int *ans:
   ans=new int[N];
   for(int i=0;i<N;i++)</pre>
```

```
if(first) ans[i]=xy[i];
    else ans[i]=yx[i];
}
return ans;
}
int hungarian(){
    int ret = 0; //weight of the optimal matching
    max_match = 0; //number of vertices in current matching
    for(int x=0;x<n;x++) xy[x]=-1;
    for(int y=0;y<n;y++) yx[y]=-1;
    init_labels(); //step 0
    augment(); //steps 1-3
    for (int x = 0; x < n; x++) //forming answer there
        ret += cost[x][xy[x]];
    return ret;
};</pre>
```

|3.4 Maximum Bipartite Matching

```
// A class to represent Bipartite graph for
// Hopcroft Karp implementation
class BGraph{
   // m and n are number of vertices on left
   // and right sides of Bipartite Graph
   // adj[u] stores adjacents of left side
   // vertex 'u'. The value of u ranges from 1 to m.
   // 0 is used for dummy vertex
   std::list<int> *adj;
   // pointers for hopcroftKarp()
   int *pair_u, *pair_v, *dist;
public:
   BGraph(int m, int n); // Constructor
   void addEdge(int u, int v); // To add edge
   // Returns true if there is an augmenting path
   // Adds augmenting path if there is one beginning
   // with u
   bool dfs(int u);
   // Returns size of maximum matching
   int hopcroftKarpAlgorithm();
// Returns size of maximum matching
int BGraph::hopcroftKarpAlgorithm(){
   // pair_u[u] stores pair of u in matching on left side
   // of Bipartite Graph.
   // If u doesn't have any pair, then pair_u[u] is NIL
   pair_u = new int[m + 1];
   // pair_v[v] stores pair of v in matching on right
   // side of Biparite Graph.
   // If v doesn't have any pair, then pair_u[v] is NIL
   pair_v = new int[n + 1];
   // dist[u] stores distance of left side vertices
   dist = new int[m + 1];
   // Initialize NIL as pair of all vertices
   for (int u = 0; u <= m; u++)
       pair_u[u] = NIL;
   for (int v = 0; v \le n; v++)
       pair_v[v] = NIL;
   // Initialize result
   int result = 0;
   // Keep updating the result while there is an
   // augmenting path possible.
   while (bfs())
       // Find a free vertex to check for a matching
       for (int u = 1: u <= m: u++)
```

```
// If current vertex is free and there is
           // an augmenting path from current vertex
           // then increment the result
           if (pair_u[u] == NIL && dfs(u))
              result++;
   return result;
// Returns true if there is an augmenting path available,
//else returns false
bool BGraph::bfs(){
 std::queue<int> q; //an integer queue for bfs
 // First layer of vertices (set distance as 0)
 for (int u = 1; u <= m; u++){
     // If this is a free vertex, add it to queue
     if (pair_u[u] == NIL)
         // u is not matched so distance is 0
         dist[u] = 0;
         q.push(u);
     // Else set distance as infinite so that this
     //vertex is considered next time for availibility
         dist[u] = INF;
 // Initialize distance to NIL as infinite
 dist[NIL] = INF;
 // q is going to contain vertices of left side only.
 while (!q.empty()){
   // dequeue a vertex
   int u = q.front();
   q.pop();
// If this node is not NIL and can provide a
   //shorter path to NIL then
   if (dist[u] < dist[NIL]){</pre>
// Get all the adjacent vertices of the dequeued vertex u
     std::list<int>::iterator it;
     for (it = adj[u].begin(); it != adj[u].end(); ++it)
       int v = *it;
       // If pair of v is not considered so far
       // i.e. (v, pair_v[v]) is not yet explored edge.
if (dist[pair_v[v]] == INF)
           // Consider the pair and push it to queue
           dist[pair_v[v]] = dist[u] + 1;
           q.push(pair_v[v]);
     }
 // If we could come back to NIL using alternating path
 // of distinct
 // vertices then there is an augmenting path available
 return (dist[NIL] != INF);
// Returns true if there is an augmenting path beginning
// with free vertex u
bool BGraph::dfs(int u){
 if (u != NIL){
   std::list<int>::iterator it;
   for (it = adj[u].begin(); it != adj[u].end();++it)
     // Adjacent vertex of u
     int v = *it;
     // Follow the distances set by BFS search
     if (dist[pair_v[v]] == dist[u] + 1)
```

```
// If dfs for pair of v also returnn true then
         if (dfs(pair_v[v]) == true)
         { // new matching possible, store the matching
            pair_v[v] = u;
            pair_u[u] = v;
            return true;
  //If there is no augmenting path beginning with u then.
     return false;
  return true;
// Constructor for initialization
BGraph::BGraph(int m, int n){
   this->m = m;
   this -> n = n;
   adj = new std::list<int>[m + 1];
// function to add edge from u to v
void BGraph::addEdge(int u, int v){
   adj[u].push_back(v); // Add v to us list.
```

13.5 Minimum Cost Maximum Flow

```
struct Edge{
   int from, to, capacity, cost;
vector<vector<int>> adj, cost, capacity;
const int INF = 1e9;
void shortest_paths(int n, int v0, vector<int>& d,
                  vector<int>& p) {
 d.assign(n, INF);
 d[v0] = 0:
 vector<bool> inq(n, false);
 queue<int> q;
 q.push(v0);
 p.assign(n, -1);
 while (!q.empty()) {
   int u = q.front();
   q.pop();
   inq[u] = false;
   for (int v : adj[u]) {
     if (capacity[u][v] > 0 && d[v] > d[u] +cost[u][v]){
         d[v] = d[u] + cost[u][v];
         p[v] = u;
         if (!inq[v]) {
             inq[v] = true;
             q.push(v);
int min_cost_flow(int N, vector < Edge > edges, int K,
                int s,int t){
 adj.assign(N, vector<int>());
 cost.assign(N, vector<int>(N, 0));
 capacity.assign(N, vector<int>(N, 0));
 for (Edge e : edges) {
   adj[e.from].push_back(e.to);
   adj[e.to].push_back(e.from);
   cost[e.from][e.to] = e.cost;
   cost[e.to][e.from] = -e.cost;
   capacity[e.from][e.to] = e.capacity;
```

```
int flow = 0;
int cost = 0:
vector<int> d, p;
while (flow < K) {
 shortest_paths(N, s, d, p);
 if (d[t] == INF)
     break;
  // find max flow on that path
 int f = K - flow;
 int cur = t;
 while (cur != s) {
     f = min(f, capacity[p[cur]][cur]);
     cur = p[cur];
 // apply flow
flow += f;
 cost += f * d[t];
 cur = t;
 while (cur != s) +
     capacity[p[cur]][cur] -= f;
     capacity[cur][p[cur]] += f;
     cur = p[cur];
if (flow < K)</pre>
   return -1;
else
   return cost;
```

3.6 notes

/* Partial Ordered Set Let S be a set of elements and be a partial ordering on the set. That is, for some elements x and y in S we may . The only properties that must satisfy a reflexivity (xx), antisymmetry (if then x = y), and transitivity (if xy and). Note that because it is a partial ordering, not all x and y are comparable. An example of a partially ordered set (poset) is the set of points (x,y) in the Cartesian plane with the operator (x1,y1) (x2, y2) if: y1y2 (e.g. NCPC 2007 problem). and

We define a chain C to be a subset of S such that the elements of C can be labelled $x1,x2,\ldots,xn$ such that $x1x2\ldots xn$. A partition P is a set of each element occurs in exactly one chain.

We define an antichain A to be a subset of S such that for any x and y in A we have neither xy nor yx. That no two elements of an antichain are comparable.

Dilworth's Theorem We can define the width of a poset in two ways, which is the result of Dilworth's Theorem. One definition is the size of the maximum antichain; the other is the size of the minimum partition. It is easy to see that any partition must have at least as many chains as the size of the maximum antichain because every element of the antichain must be in a different chain. Dilworth's Theorem tells us that there exists a partition of exactly that size and can be proved inductively.

Calculating the Width So how does one calculate the width of a poset? To solve this problem in general, we can use maximum matching on a bipartite graph. Duplicate each element x in S as ux and vx and if xy (for xy) then add the edge .If yo compute the maximum matching on this graph, this is equivalent to partitioning the elements into chains, where if the edge is chosen then x and y are in the same

chain. In particular, if the size of the matching is m and minimum cost will be the answer. Otherwise it will not be |S| = n , then the number of partitions is n-m . Notice possible. that this gives a bijection between partitions and matchings, so the maximum matching gives the minimum partition, which we know is equal to the width as given by Dilworth's Theorem.

MAX/MIN Closure :

Closure is a directed subgraph with no outgoing edges. MAX/Min Closure means a closure with max/min sum of weighted nodes outside the graph. For MAX, join source with positive nodes, sink with negative nodes and capacities are absolute value, infinite capacity between existing edges. For min, source & sink is reversed. Ans = sum of positive nodes - min cut (For max) Ans = sum of negative nodes + min cut (For min)

Demand Capacity :

We add a new source s and a new sink t, a new edge from the source s to every other vertex, a new edge for every to sink of capacity p. Ans = sum of w - maxflow vertex to the sink t, and one edge from t to s. Additionally we define the new capacity function c as:

```
c ((s,v)) = uVd ((u,v)) for each edge (s,v).
c((v, t)) = wVd((v, w)) for each edge (v, t).
c((u,v))=c((u,v)) d((u,v)) for each edge (u,v) in the
       old network.
c((t,s))=
```

If the new network has a saturating flow (a flow where each edge outgoing from s is completely filled, which is equivalent to every edge incoming to t is completely filled), then the network with demands has a valid flow, and the actual flow can be easily reconstructed from the new network. Otherwise there doesn't exist a flow that satisfies all conditions.

Knig - Egevry Theorem

the maximum matching is equal to the vertex cover number for a bipartite graph.

Minimum Vertex Cover in bipartite graph Let's say the maximum matching is M.

And orientation of edges will be : Those edges that belong to M will go from right to left, all other edges will go from left to right. Now run DFS starting at all left vertices that are not incident to edges in M. Some vertices of the graph will become visited during this DFS and some not-visited. To get minimum vertex cover you take all visited right vertices of M, and all not-visited left vertices of M.

Minimum Path cover in a Directed Acyclic Graph : find the minimum number of vertex-disjoint paths to cover each vertex in V. make a bipartite graph G⁵ = {Vout X Vin, E } Vout = nodes having positive out-degree Vin = nodes have positive in-degree G' has a matching of size m iff there exists n-m vertexdisjoint paths that cover each vertex in G.

In a DAG delete some edges so that each nodes indegree=outdegree=k and maximise the sum of the deleted edges cost:

k. cost 0 From each node j (n<j<=n+n) to Sink, add edge of capacity compute the Grundy number. O(m + n). */ For all i,j (1<=i,j<=n) If in the main graph from i to j there is an edge, then add from each node i to n+j of capacity 1 and cost = original edges cost. Then run the min cost max flow algorithm.if max flow = n*k , then

Given a graph where each node has value and each edges cost = value[y] xor value[x]. Initially some nodes values are given and we have to find other nodes value such that sum of all edge is minimum.

We make LOG(MAX_VAL) graphs for each bit. For ith bit, for each known node u,

add edge with capacity INF from source to u whose i bit is set. otherwise add edge with capacity INF from node to sink. add the regular edge with capacity 1. Then after finding the MinCut , let U are those node which are in source's component. Then their bit will be set.

find a subgraph such that cost(edge)-cost(node) is maximum Consider each EDGE & NODE as node. Add edge from Source to each EDGE of capacity w & from each EDGE to both nodes of capacity INF & from each node

Finding an euler circuit in a graph G with both directed and undirected edges where undirected edges can be used

Euler circuit exists iff every vettex's indegree=outdegree

make a bipartite graph. One partition contains all edges and other contains all nodes. For each directed edge, connect this edge to its head endpoint. For each undirected edge, connect this edge to its both endpoints Now from Source node, add edges to each G's edge. all of these edges' capacity will be 1. Now from each G's node, add edges to sink. Edge i's capacity will be total degree/2 of node i. Then check the flow and so on.

Find a subgraph where |edge|/|node] is maximum. The Knig-Egevry theorem (Knig's theorem), asserts that make a bipartite graph. One partition contains all edges and other contains all nodes. For each G's edge, connect this edge to its both endpoints } Now from Source node , add edges to each G's edge. all of these edges' capacity will be 1. Now from each G's node, add edges to sink. Edge i's capacity will be X.

> we do binary search on X's value. if(Flow < Number of Edges) left = mid otherwise right = mid.

/* tree case: g[u] = for all v : XOR[g[v] + 1]

Games

HackenBush

```
lose if no moves available
                                                          1. Colon Principle: Grundy number of a tree is the
                                                                   xor of Grundy number of child subtrees.
                                                          2. Fusion Principle: Consider a pair of adjacent
                                                              vertices u, v that has another path (i.e.,
                                                              they are in a cycle). Then, we can contract u
                                                              and v without changing Grundy number.
                                                          We first decompose graph into two-edge connected
                                                          components. Then, by contracting each components by
                                                          using Fusion Principle, we obtain a tree (and many
From source to each node i (1<=i<=n), add edge of capacityself loops) that has the same Grundy number to the
                                                          original graph. By using Colon Principle, we can
                                                           struct hackenbush {
                                                            int n; vector<vector<int>> adj;
                                                            hackenbush(int n) : n(n), adj(n) { }
                                                            void add_edge(int u, int v) {
                                                              adj[u].push_back(v);
```

```
if(u!=v) adj[v].push_back(u);
int grundy(int r) {
 vector<int> num(n), low(n); int t = 0;
 function<int(int, int)> dfs=[&](int p,int u) {
   num[u] = low[u] = ++t; int ans = 0;
   for (int v : adj[u]) {
     if (v == p) { p += 2 * n; continue; }
     if (num[v] == 0) {
       int res = dfs(u, v);
low[u] = min(low[u], low[v]);
if (low[v] > num[u]) ans ^= (1+res)^1;
        else ans ^= res;
                                   // non bridge
     else low[u] = min(low[u], num[v]);
   if (p > n) p = 2 * n;
   for (int v : adj[u])
     if (v != p && num[u] <= num[v]) ans ^= 1;</pre>
   return ans:
 return dfs(-1, r);
```

Geo

Convex Hull

```
struct pt {double x, y;};
bool cmp(pt a, pt b)
   return a.x < b.x \mid \mid (a.x == b.x \&\& a.y < b.y);
bool cw(pt a, pt b, pt c) {
return a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y)<0;
bool ccw(pt a, pt b, pt c) {
 return a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y)>0;
vector<pt> a;
vector<pair<double,pair<double,double> > > pp;
void convex_hull(vector<pt>& a) {
   if (a.size() == 1)
       return;
   sort(a.begin(), a.end(), &cmp);
   pt p1 = a[0], p2 = a.back();
   vector<pt> up, down;
   up.push_back(p1);
   down.push_back(p1);
   for (int i = 1; i < (int)a.size(); i++) {</pre>
       if (i == a.size() - 1 || cw(p1, a[i], p2)) {
           while (up.size() >= 2 &&
       !cw(up[up.size()-2], up[up.size()-1], a[i]))
              up.pop_back();
           up.push_back(a[i]);
       if (i == a.size() - 1 || ccw(p1, a[i], p2)) {
           while(down.size() >= 2 &&!
ccw(down[down.size()-2], down[down.size()-1], a[i]))
              down.pop_back();
           down.push_back(a[i]);
   a.clear();
for(int i=0;i<(int)up.size();i++)a.push_back(up[i]);
for(int i=down.size()-2;i>0;i--)a.push_back(down[i]);
```

Half Plane Intersection

```
#define MAX 200005
#define MOD 1009
#define SMOD 998244353
#define ROOT 318
#define GMAX 19
#define EPS 0.000000001
#define NIL 0
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
class HalfPlaneIntersection{
    static double eps, inf;
public:
    struct Point{
       double x, y;
explicit Point(double x=0,double y=0):x(x), y(y) {}
friend Point operator+(const Point&p,const Point&q){
    return Point(p.x + q.x, p.y + q.y);
friend Point operator-(const Point&p,const Point&q){
    return Point(p.x - q.x, p.y - q.y);
friend Point operator*(const Point&p,double&k){
    return Point(p.x * k, p.y * k);
friend double cross(const Point& p, const Point& q){
   return p.x * q.y - p.y * q.x;
};
// Basic half-plane struct.
    struct Halfplane
    {// 'p' is a passing point of the line and
   //'pq' is the direction vector of the line.
       Point p, pq; double angle;
Halfplane() {}
       Halfplane(const Point&a, const Point&b):p(a)
        ,pq(\bar{b}-a){}
           angle = atan21(pq.y, pq.x);
// Check if point 'r' is outside this half-plane.
// Every half-plane allows the region to the LEFT
// of its line.
       bool out(const Point& r){
           return cross(pq, r - p) < -eps;</pre>
//Comparator for sorting. If the angle of both half-
//planes is equal, the leftmost one should go first.
       bool operator < (const Halfplane& e) const{</pre>
           if (fabsl(angle - e.angle) < eps)</pre>
               return cross(pq, e.p - p) < 0;
           return angle < e.angle;
       }// We use equal comparator for std::unique
       //to easily remove parallel half-planes.
       bool operator == (const Halfplane& e) const{
           return fabsl(angle - e.angle) < eps;</pre>
       }// Intersection point of the lines of two
//half-planes. It is assumed they're never parallel.
       friend Point inter(const Halfplane& s,
                         const Halfplane& t){
double alpha=cross((t.p-s.p),t.pq)/cross(s.pq,t.pq);
return s.p + (s.pq * alpha);
vector<Point>hp_intersect(vector<Halfplane>&H)
       Point box[4] = // Bounding box in CCW order
           Point(inf, inf),
```

```
Point(-inf. inf).
           Point(-inf, -inf),
           Point(inf, -inf)
for(int i=0;i<4;i++){//Add bounding box half-planes.</pre>
           Halfplane aux(box[i], box[(i+1) \% 4]);
           H.push_back(aux);
       // Sort and remove duplicates
       sort(H.begin(), H.end());
       H.erase(unique(H.begin(), H.end()), H.end());
       deque < Halfplane > dq;
       int len = 0;
       for(int i = 0; i < int(H.size()); i++){</pre>
// Remove from the back of the deque while last
// half-plane is redundant
while(len>1&&H[i].out(inter(dq[len-1], dq[len-2]))){
               dq.pop_back();
--len:
           }// Remove from the front of the deque
       //while first half-plane is redundant
    while (len > 1 && H[i].out(inter(dq[0],dq[1]))){
               dq.pop_front();
               --len;
           } // Add new half-plane
           dq.push_back(H[i]); ++len;
       } // Final cleanup: Check half-planes at the
       //front against the back and vice-versa
while (len > 2 \& dq [0] \cdot out(inter(dq [len-1], dq [len-2])))
           dq.pop_back(); --len;
while(len > 2 && dq[len-1].out(inter(dq[0],dq[1]))){
           dq.pop_front(); --len;
       // Report empty intersection if necessary
       if (len < 3) return vector<Point>();
// Reconstruct the convex polygon from
//the remaining half-planes.
       vector<Point> ret(len);
       for(int i = 0; i+1 < len; i++){</pre>
           ret[i] = inter(dq[i], dq[i+1]);
       ret.back() = inter(dg[len-1], dg[0]);
       return ret;
double HalfPlaneIntersection::eps=1e-9;
double HalfPlaneIntersection::inf=1e9;
    vector<HalfPlaneIntersection::Halfplane> V;
    vector<HalfPlaneIntersection::Point> P;
    for(i=0; i<n; i++){
       int c;
       scanf("%d",&c);
       HalfPlaneIntersection::Halfplane h;
       HalfPlaneIntersection::Point p:
       for(j=0; j<c; j++){
    scanf("%lf %lf",&p.x,&p.y);</pre>
           P.push_back(p);
       for(j=0; j<c; j++){</pre>
h=HalfPlaneIntersection::Halfplane(P[j],P[(j+1)%c]);
           V.push_back(h);
       P.clear();
    P=HalfPlaneIntersection::hp_intersect(V);
    double ans=0;
    n=P.size():
    for(i=0; i<n; i++){</pre>
```

```
ans=ans+P[i].x*P[(i+1)%n].y-P[i].y*P[(i+1)%n].x;
}
ans=ans/2;
```

5.3 Line Segment Intersection

```
struct pt {
   double x, y;
   bool operator<(const pt& p) const</pre>
 return x < p.x-EPS \mid | (abs(x-p.x) < EPS & y < p.y-EPS);
struct line {
   double a, b, c;
line() {}
   line(pt p, pt q){
    a = p.y - q.y;
    b = q.x - p.x;
       c = -a * p.x - b * p.y;
       norm();
   void norm(){
       double z = sqrt(a * a + b * b);
       if (abs(z) > EPS) a /= z, b /= z, c /= z;
   double dist(pt p){ return a * p.x + b * p.y + c;}
double det(double a, double b, double c, double d){
   return a * d - b * c;
inline bool betw(double 1, double r, double x){
   return min(1, r) \leq x + EPS && x \leq max(1,r)+EPS;
bool intersect_1d(double a, double b, dbl c, dbl d)
   if (a > b)
       swap(a, b);
   if (c > d)
       swap(c, d);
   return max(a, c) <= min(b, d) + EPS;</pre>
}bool
intersect(pt a,pt b,pt c,pt d,pt& left,pt& right){
   if (!intersect_1d(a.x, b.x, c.x, d.x) ||
     !intersect_1d(a.y, b.y, c.y, d.y))return false;
   line m(a, b); line n(c, d);
   double zn = det(m.a, m.b, n.a, n.b);
   if (abs(zn) < EPS) {
   if (abs(m.dist(c)) > EPS || abs(n.dist(a)) > EPS)
           return false;
       if (b < a) swap(a, b);
       if (d < c) swap(c, d);
       left = max(a, c); right = min(b, d);
       return true;
left.x = right.x = -det(m.c, m.b, n.c, n.b) / zn;
left.y = right.y = -det(m.a, m.c, n.a, n.c) / zn;
return betw(a.x,b.x,left.x)&&betw(a.y,b.y,left.y) &&
   betw(c.x, d.x, left.x) && betw(c.y, d.y, left.y);
```

5.4 Minimum Perimeter Triangle

```
#include <ext/pb_ds/detail/standard_policies.hpp>
struct pt{
   double x, y;
   int id;
struct cmp_x{
   bool operator()(const pt & a, const pt &b)const{
       return a.x < b.x \mid \mid (a.x == b.x \&\& a.y < b.y);
};
struct cmp_y{
   bool operator()(const pt& a, const pt & b)const{
       return a.y < b.y;</pre>
int n; vector<pt> a; double mindist;
pair<int, pair<int, int> > best_pair;
void upd_ans(const pt & a, const pt &b,const pt&c){
   double distC = sqrt((a.x-b.x)*(a.x-b.x)
             +(a.y-b.y)*(a.y-b.y));
   double distA = sqrt((c.x - b.x)*(c.x - b.x)
             + (c.v - b.v)*(c.v - b.v));
   double distB = sqrt((a.x - c.x)*(a.x - c.x)
             + (a.y - c.y)*(a.y - c.y));
   if (distA + distB + distC < mindist){</pre>
   mindist = distA + distB + distC;
   best_pair=make_pair(a.id,make_pair(b.id,c.id));
vector<pt> t;
void rec(int 1, int r){
   if (r - 1 \le 3 \&\&r - 1 >= 2){
       for (int i = 1; i < r; ++i){
           for (int j = i + 1; j < r; ++j){
              for(int k=j+1;k<r;k++){</pre>
                  upd_ans(a[i],a[j],a[k]);
       sort(a.begin() + 1, a.begin() + r, cmp_v());
       return;
   int m = (1 + r) >> 1; int midx = a[m].x;
   rec(1, m); rec(m, r);
   merge(a.begin() + 1, a.begin() + m, a.begin()
   + m, a.begin() + r, t.begin(), cmp_y());
   copy(t.begin(), t.begin() + r - 1, a.begin()+1);
   int tsz = 0;
   for (int i = 1; i < r; ++i){
       if (abs(a[i].x - midx) < mindist/2){</pre>
           for (int j = tsz - 1; j >= 0
           && a[i].y - t[j].y < mindist/2; --j)
              if(i+1<r) upd_ans(a[i],a[i+1],t[j]);
               if(j>0) upd_ans(a[i], t[j-1], t[j]);
           t[tsz++] = a[i];
       }
   }
```

5.5 Minkowski

```
#define MAX 300005
#define BEGIN 1
#define CHAINS 18
#define NOT_VISITED 0
#define VISITING 1
#define VISITED 2
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
```

```
struct pt{
   long long x, y;
   pt() {}
   pt(long long _x, long long _y):x(_x), y(_y) {}
   pt operator+(const pt & p) const{
       return pt(x + p.x, y + p.y);
   pt operator-(const pt & p) const{
       return pt(x - p.x, y - p.y);
    long long cross(const pt & p) const{
       return x * p.y - y * p.x;
   long long dot(const pt & p) const{
       return x * p.x + y * p.y;
    long long cross(const pt & a, const pt&b)const{
       return (a - *this).cross(b - *this);
    long long dot(const pt & a, const pt & b) const{
       return (a - *this).dot(b - *this);
   long long sqrLen() const{
       return this->dot(*this);
class pointLocationInPolygon{
   bool lexComp(const pt & 1, const pt & r){
       return 1.x < r.x \mid | (1.x == r.x && 1.y < r.y);
   int sgn(long long val){
       return val > 0 ? 1 : (val == 0 ? 0 : -1);
   vector<pt> seq; int n; pt translate;
   bool pointInTriangle(pt a, pt b, pt c,pt point){
       long long s1 = abs(a.cross(b, c));
       long long s2 = abs(point.cross(a, b)) +
   abs(point.cross(b, c)) + abs(point.cross(c, a));
       \overline{return} s1 == s2;
public:
   pointLocationInPolygon(){}
   pointLocationInPolygon(vector<pt> & points){
       prepare(points);
   void prepare(vector<pt> & points){
       seq.clear();
       n = points.size();
       int pos = 0;
       for(int i = 1; i < n; i++){
           if(lexComp(points[i], points[pos]))
              pos = i;
       translate.x=points[pos].x;
       translate.y=points[pos].y;
       rotate(points.begin(), points.begin() +
        pos, points.end());
       seq.resize(n);
       for(int i = 0; i < n; i++)</pre>
           seq[i] = points[i + 1] - points[0];
   bool pointInConvexPolygon(pt point){
       point.x-=translate.x;
       point.y-=translate.y;
       if(seq[0].cross(point) != 0 && sgn(seq[0].
       cross(point))!=sgn(seq[0].cross(seq[n-1])))
           return false:
       if(seq[n-1].cross(point)!=0&&sgn(seq[n-1]
```

```
.cross(point))!=sgn(seq[n-1].cross(seq[0])))
           return false;
       if(seq[0].cross(point) == 0)
           return seq[0].sqrLen()>= point.sqrLen();
       int 1 = 0, r = n - 1;
       while (r - 1 > 1)
           int mid = (1 + r)/2; int pos = mid;
           if(seq[pos].cross(point) >= 0)1 = mid;
           else r = mid;
       int pos = 1;
       return pointInTriangle(seq[pos],
       seq[pos+1],pt(0,0),point);
    pointLocationInPolygon(){
       seq.clear();
class Minkowski{
   static void reorder_polygon(vector<pt> & P){
       size_t pos = 0;
       for(size_t i = 1; i < P.size(); i++){</pre>
           if(P[i].y < P[pos].y | |
       (P[i].y == P[pos].y \&\& P[i].x < P[pos].x))
              pos = i;
       rotate(P.begin(), P.begin() + pos, P.end());
public:
   vector<pt> minkowski(vector<pt> P,vector<pt>Q){
       // the first vertex must be the lowest
       reorder_polygon(P);
       reorder_polygon(Q);
       // we must ensure cyclic indexing
       P.push back(P[0]):
       P.push_back(P[1]);
       Q.push_back(Q[0]);
       Q.push_back(Q[1]);
       // main part
       vector<pt> result;
       size_t i = 0, j = 0;
       while(i < P.size() - 2 || j < Q.size() - 2){</pre>
           result.push_back(P[i] + Q[j]);
auto cross = (P[i + 1] - P[i]).cross(Q[j+1]-Q[j]);
           if(cross >= 0) ++i;
           if(cross <= 0) ++j;
       return result;
   }
```

5.6 Pair of Intersecting Segments

```
#define MAX 100009
#define MAX_NODES 100005
struct pt {
    double x, y;
};
struct seg {
    pt p, q; int id;
    double get_y(double x) const {
        if (abs(p.x - q.x) < EPS)
            return p.y;
        return p.y+(q.y-p.y)*(x-p.x)/(q.x - p.x);
};
bool intersect1d(double 11, dbl r1, dbl l2, dbl r2){
    if (11 > r1) swap(11, r1);
    if (12 > r2) swap(12, r2);
    return max(11, 12) <= min(r1, r2) + EPS;</pre>
```

```
int vec(const pt& a, const pt& b, const pt& c) {
 double s = (b.x-a.x)*(c.y-a.y)-(b.y-a.y)*(c.x-a.x);
   return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
bool intersect(const seg& a, const seg& b){
   return intersect1d(a.p.x,a.q.x,b.p.x,b.q.x) &&
          intersect1d(a.p.y,a.q.y,b.p.y,b.q.y) &&
          vec(a.p,a.q,b.p)*vec(a.p,a.q,b.q) <= 0 &&
          vec(b.p, b.q, a.p)*vec(b.p,b.q,a.q) <= 0;
bool operator<(const seg& a, const seg& b){
   double x=\max(\min(a.p.x,a.q.x),\min(b.p.x,b.q.x));
   return a.get_y(x) < b.get_y(x) - EPS;</pre>
struct event {
   double x; int tp, id;
event() {}
event(double x, int tp, int id):x(x),tp(tp),id(id){}
   bool operator<(const event& e) const {
       if (abs(x - e.x) > EPS) return x < e.x;
return tp > e.tp;
set<seg> s; vector<set<seg>::iterator> where;
set<seg>::iterator prev(set<seg>::iterator it) {
   return it == s.begin() ? s.end() : --it;
set<seg>::iterator next(set<seg>::iterator it) {
   return ++it;
pair<int, int> solve(const vector<seg>& a) {
   int n = (int)a.size(); vector<event> e;
   for (int i = 0; i < n; ++i) {
   e.push_back(event(min(a[i].p.x,a[i].q.x),+1,i));
   e.push_back(event(max(a[i].p.x,a[i].q.x),-1,i));
   sort(e.begin(), e.end()); s.clear();
   where.resize(a.size());
   for (size_t i = 0; i < e.size(); ++i) {</pre>
       int id = e[i].id;
       if (e[i].tp == +1) {
           set<seg>::iterator nxt=
           s.lower_bound(a[id]),prv=prev(nxt);
           if(nxt!=s.end()&&intersect(*nxt,a[id]))
              return make_pair(nxt->id, id);
           if(prv!=s.end()&&intersect(*prv,a[id]))
              return make_pair(prv->id, id);
           where[id] = s.insert(nxt, a[id]);
           set<seg>::iterator nxt = next(where[id])
           , prv = prev(where[id]);
           if (nxt != s.end() && prv != s.end() &&
           intersect(*nxt, *prv))
              return make_pair(prv->id, nxt->id);
           s.erase(where[id]);
   return make_pair(-1, -1);
```

5.7 Vertical Decomposition

```
#define MAX 300005
#define MDD 1000000007
#define GMAX 19
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/detail/standard_policies.hpp>
```

```
typedef double dbl:
const dbl eps = 1e-9;
inline bool eq(dbl x, dbl y){
    return fabs(x - y) < eps;</pre>
inline bool lt(dbl x, dbl y){
    return x < y - eps;
inline bool gt(dbl x, dbl y){
    return x > y + eps;
inline bool le(dbl x, dbl y){
   return x < y + eps;
inline bool ge(dbl x, dbl y){
  return x > y - eps;
struct pt{
    dbl x, y;
    inline pt operator - (const pt & p)const{
       return pt\{x - p.x, y - p.y\};
    inline pt operator + (const pt & p)const{
       return pt\{x + p.x, y + p.y\};
    inline pt operator * (dbl a)const{
       return pt\{x * a, v * a\};
    inline dbl cross(const pt & p)const{
  return x * p.y - y * p.x;
    inline dbl dot(const pt & p)const{
        return x * p.x + \bar{y} * p.\bar{y};
    inline bool operator == (const pt & p)const{
        return eq(x, p.x) && eq(y, p.y);
struct Line{
   pt p[2];
Line(){} Line(pt a, pt b):p{a, b}{}
    pt vec()const{
       return p[1] - p[0];
    pt& operator [](size_t i){
       return p[i];
inline bool lexComp(const pt & 1, const pt & r){
    if(fabs(1.x - r.x) > eps){
       return 1.x < r.x;</pre>
    else return l.y < r.y;</pre>
vector<pt> interSegSeg(Line 11, Line 12){
    if(eq(11.vec().cross(12.vec()), 0)){
        if(!eq(11.vec().cross(12[0] - 11[0]), 0))
            return {};
        if(!lexComp(l1[0], l1[1]))
            swap(11[0], 11[1]);
        if(!lexComp(12[0], 12[1]))
            swap(1\bar{2}[0], 12[1]);
        pt 1 = lexComp(11[0], 12[0]) ? 12[0] : 11[0];
        pt r = lexComp(11[1], 12[1]) ? 11[1] : 12[1];
        if(1 == r)
           return {1};
  return lexComp(l,r)?vector<pt>{1,r}:vector<pt>();
```

```
dbl s = (12[0] - 11[0]).cross(12.vec()) /
       11.vec().cross(12.vec());
       pt inter = 11[0] + 11.vec() * s;
       if(ge(s, 0) && le(s, 1) && le((12[0] -
       inter).dot(12[1] - inter), 0))
           return {inter};
       else
           return {};
char get_segtype(Line segment,pt other_point){
   if(eq(segment[0].x, segment[1].x))
       return 0;
   if(!lexComp(segment[0], segment[1]))
       swap(segment[0], segment[1]);
   return (segment[1]-segment[0]).cross(other_point
    - segment [0] > 0 ? \bar{1} : -1:
dbl union_area(vector<tuple<pt, pt,pt> > triangles){
   vector<Line> segments(3 * triangles.size());
   vector<char> segtype(segments.size());
   for(size_t i = 0; i < triangles.size(); i++){</pre>
       pt a, b, c; tie(a, b, c) = triangles[i];
 segments[3*i]=lexComp(a,b)?Line(a,b):Line(b,a);
 segtype[3 * i] = get_segtype(segments[3 * i], c);
 segments[3*i+1]=lexComp(b,c)?Line(b,c):Line(c,b);
 segtype[3*i+1] = get_segtype(segments[3*i+1],a);
 segments[3*i+2]=lexComp(c,a)?Line(c,a):Line(a,c);
 segtype[3*i+2] = get_segtype(segments[3*i+2],b);
 vector<dbl>k(segments.size()),b(segments.size());
   for(size_t i = 0; i < segments.size(); i++){</pre>
       if(segtype[i]){
           k[i]=(segments[i][1].y-segments[i][0].y)
            /(segments[i][1].x-segments[i][0].x);
       b[i]=segments[i][0].y-k[i]*segments[i][0].x;
   dbl ans = 0;
   for(size_t i = 0; i < segments.size(); i++){</pre>
       if(!segtype[i]) continue;
       dbl l = segments[i][0].x,r=segments[i][1].x;
       vector<pair<dbl, int> > evts;
       for(size_t j = 0; j < segments.size(); j++){</pre>
           if(!segtype[j] || i == j)
 dbl l1 = segments[j][0].x, r1 = segments[j][1].x;
           if(ge(l1, r) || ge(l, r1))
 dbl common_l = max(l, l1), common_r = min(r, r1);
 auto pts = interSegSeg(segments[i], segments[j]);
           if(pts.empty()){
              dbl yl1 = k[j] * common_l + b[j];
              dbl yl = k[i] * common_l + b[i];
            if(lt(vl1, vl) == (segtype[i] == 1)){
       int evt_type = -segtype[i] * segtype[j];
       evts.emplace_back(common_l, evt_type);
       evts.emplace_back(common_r, -evt_type);
           else if(pts.size() == 1u){
 dbl yl=k[i]*common_l+b[i],yl1=k[j]*common_l+b[j];
 int evt_type = -segtype[i] * segtype[j];
 if(lt(y11, y1) == (segtype[i] == 1)){
     evts.emplace_back(common_l, evt_type);
     evts.emplace_back(pts[0].x, -evt_type);
```

```
yl =k[i]*common_r+b[i],yl1=k[j]*common_r+b[i];
  if(lt(yl1, yl) == (segtype[i] == 1)){
     evts.emplace_back(pts[0].x, evt_type);
     evts.emplace_back(common_r, -evt_type);
           else{
               if(segtype[j] != segtype[i] || j>i){
                   evts.emplace_back(common_1, -2);
                   evts.emplace_back(common_r, 2);
       evts.emplace_back(1, 0);
       sort(evts.begin(), evts.end());
size_t j = 0; int balance = 0;
       while(j < evts.size()){</pre>
           siže_t ptr = j;
           while(ptr < evts.size() &&</pre>
           eq(evts[j].first, evts[ptr].first)){
               balance += evts[ptr].second;
           if(!balance && !eq(evts[j].first, r)){
dbl next_x = ptr == evts.size() ?r:evts[ptr].first;
               ans -= segtype[i] * (k[i] * (next_x
+evts[j].first)+2*b[i])*(next_x-evts[j].first);
           j = ptr;
   return ans/2;
```

area of simple polygon

```
double area(const vector<point>& fig) {
   double res = 0;
   for (unsigned i = 0; i < fig.size(); i++) {</pre>
       point p = i ? fig[i - 1] : fig.back();
       point q = fig[i];
       res += (p.x - q.x) * (p.y + q.y);
   return fabs(res) / 2;
```

circle line intersection

```
double r, a, b, c; // given as input
double x0 = -a*c/(a*a+b*b), y0 = -b*c/(a*a+b*b);
if (c*c > r*r*(a*a+b*b)+EPS)
   puts ("no points");
else if (abs (c*c - r*r*(a*a+b*b)) < EPS) {
   puts ("1 point");
   cout << x0 << ', ' << y0 << '\n';
else {
   double d = r*r - c*c/(a*a+b*b);
   double mult = sqrt (d / (a*a+b*b)):
   double ax, ay, bx, by;
   ax = x0 + b * mult;
   bx = x0 - b * mult;
   ay = y0 - a * mult;
   bv = v0 + a * mult;
   puts ("2 points");
   cout<<ax<<' '<< ay << '\n' << bx << ' ' << by;
```

circle union area-hellbent

```
struct Point {
 LD x, y ;
 LD operator*(const Point &a)const {
   return x*a.y-y*a.x;}
 LD operator/(const Point &a)const {
   return sqrt((a.x-x)*(a.x-x)+(a.y-y)*(a.y-y));}
}po[N];
LD r[N]
int sgn(LD x) {return fabs(x) < EPS?0:(x > 0.0?1:-1);}
pair<LD,bool> ARG[2*N]
LD cir_union(Point c[],LD r[],int n) {
   LD sum = 0.0 , sum1 = 0.0 ,d,p1,p2,p3 ;
 for(int i = 0 ; i < n ; i++) {
    bool f = 1;
    for(int j = 0; f \& k j < n : j++)
     if(i!=j) \&\& sgn(r[j]-r[i]-c[i]/c[j])!=-1)f=0;
    if(!f) swap(r[i],r[--n]),swap(c[i--],c[n]);
 for(int i = 0; i < n; i++) {</pre>
      int k = 0, cnt = 0;
      for(int j = 0; j < n; j++) {
       if(i!=j&&sgn((d=c[i]/c[j])-r[i]-r[j])<=0){
          p3=acos((r[i]*r[i]+d*d-r[j]*r[j])/
                                    (2.0*r[i]*d));
          p2=atan2(c[j].y-c[i].y,c[j].x-c[i].x);
          p1 = p2-p3; p2' = p2+p3;
          if(sgn(p1+PI)==-1) p1+=2*PI,cnt++;
          if(sgn(p2-PI)==1) p2-=2*PI,cnt++;
          ARG[k++] = make_pair(p1,0);
          ARG[k++] = make_pair(p2,1);
      if(k) {
        sort(ARG,ARG+k) ;
       p1 = ARG[k-1].first-2*PI;
        p3 = r[i]*r[i];
       for(int j = 0 ; j < k ; j++) {
   p2 = ARG[j].first;</pre>
          if(cnt==0) {
           sum+=(p2-p1-sin(p2-p1))*p3 ;
sum1+=(c[i]+Point(cos(p1),sin(p1))*
                    r[i])*(c[i]+
                    Point(cos(p2),sin(p2))*r[i]);
          p1 = p2;
          ARG[j].second ? cnt--:cnt++;
      else sum += 2*PI*r[i]*r[i];
 return (sum+fabs(sum1))*0.5;
```

5.11common tangent

```
struct pt {
   double x, y;
   pt operator- (pt p) {
       pt res = {x-p.x, y-p.y};
       return res;
struct circle : pt {
   double r;
struct line {
   double a, b, c;
const double EPS = 1E-9;
double sqr (double a) {
```

```
return a * a:
void tangents (pt c, double r1, double r2,
         vector<line> & ans) {
   double r = r2 - r1;
   double z = sqr(c.x) + sqr(c.y);
   double d = z - sqr(r);
   if (d < -EPS) return;</pre>
   d = sqrt (abs (d));
   line \bar{1}; l.a = (c.x * r + c.y * d) / z;
   1.b = (c.y * r - c.x * d) / z; 1.c = r1;
   ans.push_back (1);
vector<line> tangents (circle a, circle b) {
   vector<line> ans:
   for (int i=-1; i<=1; i+=2)
       for (int j=-1; j<=1; j+=2)
           tangents (b-a, a.r*i, b.r*j, ans);
   for (size_t i=0; i<ans.size(); ++i)</pre>
       ans[i].c = ans[i].a * a.x + ans[i].b * a.y;
```

5.12notes

Picks Theorem: S = I + B / 2 - 1, I = internal point, B = border point, S =area

5.13 radial sweep - anikda unga bunga

```
struct Point{
   ll x; ll y; ll val;
   bool operator < (const Point& other) const {</pre>
       if(x * other.x > 0) return y*other.x < x*other.y;</pre>
       else return y*other.x > x*other.y;
Point p[MAX];
int main(){
    int n;
   scanf("%d",&n);
ll Plus=0, Minus=0;
    for(int i=1;i<=n;i++) cin>>p[i].x>>p[i].y>>p[i].val;
    //Precalculate w.r.t. Y-axis
    for(int i=1;i<=n;i++){</pre>
       if(p[i].x>=0) Plus+=p[i].val;
       else Minus+=p[i].val;
    sort(p+1,p+n+1);
    //Calculate w.r.t. all slopes
   for(int i=1;i<=n;i++){</pre>
       if(p[i].x>=0) Plus-=p[i].val, Minus+=p[i].val;
       else Minus-=p[i].val, Plus+=p[i].val;
   }
```

Graph

6.1 Articulation Vertex

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, low;
int timer;
void dfs(int v, int p = -1) {
   visited[v] = true;
   tin[v] = low[v] = timer++;
   int children=0;
   for (int to : adj[v]) {
```

```
if (to == p) continue:
       if (visited[to]) {
           low[v] = min(low[v], tin[to]);
           dfs(to, v);
           low[v] = min(low[v], low[to]);
           if (low[to] >= tin[v] && p!=-1)
              IS_CUTPOINT(v);
           ++children;
       }
   if(p == -1 \&\& children > 1)
       IS_CUTPOINT(v);
void find_cutpoints() {
   timer = 0:
   visited.assign(n, false);
   tin.assign(n, -1);
   low.assign(n, -1);
   for (int i = 0; i < n; ++i) {
       if (!visited[i])
           dfs (i);
```

6.2 Bridge offline

```
int n; // number of nodes
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> tin, low;
int timer;
void dfs(int v, int p = -1) {
   visited[v] = true:
   tin[v] = low[v] = timer++;
   for (int to : adj[v]) {
       if (to == p) continue;
       if (visited[to]) {
           low[v] = min(low[v], tin[to]);
       } else {
           dfs(to, v);
           low[v] = min(low[v], low[to]);
           if (low[to] > tin[v])
              IS_BRIDGE(v, to);
       }
   }
void find_bridges() {
   timer = 0;
   visited.assign(n, false);
   tin.assign(n, -1);
   low.assign(n, -1);
   for (int i = 0; i < n; ++i) {
       if (!visited[i])
           dfs(i):
}
```

6.3 Eular Path-hellbent

```
int c[maxn] , d[maxn] ;
map< int , multiset<int> > g ; map <int,int> vis ;
void dfs1(int u){
    vis[u] = 1 ;
    for( auto v : g[u] )
        if( vis.find(v) == vis.end() ) dfs1(v) ;
}
//just call dfs2 with the node you want to start
//your path at first you need to make sure,
//the graph is connected and euler path exists
```

```
vector <int> ans :
void dfs2(int u){
 while( (int)g[u].size() !=0 ){
   int_v = *g[u].begin()
   g[u].erase(g[u].find(v));
   g[v].erase(g[v].find(u)); dfs2(v);
 ans.pb(u);
int main(){
 int n; scanf("%d",&n);
 for(int i=1; i<n; i++) scanf("%d",&c[i]);</pre>
 for(int i=1; i<n; i++) scanf("%d",&d[i]);</pre>
 for(int i=1 ; i<n ; i++) {</pre>
   if( c[i] > d[i] ){
     printf("-1\n"); return 0 ;
   g[c[i]].insert(d[i]); g[d[i]].insert(c[i]);
 int src = c[1] , cnt = 0 ;
 for( auto it : g ){
   if( (int)it.second.size() & 1 ){
     cnt++ ; src = it.first ;
 dfs1( src )
 if(vis.size()!=g.size()||(cnt!=0&&cnt!=2)){
   printf("-1\n");
   return 0;
 //call for printing euler path
 dfs2(src);
 for(int i=0 ; i<ans.size() ; i++){</pre>
   printf("%d",ans[i]);
   if(i == (int)ans.size() - 1) printf("\n");
   else printf(" ");
```

6.4 Strongly Connected Components

```
vector<vector<int>> adj, adj_rev;
vector<bool> used:
vector<int> order, component;
void dfs1(int v) {
   used[v] = true;
   for (auto u : adj[v])
       if (!used[u])
           dfs1(u);
   order.push_back(v);
void dfs2(int v) {
   used[v] = true;
   component.push_back(v);
   for (auto u : adj_rev[v])
       if (!used[u])
           dfs2(u):
int main() {
   int n;
   // ... read n ...
   for (;;) {
       int a, b;
       // ... read next directed edge (a,b) ...
       adj[a].push_back(b);
```

```
adj_rev[b].push_back(a);
}
used.assign(n, false);
for (int i = 0; i < n; i++)
    if (!used[i])
        dfs1(i);
used.assign(n, false);
reverse(order.begin(), order.end());
for (auto v : order)
    if (!used[v]) {
        dfs2 (v);

        // ... processing next component ...
        component.clear();
}</pre>
```

|6.5> eular formula

If a finite, connected, planar graph is drawn in the plane without any edge intersections, and v is the number of vertices, e is the number of edges and f is the number of faces (regions bounded by edges, including the outer, infinitely large region), then v - e + f = 2

$|6.6 \quad \text{scc} + 2\text{sat}|$

```
/*at first take a graph of size 2*n(for each vari
able two nodes). for each clause of type (a or b)
add two diredge !a-->b and !b-->a. if both x_i and
!x_i is in same connected component for some i, then this equations are unsatisfiable . Otherwise
there is a solution. Assume, f is satisfiable. Now
we want to give values to each var in order to satisfy f. It can be done with a top sort of vertices of the graph we made. If !x_i is after
x_i in topological sort, x_i should be FALSE. It
should be TRUE otherwise. say we have equation
with three var x1, x2, x3.(x1 \text{ or } !x2) and (x2 \text{ or } x3)
= 1. so we addx1, x2, x3 and x4(as !x1), x5(!x2)
and x6(!x3). Add edge x4-->x2,x2-->x1, x5-->x3,
x6-->x2.
you need to pass array to the function findSCC
in which result will be returned every node will
be given a number, for nodes of a single connected component the number will be same this number
representing nodes willbe topsorted*/
class SCC{
public:
 vector<int> *g1, *g2; int maxNode, *vis1, *vis2;
  stack<int> st;
  SCC(int MaxNode){
    maxNode = MaxNode ; vis1 = new int[maxNode+2];
    vis2 = new int[maxNode+2] ;
    g1 = new vector<int>[maxNode+2]
    g2 = new vector<int>[maxNode+2] ;
  void addEdge(int u,int v){
    g1[u].push_back(v); g2[v].push_back(u);
  void dfs1(int u){
    if(vis1[u]==1) return; vis1[u]=1;
    for(int i=0;i<g1[u].size();i++)dfs1(g1[u][i]);</pre>
    st.push(u); return;
  void dfs2(int u, int cnt , int *ans){
    if(vis2[u]==1) return ; vis2[u] = 1 ;
```

```
for(int i=0:i<g2[u].size():i++)</pre>
     dfs2(g2[u][i],cnt,ans);
   ans[u] = cnt;
  int findSCC( int *ans ) {
   for(int i=1; i<=maxNode; i++) vis1[i] = 0;</pre>
   for(int i=1 ; i<=maxNode ; i++)</pre>
     if(vis1[i]==0) dfs1(i);
    int cnt = 0;
    for(int i=1; i<=maxNode; i++) vis2[i] = 0;</pre>
    while( !st.empty() ) {
     int u = st.top() ;
     if(vis2[u]==0) {++cnt; dfs2(u, cnt, ans);}
     st.pop();
    for(int i=1 ; i<=maxNode ; i++) {</pre>
     g1[i].clear(); g2[i].clear();
    delete vis1; delete vis2; return cnt;
};
```

Math

Crt

```
PII CRT(int x,int a,int y, int b) {
 int s, t;
 int d = extended_euclid(x, y, s, t);
 if (a % d != b % d) return make_pair(0, -1);
 return make_pair(mod(s*b*x+t*a*y,x*y)/d,x*y/d);
```

Discrete Root

```
#define MAX 100000
int prime[MAX+1],Phi[MAX+1];
void sieve(){
   int i,j;
for(i=2; i*i<=MAX; i++){</pre>
       if(prime[i]) continue;
       for(j=i; j*i<=MAX; j++){</pre>
           if(prime[i*j]==0) prime[i*j]=i;
void PhiWithSieve(){
    int i;
   for(i=2; i<=MAX; i++){</pre>
       if(prime[i]=0){
           Phi[i]=i-1;
       else if((i/prime[i])%prime[i]==0){
           Phi[i]=Phi[i/prime[i]]*prime[i];
       else{
           Phi[i]=Phi[i/prime[i]]*(prime[i]-1);
   }
int gcd(int a,int b){
    if(b==0) return a;
    else return gcd(b,a%b);
int powmod (int a, int b, int p) {
    int res = 1;
    while (b)
       if (b & 1)
           res = int (res * 111 * a % p), --b;
```

```
a = int (a * 111 * a % p), b >>= 1:
   return res;
int PrimitiveRoot(int p){
   vector<int>fact;
   int phi=Phi[p];
   int n=phi;
   while (n>1) {
       if(prime[n]==0){
           fact.push_back(n);
       else{
           int f=prime[n];
           while (n\%f==0) {
              n=n/f;
           fact.push_back(f);
       }
   int res;
   for(res=p-1; res>1; --res){
       for(n=0; n<fact.size(); n++){</pre>
           if(powmod(res,phi/fact[n],p)==1){
              break;
       if(n>=fact.size()) return res;
   return -1;
int DiscreteLog(int a, int b, int m) {
   a \% = m, b \% = m;
   int n = sqrt(m) + 1;
   map<int, int> vals;
   for (int p = 1; p \le n; ++p)
       vals[powmod(a, (int) (111* p * n) %m , m)]=p;
   for (int q = 0; q <= n; ++q) {
       int cur = (powmod(a, q, m) * 111 * b) % m;
       if (vals.count(cur)) +
           int ans = vals[cur] * n - q;
           return ans;
   return -1;
vector<int> DiscreteRoot(int n,int a,int k){
   int g = PrimitiveRoot(n);
   vector<int> ans;
   int any_ans = DiscreteLog(powmod(g,k,n),a,n);
   if (any_ans == -1){
       return ans;
   int delta = (n-1) / gcd(k, n-1);
   for(intcur=any_ans%delta;cur<n-1;cur+=delta)</pre>
       ans.push_back(powmod(g, cur, n));
   sort(ans.begin(), ans.end());
   return ans;
```

7.3 Fast Fourier Transform

```
#define MOD 1000000007
#define MAX 200005
#define PMAX 55
#define PRECISION 0.000001
#define INF 2000000000
using cd = complex<double>
const double PI = acos(-1);
void fft(vector<cd>& a, bool invert){
   int n = a.size();
```

```
for(int i = 1, j = 0; i < n; i++){
       int bit = n > 1:
       for(; j&bit; bit>>=1){
           j^=ੱbit;
       j ^= bit;
       if(i < j) swap(a[i], a[j]);</pre>
   for(int len = 2; len <= n; len <<= 1){</pre>
       double ang = 2*PI/len*(invert ? -1 : 1);
       cd wlen(cos(ang), sin(ang));
       for(int i = 0; i < n; i += len){</pre>
           cd w(1);
           for(int j = 0; j < len/2; j++){
               cd u = a[i+j], v = a[i+j+len/2]*w;
               a[i+i] = u+v:
               a[i+j+len/2] = u-v;
               w *= wlen;
   if(invert){
       for(cd &x: a) x /= n;
vector<int> multiply(vector<int>&a,vector<int>&b){
   vector<cd> fa(a.begin(), a.end());
   vector<cd> fb(b.begin(), b.end()); int n = 1;
   while(n < a.size()+b.size()) n <<= 1;
   fa.resize(n); fb.resize(n);
   fft(fa, false);
   fft(fb, false);
   for(int i = 0; i < n; i++)
       fa[i] *= fb[i];
   fft(fa, true);
   vector<int> result(n);
   for(int i = 0; i < n; i++)</pre>
       result[i] = round(fa[i].real());
   return result;
//Number Theoretic Transformation
ll gcd(ll a,ll b){
   if(b==0) return a;
   else return gcd(b,a%b);
ll egcd(ll a, ll b, ll & x, ll & y) {
   if (a == 0) {
       x = 0; y = 1;
       return b;
   ll x1, y1; ll d = egcd(b \% a, a, x1, y1);
   x = y1 - (b / a) * x1; y = x1;
   return d;
|ll ModuloInverse(ll a,ll n){
   11 x,y; x=gcd(a,n);
   a=a/x; n=n/x;
   ll res = egcd(a,n,x,y); x=(x^n+n)^n;
   return x;
const int mod = 998244353:
const int root = 15311432
const int root_1 = 469870224;
const int root_pw = 1 << 23;</pre>
void fft(vector<int> & a, bool invert) {
   int n = a.size();
   for (int i = 1, j = 0; i < n; i++) {
       int bit = n \gg 1:
       for (; j & bit; bit >>= 1)
    j ^= bit;
```

```
i ^= bit:
       if (i < j)
           swap(a[i], a[j]);
   for (int len = 2; len <= n; len <<= 1) {
       int wlen = invert ? root_1 : root;
       for (int i = len; i < root_pw; i <<= 1)</pre>
           wlen = (int)(1LL * wlen * wlen % mod);
       for (int i = 0: i < n: i += len) {
           int w = 1;
           for (int j = 0; j < len / 2; j++) {
int u=a[i+j], v=(int)(1LL*a[i+j+len/2]*w%mod);
a[i+j] = u + v < mod ? u + v : u + v - mod;
a[i+j+len/2] = u - v >= 0 ? u - v : u - v + mod;
w = (int)(1LL * w * wlen % mod);
   if (invert) {
       int n_1 = (int) ModuloInverse(n, mod);
       for (int & x : a)
           x = (int)(1LL * x * n_1 \% mod);
vector<int> multiply(vector<int>&a, vector<int> &b){
   vector<int> fa(a.begin(), a.end());
   vector<int> fb(b.begin(), b.end());
   int n = 1;
   while(n < a.size()+b.size())</pre>
       n <<= 1:
   fa.resize(n); fb.resize(n);
   fft(fa, false); fft(fb, false);
   for(int i = 0; i < n; i++)</pre>
       fa[i] = (int) (1LL*fa[i]*fb[i]%mod);
   fft(fa. true):
   vector<int> result(n);
   for(int i = 0; i < n; i++)</pre>
       result[i] = fa[i];
   return result:
```

7.4 Ncr anymod

```
bitset<MAXN+7>siv;
vector<ll>primes;
void Sieve(){
 for(11 i=2; i<=MAXN; i++){</pre>
   if(!siv.test(i)){
     primes.push_back(i);
     for(ll j=i+i; j<=MAXN; j+=i) siv.set(j);</pre>
void PrimePowerFactorization(ll x
                   vector<pair<11.11>>&factor){
  if(x<=MAXN*MAXN){</pre>
   for(int i=0; i<primes.size(); i++){</pre>
     11 cnt=0;
     while(x%primes[i]==0){
       cnt++;
       x/=primes[i];
    if(cnt!=0)factors.push_back({primes[i],cnt});
   if(x!=1) factors.push_back(make_pair(x,1));
  else{
     //Complexity O(Number of Primes Upto MAXN)
```

```
This is the weak form. Mods contains all ai.mi
pairs such that x=ai mod mi, all mi are co-prime
and product all mi<10^18.
Complexity O(number of mods*log of their Product)
11 CRT(vector<pair<11,11> >&mods,11 Product){
    ll ret=0:
   for(int i=0; i<mods.size(); i++){</pre>
       11 cur=mods[i].first;
       11 cof=Product/mods[i].second;
       11 invcof=InvMod(cof,mods[i].second);
       cur*=cof;
       cur%=Product;
//cur=SafeMul(cur,cof,Product);//handles overflow
       cur*=invcof;
       cur%=Product;
       //cur=SafeMul(cur,invcof,Product);
       ret+=cur;
       ret%=Product;
    return ret;
//calculates N!(ignoring P's POWERS) % mod.Here
//mod=P^M(P is a Prime).
//O(mod) preprocessing and O(logP N) query using
//Gauss Generalization of Wilson's Theorem and
//also returns maximum power of P that divides N!.
//fact[i] holds i!(ignoring P's MULTIPLES) % mod
//for all i upto P^M-1 helps in calculating NCR %
//MOD where MOD is small prime power. i.e. P<=10^6
FactorialModPrimePowerIgnoringPrimeFactors=FMPPIPF
pair<11,11>FMPPIPF(11 N,11 mod,11 P,11 M,11*fact){
 ll ret1=1; ll ret2=0;
  while(N){
     11 res=N\( mod; ret1*=fact[res]; ret1\( =mod; \)
     if(P!=2||M<3){}
         if((N/mod)%2) ret1*=-1;
     N/=P: ret2+=N:
  ret1=(ret1+mod)%mod;return make_pair(ret1,ret2);
//N<10^6 and mod is any PRIME. Complexity: O(N)
//preprocessing. O(1) query.
ll fact[MAXN+7]; ll inv[MAXN+7];
11 NCRforSmallNPrimeMod(ll N,ll R,ll mod){
 ll ret=fact[N]; ret*=inv[fact[R]]; ret%=mod;
 //ret=SafeMul(ret,inv[fact[R]],mod);
ret*=inv[fact[N-R]]; ret%=mod;
  //ret=SafeMul(ret,inv[fact[N-R]],mod);
  return ret;
voidPreprocessNCRforSmallNPrimeMod(11 lim,11 mod){
   fact[0]=1;
   for(ll i=1; i<=lim; i++){</pre>
       fact[i]=i*fact[i-1]; fact[i]%=mod;
    genAllInvUptoPrime(lim,mod,inv);
//N<10^18 and mod=P^R where P is a mod<10^6.
//Complexity:O(mod) preprocessing. O(logPN) query.
//factIgnoringP[i] holds i!(ignoring P's MULTIPLES
//) % mod for all i upto P^R-1
ll NCRforPrimePower(ll N,ll R,ll mod,ll P,ll M,
                   11* factIgnoringP){
  if(R>N) return 0;
  11 up,down1,down2 ,primePowerFactor=0 , ret;
  pair<ll, ll>tmp;
  tmp=FMPPIPF(N,mod,P,M,factIgnoringP);
  up=tmp.first; primePowerFactor+=tmp.second;
  tmp=FMPPIPF(R,mod,P,M,factIgnoringP);
```

```
down1=InvMod(tmp.first.mod):
 primePowerFactor-=tmp.second;
 tmp=FMPPIPF(N-R,mod,P,M,factIgnoringP);
 down2=InvMod(tmp.first.mod):
 primePowerFactor-=tmp.second;
 if(primePowerFactor>=M) return 0:
 ret=up; ret*=down1; ret%=mod; ret*=down2;
 ret%=mod; ll pFactor=1;
 for(ll i=1; i<=primePowerFactor;i++) pFactor*=P;</pre>
 ret*=pFactor; ret%=mod;
 ret=(ret+mod)%mod:
 return ret;
//N<10^18 and mod is any INTEGER so that all of
//its prime divisors are <10^6.Complexity:</pre>
//O(MOD) preprocessing. O(log N) query.
vector<pair<11,11> >PFactors; vector<11>Pmods;
11 factIgnoringP[22][MAXN];
11 NCRforAnyMOD(11 N,11 R,11 mod){
 vector<pair<11,11> >crt;
 for(int i=0; i<PFactors.size(); i++){</pre>
   11 m=Pmods[i]; 11 prm=PFactors[i].first;
   11 pwr=PFactors[i].second;
   11 a=NCRforPrimePower(N,R,m
                     ,prm,pwr,factIgnoringP[i]);
   crt.push_back(make_pair(a,m));
 return CRT(crt,mod);
void preporcessNCRforAnyMOD(11 mod){
 PFactors.clear(); Pmods.clear();
 PrimePowerFactorization(mod, PFactors);
 for(int i=0; i<PFactors.size(); i++){</pre>
   11 prm=PFactors[i].first;
   11 pwr=PFactors[i].second; 11 m=1;
   for(ll j=1; j<=pwr; j++) m*=prm;
Pmods.push_back(m); factIgnoringP[i][0]=1;</pre>
   for(ll j=1; j<m; j++){
    ll now=j;</pre>
       if(now%prm==0) now=1;
factIgnoringP[i][j]=(factIgnoringP[i][j-1]*now)%m;
/* Sieve(); , preporcessNCRforAnyMOD(M);
CRforAnyMOD(N,R,M) */
```

7.5 Pollard Rho (kundu vai)

```
return r < 0? r + mod : r:
LL power(LL x, LL p, LL mod){
  LL s=1, m=x;
  while(p) {
   if(p\&1) s = mult(s, m, mod);
   p>>=1;
   m = mult(m, m, mod);
 return s;
bool witness(LL a, LL n, LL u, int t){
 LL x = power(a,u,n);
 for(int i=0; i<t; i++) {</pre>
   LL nx = mult(x, x, n);
   if (nx==1 && x!=1 && x!=n-1) return 1;
x = nx;
 return x!=1;
vector<LL> bases = {2, 325, 9375, 28178, 450775, 9780504,
1795265022};
bool miller_rabin(LL n) {
  if (n<2) return 0;
  if (n%2==0) return n==2;
  LL u = n-1;
  int t = 0;
  while (u\%2==0) u/=2, t++; // n-1 = u*2^t
  for (LL v: bases) {
   LL a = v\%(n-1) + 1;
   if(witness(a, n, u, t)) return 0;
 return 1;
LL gcd(LL u, LL v) {
   if (u == 0) return v;
    if (v == 0) return u;
    int shift = __builtin_ctzll(u | v);
   u >>= __builtin_ctzll(u);
       v >>= __builtin_ctz(v);
       if (u > v) swap(u, v);
       v = v - u;
   } while (v):
   return u << shift;</pre>
mt19937_64 rng(chrono::steady_clock::now().
             time_since_epoch().count());
LL pollard_rho(LL n) {
    if (n==1)
                          return 1;
    if (n\%2==0)
                          return 2;
  if (miller_rabin(n)) return n;
    while (true) {
       LL x = uniform_int_distribution < LL > (1, n-1)(rng);
       LL y = 2, res = 1;
       for (int sz=2; res == 1; sz*=2) {
           for (int i=0; i<sz && res<=1; i++) {</pre>
               x = mult(x, x, n) + 1;
               res = gcd(abs(x-y), n)
       if (res!=0 && res!=n) return res;
```

7.6 Polynomial Algebra

```
namespace algebra {
  const int inf = 1e9; const int magic = 500;
// threshold for sizes to run the naive algo
```

```
namespace fft {
 const int maxn = 1 << 18;</pre>
 typedef double ftype;
 typedef complex<ftype> point;
 point w[maxn];
 const ftype pi = acos(-1);
 bool initiated = 0;
 void init() {
   if(!initiated) {
     for(int i = 1; i < maxn; i *= 2) {</pre>
       for(int j = 0; j < i; j++) {
         w[i + j] = polar(ftype(1), pi * j / i);
     initiated = 1;
 template<typename T>
 void fft(T*in, point *out, int n, int k = 1) {
   if(n == 1) {
     *out = *in;
   } else {
     n /= 2;
     fft(in, out, n, 2 * k);
     fft(in + k, out + n, n, 2 * k);
     for(int i = 0; i < n; i++) {
        auto t = out[i + n] * w[i + n];
        out[i + n] = out[i] - t;
       out[i] += t;
 template<typename T>
 void mul_slow(vector<T> &a, const vector<T>&b){
   vector<T> res(a.size() + b.size() - 1);
   for(size_t i = 0; i < a.size(); i++) {</pre>
     for(size_t j = 0; j < b.size(); j++) {</pre>
        res[i + j] += a[i] * b[j];
   \dot{a} = res;
 template<typename T>
 void mul(vector<T> &a, const vector<T> &b) {
   if(min(a.size(), b.size()) < magic) {</pre>
     mul_slow(a, b);
   init();
   static const int shift= 15,mask=(1<<shift)-1;</pre>
   size_t n = a.size() + b.size() - 1;
   while(__builtin_popcount(n) != 1) {
   a.resize(n);
   static point A[maxn], B[maxn];
   static point C[maxn], D[maxn];
   for(size_t i = 0; i < n; i++) {
   A[i] = point(a[i] & mask, a[i] >> shift);
     if(i < b.size()) {
       B[i] = point(b[i] & mask, b[i] >> shift);
     } else {
       B[i] = 0;
   fft(A, C, n); fft(B, D, n);
   for(size_t i = 0; i < n; i++) {
  point c0 = C[i] + conj(C[(n - i) % n]);
  point c1 = C[i] - conj(C[(n - i) % n]);</pre>
     point d0 = D[i] + conj(D[(n - i) \% n]);
```

```
point d1 = D[i] - coni(D[(n - i) \% n]):
     A[i] = c0 * d0 - point(0, 1) * c1 * d1;
     B[i] = c0 * d1 + d0 * c1:
   fft(A, C, n); fft(B, D, n);
   reverse(C + 1, C + n);
   reverse(D + 1, D + n);
   int t = 4 * n;
   for(size_t i = 0; i < n; i++) {</pre>
     int64_t A0 = llround(real(C[i]) / t);
     T A1 = llround(imag(D[i]) / t);
     T A2 = llround(imag(C[i]) / t);
     a[i] = A0 + (A1 << shift) + (A2 << 2*shift);
   return;
template<typename T>
T bpow(T x, size_t n) {
 return n?n\%2?x*bpow(x,n-1):bpow(x*x,n/2):T(1);
template<typename T>
T bpow(T x, size_t n, T m) {
 return n?n\%2?x*bpow(x,n-1,m)\%m:
         bpow(x*x/m,n/2,m):T(1);
template<typename T>
T gcd(const T &a, const T &b) {
 return b == T(0) ? a : gcd(b, a \% b);
template<typename T>
T nCr(T n, int r) { // runs in O(r)}
 T res(1);
 for(int i = 0; i < r; i++) {
   res *= (n - T(i));
   res /= (i + 1);
 return res;
template<int m>
struct modular {
 int64_t r;
 modular() : r(0) {}
 modular(int64_t rr) : r(rr) {
 if(abs(r) >= m) r \% = m; if(r < 0) r += m;}
 modular inv() const {return bpow(*this, m - 2);}
 modular operator*(modular&t){return(r*t.r)%m;}
modular operator/(modular&t){return*this*t.inv();}
 modular operator+=(modular&t){r+=t.r;
 if(r>=m) r-=m; return *this;}
 modular operator -= ( modular &t) {
     r \rightarrow t.r; if(r < 0) r + m; return *this;}
 modular operator + ( modular &t) {
     return modular(*this) += t;}
 modular operator - ( modular &t) {
     return modular(*this) -= t;}
 modular operator *= ( modular &t) {
     return *this = *this * t;}
 modular operator /= ( modular &t) {
     return *this = *this / t;}
 bool operator==(modular&t){return r==t.r;}
 bool operator!=(modular&t){return r != t.r;}
 operator int64_t() const {return r;}
template<int T>
istream& operator>> (istream &in, modular<T> &x){
 return in >> x.r;
template<typename T>
struct poly {
```

```
vector<T> a:
void normalize() { // get rid of leading zeroes
 while(!a.empty() && a.back() == T(0)) {
   a.pop_back();
poly(){}
poly(T a0) : a{a0}{normalize();}
poly(vector<T> t) : a(t){normalize();}
poly operator += (const poly &t) {
 a.resize(max(a.size(), t.a.size()));
 for(size_t i = 0; i < t.a.size(); i++) {</pre>
   a[i] += t.a[i];
 normalize();
 return *this;
poly operator -= (const poly &t) {
 a.resize(max(a.size(), t.a.size()));
 for(size_t i = 0; i < t.a.size(); i++) {</pre>
   a[i] = t.a[i];
 normalize();
 return *this:
poly operator+(poly &t){return poly(*this)+=t;}
poly operator-(poly &t){return poly(*this)-=t;}
poly mod_xk(size_t k){
// get same polynomial mod x^k
 k = min(k, a.size());
 return vector<T>(begin(a), begin(a) + k);
poly mul_xk(size_t k) const { // multiply by x^k
 poly res(*this);
 res.a.insert(begin(res.a), k, 0);
 return res;
poly div_xk(size_t k) const {
// divide by x^k, dropping coefficients
 k = min(k, a.size());
 return vector<T>(begin(a) + k, end(a));
poly substr(size_t 1, size_t r) const {
// return mod_xk(r).div_xk(1)
 1 = min(1, a.size());
 r = min(r, a.size());
 return vector<T>(begin(a) + 1, begin(a) + r);
poly inv(size_t n) const {
// get inverse series mod x^n
 assert(!is zero()):
 poly ans = a[0].inv();
 size_t a = 1;
 while (a < n) {
   poly C = (ans*mod_xk(2*a)).substr(a, 2 * a);
   ans -= (ans * C).mod_xk(a).mul_xk(a);
   a *= 2;
 return ans.mod_xk(n);
poly operator *= (const poly &t)
 {fft::mul(a, t.a); normalize(); return *this;}
poly operator * (const poly &t) const
 {return poly(*this) *= t;}
poly reverse(size_t n, bool rev = 0) const {
// reverses and leaves only n terms
 poly res(*this);
 if(rev) { // If rev = 1 then tail goes to head
   res.a.resize(max(n, res.a.size()));
```

```
std::reverse(res.a.begin(), res.a.end());
  return res.mod_xk(n);
pair<poly, poly> divmod_slow(const poly &b)
const { // when divisor or quotient is small
  vector<T> A(a);
  vector<T> res;
  while(A.size() >= b.a.size()) {
   res.push_back(A.back() / b.a.back());
   if(res.back() != T(0)) {
     for(size_t i = 0; i < b.a.size(); i++) {</pre>
       A[A.size() - i - 1] -= res.back() *
                 b.a[b.a.size() - i - 1];
   A.pop_back();
  std::reverse(begin(res), end(res));
 return {res, A};
pair<poly, poly> divmod(const poly &b) const
{ // returns quotiend and remainder of a mod b
  if(deg() < b.deg()) {
   return {poly{0}, *this};
  int d = deg() - b.deg();
  if(min(d, b.deg()) < magic) {</pre>
   return divmod_slow(b);
  poly D = (reverse(d + 1) * b.reverse(d + 1).
  inv(d + 1). mod_xk(d + 1). reverse(d + 1, 1);
 return \{D, *this - D * b\};
poly operator/(poly&t){return divmod(t).first;}
poly operator%(poly&t){return divmod(t).second;}
poly operator *= (T &x) {
 for(auto &it: a) {
   it *= x;
  normalize();
 return *this;
poly operator /= (const T &x) {
 for(auto &it: a) {
   it /= x;
 normalize();
 return *this;
poly operator *(T &x){return poly(*this) *= x;}
poly operator /(T &x){return poly(*this) /= x;}
void print() const {
 for(auto it: a) {
   cout << it << ''';
  cout << endl;
T \text{ eval}(T \text{ x}) \{ // \text{ evaluates in single point } x \}
 T res(0);
 for(int i = int(a.size()) - 1; i >= 0; i--) {
  res *= x;
   res += a[i];
 return res;
T& lead() { // leading coefficient
 return a.back():
int deg() const { // degree
```

```
return a.empty() ? -inf : a.size() - 1;
bool is_zero() const { // is polynomial zero
 return a.empty();
T operator [](int idx) const {
 return idx>=(int)a.size()||idx<0?T(0):a[idx];</pre>
T& coef(size_t idx) {
// mutable reference at coefficient
 return a[idx];
bool operator == ( poly &t) {return a == t.a;}
bool operator != ( poly &t) {return a != t.a;}
poly deriv() { // calculate derivative
 vector<T> res;
 for(int i = 1; i <= deg(); i++) {</pre>
   res.push_back(T(i) * a[i]);
poly integr() { // calculate integral with C = 0
 vector<T> res = {0}:
 for(int i = 0; i <= deg(); i++) {</pre>
   res.push_back(a[i] / T(i + 1));
 return res;
size_t leading_xk() const {
// Let p(x) = x^k * t(x), return k
 if(is_zero()) {
   return inf;
 int res = 0;
 while(a[res] == T(0)) {
   res++;
 return res;
poly log(size_t n){//calculate log p(x) mod x^n
 assert(a[0] == T(1));
 return (deriv().mod_xk(n)*inv(n))
         .integr().mod_xk(n);
poly exp(size_t n){//calculate exp p(x) mod x^n
 if(is_zero()) {
   return T(1);
 assert(a[0] == T(0));
 poly ans = T(1); size_t a = 1;
 while(a < n) {
   poly C=ans.log(2*a).div_xk(a)-substr(a,2*a);
   ans -= (ans * C).mod_xk(a).mul_xk(a);
   a *= 2;
 return ans.mod_xk(n);
poly pow_slow(size_t k, size_t n){//if k is small
 return k?k%2?(*this*pow_slow(k-1,n)).mod_xk(n)
:(*this* *this).mod_xk(n).pow_slow(k/2,n):T(1);
poly pow(size_t k,size_t n){
//calculate p^k(n) mod x^n
 if(is_zero()) return *this;
 if(k < magic) return pow_slow(k, n);</pre>
 int i = leading_xk();
 T j = a[i];
 poly t = div_xk(i) / j;
 return bpow(j,k)*(t.log(n)*T(k))
         .exp(n).mul_xk(i*k).mod_xk(n);
```

```
poly mulx(T x){
//component-wise multiplication with x^k
 T cur = 1; poly res(*this);
 for(int i = 0; i <= deg(); i++) {</pre>
   res.coef(i) *= cur;
   cur *= x:
 return res;
poly mulx_sq(T x){
//component-wise multiplication with x^{k^2}
 T \text{ cur} = x; T \text{ total} = 1; T xx = x * x;
 polv res(*this);
 for(int i = 0; i <= deg(); i++) {
   res.coef(i) *= total;
   total *= cur;
   cur *= xx;
 return res;
vector<T> chirpz_even(T z, int n) {
// P(1), P(z^2), P(z^4), ..., P(z^2(n-1))
 int m = deg();
 if(is_zero()) {
   return vector<T>(n, 0);
 vector<T> vv(m + n);
 T zi = z.inv(); T zz = zi * zi;
 T cur = zi; T total = 1;
 for(int i = 0; i \le max(n - 1, m); i++) {
   if(i <= m) {vv[m - i] = total;}</pre>
   if(i < n) \{vv[m + i] = total;\}
   total *= cur;
   cur *= zz;
poly w=(mulx_sq(z)*vv).substr(m,m+n).mulx_sq(z);
 vector<T> res(n);
 for(int i = 0; i < n; i++) {</pre>
   res[i] = w[i]:
 return res;
}
vector<T> chirpz(T z, int n) {
// P(1), P(z), P(z^2), ..., P(z^{(n-1)})
 auto even = chirpz_even(z, (n + 1) / 2);
 auto odd = mulx(z).chirpz_even(z, n / 2);
 vector<T> ans(n):
 for(int i = 0; i < n / 2; i++) {
  ans[2 * i] = even[i];</pre>
   ans[2 * i + 1] = odd[i];
 if(n % 2 == 1) {
   ans[n-1] = even.back():
 return ans;
template<typename iter>
vector<T> eval(vector<poly>&tree,int v,iter 1
,iter r) { // auxiliary evaluation function
 if(r - 1 == 1) {
   return {eval(*1)};
 } else {
   auto m = 1 + (r - 1) / 2;
   auto A = (*this % tree[2 * v]).eval(tree,
             2 * v, 1, m);
   auto B = (*this \% tree[2 * v + 1]).eval
                 (tree, 2 * v + 1, m, r);
   A.insert(end(A), begin(B), end(B));
   return A:
```

```
vector<T> eval(vector<T> x) {
     // evaluate polynomial in (x1, ..., xn)
     int n = x.size():
     if(is_zero()) {
       return vector<T>(n, T(0)):
     vector<poly> tree(4 * n);
     build(tree, 1, begin(x), end(x));
     return eval(tree, 1, begin(x), end(x));
   template<typename iter>
   poly inter(vector<poly>&tree,int v,iter 1,iter
r, iter ly, iter ry) {//axiliary interpolation function
     if(r - 1 = 1) {
       return {*ly / a[0]};
     } else {
       auto m = 1 + (r - 1) / 2;
       auto my = ly + (ry - ly) / 2;
       auto A = (*this \% tree[2 * v]).
       inter(tree, 2 * v, 1, m, 1y, my);
auto B = (*this % tree[2 * v + 1]).
inter(tree, 2 * v + 1, m, r, my, ry);
       return A * tree[2 * v + 1] + B * tree[2*v];
  témplate<typename T>
  poly<T> operator * (const T& a,const poly<T>&b){
   return b * a;
  template<typename T>
  poly<T> xk(int k) { // return x^k
   return poly<T>{1}.mul_xk(k);
  template<typename T>
 T resultant(poly<T>a,poly<T>b){
   //computes resultant of a and b
   if(b.is zero()) {
     return 0;
   } else if(b.deg() == 0) {
     return bpow(b.lead(), a.deg());
   } else {
     int pw = a.deg();
     a %= b; pw -= a.deg();
     T mul=bpow(b.lead(),pw)*T((b.deg()&a.deg()&1)
     ?-1:1); T ans= resultant(b, a);
     return ans * mul;
  template<typename iter>
  poly<typename iter::value_type> kmul(iter L,
                          iter R) { // computes
 //(x-a1)(x-a2)...(x-an) without building tree
   if(R - L == 1) {
   return vector<typename iter::value_type>{-*L,1};
   } else {
     iter M = L + (R - L) / 2;
     return kmul(L, M) * kmul(M, R);
  template<typename T, typename iter>
  poly<T> build(vector<poly<T>> &res, int v, iter L,
                          iter R){//
  builds evaluation tree for (x-a1)(x-a2)...(x-an)
   if(R - L == 1) {
     return res[v] = vector<T>{-*L, 1};
     iter M = L + (R - L) / 2;
```

```
return res[v] = build(res. 2 * v. L. M)
             * build(res, 2 * v + 1, M, R);
 template<typename T>
 poly<T> inter(vector<T>x, vector<T>y){//interpo-
 //lates minimum polynomial from (xi, yi) pairs
   int n = x.size();
   vector<poly<T>> tree(4 * n);
   return build(tree, 1, begin(x), end(x)).deriv()
.inter(tree, 1, begin(x), end(x), begin(y), end(y));
using namespace algebra;
const int mod = 1e9 + 7;
typedef modular<mod> base:
typedef poly<br/>base> polyn;
using namespace algebra;
signed main() {
 int n = 100000;
 polyn a; vector<base> x;
 for(int i = 0; i <= n; i++) {
   a.a.push_back(1 + rand() % 100);
   x.push_back(1 + rand() \% (2 * n));
 sort(begin(x), end(x));
 x.erase(unique(begin(x), end(x)), end(x));
 auto b = a.eval(x);
cout << clock() / double(CLOCKS_PER_SEC) << endl;</pre>
 auto c = inter(x, b);
 polyn md = kmul(begin(x), end(x));
 cout << clock() / double(CLOCKS_PER_SEC) << endl;</pre>
 assert(c == a % md);
```

7.7 all comb

```
vector<int> ans;
void gen(int n, int k, int idx, bool rev) {
   if (k > n \mid \mid k < 0) return;
   if (!n) {
       for (int i = 0; i < idx; ++i) {</pre>
           if (ans[i])
              cout << i + 1;
       cout << "\n":
       return;
   ans[idx] = rev;
   gen(n-1, k-rev, idx + 1, false);
   ans[idx] = !rev:
   gen(n-1, k-!rev, idx + 1, true);
void all_combinations(int n, int k) {
   ans.resize(n);
   gen(n, k, 0, false);
```

7.8 gauss elimination

```
Det = 1.0. Rank = 0:
  memset(where, -1, sizeof(where));
  for(int col=0,row = 0;col<m&&row < n; ++col) {
   int sel = row;
   for(int i = row+1; i < n; ++i)</pre>
     if(fabs(a[i][col])>fabs(a[sel][col])) sel=i;
   if(fabs(a[sel][col]) < EPS) {Det=0.0; continue;}</pre>
  for(int j=0;j<=m;++j)swap(a[sel][j],a[row][j]);</pre>
   if(row != sel) Det = -Det;
   Det *= a[row][col];
   where[col] = row;
   double s = (1.0 / a[row][col]);
   for(int j = 0; j <= m; ++j) a[row][j] *= s;
   for(int i = 0; i < n; ++i) if (i != row &&
         fabs(a[i][col]) > EPS) {
      double t = a[i][col];
      for(int j = 0; j \le m; ++j)
       a[i][j] -= a[row][j] * t;
   ++row, ++Rank;
  for(int_i = 0; i \leq m; ++i)
   ans[i] = (where[i] == -1) ?0.0:a[where[i]][m];
 for(int i = Rank; i < n; ++i)</pre>
   if(fabs(a[i][m]) > EPS) return 0;
 for(int i = 0; i < m; ++i)</pre>
   if(where[i] == -1) return INF;
// calculates gauss modulo a prime
long long Det;
long long bigmod(long long x,
               long long pow, long long mod) {
  long long ret = 1;
  while(pow > 0) {
     if(pow & 1) ret = (ret * x) % mod;
     x = (x * x) \% mod;
     pow >>= 1;
  return ret;
#define INVERSE(a, m) bigmod(a, m-2, m)
int gauss(long long a[MAX][MAX],
 long long ans[MAX],int n,int m,long long mod){
Det = 1, Rank = 0;
memset(where, -1, sizeof(where));
for(int col = 0, row = 0; col<m&&row<n;++col){
 int sel = row;
  for(int i = row+1; i < n; ++i)</pre>
    if(abs(a[i][col]) > abs(a[sel][col])) sel=i;
  if(!a[sel][col]) { Det = 0; continue; }
  for(int j=0; j<=m;++j) swap(a[sel][j],a[row][j]);</pre>
 if(row != sel) Det = -Det;
  Det = (Det * a[row][col]) % mod;
  where [col] = row;
                     // inverse of a[row][col]
  long long s = INVERSE(a[row][col], mod);
  for(int j = 0; j \le m; ++j)
   a[row][j] = (\tilde{a}[row][j] * s) % mod;
  for(int i = 0; i < n; ++i) if (i != row &&
                               a[i][col] > 0) {
    long long t = a[i][col];
    for(int j = 0; j <= m; ++j) a[i][j] =
     (a[i][j] - (a[row][j]*t)^{\%} \mod + mod)^{\%}
 ++row, ++Rank;
for(int i = 0; i < m; ++i)
 ans[i] = (where[i] = -1) ? 0 : a[where[i]][m];
for(int i = Rank; i < n; ++i)
```

```
if(a[i][m]) return 0:
for(int i = 0; i < m; ++i)</pre>
 if(where[i] == -1) return INF;
// calculates 32 times faster for modulo 2
int Det; // number of variables (must be defined)
int gauss(vector < bitset<MAX> > &a,
         bitset<MAX> &ans, int n, int m) {
 Det = 1, Rank = 0;
  memset(where, -1, sizeof(where));
  for(int col=0,row=0; col < m && row < n;++col){</pre>
     int sel = row;
     for(int i = row; i < n; ++i)</pre>
     if(a[i][col]) { sel = i; break; }
if(!a[sel][col]) { Det = 0; continue; }
     swap(a[sel], a[row]);
     if(row != sel) Det = -Det;
     Det &= a[row][col];
     where [col] = row;
     for(int i = 0; i < n; ++i)</pre>
       if (i != row&&a[i][col] > 0)a[i]^=a[row];
     ++row, ++Rank;
  for(int i = 0; i < m; ++i)</pre>
     ans[i] = (where[i] == -1)?0:a[where[i]][m];
  for(int i = Rank;i<n;++i)if(a[i][m]) return 0;</pre>
  for(int i = 0; i < m; ++i)
   if(where[i] == -1) return INF;
  return 1;
```

7.9 linear sieve

```
vector<int> lp(N+1);
vector<int> pr;
for (int i=2; i <= N; ++i) {
   if (lp[i] == 0) {
      lp[i] = i;pr.push_back(i);}
for (int j=0; j < (int)pr.size()
   && pr[j] <= lp[i] && i*pr[j] <= N; ++j)
      lp[i * pr[j]] = pr[j];
}</pre>
```

7.10 matrix inverse

```
#define MOD 100000007
#define MAX 100
11 BigMod(ll a,ll r,ll Mod) {
 ll ret=1;
 while(r) {
   if(r&1)ret=ret*a,ret=ret%Mod;
   a=a*a; a=a\%Mod; r>>=1;
 return ret;
ll InverseMod(ll a,ll Mod){
 return BigMod(a,Mod-2,Mod);
11 Mul(11 x,11 y) { return (x*y)%MOD; }
ll Div(ll x,ll y){
 return(x*InverseMod(y,MOD))%MOD;
//1-based
struct Matrix{
 int row, col;
ll m[MAX][MAX];
 Matrix() {memset(m,0,sizeof(m));}
 void Set(int r,int c) {row = r; col = c;}
 Matrix(int r,int c){
     memset(m,0,sizeof(m)); Set(r,c);
```

```
void normalize(){
   for(int i=1; i<=row; i++){</pre>
     for(int j=1; j<=col; j++){</pre>
       m[i][j]\%=MOD; if(m[i][j]<0)m[i][j]+=MOD;
ll Det(Matrix mat){
 assert(mat.row == mat.col):
 int n = mat.row; mat.normalize(); ll ret = 1;
 for(int i = 1; i <= n; i++){
   for(int j = i + 1; j \le n; j++){
     while(mat.m[j][i]){
       11 t = Div(mat.m[i][i], mat.m[j][i]);
       for(int k = i; k <= n; ++k){
       mat.m[i][k] -= Mul(mat.m[j][k] , t);
       if(mat.m[i][k] < 0) mat.m[i][k] += MOD;</pre>
       swap(mat.m[j][k], mat.m[i][k]);
       ret = MOD - ret;
   if(mat.m[i][i] == 0) return 0;
   ret = Mul(ret, mat.m[i][i]);
 if(ret < 0) ret += MOD; return ret;</pre>
ll Tmp[MAX<<1][MAX<<1];
Matrix Inverse(Matrix mat){
 assert(mat.row==mat.col);assert(Det(mat) != 0);
 int n = mat.row; mat.normalize();
 for(int i=1;i<=n;i++){</pre>
   for(int j=1;j<=n;j++) Tmp[i][j] = mat.m[i][j];</pre>
   for(int j=n+1; j \le 2*n; j++) Tmp[i][j] = 0;
   Tmp[i][i+n] = 1;
 for(int i=1; i<=n; i++){</pre>
   assert(Tmp[i][i] != 0);
   for(int j=1; j<=n; j++){
     if(i == j) continue
     11 c = Div(Tmp[j][i], Tmp[i][i]);
     for(int k=i; k<=2*n; k++){
Tmp[j][k] = Tmp[j][k] - Mul(Tmp[i][k], c);</pre>
     if(Tmp[j][k] < 0) Tmp[j][k] += MOD;</pre>
 Matrix Inv(n,n);
 for(int i=1; i<=n; i++){
   for(int j = 1; j <= n; j++){
  Inv.m[i][j] = Div(Tmp[i][j+n], Tmp[i][i]);</pre>
 return Inv;
```

3 String

8.1 Aho Corasick

```
const int K = 26;
struct Vertex {
  int next[K];
  bool leaf = false;
  int p = -1;
  char pch;
  int link = -1;
  int go[K];
```

```
Vertex(int p=-1, char ch='$') : p(p), pch(ch) {
   fill(begin(next), end(next), -1);
   fill(begin(go), end(go), -1);
vector<Vertex> t(1);
void add_string(string const& s) {
 int v = 0;
 for (char ch : s) {
     int c = ch - 'a';
     if (t[v].next[c] == -1) {
         t[v].next[c] = t.size();
         t.emplace_back(v, ch);
     v = t[v].next[c];
 t[v].leaf = true;
int go(int v, char ch);
int get_link(int v) {
  if (t[v].link == -1) {
   if (v == 0 || t[v].p == 0)
       t[v].link = 0:
   else
       t[v].link = go(get_link(t[v].p), t[v].pch);
 return t[v].link;
int go(int v, char ch) {
 int c = ch - 'a';
 if (t[v].go[c] == -1) {
   if (t[v].next[c] != -1)
       t[v].go[c] = t[v].next[c];
       t[v].go[c] = v == 0?0: go(get_link(v), ch);
 return t[v].go[c];
```

8.2Hashing

```
const int MAX = 3000009;
11 \mod [2] = \{1000000007, 1000000009\};
//Some back-up primes: 1072857881, 1066517951, 1040160883 vector<int> d1(n);
ll bases[2] = \{137, 281\};
11 pwbase[3][MAX];
void Preprocess(){
  pwbase[0][0] = pwbase[1][0] = 1;
  for(ll i = 0; i < 2; i++){
   for(11 j = 1; j < MAX; j++){}
     pwbase[i][j]=(pwbase[i][j-1]*bases[i])%mods[i];
11 fmod(l1 a, l1 b, int md=mods[0]) {
 unsigned long long x = (long long) a * b;
  unsigned xh=(unsigned)(x >> 32),xl=(unsigned) x, d, m;
   "divl %4; \n\t"
   : "=a" (d), "=d" (m)
: "d" (xh), "a" (xl), "r" (md)
 réturn m;
struct Hashing{
 vector<vector<ll>>> hsh;
  Hashing(){}
  Hashing(string &_str) {
     hsh.push_back(vector<ll>(_str.size()+5, 0));
     hsh.push_back(vector<ll>(_str.size()+5, 0));
```

```
Build1(str):
void Build(string &str){Build1(str);Build2(str);}
void Build1(const string &str) {
 for(Il i = str.size() - 1: i >= 0: i--){}
   hsh[j][i] = fmod(hsh[j][i + 1],
                       bases[j], mods[j]) + str[i];
   if (hsh[j][i] > mods[j])
       hsh[j][i] -= mods[j];
void Build2(const string &str) {
 for(Il i = str.size() - 1; i >= 0; i--){
     hsh[j][i] = fmod(hsh[j][i + 1]),
                      bases[j], mods[j]) + str[i];
     if (hsh[j][i] > mods[j])
         hsh[j][i] -= mods[j];
pair<ll, ll> GetHash(ll i, ll j){
   assert(i <= j);</pre>
   11 \text{ tmp1} = (hsh[0][i] -
         fmod(hsh[0][j + 1] ,pwbase[0][j - i + 1]));
   ll tmp2 = (hsh[1][i] -
fmod(hsh[1][j + 1], pwbase[1][j - i + 1], mods[1]));
   if(tmp1 < 0) tmp1 += mods[0];</pre>
   if(tmp2 < 0) tmp2 += mods[1];
   return make_pair(tmp1, tmp2);
ll getSingleHash(ll i, ll j) {
  assert(i <= j);
  11 \text{ tmp1} = (hsh[0][i] - fmod(hsh[0][j + 1]),
                            pwbase[0][j - i + 1]));
  if(tmp1 < 0) tmp1 += mods[0];
 return tmp1;
```

Manacher's Algorithm

```
for (int i = 0, l = 0, r = -1; i < n; i++){
 int k = (i > r)?1: min(d1[1 + r - i], r - i + 1);
  while (0 \le i - k \& i + k \le k \le [i - k] = s [i + k]) k + +;
  d1[i] = k--;
  if^{(i+k)}
    l = i - k; r = i + k;
vector<int> d2(n);
for (int i = 0, 1 = 0, r = -1; i < n; i++){
  int k=(i>r)?0:min(d2[l+r-i+1],r-i+1);

  while (0 \le i - k - 1 \& \& i + k \le n \& \& s [i - k - 1] == s [i + k]) k++;
  d2[i]=k--:
  if(i+k>r){
    l=i-k-1;
    r=i+k;
```

Palindromic Tree

```
int tree[N][26], idx;
int len[N], link[N], t;
char s[N]; // 1-indexed
len[1] = -1, link[1] = 1,;
len[2] = 0, link[2] = 1;
idx = t = 2;
```

```
roid extend(int p) {
while (s[p - len[t] - 1] != s[p]) t = link[t];
int x = link[t], c = s[p] - 'a';
while (s[p - len[x] - 1]] = s[p]) x = link[x];
if(!tree[t][c]) {
  tree[t][c] = ++idx;
  len[idx] = len[t] + 2;
  link[idx] = len[idx] == 1 ? 2 : tree[x][c];
} t = tree[t][c]:
```

```
8.5 Suffix Array
vector<int> sort_cyclic_shifts(char *s) {
   int n = strlen(s);
   const int alphabet = 256;
   vector<int> p(n), c(n), cnt(max(alphabet, n),0);
   for (int i = 0; i < n; i++) cnt[s[i]]++;
   for (int i=1;i<alphabet;i++) cnt[i]+=cnt[i-1];</pre>
   for (int i = 0; i < n; i++) p[--cnt[s[i]]] = i;
   c[p[0]] = 0;
   int classes = 1;
   for (int i = 1; i < n; i++) {
       if (s[p[i]] != s[p[i-1]]) classes++;
       c[p[i]] = classes - 1;
   vector<int> pn(n), cn(n);
   for (int h = 0; (1 << h) < n; ++h) {
       for (int i = 0; i < n; i++) {</pre>
           pn[i] = p[i] - (1 << h);
           if (pn[i] < 0)
               pn[i] += n;
       fill(cnt.begin(), cnt.begin() + classes, 0);
       for (int i = 0; i<n;i++) cnt[c[pn[i]]]++;</pre>
     for(int i=1;i<classes;i++) cnt[i]+=cnt[i-1];</pre>
       for (int i = n-1; i >= 0; i--)
           p[--cnt[c[pn[i]]]] = pn[i];
       cn[\overline{p}[0]] = 0:
       classes = 1:
       for (int i = 1; i < n; i++) {
   int ind=p[i] + (1 << h);</pre>
           if(ind>=n) ind=ind-n;
           pair<int, int> cur = {c[p[i]], c[ind]};
ind=p[i-1] + (1 << h);</pre>
           if(ind>=n) ind=ind-n;
           pair<int, int> prev = {c[p[i-1]], c[ind]};
           if (cur != prev)
               ++classes;
           cn[p[i]] = classes - 1;
       c.swap(cn);
   return p;
vector<int> suffix_array_construction(char *s) {
   int n=strlen(s);
   s[n]='#';
   vector<int> sorted_shifts=sort_cyclic_shifts(s);
   sorted_shifts.erase(sorted_shifts.begin());
   return sorted_shifts;
vector<int> lcp_construction(char *s,
                   vector<int> const& p) {
   int n = strlen(s);
   vector<int> rank(n, 0);
   for (int i = 0; i < n; i++)
       rank[p[i]] = i;
   int k = 0;
```

```
vector < int > lcp(n-1, 0):
   for (int i = 0; i < n; i++) {
       if (rank[i] == n - 1) {
           k = 0;
           continue;
       int j = p[rank[i] + 1];
       while (i + k < n \&\& j + k < n \&\& s[i+k] == s[j+k])
           k++;
       lcp[rank[i]] = k;
       if (k)
           k--:
   return lcp;
int lcp(int i, int j) {
   int ans = 0:
   for (int k = log_n; k >= 0; k--) {
       if (c[k][i] == c[k][j]) {
           ans += 1 << k;
           i += 1 << k;
           j += 1 << k;
   return ans;
```

8.6 Suffix Automaton

```
class SuffixAutomaton{
 bool complete; int last; set<char> alphabet;
 struct state{
     // shortest_non_appearing_string -> snas
     // length_of_substrings -> los
     int len, link, endpos, first_pos, snas, height;
     ll substrings, los;
     bool is_clone;
     map<char, int> next; vector<int> inv_link;
     state(int leng=0,int li=0){
       len=leng; link=li; first_pos=-1;
       substrings=0; los=0;
       endpos=1; snas=0;
       is_clone=false; height=0;
 véctor<state> st;
 void process(int node){
   map<char, int> ::iterator mit;
   st[node].substrings=1;st[node].snas=st.size();
   if((int) st[node].next.size()
       <(int) alphabet.size()) st[node].snas=1;
   for(mit=st[node].next.begin();
         mit!=st[node].next.end(); ++mit){
       if(st[mit->second].substrings==0)
           process(mit->second);
       st[node].height=max(st[node].height,
              1+st[mit->second].height);
       st[node].substrings=st[node].substrings
              +st[mit->second].substrings;
       st[node].los=st[node].los
 +st[mit->second].los+st[mit->second].substrings;
       st[node].snas=min(st[node].snas
       ,1+st[mit->second].snas);
   if(st[node].link!=-1){
     st[st[node].link].inv_link.push_back(node);
 void set_suffix_links(int node){
   for(i=0; i<st[node].inv_link.size(); i++){</pre>
```

```
set suffix links(st[node].inv link[i]):
     st[node].endpos=st[node].endpos+st[st[node].
                inv_link[i]].endpos;
 }
  void output_all_occurrences(int v,
  int P_length, vector<int> &pos){
   if (!st[v].is_clone)
       pos.push_back(st[v].first_pos - P_length+1);
   for (int u : st[v].inv_link)
       output_all_occurrences(u, P_length, pos);
void kth_smallest(int node,int k,vector<char> &str){
 if(k==0) return;
map<char, int> ::iterator mit;
 for(mit=st[node].next.begin();
   mit!=st[node].next.end(); ++mit){
   if(st[mit->second].substrings<k)</pre>
         k=k-st[mit->second].substrings;
   else{
       str.push_back(mit->first);
       kth_smallest(mit->second,k-1,str);
       return;
int find_occurrence_index(int node,int index,
                    vector<char> &str){
   if(index==str.size()) return node;
   if(!st[node].next.count(str[index])) return -1:
   else return find_occurrence_index(st[node].
             next[str[index]],index+1,str);
void klen_smallest(int node,int k,vector<char> &str)
 if(k==0) return;
map<char, int> ::iterator mit;
 for(mit=st[node].next.begin();
   mit!=st[node].next.end(); ++mit){
     if(st[mit->second].height>=k-1){
         str.push_back(mit->first);
         klen_smallest(mit->second,k-1,str);
 }
void minimum_non_existing_string(int node,
                        vector<char> &str){
  map<char, int> ::iterator mit;
  set<char>::iterator sit;
 for(mit=st[node].next.begin(),sit=
alphabet.begin();sit!=alphabet.end();++sit,++mit){
   if (mit==st[node].next.end() | |mit->first!=(*sit))
       str.push_back(*sit);
       return;
   else if(st[node].snas==1+st[mit->second].snas)
       str.push back(*sit):
     minimum_non_existing_string(mit->second,str);
void find_substrings(int node,int index,vector<char>
&str,vector<pair<l1,11> > &sub_info){
  sub_info.push_back(make_pair(st[node].substrings,
  st[node].los+st[node].substrings*index));
 if(index==str.size()) return;
```

```
if(st[node].next.count(str[index]))
   find_substrings(st[node].next[str[index]],
                       index+1,str,sub_info);
   return;
 else
   sub_info.push_back(make_pair(0,0));
void check()
   if(!complete)
       process(0);
       set_suffix_links(0);
       compléte=true;
SuffixAutomaton(set<char> &alpha)
 st.push_back(state(0,-1));
 last=0:
 complete=false;
 set<char>::iterator sit;
 for(sit=alpha.begin(); sit!=alpha.end(); sit++)
   alphabet.insert(*sit);
 st[0].endpos=0;
void sa_extend(char c){
 int cur = st.size();
//printf("New node (%d,%c)\n",cur,c);
 st.push_back(state(st[last].len + 1));
 st[cur].first_pos=st[cur].len-1;
 int p = last;
 while (p != -1 \&\& !st[p].next.count(c)){
     st[p].next[c] = cur;
     //printf("Set edge %d -> %d (%c)\n",p,cur,c);
     p = st[p].link;
 if (p == -1){
     st[cur].link = 0;
     //printf("Set link %d -> %d\n",cur,0);
 else{
     int q = st[p].next[c];
     if (st[p].len + 1 == st[q].len){
         st[cur].link = q;
         //printf("Set link %d -> %d\n",cur,q);
     else{
         int clone = st.size();
//printf("Create clone node %d from %d\n",clone,q);
//printf("Set link %d -> %d\n",clone,st[q] link);
       st.push_back(state(st[p].len+1,st[q].link));
         st[clone].next = st[q].next;
         st[clone].is_clone=true;
         st[clone].endpos=0;
         st[clone].first_pos=st[q].first_pos;
         while (p != -1 \&\& st[p].next[c] == q){
//printf("Change transition %d -> %d : %d -> %d (%c)
// \ln",p,q,p,clone,c);
             st[p].next[c] = clone;
             p = st[p].link;
```

```
//printf("Change link %d \rightarrow %d : %d \rightarrow %d\n",q,st[q]
// .link,q,clone);
//printf("Set link %d -> %d\n",cur,clone);
         st[q].link = st[cur].link = clone;
  last = cur;
  complete=false;
SuffixAutomaton(){
   int i;
   for(i=0; i<st.size(); i++){
    st[i].next.clear();</pre>
       st[i].inv_link.clear();
   st.clear();
   alphabet.clear();
  void kth_smallest(int k,vector<char> &str){
     check():
     kth_smallest(0,k,str);
int FindFirstOccurrenceIndex(vector<char> &str){
    check();
    int ind=find_occurrence_index(0,0,str);
    if(ind==0) return -1;
    else if(ind==-1) return st.size();
    else return st[ind].first_pos+1-(int)str.size();
void FindAllOccurrenceIndex(vector<char> &str,
       vector<int> &pos){
    check():
    int ind=find_occurrence_index(0,0,str);
   if(ind!=-1)
       output_all_occurrences(ind,str.size(),pos);
int Occurrences(vector<char> &str){
    check();
    int ind=find occurrence index(0.0.str):
    if(ind==0) return 1;
    else if(ind==-1) return 0;
    else return st[ind].endpos:
```

```
void klen_smallest(int k,vector<char> &str){
   check():
   if(st[0].height>=k) klen_smallest(0,k,str);
void minimum non existing string(vector<char> &str){
   int ind=find_occurrence_index(0,0,str);
 if(ind!=-1) minimum_non_existing_string(ind,str);
ll cyclic_occurrence(vector<char> &str){
 check();
 int i,j,len; ll ans=0;
 int n=str.size(); set<int> S;
 set<int>::iterator it;
 for (i=0, j=0, len=0; i < n*2-1; i++){
   //printf("%d->%c\n",i,str[i%n]);
   if(st[i].next.count(str[i%n])){
       len++:
       j=st[j].next[str[i%n]];
   else{
     while(j!=-1&&(!st[j].next.count(str[i%n])))
         j=st[j].link;
     if(j!=-1){\{}
         len=st[j].len+1;
         j=st[j].next[str[i%n]];
     else{
       len=0:
       j=0;
   while (st[j].link!=-1\&\&st[st[j].link].len>=n){
      j=st[j].link; len=st[j].len;
   if(len>=n) S.insert(j);
 for(it=S.begin();it!=S.end();++it){
     ans=ans+st[*it].endpos;
 return ans:
```

```
}; // main
vector<char> X;
 int i;
 set<char> alpha;
 for(i=0; i<26; i++){
     alpha.insert('a'+i);
 SuffixAutomaton sa(alpha);
 for(i=0;; i++){
     scanf("%c",&c);
     if(!('a'<=c&&c<='z')) break;
     sa.sa_extend(c);
 int n,j;
scanf("%d ",&n);
 for(j=0; j<n; j++){</pre>
     for(i=0;; i++){
         scanf("%c",&c);
         if(!('a'<=c&&c<='z')) break;
         X.push_back(c);
     11 ans=sa.cyclic_occurrence(X);
     X.clear();
     printf("%164d\n",ans);
```

```
vector<int> z_function(string s) {
   int n = (int) s.length();
   vector<int> z(n);
   for (int i = 1, \dot{1} = 0, r = 0; i < n; ++i) {
       if (i <= r)
           z[i] = \min_{x \in A} (x - i + 1, z[i - 1]);
       while (i + z[i] < n \&\& s[z[i]] == s[i + z[i]])
       ++z[i];
if (i + z[i] - 1 > r)
           l = i, r = i + z[i] - 1;
   return z;
```