

Shortest Common Supersequence

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LCS: Longest Common Subsequence

Given two sequences, print all longest subsequence present in both of them.

X = "opengenus"
Y = "operagenes"

Z = "opegens"



X = "AGGTAB"
Y = "GXTXAYB"

Z = "GTAB"

SCS: Shortest Common Supersequence

Given two strings str1 and str2, the task is to find the length of the shortest string that has both str1 and str2 as subsequences.

X = "open**g**enus"
Y = "oper**a**gen**e**s"

z = "open**r**agen**u**e**s**"



X = "AGG**T**AB"
Y = "GXT**X**A**Y**B"

Z = "AGXG**T**X**A**Y**B**"

Considerations:

- Difference between subsequence and substring

String 1 = Frankenstein

Substring= rank, enstein

Subsequence = rnk, esi

X = "AGGTAB"

Y = "GXTXAYB"

Z = "AGXGTXAYB"



$$\text{SCSthLeng}(X, Y) = m + n - \text{LCSLength}(X, Y)$$

-Recursive (Brute Force) approach - TC: $O(2^{\min(m,n)})$

Case 1.

If either X or Y is empty, then return the length of the other sequence.

Case 2.

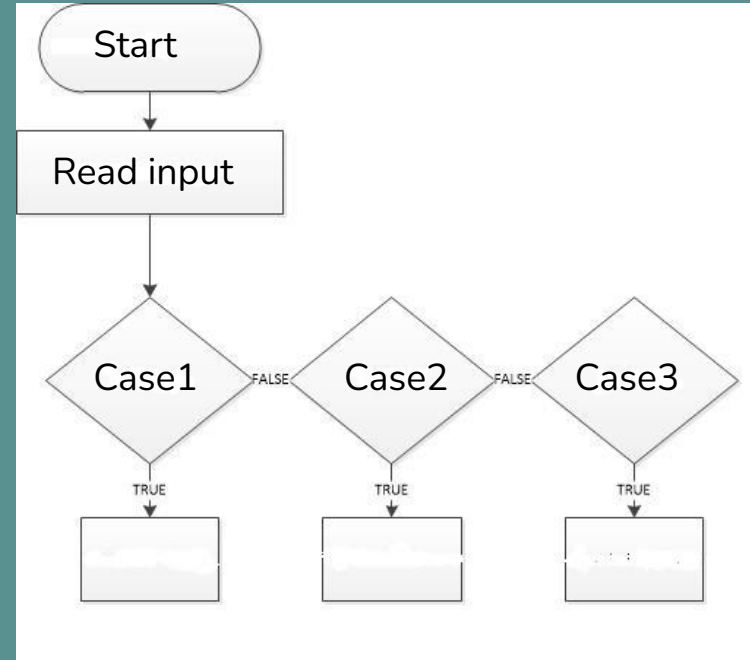
If both sequences X and Y start with the same character, then return $char[0] + SCS(X_{n-1}, Y_{m-1})$.

Case 3.

If X and Y start with a different character, then return $char + \text{the smallest of } SCS(X_{n-1}, Y_m) \text{ and } SCS(X_n, Y_{m-1})$.

X = "AGGTAB"
Y = "GXTXAYB"

Z = "AGXGTXAYB"



```

1 package presentation;
2 import java.util.Scanner;
3 public class Recursive {
4
5     private static boolean isEmpty(String s) {
6         return null == s || s.isEmpty();
7     }
8
9     private static String scs(String x, String y) {
10         if (isEmpty(x)) {
11             return y;
12         }
13         if (isEmpty(y)) {
14             return x;
15         }
16         if (x.charAt(0) == y.charAt(0)) {
17             return x.charAt(0) + scs(x.substring(1), y.substring(1));
18         }
19         if (scs(x, y.substring(1)).length() <= scs(x.substring(1), y).length()) {
20             return y.charAt(0) + scs(x, y.substring(1));
21         }
22         else {
23             return x.charAt(0) + scs(x.substring(1), y);
24         }
25     }
26
27     public static void main(String[] args) {
28         System.out.println(scs("aggtab", "gtxayb"));
29     }
30 }

```

AGGTAB
GXTXAYB

XAGGTXAYB
9

(A)GGTAB
GXTXAYB

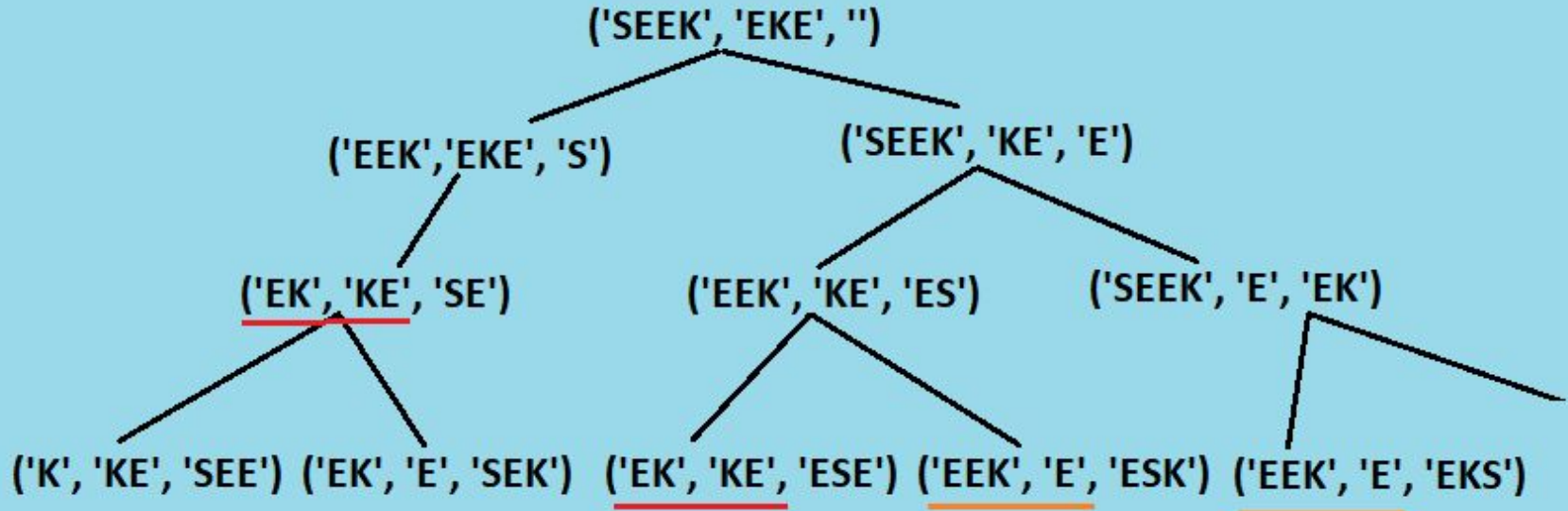
GXGTXAYB
8

X = "AGGTAB"

Y = "GXTXAYB"

Z = "AGXGTXAYB"

Recursive approach leads to overlapping problem:



Partial Recursion Tree of "SEEK" and "EKE" with repeated subproblems

-Dynamic Programming approach - TC: $O(m*n)$

S1 = "A**GGTAB**"



S2 = "**G**XT**X**A**Y****B**"



SCS = "A**G**X**G**T**X**A**Y****B**"



LCS = "**G**T**A****B**"



- Add LCS chars only one
- The first common char belongs to SCS
- Add non-lcs chars in order



	Y	A	B	C
X	0 (00)	1 (01)	2 (02)	3 (03)
A	1 (10)	1 (11)	2 (12)	3 (13)
D	2 (20)	2 (21)	3 (22)	4 (23)

X = "AD"

Y = "ABC"

Output = "ADBC"

Algorithm 2: SCS Dynamic Programming

Data: X, Y : arrays of characters, n : integer, m : integer

```

1  $M[0...n][0...m]$ ;
2 for  $i \leftarrow 0$  to  $n$  do
3    $M[i][0] \leftarrow i$ ;
4 end
5 for  $j \leftarrow 0$  to  $m$  do
6    $M[0][j] \leftarrow j$ ;
7 end
8 for  $i \leftarrow 1$  to  $n$  do
9   for  $j \leftarrow 1$  to  $m$  do
10    if  $X[i] == Y[j]$  then
11       $M[i][j] \leftarrow 1 + M[i-1][j-1]$ ;
12    else
13       $M[i][j] \leftarrow \min\{M[i-1][j], M[i][j-1]\} + 1$ ;
14    end
15  end
16 end
17 return  $M[n][m]$ ;

```

-Dynamic Programming approach - TC: $O(m*n)$

	Y	G	X	T	X	A	Y	B
X 	0	1	2	3	4	5	6	7
A	1	2	3	4	5	5	6	7
G	2	2	3	4	5	6	7	8
G	3	3	4	5	6	7	8	9
T	4	4	5	5	6	7	8	9
A	5	5	6	6	7	7	8	9
B	6	6	7	7	8	8	9	9

X = "AGGTAB"

Y = "GXTXAYB"

Z = "AGXGTXAYB"

```
1 package presentation;
2 import java.util.Scanner;
3 public class DP{
4
5     private static String SCS(String x, String y, int n, int m) {
6         int[][] t = new int[n+1][m+1];
7
8         for (int i = 1; i <= n; i++) {
9             for (int j = 1; j <= m; j++) {
10                 if (x.charAt(i-1) == y.charAt(j-1)) {
11                     t[i][j] = 1 + t[i-1][j-1];
12                 } else {
13                     t[i][j] = Math.max(t[i-1][j], t[i][j-1]);
14                 }
15             }
16         }
17     }
18 }
```

→ Creating 2D Array

```
15     StringBuilder sb = new StringBuilder();
16     int i = n, j = m;
17     while (i > 0 && j > 0) {
18         if (x.charAt(i-1) == y.charAt(j-1)) {
19             sb.insert(0, x.charAt(i-1));
20             i--;
21             j--;
22         } else {
23             if (t[i][j-1] > t[i-1][j]) {
24                 sb.insert(0, y.charAt(j-1));
25                 j--;
26             } else {
27                 sb.insert(0, x.charAt(i-1));
28                 i--;
29             }
30         }
31         while (i > 0) {
32             sb.insert(0, x.charAt(i-1));
33             i--;
34         }
35         while (j > 0) {
36             sb.insert(0, y.charAt(j-1));
37             j--;
38         }
39         return sb.toString();
40     }
41 }
```

```
37
38     public static void main(String[] args) {
39         Scanner input = new Scanner(System.in);
40         String x = input.next();
41         String y = input.next();
42         int n = x.length();
43         int m = y.length();
44         System.out.println(SCS(x,y,n,m));
45     }
46 }
```



Finding SCS of more than two strings:

Input Subsequences

str1 = FWCA
str2 = LOWRTS
str3 = FLAR
str4 = CHAS

Output

FLOWCHARTS

F WC A

LOW RTS

FL AR

CHA S



Thank you for your ATTENTION.