

Custom Search

COURSES

HIRE WITH US

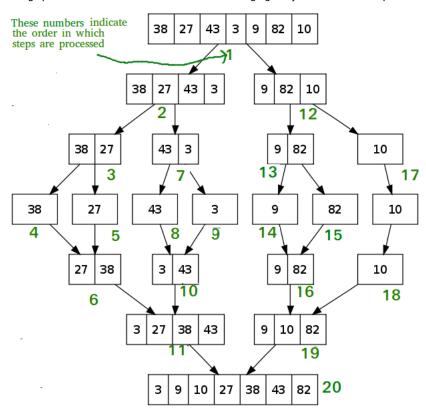


0

Merge Sort

Like QuickSort, Merge Sort is a Divide and Conquer algorithm. It divides input array in two halves, calls itself for the two halves and then merges the two sorted halves. **The merge() function** is used for merging two halves. The merge(arr, I, m, r) is key process that assumes that arr[I..m] and arr[m+1..r] are sorted and merges the two sorted sub-arrays into one. See following C implementation for details.

The following diagram from wikipedia shows the complete merge sort process for an example array {38, 27, 43, 3, 9, 82, 10}. If we take a closer look at the diagram, we can see that the array is recursively divided in two halves till the size becomes 1. Once the size becomes 1, the merge processes comes into action and starts merging arrays back till the complete array is merged.





Recommended: Please solve it on "PRACTICE" first, before moving on to the solution.

```
C/C++
/* C program for Merge Sort */
#include<stdlib.h>
#include<stdio.h>
// Merges two subarrays of arr[].
// First subarray is arr[l..m]
// Second subarray is arr[m+1..r]
void merge(int arr[], int 1, int m, int r)
    int i, j, k;
int n1 = m - 1 + 1;
    int n2 = r - m;
    /* create temp arrays */
    int L[n1], R[n2];
    /* Copy data to temp arrays L[] and R[] */
    for (i = 0; i < n1; i++)</pre>
         L[i] = arr[l + i];
    for (j = 0; j < n2; j++)
         R[j] = arr[m + 1+ j];
    /* Merge the temp arrays back into arr[1..r]*/
    i = 0; // Initial index of first subarray
    j = 0; // Initial index of second subarray
    k = 1; // Initial index of merged subarray
    while (i < n1 && j < n2)
         if (L[i] <= R[j])</pre>
             arr[k] = L[i];
             i++;
         }
         else
         {
             arr[k] = R[j];
             j++;
         k++;
    }
    /* Copy the remaining elements of L[], if there
        are any */
    while (i < n1)
         arr[k] = L[i];
         i++;
         k++;
    }
    /* Copy the remaining elements of R[], if there
       are any */
    while (j < n2)
    {
         arr[k] = R[j];
         j++;
        k++;
    }
}
/* 1 is for left index and r is right index of the
   sub-array of arr to be sorted */
void mergeSort(int arr[], int 1, int r)
    if (1 < r)
    {
         // Same as (1+r)/2, but avoids overflow for
         // large 1 and h
         int m = 1 + (r-1)/2;
         // Sort first and second halves
        mergeSort(arr, 1, m);
        mergeSort(arr, m+1, r);
        merge(arr, 1, m, r);
    }
}
```

```
/* UTILITY FUNCTIONS */
/* Function to print an array */
void printArray(int A[], int size)
{
    int i;
    for (i=0; i < size; i++)</pre>
       printf("%d ", A[i]);
    printf("\n");
}
/* Driver program to test above functions */
int main()
{
    int arr[] = {12, 11, 13, 5, 6, 7};
    int arr_size = sizeof(arr)/sizeof(arr[0]);
    printf("Given array is \n");
    printArray(arr, arr_size);
    mergeSort(arr, 0, arr_size - 1);
    printf("\nSorted array is \n");
    printArray(arr, arr_size);
    return 0;
}
Java
/* Java program for Merge Sort */
class MergeSort
    // Merges two subarrays of arr[].
    // First subarray is arr[1..m]
    // Second subarray is arr[m+1..r]
    void merge(int arr[], int 1, int m, int r)
        // Find sizes of two subarrays to be merged
        int n1 = m - l + 1;
        int n2 = r - m;
        /* Create temp arrays */
        int L[] = new int [n1];
        int R[] = new int [n2];
        /*Copy data to temp arrays*/
        for (int i=0; i<n1; ++i)</pre>
            L[i] = arr[l + i];
        for (int j=0; j<n2; ++j)</pre>
            R[j] = arr[m + 1+ j];
        /* Merge the temp arrays */
        // Initial indexes of first and second subarrays
        int i = 0, j = 0;
        // Initial index of merged subarry array
        int k = 1;
        while (i < n1 && j < n2)
        {
            if (L[i] <= R[j])</pre>
                 arr[k] = L[i];
                 i++;
            else
                 arr[k] = R[j];
                 j++;
             k++;
        }
        /* Copy remaining elements of L[] if any */
        while (i < n1)
        {
            arr[k] = L[i];
            i++;
```

```
k++;
        }
        /* Copy remaining elements of R[] if any */
       while (j < n2)
            arr[k] = R[j];
            j++;
            k++;
        }
   }
   // Main function that sorts arr[1..r] using
   // merge()
   void sort(int arr[], int 1, int r)
        if (1 < r)
            // Find the middle point
            int m = (1+r)/2;
            // Sort first and second halves
            sort(arr, 1, m);
            sort(arr , m+1, r);
            // Merge the sorted halves
            merge(arr, 1, m, r);
       }
   }
    /* A utility function to print array of size n */
   static void printArray(int arr[])
    {
        int n = arr.length;
        for (int i=0; i<n; ++i)</pre>
           System.out.print(arr[i] + " ");
        System.out.println();
   }
    // Driver method
   public static void main(String args[])
        int arr[] = {12, 11, 13, 5, 6, 7};
        System.out.println("Given Array");
        printArray(arr);
        MergeSort ob = new MergeSort();
        ob.sort(arr, 0, arr.length-1);
        System.out.println("\nSorted array");
        printArray(arr);
/* This code is contributed by Rajat Mishra */
```

Python3

```
# Python program for implementation of MergeSort
def mergeSort(arr):
    if len(arr) >1:
        mid = len(arr)//2 #Finding the mid of the array
        L = arr[:mid] # Dividing the array elements
        R = arr[mid:] # into 2 halves
        mergeSort(L) # Sorting the first half
        mergeSort(R) # Sorting the second half
        i = j = k = 0
        # Copy data to temp arrays L[] and R[]
        while i < len(L) and j < len(R):</pre>
            if L[i] < R[j]:
                arr[k] = L[i]
                i+=1
            else:
                arr[k] = R[j]
                j+=1
            k+=1
```

```
# Checking if any element was left
        while i < len(L):</pre>
            arr[k] = L[i]
            i+=1
            k+=1
        while j < len(R):
            arr[k] = R[j]
            i+=1
# Code to print the list
def printList(arr):
    for i in range(len(arr)):
        print(arr[i],end=" ")
# driver code to test the above code
if __name__ == '__main__':
    arr = [12, 11, 13, 5, 6, 7]
   print ("Given array is", end="\n")
    printList(arr)
   mergeSort(arr)
    print("Sorted array is: ", end="\n")
    printList(arr)
# This code is contributed by Mayank Khanna
```

Output:

```
Given array is
12 11 13 5 6 7
Sorted array is
5 6 7 11 12 13
```

Time Complexity: Sorting arrays on different machines. Merge Sort is a recursive algorithm and time complexity can be expressed as following recurrence relation.

$$T(n) = 2T(n/2) + \Theta(n)$$

The above recurrence can be solved either using Recurrence Tree method or Master method. It falls in case II of Master Method and solution of the recurrence is $\Theta(nLogn)$.

Time complexity of Merge Sort is $\Theta(nLogn)$ in all 3 cases (worst, average and best) as merge sort always divides the array into two halves and take linear time to merge two halves.

Auxiliary Space: O(n)

Algorithmic Paradigm: Divide and Conquer

Sorting In Place: No in a typical implementation

Stable: Yes

Applications of Merge Sort

1. Merge Sort is useful for sorting linked lists in O(nLogn) time. In the case of linked lists, the case is different mainly due to the difference in memory allocation of arrays and linked lists. Unlike arrays, linked list nodes may not be adjacent in memory. Unlike an array, in the linked list, we can insert items in the middle in O(1) extra space and O(1) time. Therefore merge operation of merge sort can be implemented without extra space for linked lists.

In arrays, we can do random access as elements are contiguous in memory. Let us say we have an integer (4-byte) array A and let the address of A[0] be x then to access A[i], we can directly access the memory at (x + i*4). Unlike arrays, we can not do random access in the linked list. Quick Sort requires a lot of this kind of access. In linked list to access i'th index, we have to travel each and every node from the head to i'th node as we don't have a continuous block of memory. Therefore, the overhead increases for quicksort. Merge sort accesses data sequentially and the need of random access is low.

- 2. Inversion Count Problem
- 3. Used in External Sorting



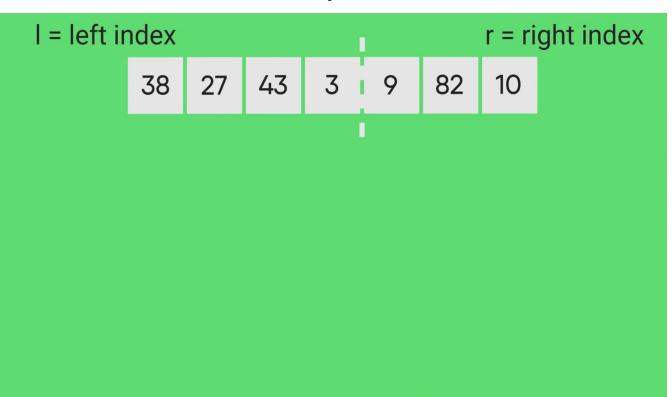
Snapshots:

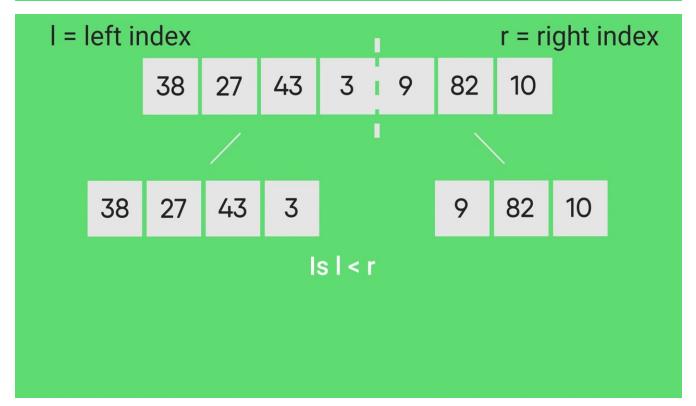
- Recent Articles on Merge Sort
- Coding practice for sorting.
- Quiz on Merge Sort

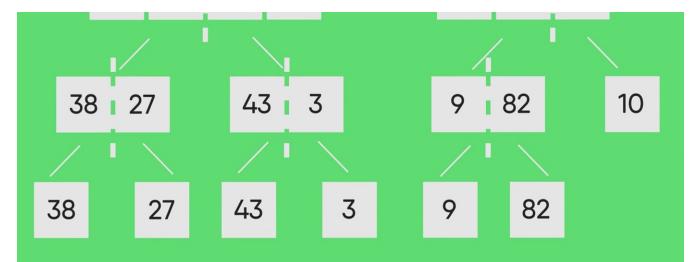
Other Sorting Algorithms on GeeksforGeeks:

3-way Merge Sort, Selection Sort, Bubble Sort, Insertion Sort, Merge Sort, Heap Sort, QuickSort, Radix Sort, Counting Sort, Bucket Sort, ShellSort, Comb Sort

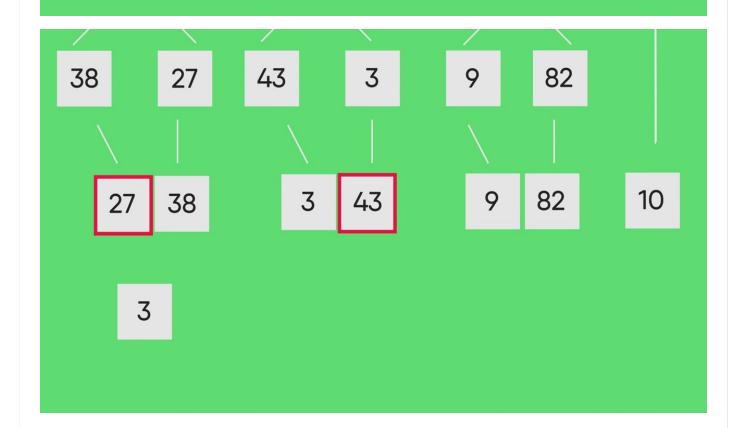
Please write comments if you find anything incorrect, or you want to share more information about the topic discussed above.

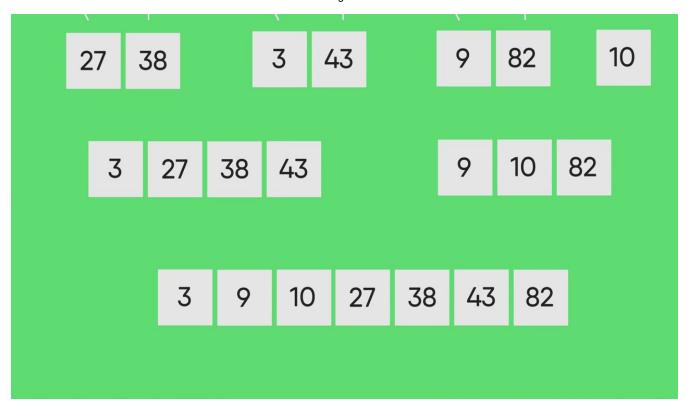






After dividing the array into smallest units merging starts, based on comparision of elements.





Recommended Posts:

Merge Sort with O(1) extra space merge and O(n lg n) time

Why Quick Sort preferred for Arrays and Merge Sort for Linked Lists?

Quick Sort vs Merge Sort

3-way Merge Sort

Iterative Merge Sort

In-Place Merge Sort

Merge Sort using Multi-threading

Merge Sort for Linked Lists

C Program for Iterative Merge Sort

Merge Sort for Linked Lists in JavaScript

Count Inversions in an array | Set 1 (Using Merge Sort)

Merge Sort for Doubly Linked List

Java Program for Iterative Merge Sort

Python Program for Iterative Merge Sort

Union and Intersection of two linked lists | Set-2 (Using Merge Sort)

Improved By: Ita_c, Mayank Khanna 2





Based on 217 vote(s)

Feedback/ Suggest Improvement



Improve Article

Please write to us at contribute@geeksforgeeks.org to report any issue with the above content.

Writing code in comment? Please use ide.geeksforgeeks.org, generate link and share the link here.

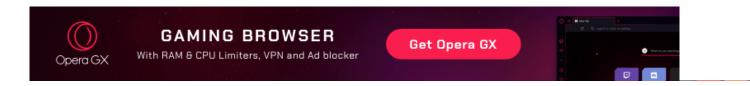
Important Update

When you log in with Disqus, we process personal data to facilitate your authentication and posting of comments. We also store the comments you post and those comments are immediately viewable and searchable by anyone around the world.

Please access our Privacy Policy to learn what personal data Disqus collects and your choices about how it is used. All users of our service are also subject to our Terms of Service.

Proceed







A computer science portal for geeks

5th Floor, A-118, Sector-136, Noida, Uttar Pradesh - 201305 feedback@geeksforgeeks.org

COMPANY

About Us Careers Privacy Policy Contact Us LEARN

Algorithms
Data Structures
Languages
CS Subjects
Video Tutorials

PRACTICE

Courses Company-wise Topic-wise How to begin? CONTRIBUTE

Write an Article Write Interview Experience Internships Videos

@geeksforgeeks, Some rights reserved