EXPERIMENT No. 4

Date of Performance:

Date of Submission:

Aim: To implement Election algorithm

Objectives/ Requirements:

- 1)Create a dynamic group
- 2) Create number of processes in the group
- 3) Calculate number of processes in the group.
- 4) Decide on the election algorithm (Optional)
- 5) Elect coordinator.

Theory:

Principle: Many distributed algorithms require that some process acts as a coordinator. The question is how to select this special process dynamically. Note: In many systems the coordinator is chosen by hand (e.g., file servers, DNS servers). This leads to centralized solutions => single point of failure.

- Doesn't matter which process does the job, just needto pick one
- Example: pick a master in Berkeley clocksynchronization algorithm
- Election algorithms: technique to pick a unique coordinator
- Assumption: each process has a unique ID
- Goal: find the non-crashed process with the highest ID

Types of election algorithm:

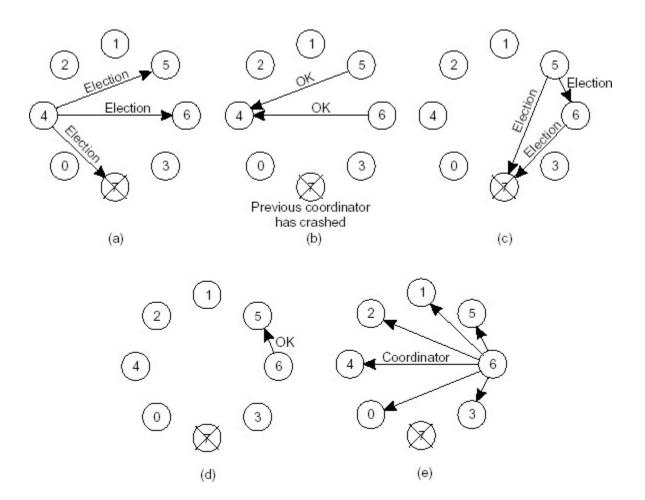
1) Bully's algorithm

Assumptions

- Each process knows the ID and address of every otherprocess
- Communication is reliable
- Need consistent result
- A process initiates an election if it just recovered from failure or it notices that the coordinator has failed
- Three types of messages: Election, OK, Coordinator.
- Several processes can initiate an election simultaneously.

Algorithm:

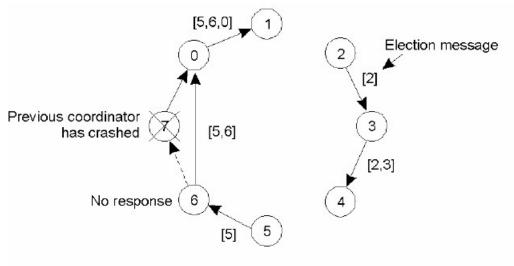
- Any process P can initiate an election
- P sends Election messages to all process with higher ID sand awaits OK messages
- $-\operatorname{If}$ no OK messages, P becomes coordinator and sends Coordinator messages to all processes with lower IDs
- If it receives an OK, it drops out and waits for an Coordinator message
- If a process receives an Election message
- Immediately sends Coordinator message if it is the process with highest ID
- Otherwise, returns an OK and starts an election
- If a process receives a Coordinator message, it treats sender as the coordinator.



2) Ring Algorithm:

Processes are arranged in a logical ring, each process knows the structure of the ring

- A process initiates an election if it just recovered from failure or it notices that the coordinator has failed
- Initiator sends Election message to closest downstream node that is alive
- Election message is forwarded around the ring
- Each process adds its own ID to the Election message
- When Election message comes back, initiator picks node with highest ID and sends a Coordinator message specifying the winner of the election
- Coordinator message is removed when it has circulated once.
- Multiple elections can be in progress



Assume n processes and one election in progress

- Bully algorithm
- Worst case: initiator is node with lowest ID
- Triggers n-2 elections at higher ranked nodes: O(n2) messages
- Best case: initiator is node with highest ID
- Immediate election: n-1 messages
- Ring algorithm
- 2n messages always

CODE:

Bully's algorithm

```
class BullyElectionAlgorithm:
    def __init__(self, nodes):
        self.nodes = nodes
        self.coordinator = None

def start_election(self, initiating_node):
        print(f"Node {initiating_node} initiates election.")
        higher nodes = [node for node in self.nodes if node > initiating_node]
```

```
if not higher nodes:
       self.coordinator = initiating node
       print(f"Node {self.coordinator} becomes the coordinator.")
    higher nodes.sort()
    for node in higher nodes:
       print(f"Node {initiating node} sends election message to Node {node}.")
    for node in higher nodes:
       print(f"Node {node} acknowledges the election.")
    for node in higher nodes:
       if self.check alive(node):
         self.coordinator = node
         print(f"Node {self.coordinator} becomes the coordinator.")
         return
  def check alive(self, node):
    # Placeholder function to check if node is alive
    return True
# Example usage:
nodes = [1, 2, 3, 4, 5]
election = BullyElectionAlgorithm(nodes)
election.start election(3)
```

Output:

```
Node 3 initiates election.

Node 3 sends election message to Node 4.

Node 3 sends election message to Node 5.

Node 4 acknowledges the election.

Node 5 acknowledges the election.

Node 4 becomes the coordinator.
```

Code:

Ring Algorithm

```
class Pro:
 def init (self, id):
  self.id = id
  self.act = True
class GFG:
 def init (self):
  self.TotalProcess = 0
  self.process = []
 def initialiseGFG(self):
  print("No of processes 5")
  self.TotalProcess = 5
  self.process = [Pro(i) for i in range(self.TotalProcess)]
 def Election(self):
  print("Process no " + str(self.process[self.FetchMaximum()].id) + " fails")
  self.process[self.FetchMaximum()].act = False
  print("Election Initiated by 2")
  initialized Process = 2
  old = initializedProcess
  newer = old + 1
  while (True):
   if (self.process[newer].act):
            print("Process " + str(self.process[old].id) + " pass Election(" +
str(self.process[old].id) + ") to" + str(self.process[newer].id))
     old = newer
   newer = (newer + 1) \% self.TotalProcess
   if (newer == initializedProcess):
     break
  print("Process " + str(self.process[self.FetchMaximum()].id) + " becomes coordinator")
  coord = self.process[self.FetchMaximum()].id
  old = coord
  newer = (old + 1) % self. Total Process
  while (True):
   if (self.process[newer].act):
      print("Process " + str(self.process[old].id) + " pass Coordinator(" + str(coord) + ")
message to process " + str(self.process[newer].id))
     old = newer
   newer = (newer + 1) \% self.TotalProcess
   if (newer == coord):
     print("End Of Election ")
     break
```

```
def FetchMaximum(self):
    maxId = -9999
    ind = 0
    for i in range(self.TotalProcess):
        if (self.process[i].act and self.process[i].id > maxId):
            maxId = self.process[i].id
        ind = i
        return ind

def main():
    object = GFG()
    object.initialiseGFG()
    object.Election()

if __name__ == "__main__":
    main()
```

Output:

```
No of processes 5
Process no 4 fails
Election Initiated by 2
Process 2 pass Election(2) to3
Process 3 pass Election(3) to0
Process 0 pass Election(0) to1
Process 3 becomes coordinator
Process 3 pass Coordinator(3) message to process 0
Process 0 pass Coordinator(3) message to process 1
Process 1 pass Coordinator(3) message to process 2
End Of Election
```

: Conclusion:

Successfully implemented Election algorithm.

SIGN AND REMARK

DATE

R1		R2	R3	R4	Total	Signature
(4 Mar	ks)	(4 Marks)	(4 Marks)	(3 Mark)	(15 Marks)	