

Testing Static Analyses for Precision and Soundness

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Static Analysis

- Sound
- Precise
- Fast

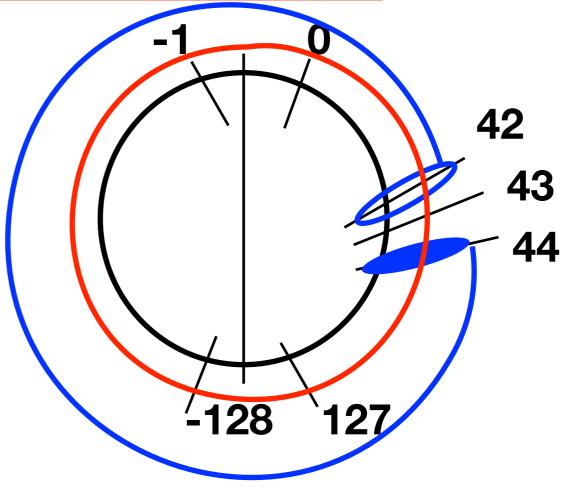
Motivation: Precision

```
%3 = select i1 %2, i8 1, i8 %x
```

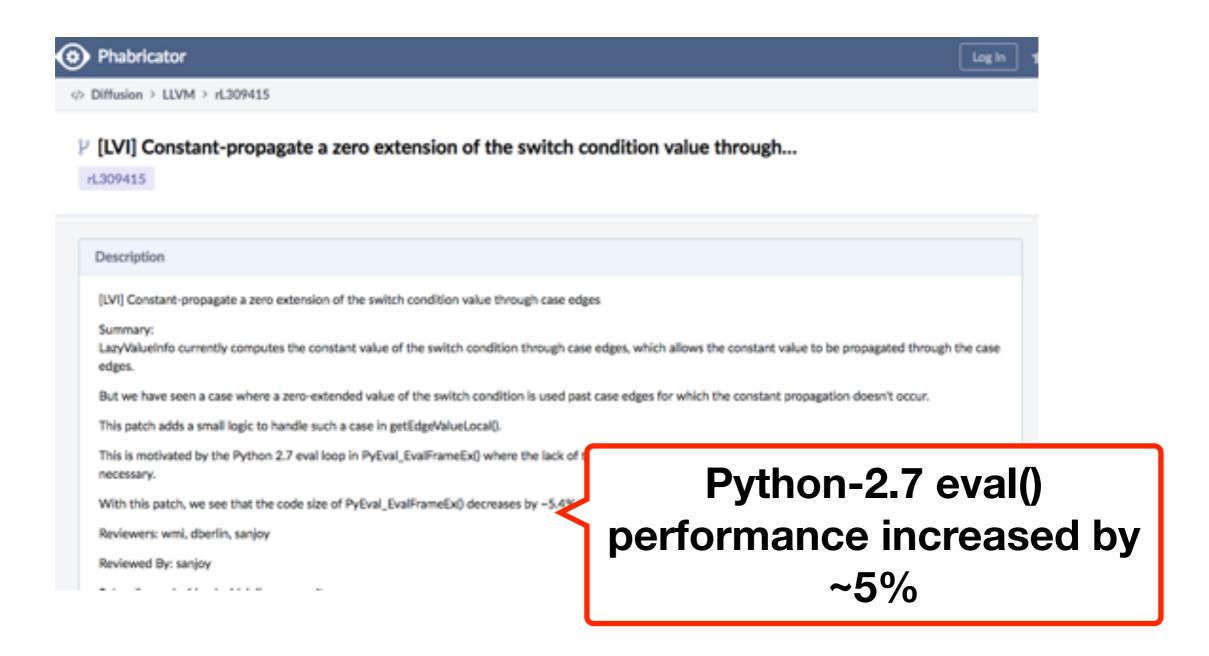
```
%4 = icmp eq i8 %3, 42 ret i1 %4
```

LLVM Result: No info

Best Result: [44, 42)



Imprecision in LLVM's Integer Range Analysis



Miscompilation Bug in LLVM

Bug 23011 - miscompile of % in loop

Status: RESOLVED FIXED

Nick Lewycky 2015-03-24 18:46:31 PDT

```
$ clang++ -v
clang version 3.7.0 (trunk 233044)
Target: x86_64-unknown-linux-gnu

Testcase:
#include <stdio.h>
#include <stdlib.h>
#include <string>

using namespace std;

int main(int argc, char **argv) {
   int r = 2;
   bool ok = true;
   while (ok) {
     string ab;
```

Motivation: Soundness

Q: Number of sign bits for %0?

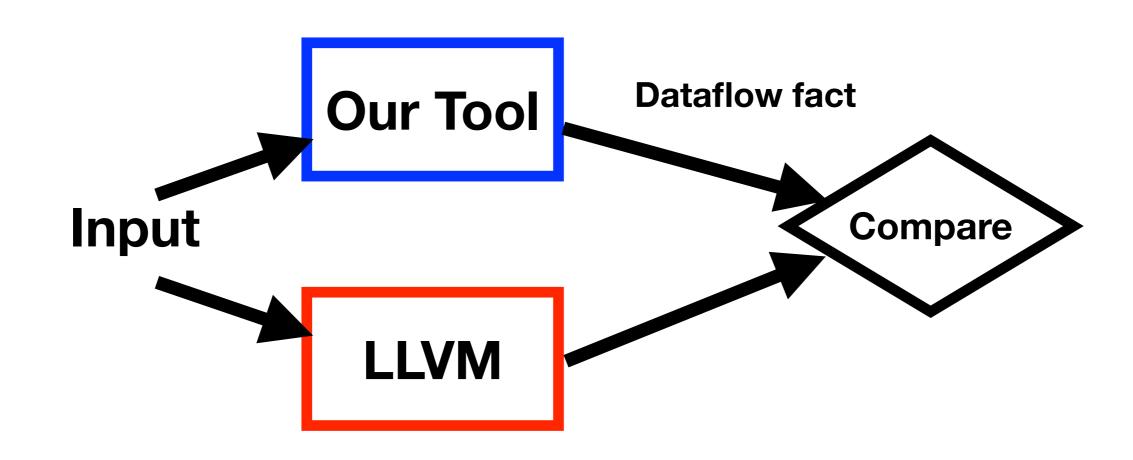
LLVM-3.7 Result: 31
Best Result: 30

Problems

- Imprecision
- Unsoundness
- Developers are manually improving the static analysis without any help of formal methods

Goal

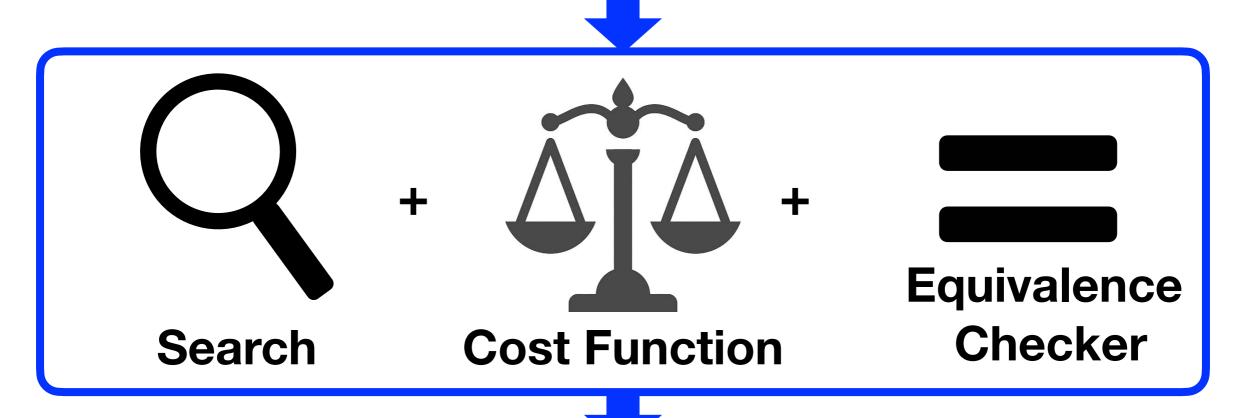
Automatic testing of LLVM's static analyses using formal methods



Souper

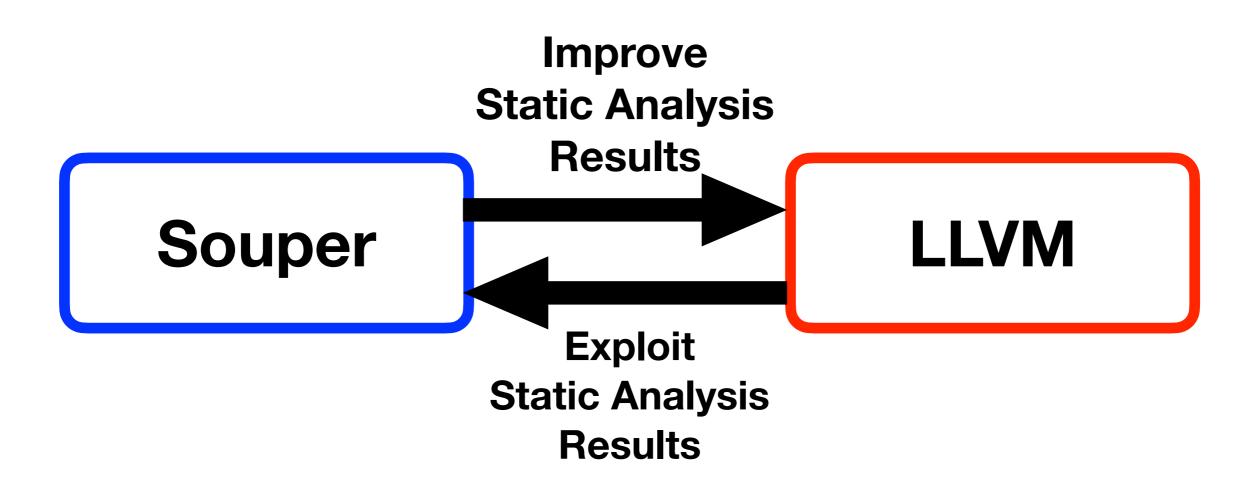
Automatically discover missing peephole optimizations

Source Code



Optimizations

Static Analysis and Souper



- Known Bits
- Integer Range
- Number of Sign Bits
- Non-zero
- Non-negative
- Negative
- Power of Two
- Demanded Bits

Forward

Backward

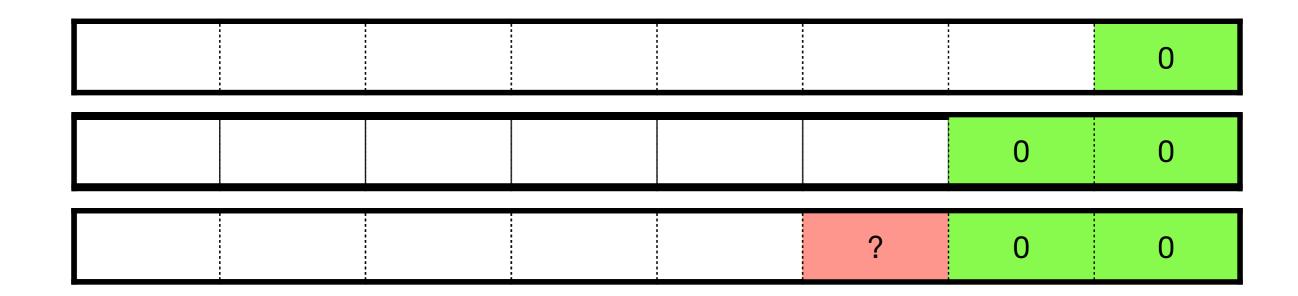
Known Bits

%0 = shl i8 %x, 4

%0 = shl i8 4, %x

? ? ? ? ? ? ?

Solver-based Algorithm to Compute Known Bits for 4 << x



more failing guesses ...

? ? ?	? ?	? 0	0
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%0 = shl i8 4, %x

LLVM:

? ? ? ? ? ? ?

Our Tool:

? ? ? ? ? 0 0

• Our algorithm uses at most | 2 * BitWidth solver calls

- Brute force algorithm uses 3^{BitWidth} calls, which is infeasible
- Computes maximally precise known bits as the lattice is separable at bit level
 - Details in the paper

Integer Range

```
define i4 @foo(i4 %y) {
  entry:
    %0 = mul i4 %y, %y
    ret i4 %0
}
```

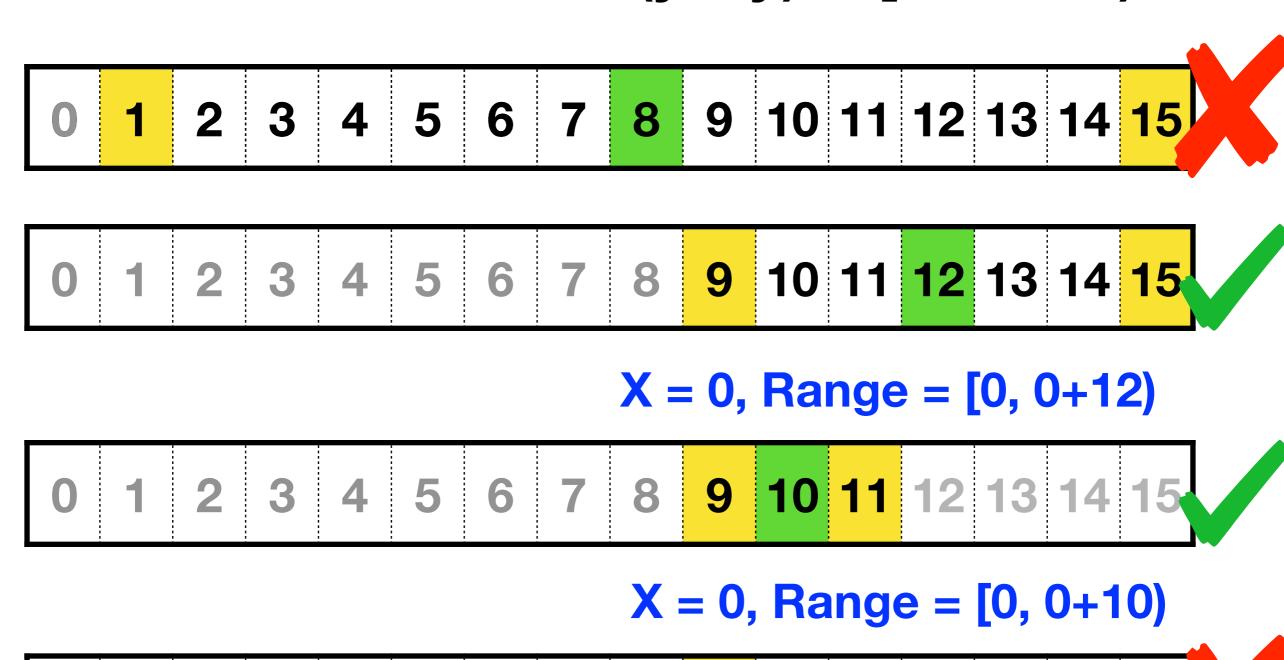
Range of %0?

LLVM: No information

Range [X, X+M)

- Algorithm for M Binary Search to find the smallest M
- Algorithm for X Solver-based
 Constant Synthesis

Find X such that $(y * y) \in [X, X+M)$



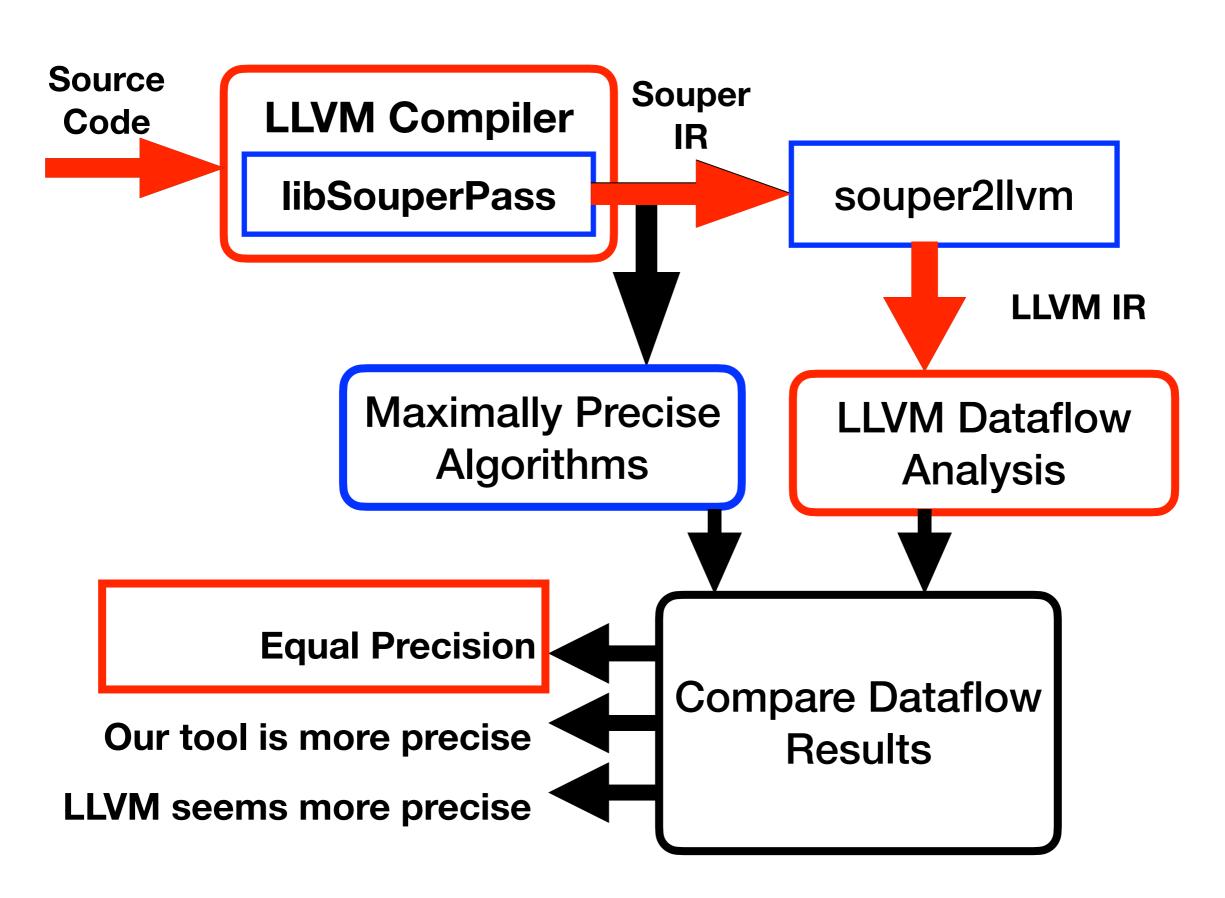
%0 = mul i4 %y, %y

LLVM: No Information

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15

Our Tool: [0, 10)

0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15



Impact on LLVM

- Five concrete precision improvements made in LLVM's static analyses (< version 8) already discussed in the paper.
- All integer range imprecisions discussed in the paper have already been fixed in LLVM 10.
- More known bits imprecisions have also been fixed in code generation phase.

What happens if LLVM calls our analyses instead of its own?

Too Slow!

Is LLVM Unsound?

- No new soundness bugs were found in LLVM+Clang-8.0
- Introduced three old soundness bugs from LLVM-2.9+ and our tool detected all of them

Conclusion

- Solver-based algorithms to compute maximally precise dataflow results to find imprecisions and unsoundness issues
- Encourage compiler developers to use formal methods based techniques to test static analyses



Backup Slides Constant Synthesis

 We use SMT Solvers to compute a constant X that satisfies the constraint.

								! !			! !				
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
	-			-			_								-

$$0*0=0$$
 $1*1=1$
 $2*2=4$
 $3*3=9$
 $4*4=16=0$
 $5*5=25=9$
 $6*6=36=4$
 $7*7=49=1$

Constant Synthesis

 Use SMT solver to compute a constant C that satisfies the given constraint.

Compute Integer Range for (y + y) & 1

