

Confidential

Development of a Versioned SQL Database System for Parameter Management in the Common Powertrain Controller

Entwicklung eines versionierten SQL-Datenbanksystems für das Parametermanagement im Common Powertrain Controller

Master's thesis

In the study program Mechanical Engineering

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Title: Development of a Versioned SQL Database System for Parameter

Management in the Common Powertrain Controller

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termanagement im Common Powertrain Controller

In current vehicles, the coordination of variant diversity plays an important role. For cost reasons, attempts are made to use common parts across model series and for different equipment variants within a model series. For control units, this approach also extends to the software, where possible. Different characteristics are then implemented using parameters. Parameter management is typically done using databases.

This master's thesis aims to develop a concept for an SQL database for managing parameters for a central control unit in vehicle applications. Various aspects shall be considered.

First, various versioning approaches shall be tested and evaluated. The goal is to document changes in a traceable manner.

Based on the requirements of software developers, various use cases shall be defined. Based on this, a concept for the tables within the database shall be developed.

Furthermore, it shall be described how the database can later be used by both software developers and development engineers. This also includes the development of a user interface. Finally, the limitations and restrictions of the developed approach should be described.

The results must be documented and presented.



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Preamble

Contents

| Acro | nyms | IX |
|----------------|---------------------------------------|----------------|
| Symb | pols | ΧI |
| List | of Figures | XIII |
| List | of Tables | ΧV |
| Kurzf | fassung | XVII |
| Abstı | ract | XIX |
| 1.6.2 | Introduction Background and Context | |
| 2 | Theoretical Background | 11 |
| 2.1 2.1.1 | Automotive Electronic Control Systems | 11 11 13 |
| 2.2 | Release and Phase Management | 13 15 15 |
| 2.3 | Non-Relational Database Systems | 17 18 18 |
| 2.3.2 2.3.3 | Entity-Relationship Modeling | 20 21 |
| 2.3.5 2.3.6 | Role-Based Access Control Models | 22 23 23 |

| 2.3.8 | Conceptual, Logical, and Physical Design Levels | 26 |
|-----------------|---|-----------|
| 3 | State of the Art | 27 |
| 3.1 | Parameter Management in Automotive Software Development | 27 |
| 3.1.1 | Evolution of Automotive Parameter Management | 27 |
| 3.1.2 | Challenges in Automotive Parameter Management | 28 |
| 3.1.3 | Current Approaches and Tools | 28 |
| 3.2 | Database Version Control Systems | 29 |
| 3.2.1 | Traditional Database Versioning Approaches | 29 |
| 3.2.2 | Temporal Database Approaches | 30 |
| 3.2.3 | Version Control for Parameter Management | 31 |
| 3.3 | Role-Based Access Control in Enterprise Systems | 32 |
| 3.3.1 | RBAC Model and Extensions | 32 |
| 3.3.2 | RBAC in Database Systems | 33 |
| 3.3.3 | Access Control for Parameter Management | 33 |
| 3.4 | Database Integration with Enterprise Systems | 34 |
| 3.4.1 | Enterprise Integration Patterns | 34 |
| 3.4.2 | Database Synchronization Approaches | 35 |
| 3.5 | Summary and Research Gaps | 35 |
| 4 | Mathadalawy and Canaant Pavalanment | 37 |
| 4 4.1 | Methodology and Concept Development | 37 |
| | Requirements Analysis | 37 |
| | Functional Requirements | 38 |
| | Integration with External Systems | 39 |
| | User Role Requirements | 39 |
| | Data Management Requirements | |
| 4.2 4.3 | Use Case Modeling | 40 41 |
| | | 41 |
| | Traditional Role-Based Approach | 43 |
| 4.3.2 4.4 | Hybrid Role-Permission Approach | 43 |
| | Parameter Synchronization Approaches | 44 |
| | Change-Based Synchronization Approach | 44 |
| | Phase-Based Synchronization Approach | |
| 4.5 | Database System Considerations | 46 |
| | Database System Requirements | 46 |
| | Comparative Analysis of Database Systems | 47 |
| 4.6 | Entity-Relationship Model | 47 |
| | Core Data Entities | 49 |
| | Relationship Structure | 50 51 |
| 4 ~ ~ | Narmalization and (Intimization | |

| 4.7 | Validation Mechanisms | 51 |
|-------|---|-----|
| 4.7.1 | Data Integrity Constraints | 51 |
| 4.7.2 | Business Rule Validation | 52 |
| 4.7.3 | Conflict Resolution Strategies | 52 |
| 4.7.4 | Audit and Traceability Mechanisms | 53 |
| 5 | Implementation | 55 |
| 5.1 | Database Structure Implementation | 55 |
| | Core Data Entities | 55 |
| | Version Control Implementation | 58 |
| 5.1.3 | Variant and Segment Management | 62 |
| 5.2 | Access Control Implementation | 68 |
| | Role-Based Permission Model | 68 |
| 5.2.2 | Module-Based Access Control | 71 |
| 5.3 | Query Optimization Implementation | 73 |
| 5.3.1 | Indexing Implementation | 73 |
| 5.3.2 | Function-Based Optimization | 74 |
| 5.4 | Change Tracking Implementation | 83 |
| 5.4.1 | Automatic Change Logging | 83 |
| 5.4.2 | Change Analysis Functions | 87 |
| 5.4.3 | Partitioning Implementation | 89 |
| 5.5 | Integration Implementation | 97 |
| 5.5.1 | Parameter Definition Database Synchronization | 97 |
| 5.5.2 | Vehicle Configuration Database Integration | 100 |
| 5.5.3 | Parameter File Generation | 104 |
| 6 | | 107 |
| 6.1 | Validation Methodology | 107 |
| 6.1.1 | Test Scenario Development | 107 |
| 6.1.2 | Performance Measurement Framework | 108 |
| 6.2 | 3 | 108 |
| 6.2.1 | User Management Validation | 108 |
| 6.2.2 | Module Access Impact on Performance | 111 |
| 6.2.3 | Release Management Validation | 112 |
| 6.2.4 | Variant Management Validation | 114 |
| 6.3 | Performance Analysis | 117 |
| 6.3.1 | Query Performance Assessment | 117 |
| 6.3.2 | Index Performance Analysis | 118 |
| 6.3.3 | Storage Requirements Analysis | 119 |
| 6.3.4 | Versioning Approach Performance | 121 |
| | | |

| 6.4 | Integration Testing | 122 |
|--------|--|-----|
| 6.4.1 | Parameter Definition Database Synchronization | 122 |
| 6.4.2 | Vehicle Configuration Integration | 124 |
| 6.5 | Feature Comparison with Excel-Based Approach | 124 |
| 6.5.1 | Data Integrity Improvements | 125 |
| 6.5.2 | Development Process Impact | 125 |
| 7 | Conclusion and Future Work | 127 |
| 7.1 | Summary of Contributions | 127 |
| 7.2 | Key Findings | 128 |
| 7.2.1 | Versioning Approach Effectiveness | 128 |
| 7.2.2 | Access Control Performance | 129 |
| 7.2.3 | Storage Distribution Insights | 129 |
| 7.2.4 | External System Integration Challenges | 130 |
| 7.3 | Limitations | 130 |
| 7.3.1 | Performance Limitations | 131 |
| 7.3.2 | Architectural Limitations | 131 |
| 7.3.3 | Integration Limitations | 132 |
| 7.4 | Future Work | 132 |
| 7.4.1 | Performance Optimizations | 132 |
| 7.4.2 | Architectural Enhancements | 134 |
| 7.4.3 | Integration Enhancements | 135 |
| 7.5 | Broader Implications | 136 |
| 7.5.1 | Implications for Database Research | 136 |
| 7.5.2 | Implications for Automotive Software Development | 137 |
| 7.6 | Conclusion | 137 |
| Biblio | ography | 139 |
| Α | User Management Test Cases | 143 |
| A.1 | Role-Based Permission Test Cases | 143 |
| A.2 | Module-Based Access Control Test Cases | 145 |
| A.3 | Direct Permission Assignment Test Cases | 146 |
| A.4 | Phase-Specific Permission Test Cases | 147 |
| A.5 | Boundary Case Test Cases | 147 |
| A.6 | Test Implementation Details | 148 |
| A.7 | Role Permission Matrix | 150 |
| В | Variant Management Test Cases | 151 |
| B.1 | Variant Creation Test Cases | 151 |
| B 2 | Segment Modification Test Cases | 152 |

| B.3 | Performance Test Cases | 155 |
|-----|--------------------------------|-----|
| B.4 | Test Implementation Details | 156 |
| B.5 | Test Environment Configuration | 161 |

Acronyms

CPC Common Powertrain Controller

ECU Electronic Control Unit

PDD Parameter Definition Database

PID Parameter ID

VCD Vehicle Configuration Database

VMAP Variant Management and Parametrization

Symbols

| Symbol | Unit | Description |
|------------|------|-------------|
| $F_{ m N}$ | N | Force |
| au | Nm | Torque |

List of Figures

| 1.1 | Gantt Chart of the planned work schedule | 9 |
|-----|--|-----|
| 2.1 | Hierarchical Organization of Automotive Electronic Systems | 12 |
| 2.2 | Automotive Parameter Release Cycle | 14 |
| 2.3 | Example of a Relational Schema [39] | 16 |
| 2.4 | Major Types of NoSQL Databases [17] | 17 |
| 2.5 | Simple Use Case Diagram for a Telephone System | 19 |
| 2.6 | Entity-Relationship Diagram for a Fitness Studio Database | 20 |
| 2.7 | Temporal Table Tracking Employee Department History | 24 |
| 4.1 | VMAP System Use Case Diagram | 41 |
| 4.2 | Traditional Role-Based Access Control Approach | 42 |
| 4.3 | Hybrid Role-Permission Access Control Approach | 43 |
| 4.4 | Change-Based Parameter Synchronization Approach | 44 |
| 4.5 | Phase-Based Parameter Synchronization Approach | 45 |
| 4.6 | Comprehensive Entity-Relationship Diagram for the Variant Management | |
| | and Parametrization (VMAP) System | 49 |
| 6.1 | Impact of Module-Specific Access Checks | 112 |
| 6.2 | Query Performance With and Without Indexes | 118 |
| 6.3 | Storage Distribution by Entity Type | 119 |
| 6.4 | Version-based vs. Phase-based Parameter Management Comparison . | 121 |
| 6.5 | Parameter Definition Database Synchronization Time Trends | 123 |

List of Tables

| 4.1 | Comparison of Database Systems for Automotive Parameter Management 48 |
|-------------|---|
| 6.1 | Sample Module Developer Role Permission Test Cases 109 |
| 6.2 | User Management Test Results |
| 6.3 | Phase Transition Test Results |
| 6.4 | Phase Freeze Protection Test Cases |
| 6.5 | Variant Operation Performance Metrics |
| 6.6 | Segment Operation Performance |
| 6.7 | Query Performance Comparison |
| 6.8 | Storage Requirements Analysis |
| 6.9 | Parameter Definition Database Synchronization Results |
| 6.10 | Feature Comparison with Excel-Based Approach |
| A.1 | Administrator Role Permission Test Cases |
| A.2 | Module Developer Role Permission Test Cases |
| A.3 | Documentation Team Role Permission Test Cases |
| A.4 | Read-Only User Role Permission Test Cases |
| A.5 | Module-Based Access Control Test Cases |
| A.6 | Direct Permission Assignment Test Cases |
| A .7 | Phase-Specific Permission Test Cases |
| 8.A | Boundary Case Test Cases |
| A.9 | Role Permission Matrix |
| B.1 | Variant Creation Test Cases |
| B.2 | Segment Creation Test Cases |
| B.3 | Segment Update Test Cases |
| B.4 | Segment Deletion Test Cases |
| B.5 | Variant and Segment Performance Test Cases |
| | |

Kurzfassung

Moderne Nutzfahrzeuge basieren auf hochentwickelten elektronischen Steuergeräten (ECUs), die tausende konfigurierbare Parameter zur Steuerung von Antriebsstrangfunktionen benötigen. Der derzeitige Excel-basierte Ansatz für das Parametermanagement in der Automobilsoftwareentwicklung birgt erhebliche Risiken, darunter Dateninkonsistenzen, Versionskonflikte und eingeschränkte Nachvollziehbarkeit [41]. Diese Arbeit adressiert diese Herausforderungen durch die Entwicklung einer umfassenden Datenbankarchitektur für Variantenmanagement und Parametrisierung (VMAP) in PostgreSQL.

Die Forschung implementiert einen phasenbasierten Versionierungsansatz, der an den Entwicklungszyklen der Automobilindustrie ausgerichtet ist und gegenüber traditionellen änderungsbasierten Versionierungsmodellen erhebliche Leistungsvorteile bietet, während die Datenintegrität gewährleistet bleibt. Empirische Tests zeigen, dass dieser Ansatz Abfrageantwortzeiten unter 100ms selbst bei Datensätzen mit über 100.000 Parametern erzielt, bei akzeptabler Speichereffizienz [4]. Das hybride Rollen-Berechtigungszugriffsmodell des Systems kombiniert rollenbasierte Sicherheit mit modulspezifischen Zugriffskontrollen, was durch umfangreiche Benutzervalidierung mit realen Engineering-Workflows verifiziert wurde.

Die Integration mit Unternehmenssystemen ermöglicht synchronisierte Parameterdefinitionen und Fahrzeugkonfigurationsdaten, entsprechend den von Hohpe und Woolf [18] beschriebenen Enterprise-Integration-Patterns. Die Leistungsanalyse ergab, dass die Datenbank interaktive Antwortzeiten für die meisten Operationen auch bei Produktionsdatensätzen mit 830 Varianten und 167.990 Segmenten beibehält, wobei bestimmte komplexe Operationen wie Phasenvergleiche eine suboptimale Skalierung aufweisen, die von weiterer Optimierung profitieren würden.

Das resultierende System bietet eine 6,5- bis 21,8-fache Leistungsverbesserung gegenüber nicht-indexierten Implementierungen und löst kritische Einschränkungen des Excel-basierten Ansatzes, einschließlich gleichzeitigem Mehrbenutzer-Zugriff, automatisierter Validierung und umfassender Änderungsverfolgung. Die Leistungsanalyse zeigt, dass, obwohl Änderungshistoriendatensätze den Speicherbedarf dominieren (60.8

Abstract

Modern commercial vehicles rely on sophisticated electronic control units (ECUs) that require thousands of configurable parameters to manage powertrain functions. The current Excel-based approach to parameter management in automotive software development introduces significant risks including data inconsistency, version conflicts, and limited traceability [41]. This thesis addresses these challenges through the development of a comprehensive database architecture for variant management and parametrization (VMAP) in PostgreSQL.

The research implements a phase-based versioning approach aligned with automotive development cycles, providing significant performance advantages over traditional change-based versioning models while maintaining data integrity. Empirical testing demonstrates that this approach yields query response times below 100ms even with datasets exceeding 100,000 parameters, while maintaining acceptable storage efficiency [4]. The system's hybrid role-permission access control model combines role-based security with module-specific access controls, verified through extensive user management validation with real-world engineering workflows.

Integration with enterprise systems enables synchronized parameter definitions and vehicle configuration data, implementing the enterprise integration patterns described by Hohpe and Woolf [18]. Performance analysis revealed that the database maintains interactive response times for most operations even with production-scale datasets containing 830 variants and 167,990 segments, though certain complex operations like phase comparison exhibit suboptimal scaling that would benefit from further optimization.

The resulting system provides a 6.5x-21.8x performance improvement over non-indexed implementations and resolves critical limitations of the Excel-based approach, including multi-user concurrent access, automated validation, and comprehensive change tracking. Performance analysis demonstrates that while change history records dominate storage requirements (60.8% of database size), the storage-performance tradeoff effectively supports the audit requirements common in regulated automotive development environments [37].

1 Introduction

Modern commercial vehicles represent a quintessential example of cyber-physical systems, where sophisticated software enables precise control over complex mechanical components. The software controlling these vehicles has grown exponentially in complexity over recent decades, evolving from simple engine management to comprehensive control of virtually all vehicle functions. At the core of this evolution is the Electronic Control Unit (ECU)—a specialized computer that executes software to manage specific vehicle functions [8]. Contemporary commercial vehicles contain dozens of interconnected Electronic Control Units (ECUs) working in concert to ensure optimal performance, efficiency, and safety across diverse operating conditions.

1.1 Background and Context

The automotive industry has undergone a profound transformation over the past decades, evolving from predominantly mechanical systems to highly sophisticated mechatronic platforms [29]. This evolution has been particularly pronounced in the commercial vehicle sector, where modern trucks rely on complex networks of Electronic Control Units to manage everything from engine performance to safety systems [8]. These systems must adapt to a wide range of operational conditions, regulatory requirements, and market-specific configurations, creating a significant challenge in managing software variability.

At the heart of this variability management lies the Common Powertrain Controller (Common Powertrain Controller (CPC))—a central ECU managing critical functions related to engine and transmission control. The CPC's operation is governed by thousands of configurable parameters that determine how the powertrain behaves under specific conditions [37]. These parameters influence everything from basic engine timing to sophisticated emission control strategies, making their precise configuration essential for vehicle performance, efficiency, and regulatory compliance.

The parameter management challenge is further complicated by the global nature of modern vehicle development. Commercial vehicles must conform to different emissions regulations, operate in diverse environmental conditions, and meet varying customer expectations across global markets. Consequently, a single vehicle model may require numerous parameter configurations, each tailored to specific combinations of market requirements, hardware configurations, and customer specifications [41].



1.2 Problem Statement

The current approach to parameter management in commercial vehicle development relies predominantly on distributed Excel spreadsheets, a methodology that emerged during a period when parameter counts were manageable and development teams were smaller [41]. However, as software complexity has increased exponentially, this fragmented approach has introduced significant limitations and risks to the development process.

Development teams distributed across different locations must coordinate changes to thousands of parameters, track their versions, and ensure consistency across multiple vehicle platforms. The absence of a centralized version control system makes it exceptionally difficult to track changes effectively and manage releases. This situation becomes particularly critical when dealing with safety-critical parameters that directly influence vehicle performance and regulatory compliance.

The manual nature of current processes, combined with the lack of automated validation mechanisms, introduces substantial risks of data inconsistency, version conflicts, and delayed implementation of critical parameter updates. Parameter changes are not consistently verified against established rules and constraints, potentially leading to incompatible configurations or non-compliant behavior [37].

Integration with critical enterprise systems presents another significant challenge. The current process of synchronizing data with internal database systems involves several manual steps, consuming valuable development resources and introducing potential points of failure in the configuration management workflow. The absence of automated data validation and synchronization mechanisms creates additional risks for data integrity and consistency across these interconnected systems.

Furthermore, the increasing emphasis on rapid development cycles and continuous integration in the automotive industry demands a more sophisticated approach to parameter management [8]. The existing system's limitations become particularly apparent when considering the need for simultaneous development of multiple vehicle variants, each requiring specific parameter configurations for different markets and regulatory environments.

These challenges collectively underscore the urgent need for a modern, databasedriven solution that can address the complexities of contemporary automotive software development while providing a scalable foundation for future growth and adaptation.



1.3 Research Objectives

This thesis aims to address the fundamental challenges in automotive parameter management through the development of database architecture for Variant Management and Parametrization (VMAP) (Variant Management and Parametrization), a web-based application for powertrain parameter configuration. The research objectives encompass both theoretical foundations and practical implementation considerations, focusing on creating a robust solution that meets the complex demands of modern vehicle development processes.

The primary research objective centers on developing a centralized database architecture that can effectively manage the complexity of powertrain parameters while maintaining data integrity and traceability [42]. This architecture must support sophisticated version control mechanisms that can handle parameter variations across different development stages and vehicle variants. The system should provide comprehensive audit trails and change history, enabling development teams to track modifications and understand the evolution of parameter configurations over time.

A second crucial objective focuses on the implementation of a sophisticated version control system that addresses the unique requirements of parameter management in automotive software development. This system must go beyond traditional source code version control approaches to handle the complex relationships between parameters, their variants, and their applications across different vehicle platforms [37]. The version control mechanism should support parallel development streams while maintaining consistency and preventing conflicts in parameter configurations.

The research also aims to establish a comprehensive role-based access control system that supports the diverse needs of different user groups within the development process. This includes creating specialized interfaces and permissions for Module Developers, Documentation Team members, Administrators, and Read-only Users, each with specific capabilities and restrictions aligned with their responsibilities [32]. The access control system must balance security requirements with the need for efficient collaboration among development teams.

Integration with existing enterprise systems represents another critical objective of this research. TheVMAP system must establish seamless data exchange mechanisms with internal database systems, ensuring consistent information flow while minimizing manual intervention [8]. This integration should support automated validation of parameter changes and provide mechanisms for maintaining data consistency across different systems.

A final key objective involves the development of database interfaces and query optimization strategies that will support the web-based interface implementation. While the actual User Interface (UI) development falls outside the scope of the thesis, the research will focus on designing efficient database structures, stored procedures, and APIs that enable seamless integration with the planned web interface [29]. This includes developing optimized query patterns for complex operations such as parameter comparison, variant management, and release workflows, while ensuring robust data validation and business rule enforcement.

1.4 Significance of the Study

The significance of this research extends beyond addressing immediate technical challenges in parameter management. By developing a comprehensive database solution for variant management and parametrization, this work contributes to the broader field of automotive software engineering in several important ways.

First, the research advances the understanding of version control in parameter-centric systems, extending traditional concepts of software versioning to accommodate the unique characteristics of automotive parameter configurations. While considerable research has been conducted on code versioning, the versioning of parameter data presents distinct challenges that require specialized approaches [4]. This thesis contributes to closing this gap by developing and evaluating new methods for parameter versioning in complex automotive systems.

Second, the work addresses critical industry needs for improved quality and efficiency in vehicle development. Commercial vehicle manufacturers face increasing pressure to reduce development time while managing growing software complexity and ensuring regulatory compliance across global markets [8]. By providing a more robust and efficient parameter management solution, this research directly contributes to these industry priorities, potentially reducing development costs and improving vehicle quality through more consistent parameter configurations.

Third, the research advances the integration of database technology with domainspecific engineering processes. By developing specialized database structures and functions tailored to the unique requirements of automotive parameter management, this work demonstrates how database technology can be adapted to support complex engineering workflows [14]. This integration perspective is valuable not only for automotive applications but also for other engineering domains facing similar challenges in managing complex, highly variable system configurations.



Finally, the research contributes to the growing field of model-based systems engineering by providing a structured approach to managing the parametric aspects of system models. As the automotive industry continues to adopt model-based approaches for system development, the management of parameter configurations becomes increasingly critical for maintaining model integrity and traceability [37]. This thesis provides insights and solutions that support this evolution toward more systematic model-based development practices.

1.5 Thesis Structure

The thesis is organized into six chapters that systematically address the research objectives and present a comprehensive solution for automotive parameter management. The structure follows a logical progression from theoretical foundations through practical implementation, ensuring thorough coverage of both academic and industry perspectives.

Following this introduction, Chapter 2 presents a comprehensive review of the theoretical background relevant to automotive electronic control systems and database management approaches. This chapter examines the hierarchical organization of automotive electronic systems, explores database management systems with a focus on relational databases, investigates database design methodologies including entity-relationship modeling and normalization, and evaluates access control models and version control concepts applicable to parameter management.

Chapter 3 details the methodology and concept development, beginning with a thorough requirements analysis based on industry needs and academic best practices. This chapter explores the conceptual architecture design through use case modeling, compares different approaches for user management and parameter synchronization, and presents a comprehensive entity-relationship model for the system. Particular attention is given to validation mechanisms and integration approaches with existing enterprise systems.

Chapter 4 presents the implementation strategy and technical design of theVMAP system. This chapter describes the practical realization of the database structure, detailing the core data entities, version control implementation, variant and segment management, access control mechanisms, query optimization approaches, change tracking implementation, and integration with external systems. The chapter provides concrete code examples showing the translation of conceptual designs into functional database components.



Chapter 5 focuses on system evaluation and validation, presenting a comprehensive assessment of theVMAP system against the defined research objectives. This chapter describes the validation methodology used to evaluate the system, presents detailed functional testing results for user management, release management, and variant management capabilities, analyzes the system's performance characteristics including query optimization and storage requirements, and evaluates the versioning approach and integration capabilities.

The thesis concludes with Chapter 6, which summarizes the research findings, discusses the limitations of the current implementation, and presents recommendations for future development. This chapter reflects on the contributions of the research to both academic knowledge and industry practice, presenting both technical optimizations for future enhancement and broader implications for database research and automotive software development.

Throughout these chapters, the research methodology combines theoretical analysis with practical implementation, ensuring that the resulting system meets both academic standards and industry requirements. Special attention is given to database versioning approaches, user role management, and integration strategies with existing systems, addressing the unique challenges of automotive software configuration management [37].

1.6 Project Plan

The research project follows a structured approach spanning six months from November 2024 to April 2025, organized into three distinct phases: Exposé, Implementation, and Finalization. The comprehensive timeline ensures systematic progression through all research objectives while maintaining academic rigor and quality standards.

1.6.1 Exposé Phase (November - January)

The initial phase focuses on establishing strong theoretical foundations and gathering comprehensive requirements. Literature review constitutes a significant portion of this phase, extending over six weeks to ensure thorough coverage of current database versioning approaches, parameter management systems, and industry practices in automotive applications. This review encompasses analysis of existing version control systems, examination of industry standards for software configuration management, and evaluation of current parameter management solutions.



Requirements analysis follows the literature review, spanning three weeks to capture detailed system specifications. This phase involves extensive stakeholder consultation to document system requirements, analyze existing Excel-based workflows, define integration requirements with internal database systems, and establish user roles and access control specifications. The Exposé phase concludes with the submission of a comprehensive research proposal at the end of Week 3 in January.

1.6.2 Implementation Phase (December - March)

The implementation phase encompasses four major components, each allocated four weeks for development and refinement. Database design initiates this phase, focusing on developing the schema for parameter management, designing version control mechanisms, creating data models for user management, and planning integration interfaces with existing systems.

System architecture development follows, concentrating on overall system design, version control workflows, user management frameworks, and validation mechanisms. This stage establishes the foundational structure for the entire system while ensuring alignment with identified requirements and industry standards.

Prototype development constitutes the third component, involving implementation of core database functionality, development of version control features, creation of user management interfaces, and construction of system integration components. This stage transforms theoretical designs into practical implementations while maintaining focus on system usability and performance.

The final component of this phase involves comprehensive testing and validation, including database performance testing, validation of version control mechanisms, testing of user management functions, and verification of system integration capabilities. This stage ensures all implemented features meet specified requirements and performance standards.

1.6.3 Finalization Phase (April)

The concluding phase focuses on documentation and thesis preparation over four weeks. The first two weeks are dedicated to comprehensive documentation, including compilation of implementation details and system architecture documentation.

The subsequent two weeks concentrate on thesis writing, involving comprehensive documentation of research findings, inclusion of test results and analysis, preparation

of conclusions and recommendations, and thorough content review and refinement. The phase concludes with thesis submission in Week 16 and final project presentation in Week 17.

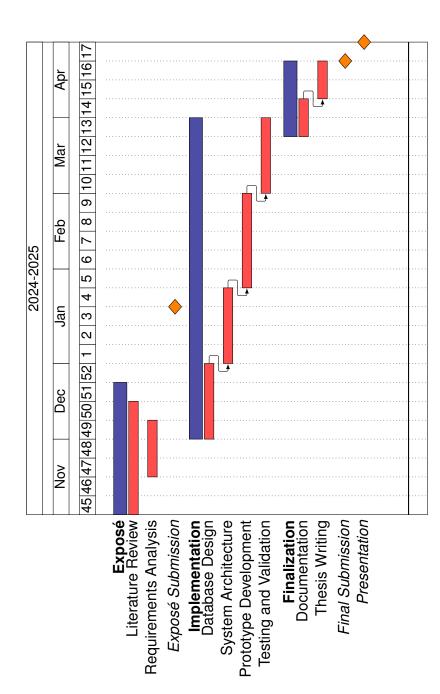


Figure 1.1: Gantt Chart of the planned work schedule.

2 Theoretical Background

This chapter establishes the theoretical foundation necessary for understanding the design and implementation of the Variant Management and Parametrization (VMAP) database system. It begins with an overview of automotive electronic control systems and parameter management, explaining the fundamental concepts that drive the requirements for the VMAP system. The chapter then explores database management systems, database design methodologies, and access control models relevant to the implementation. Finally, it discusses version control concepts and temporal database management approaches, which are critical for the parameter versioning requirements in automotive software development.

2.1 Automotive Electronic Control Systems

Modern commercial vehicles contain dozens of Electronic Control Units (ECUs) that manage various vehicle subsystems. Each ECU is a specialized computing device that controls specific functions through software parameters [37]. Understanding the structure and organization of these systems is essential for designing an effective parameter management solution.

2.1.1 ECU Hierarchy and Parameter Organization

Automotive electronic systems follow a hierarchical organization that structures parameters into logical groupings. Figure 2.1 illustrates this hierarchical structure, showing how parameters are organized in automotive electronic systems. At the top level, Electronic Control Units (ECUs) represent distinct hardware components controlling specific vehicle functions such as engine management, transmission control, or brake systems [22]. Within each ECU, modules represent functional software units that implement specific capabilities such as cruise control, adaptive power steering, diagnosis. Each module contains Parameter IDs (PIDs) that group related parameters, and finally, individual parameters define specific configuration values that affect system behavior [38].

As shown in Figure 2.1, parameters can either use their default values or be customized through variants and segments. When a parameter is used without customization, its default value is applied. However, when specific vehicle configurations require different

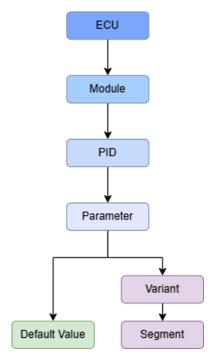


Figure 2.1: Hierarchical Organization of Automotive Electronic Systems

parameter settings, variants are created, and segments define the modified parameter values within these variants. This hierarchical structure is not merely organizational but reflects the actual architecture of automotive electronic systems, where software components are modularized for maintainability, reusability, and functional separation. The Common Powertrain Controller (CPC), a central ECU in modern trucks, manages critical powertrain functions through thousands of configurable parameters organized into this hierarchical structure [37].

Parameters themselves have complex characteristics beyond simple values. They can be scalar values, one-dimensional arrays (curves), or multi-dimensional arrays (maps or tables). Each parameter has specific attributes defining its data type, valid range, engineering units, scale factors, and default values [29]. For instance, an engine timing map might be represented as a two-dimensional array where engine speed and load are the independent variables, and ignition timing angle is the dependent variable. This complexity in parameter structure creates specific requirements for the database system designed to manage them.



2.1.2 Parameter Variants and Customization

A fundamental challenge in automotive parameter management is supporting multiple parameter configurations for different vehicle variants, regional requirements, and operating conditions. Rather than maintaining separate complete parameter sets for each configuration, which would lead to significant redundancy, automotive systems implement a variant mechanism that allows selective overriding of parameter values based on specific conditions [37].

In this approach, each parameter has a default value defined in the baseline configuration. Variants are created to represent specific vehicle configurations or conditions, and segments define modified parameter values within these variants. If no segment exists for a particular parameter in an applicable variant, the default value is used. This approach minimizes redundancy by storing only the modified values rather than complete parameter sets for each configuration [8].

The lower part of Figure 2.1 illustrates this concept, showing how parameters can follow two paths: either using their default values (left branch) or being customized through variants and segments (right branch). Variants are associated with code rules—boolean expressions that determine when a variant applies based on vehicle configuration codes. For example, a variant might apply only to vehicles with a specific engine type and transmission combination, or to vehicles destined for a particular market with unique regulatory requirements. The code rule evaluation process selects the appropriate variants for a specific vehicle configuration during parameter file generation [38].

This variant approach creates specific requirements for the database system, which must efficiently store and retrieve variant definitions and segment values while maintaining the relationships between parameters, variants, and segments. The system must also implement a parameter resolution process that correctly applies variants based on vehicle configuration codes, ensuring that the right parameter values are used for each specific vehicle.

2.1.3 Release and Phase Management

Automotive software development follows a structured release process with well-defined phases representing increasing levels of maturity and stability [8]. For parameter management, this translates into a phase-based development process where parameter configurations evolve through sequential stages before being released for production.

Figure 2.2: Automotive Parameter Release Cycle

Figure 2.2 illustrates the typical release cycle in automotive parameter development, consisting of bi-annual releases (e.g., "24.1" and "24.3" for first and third quarters of 2024), with each release progressing through four sequential phases: Phase1 (Initial definition & Configuration), Phase2 (Initial Testing), Phase3 (Refined Testing), and Phase4 (Ready for Production). The diagram shows the linear progression through these phases, with each phase building upon the work completed in the previous phase. Different ECUs may progress through these phases at different rates, requiring the parameter management system to support concurrent work on multiple phases.

Each phase represents a milestone in the development process with specific activities and quality gates. The Phase1 involves the creation of new parameters and initial configuration. Phase2 and Phase3 involve refinement based on testing feedback, with increasing levels of validation. The Phase4 represents the completed configuration ready for production release [37].

When a phase transitions to the next stage, parameter configurations are copied forward, establishing a new baseline for continued development. Changes made in earlier phases should propagate to later phases unless explicitly overridden, creating a complex versioning requirement for the parameter management system [29]. Additionally, at specific development milestones, phases may be "frozen" to create stable reference points for documentation and testing, requiring the parameter management system to enforce read-only access to frozen phases while still allowing continued development in active phases.

This phase-based release process establishes specific requirements for the database system's versioning model, which must maintain distinct parameter configurations for each phase while supporting phase transitions, change propagation, and selective freezing. The versioning approach must align with this development process rather than implementing a generic temporal model, ensuring that the system supports the actual workflows used in automotive parameter development.



2.2 Database Management Systems

Database management systems (DBMS) serve as the foundation for structured information storage and retrieval. They provide mechanisms for storing, organizing, and accessing data while ensuring integrity, security, and concurrent access [14]. For the VMAP system, selecting an appropriate database approach is critical for meeting the complex requirements of automotive parameter management.

2.2.1 Relational Database Management Systems

Relational Database Management Systems (RDBMS) organize data into structured tables composed of rows and columns, based on the relational model proposed by E.F. Codd in 1970 [10]. The relational model establishes a mathematical foundation for representing data as relations (tables) with well-defined operations for data manipulation. This approach has dominated database technology for decades due to its solid theoretical foundation and practical advantages for structured data management.

Figure 2.3 illustrates a practical example of a relational schema with two related tables. The Product Table contains product information with ModelNumber as the primary key, while the Sales Table captures order information with OrderNumber as its primary key. The relationship between these tables is established through the ModelNumber field in the Sales Table, which serves as a foreign key referencing the Product Table. This one-to-many relationship indicates that a single product can appear in multiple sales orders. The diagram clearly shows how primary and foreign keys establish relationships between tables, implementing the referential integrity that ensures consistency across related data.

In relational databases, tables adhere to predefined schemas that specify the structure, data types, and constraints applicable to the data. Each table typically includes a primary key that uniquely identifies each row, while foreign keys establish relationships between tables, implementing the referential integrity that ensures consistency across related data [14].

A key strength of relational databases is their adherence to ACID properties (Atomicity, Consistency, Isolation, Durability), which ensure reliable transaction processing. Atomicity guarantees that transactions are treated as indivisible units that either complete entirely or have no effect. Consistency ensures that transactions maintain database integrity by transforming the database from one valid state to another. Isolation prevents interference between concurrent transactions, making them appear as if executed

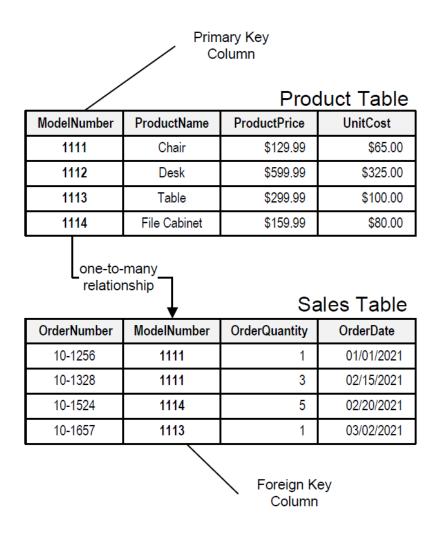


Figure 2.3: Example of a Relational Schema [39]

sequentially. Durability ensures that committed transactions persist even after system failures [14].

ACID compliance makes relational databases particularly suitable for automotive parameter management, where data integrity and consistency are paramount. Incorrect parameter values could potentially affect vehicle safety and performance, making the strong consistency guarantees of relational databases essential for maintaining data integrity [37]. Additionally, the hierarchical structure of automotive parameter systems—with well-defined relationships between ECUs, modules, Parameter IDs (PIDs), and parameters—aligns naturally with the relational model's representation of structured data and relationships.



2.2.2 Non-Relational Database Systems

Non-relational databases, often referred to as NoSQL (Not Only SQL) databases, emerged as alternatives to the relational model, particularly for use cases involving large-scale distributed systems, unstructured data, or schema flexibility requirements. Unlike relational databases, NoSQL systems typically sacrifice some aspects of ACID compliance in favor of scalability, flexibility, and performance characteristics suited to specific application domains [4].

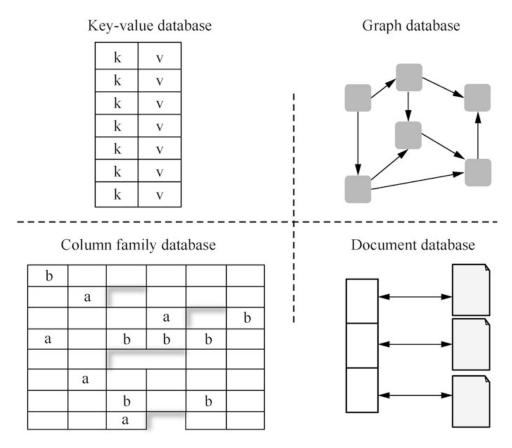


Figure 2.4: Major Types of NoSQL Databases [17]

Figure 2.4 illustrates the four major types of NoSQL databases, each with distinct data models and storage approaches. Key-value databases (top-left) store data as simple key-value pairs, offering fast lookups but limited query functionality. Graph databases (top-right) organize data as networks of interconnected nodes and relationships, excelling at representing complex relationships. Column-family databases (bottom-left) store data in column-oriented structures that can efficiently handle sparse data. Document databases (bottom-right) store semi-structured data as documents, typically in formats like JSON or XML, providing schema flexibility while maintaining query capabilities.



NoSQL databases can be categorized into these four types based on their data models: document databases (MongoDB, CouchDB), key-value stores (Redis, DynamoDB), column-family stores (Cassandra, HBase), and graph databases (Neo4j, Amazon Neptune). Many NoSQL systems follow the BASE principle (Basically Available, Soft state, Eventually consistent) rather than ACID, prioritizing availability and partition tolerance over immediate consistency [7].

While NoSQL databases excel in specific domains such as high-volume web applications, real-time analytics, and social networks, they present challenges for applications requiring complex transactions, strict data integrity, or sophisticated query capabilities across related entities [23]. For automotive parameter management, these limitations make NoSQL systems generally less suitable than relational databases.

The potential for eventual consistency rather than immediate consistency in many NoSQL systems could lead to incorrect parameter configurations being used during development or testing, creating significant risks for vehicle performance and safety. Additionally, the hierarchical nature of automotive electronic systems, with well-defined relationships between entities, aligns naturally with the relational model's approach to representing structured data and relationships. The ability to enforce these relationships through foreign key constraints provides important safeguards against data inconsistency that would be more difficult to implement in many NoSQL systems.

2.3 Database Design Methodologies

Database design methodologies provide structured approaches to creating efficient, reliable database systems. These methodologies help translate real-world information needs into technical implementations that can store and manage data effectively. This section explores fundamental approaches that form the theoretical foundation for database design, presented in a sequence that follows the natural progression from user requirements to technical implementation.

2.3.1 Use Case Modeling

Before designing a database structure, it is essential to understand how users will interact with the system. Use case modeling provides a technique for capturing user requirements by identifying who will use the system (actors) and what they need to accomplish (use cases). Developed by Ivar Jacobson, use case modeling has become a cornerstone of requirements analysis in system development [20].



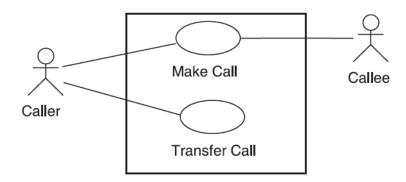


Figure 2.5: Simple Use Case Diagram for a Telephone System

Figure 2.5 illustrates a simple use case diagram for a telephone system. The diagram shows two actors, the "Caller" and "Callee," and two use cases, "Make Call" and "Transfer Call." The lines connecting the Caller to both use cases indicate that the Caller can initiate either a direct call or a call transfer, while the Callee is connected only to the "Make Call" use case, indicating their role as the recipient of calls. The rectangle surrounding the use cases represents the system boundary, clearly delineating the functional scope of the telephone system. This visual representation effectively communicates the basic functionality that the system must support, providing a foundation for further design and implementation decisions.

A use case represents a specific goal that an actor wishes to achieve using the system. Actors can be human users with different roles (such as administrators or regular users) or external systems that interact with the database. The collection of all use cases defines the system's functional boundaries—what it must do to satisfy user needs [20].

Use case diagrams provide a visual representation of these relationships, showing actors as stick figures and use cases as ovals, with lines connecting actors to their associated use cases. This visual format makes the system's purpose accessible to non-technical stakeholders, facilitating communication between developers and users. As noted by Jacobson, "Use cases bridge the gap between the users' and the developers' views of the system" [20].

For automotive parameter management systems, use case modeling helps identify the different ways in which engineers, documentation specialists, administrators, and other stakeholders need to interact with parameter data. These use cases then inform the database design, ensuring that the resulting structure effectively supports all required operations.



2.3.2 Entity-Relationship Modeling

After understanding user requirements through use cases, the next step is to model the data itself. Entity-Relationship (ER) modeling provides a conceptual framework for representing the data structure needed to support the identified use cases. Introduced by Peter Chen in 1976, ER modeling has become the most widely used approach for conceptual database design [9].

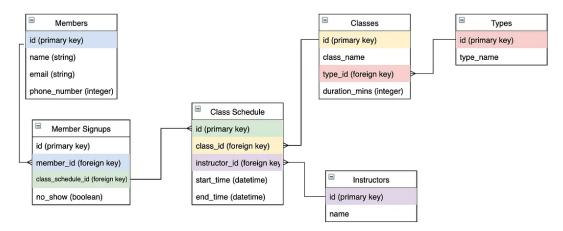


Figure 2.6: Entity-Relationship Diagram for a Fitness Studio Database

Figure 2.6 presents a detailed entity-relationship diagram for a fitness studio database. The diagram shows six main entities: Members, Member Signups, Class Schedule, Classes, Types, and Instructors. Each entity is represented as a rectangle containing its attributes, with primary keys clearly labeled at the top of each entity. The diagram uses color coding to visually distinguish different entity types and their related attributes. Foreign key relationships are explicitly shown with connecting lines and arrows pointing to the referenced primary keys, illustrating the database's referential integrity constraints.

The diagram demonstrates several important relational concepts. The Members entity contains basic member information including a unique identifier, name, email, and phone number. The Class Schedule entity serves as a central hub, connecting classes with their instructors and time slots. The Member Signups entity implements a many-to-many relationship between members and scheduled classes, with an additional attribute to track attendance (no_show). The Classes entity includes a foreign key reference to the Types entity, establishing a categorization system for different class offerings. This comprehensive diagram provides a clear blueprint for database implementation, showing both the structural components and their interrelationships in the fitness studio domain.



ER modeling identifies three main components: Entities represent the objects or concepts about which information needs to be stored. In an automotive context, these might include vehicles, electronic control units (ECUs), parameters, and users. Entities are represented as rectangles in ER diagrams. Attributes describe the specific properties or characteristics of each entity. For example, a parameter entity might have attributes like name, value, unit, and description. Attributes are shown as ovals connected to their entity. Relationships describe the associations between entities. For instance, "ECUs contain parameters" expresses a relationship between ECU and parameter entities. Relationships are shown as diamonds connecting the related entities, with cardinality notations indicating how many instances of each entity can participate in the relationship [14].

ER modeling is particularly valuable for complex domains like automotive systems because it provides a visual representation that stakeholders can understand while being precise enough to guide database implementation. Chen explains that "the entity-relationship model adopts the more natural view that the real world consists of entities and relationships" [9], making it an intuitive approach for modeling real-world systems.

2.3.3 Database Normalization

Once the conceptual model is established through ER modeling, database normalization helps refine this model into an efficient, consistent structure. Normalization is a systematic process developed by E.F. Codd that organizes data to minimize redundancy and avoid update anomalies [10].

Normalization proceeds through several "normal forms," each addressing specific types of data inconsistencies:

First Normal Form (1NF) requires that each cell in a table contains only a single value, not a list of values. For example, storing multiple phone numbers in a single field would violate 1NF. This ensures that data is atomic (indivisible) and can be manipulated consistently.

Second Normal Form (2NF) builds on 1NF by requiring that all non-key attributes depend on the entire primary key, not just part of it. This prevents situations where changing one piece of data requires multiple updates in different places.

Third Normal Form (3NF) further refines the structure by requiring that non-key attributes depend only on the primary key, not on other non-key attributes. This eliminates transitive dependencies that can lead to update anomalies [14].



For most practical applications, achieving 3NF provides a good balance between data integrity and system performance. As explained by Date, "Third normal form is considered adequate for most practical purposes; further normalization is usually performed only when necessary" [12].

In automotive parameter management, normalization helps organize complex data about ECUs, modules, and parameters into a structure that maintains consistency while supporting efficient access. For example, normalizing parameter data ensures that when a parameter value changes, that change only needs to be recorded in one place, eliminating the risk of inconsistent values across the database.

2.3.4 Role-Based Access Control Models

Database systems often contain sensitive information that should not be accessible to all users. Role-Based Access Control (RBAC) provides a structured approach to managing permissions within a database system. Introduced by David Ferraiolo and Richard Kuhn in the 1990s, RBAC has become the predominant model for access control in enterprise systems due to its balance of security and administrative simplicity [32].

The core concept of RBAC is that permissions are associated with roles, and users are assigned to appropriate roles rather than being granted permissions directly. A role represents a specific function within an organization, such as "administrator," "engineer," or "analyst." Each role is granted a set of permissions that allow users assigned to that role to perform specific operations on database objects like reading, creating, updating, or deleting records [15].

This structure provides several important advantages over direct permission assignment. First, it simplifies administration by allowing permissions to be managed at the role level rather than the individual user level. When a new user joins the organization, they can simply be assigned to the appropriate roles rather than requiring configuration of individual permissions. Second, it improves security by implementing the principle of least privilege, ensuring that users have only the permissions necessary for their specific responsibilities which reduces the risk of unauthorized access or accidental data modifications [32].

The theoretical foundation of RBAC includes several key components: users (individuals who need access to the system), roles (collections of permissions that correspond to job functions), permissions (defined operations on specific resources), and sessions (temporary bindings between users and their assigned roles). These components



provide a flexible framework for implementing access control policies tailored to specific organizational needs while maintaining a clear separation between users and permissions through the role abstraction [31].

2.3.5 Version Control for Databases

Version control for databases addresses the challenge of tracking changes to data structures and content over time. Unlike traditional file-based version control systems designed for source code, database versioning must maintain complex relationships between entities while preserving historical states and supporting evolution through distinct development stages [4].

Several approaches have emerged for implementing version control in database systems. The snapshot approach captures complete database states at specific points in time, providing simple retrieval of historical states but potentially consuming significant storage resources. The change-based approach records only modifications to database content, reducing storage requirements but requiring reconstruction of historical states through the application of change records [4].

The temporal approach extends traditional database structures with time dimensions, enabling direct querying of historical states through time-based predicates. This approach typically introduces valid time (when facts are true in the real world) and transaction time (when facts are recorded in the database) dimensions, allowing sophisticated historical analysis but adding complexity to schema design and query formulation [35].

Phase-based versioning represents a domain-specific approach that aligns database versioning with development phases rather than continuous time. This approach explicitly models development stages as first-class entities in the database schema, associating data with specific phases rather than temporal timestamps. According to Bhattacherjee et al., "Domain-specific versioning approaches often provide better performance and usability than generic temporal database techniques when tailored to specific application requirements" [4].

2.3.6 Temporal Database Concepts

Many applications, including automotive parameter management, need to track how data changes over time. Temporal database concepts address these requirements by providing mechanisms for managing time-varying data. Unlike traditional databases



that store only the current state, temporal databases maintain historical states and support queries based on time dimensions [35].

| ENo | EStart | EEnd | EDept |
|-------|------------|------------|-------|
| 22217 | 2010-01-01 | 2011-02-03 | 3 |
| 22217 | 2011-02-03 | 2011-09-10 | 4 |
| 22217 | 2011-09-10 | 2011-11-12 | 3 |

Figure 2.7: Temporal Table Tracking Employee Department History

Figure 2.7 illustrates a temporal database table that tracks an employee's department history over time. The table contains four columns: ENo (Employee Number), EStart (employment start date in a department), EEnd (employment end date in a department), and EDept (department identifier). The example shows the complete employment history of employee number 22217, who moved between departments over time. The first row indicates that the employee started in department 3 on January 1, 2010, and remained there until February 3, 2011. The second row shows that on February 3, 2011, the employee transferred to department 4, where they stayed until September 10, 2011. The third row demonstrates that on September 10, 2011, the employee returned to department 3, where they remained until November 12, 2011.

This temporal table implementation enables tracking of the complete history of the employee's department assignments without overwriting previous records. Each row represents a specific period during which the employee belonged to a particular department, with non-overlapping time intervals. The start date of each subsequent record matches the end date of the previous record, ensuring continuity in the historical record. This approach to temporal data storage supports historical queries like "Which department did employee 22217 belong to on March 15, 2011?" by evaluating which row's time interval (EStart to EEnd) contains the query date. Such temporal structures are essential for maintaining comprehensive historical records in domains where tracking changes over time is critical, such as automotive parameter management.

Temporal databases typically support two key time dimensions. Valid time represents when facts are true in the modeled reality—for example, when a particular parameter configuration becomes active in a vehicle. Transaction time represents when facts are recorded in the database—for example, when a parameter value was updated in the system. Databases that support both dimensions are known as bi-temporal databases [24].



Temporal database implementations often use specialized table structures called temporal tables. These tables extend traditional table structures with additional timestamp columns that define the time periods during which each record is valid. For example, a temporal parameter table might include ValidFrom and ValidTo columns that define when each parameter value is applicable, allowing the database to maintain a complete history of parameter changes [30].

Kulkarni and Michels explain that "temporal tables provide a systematic way to track and query historical data without requiring application-level version management" [24]. This capability is particularly valuable in regulated industries like automotive development, where traceability and auditability of parameter changes are essential for compliance and quality assurance.

In automotive parameter management, temporal database concepts can support critical requirements such as tracking parameter evolution throughout the development lifecycle, maintaining historical records for diagnostic and compliance purposes, and enabling historical analysis to understand how parameter configurations have evolved over time. These capabilities form an important theoretical foundation for designing systems that manage time-sensitive data in complex domains.

2.3.7 Strategic Denormalization

While normalization provides a theoretical foundation for database integrity, practical database design often requires balancing normalization principles with performance considerations. Strategic denormalization involves deliberately introducing controlled redundancy to improve performance for specific operations [4].

Consider a fully normalized database where information about parameters, their modules, and their ECUs is stored in separate tables. To retrieve a parameter with its associated module and ECU information would require joining all three tables—an operation that becomes increasingly expensive as the database grows. In cases where this retrieval happens frequently, storing the module and ECU names directly in the parameter table (introducing controlled redundancy) could significantly improve performance [33].

Molinaro emphasizes that "denormalization is not about abandoning normalization principles, but about making strategic exceptions for performance reasons" [25]. These exceptions should be carefully documented and justified based on specific performance requirements.



For automotive systems, where both data integrity and query performance are critical, finding the right balance between normalization and strategic denormalization is essential. This balance ensures that parameter data maintains consistency while providing the performance needed for engineering workflows.

2.3.8 Conceptual, Logical, and Physical Design Levels

Database design typically proceeds through three levels of abstraction, allowing designers to manage complexity by focusing on different aspects at each stage:

Conceptual design focuses on what data needs to be stored, without concern for implementation details. The ER model created at this stage captures entities, attributes, and relationships from a business perspective, providing a foundation that both technical and non-technical stakeholders can understand [14].

Logical design transforms the conceptual model into structures specific to the chosen database model (typically relational), defining tables, columns, keys, and relationships. This stage applies normalization principles to refine the structure, independent of any specific database system [14].

Physical design addresses how the logical design will be implemented in a specific database management system, considering factors like storage structures, indexing strategies, and access methods. This stage optimizes the design for performance based on anticipated usage patterns [28].

Moving through these levels allows database designers to progressively refine the database structure, addressing different concerns at each stage. As Elmasri and Navathe observe, "The separation of conceptual, logical, and physical design allows database designers to focus on the appropriate level of abstraction at each stage" [14].

For automotive parameter management, this layered approach helps manage the complexity of the domain, ensuring that the resulting database effectively supports both the business requirements (storing and managing parameter configurations) and the technical requirements (performance, scalability, and maintainability).

3 State of the Art

This chapter examines the current state of the art in database version control systems and automotive parameter management. It begins by analyzing existing approaches to software configuration management in the automotive industry, followed by an evaluation of database versioning techniques and their applicability to parameter management systems. The chapter also explores role-based access control models and integration strategies for enterprise systems, establishing the theoretical foundation for the VMAP system design.

3.1 Parameter Management in Automotive Software Development

The complexity of automotive software has grown exponentially in recent decades, with modern vehicles containing up to 100 million lines of code distributed across dozens of electronic control units (ECUs) [29]. This growth has significantly increased the importance and complexity of parameter management in automotive development.

3.1.1 Evolution of Automotive Parameter Management

Parameter management in automotive systems has evolved from simple calibration tables to sophisticated configuration frameworks managing thousands of parameters across multiple vehicle variants. Broy [8] describes the fundamental challenges in automotive software engineering, highlighting that software complexity is driven by the need to address multiple variants, market requirements, and technical functions. The parameter configuration problem is specifically identified as one of the key challenges in this domain.

Early approaches to parameter management relied on specialized tools provided by ECU suppliers, which typically stored parameters in proprietary formats with limited version control capabilities. Pretschner et al. [29] note that these tools evolved from simple memory editors to more sophisticated calibration environments, but remained focused on individual ECUs rather than system-wide parameter management.

Staron [37] describes how AUTOSAR (Automotive Open System Architecture) has contributed to more structured parameter management by defining standard interfaces



and component models that separate parameters from implementation. However, the practical implementation of these standards varies across organizations and ECU suppliers, creating integration challenges for comprehensive parameter management.

3.1.2 Challenges in Automotive Parameter Management

The management of parameters in automotive software development presents specific challenges that distinguish it from general software configuration management. Pretschner et al. [29] identify several key challenges related to variability management in automotive software, including the need to maintain multiple parameter configurations for different vehicle variants, markets, and operating conditions.

Broy [8] emphasizes the challenge of managing interdependencies between parameters, noting that changes to one parameter often require coordinated changes to related parameters to maintain system consistency. This creates a need for sophisticated dependency tracking mechanisms that go beyond traditional version control systems.

Another significant challenge relates to validation requirements for parameter changes. Unlike source code, which can be validated through compilation and static analysis, parameters require functional testing to verify their correctness. Pretschner et al. [29] describe how this validation often involves specialized hardware-in-the-loop or vehicle-level testing, creating a significant gap between parameter modification and validation.

Kiencke and Nielsen [22] discuss the specific challenges related to powertrain control parameters, noting the complex interactions between engine control parameters and their effects on vehicle performance, emissions, and fuel economy. These interactions create a need for sophisticated parameter testing and validation processes beyond simple version control.

3.1.3 Current Approaches and Tools

Current parameter management solutions in the automotive industry span a spectrum from general-purpose tools to specialized automotive calibration systems.

Staron [37] discusses how AUTOSAR tools provide standardized interfaces for parameter management in modern automotive systems, but notes that these tools focus primarily on the technical aspects of parameter definition rather than the organizational



processes of parameter development and validation through multiple development phases.

Broy [8] identifies the challenges of integrating parameter management into broader software development processes, noting that many organizations maintain separate workflows for software development and parameter calibration. This separation creates coordination challenges, particularly when parameter changes affect multiple software components or require software modifications.

Pretschner et al. [29] discuss how model-based development approaches are increasingly used in automotive development, with parameters linked to model elements to provide traceability and support automated validation. However, they note that the integration between parameter management tools and modeling environments remains incomplete in many organizations.

For database-oriented approaches to parameter management, Bhattacherjee et al. [4] provide a theoretical foundation by examining the principles of dataset versioning. They describe the fundamental trade-offs between storage efficiency and reconstruction performance, which are particularly relevant for systems that must maintain multiple parameter configurations across different development phases.

3.2 Database Version Control Systems

Version control for database content presents distinct challenges compared to traditional source code version control. While source code version control focuses on tracking changes to text files, database version control must address structured data with complex relationships and constraints [4]. This section examines current approaches to database version control and their applicability to automotive parameter management.

3.2.1 Traditional Database Versioning Approaches

Traditional approaches to database versioning fall into several categories, each addressing different aspects of the versioning challenge. Schema evolution tools focus on tracking and managing changes to database structure through migration scripts or schema manipulation languages. Curino et al. [11] describe an approach for automating database schema evolution in information system upgrades, focusing on maintaining data integrity during schema transitions.



Bhattacherjee et al. [4] provide a comprehensive analysis of dataset versioning approaches, identifying a fundamental trade-off between storage and recreation costs. They categorize versioning strategies into several approaches:

- 1. Version-first approaches maintain complete snapshots of datasets at specific version points, providing simple retrieval of historical states but requiring substantial storage space.
- 2. Delta-based approaches store only the changes between versions, reducing storage requirements but increasing the computational cost of reconstructing historical states.
- 3. Hybrid approaches combine elements of both strategies, typically storing periodic full snapshots with incremental deltas between snapshots.

The authors note that the optimal strategy depends on specific usage patterns, particularly the ratio between storage costs and the frequency and complexity of historical data access operations.

Mueller and Müller [26] describe a practical implementation of database versioning between research institutes, highlighting the challenges of maintaining consistency across systems with different update cycles. Their approach uses a combination of schema versioning and data synchronization mechanisms to maintain consistency while supporting independent evolution.

3.2.2 Temporal Database Approaches

Temporal database approaches provide a theoretical foundation for managing time-varying data in database systems. Kulkarni and Michels [24] describe the temporal features introduced in SQL:2011, which formalized support for period data types and temporal tables in the SQL standard. These features enable tracking of both valid time (business time) and transaction time (system time) dimensions, supporting bi-temporal data management.

The valid time dimension represents when facts are true in the modeled reality, independent of when they are recorded in the database. This dimension supports business-oriented temporal queries such as "What was the value of this parameter in a specific phase?" or "When did this parameter change from value A to value B?" [6]. The transaction time dimension represents when facts are recorded in the database, supporting auditability through questions like "Who changed this parameter, and when did they change it?" [24].



Bi-temporal databases combine both dimensions, providing a comprehensive framework for tracking both when changes occurred in the system and when they became effective in the real world [6]. This approach is particularly valuable for regulated industries like automotive development, where both historical accuracy and change auditability are essential for compliance and quality assurance.

Snodgrass [35] provides a comprehensive guide to developing time-oriented database applications in SQL, describing practical techniques for implementing temporal functionality in relational database systems. The author presents various approaches to tracking historical data, including transaction-time tables, valid-time tables, and bi-temporal tables, with practical implementation guidance for each approach.

Biriukov [5] examines practical implementation aspects of bi-temporal databases, highlighting the challenges of schema design, query formulation, and performance optimization. The author notes that domain-specific temporal approaches often provide more practical solutions than generic bi-temporal frameworks, particularly for applications with specialized temporal requirements.

3.2.3 Version Control for Parameter Management

Version control for automotive parameter management presents specific requirements that differ from general database versioning needs. Drawing from the literature, several key requirements can be identified:

Broy [8] discusses the need for version control approaches that align with automotive development processes, which typically follow a structured progression through predefined development phases. Unlike source code versioning, which often follows continuous development with arbitrary version points, parameter versioning must support specific phase-based workflows.

Bhattacherjee et al. [4] examine the trade-offs between different versioning strategies, which are particularly relevant for parameter management systems that must maintain multiple configurations across different development phases. The authors' analysis of storage versus reconstruction costs provides a theoretical foundation for designing efficient parameter versioning systems.

Snodgrass [35] describes techniques for tracking valid-time information in database systems, which aligns with the need to maintain parameter configurations that are valid for specific development phases or vehicle configurations. However, the author's focus on general temporal database approaches does not address the specific requirements of phase-based development.



Bhattacherjee et al. [4] note that domain-specific versioning systems often provide more effective solutions than generic versioning frameworks, particularly for domains with structured development processes and complex entity relationships. This observation supports the development of specialized versioning approaches tailored to automotive parameter management rather than adopting generic temporal database techniques.

3.3 Role-Based Access Control in Enterprise Systems

Role-Based Access Control (RBAC) has become a dominant paradigm for managing access rights in enterprise systems, providing a structured approach to security management that aligns with organizational responsibilities [32]. For automotive parameter management, where different user roles have distinct responsibilities and access requirements, RBAC provides a foundation for implementing appropriate security controls.

3.3.1 RBAC Model and Extensions

The core RBAC model, as defined by Sandhu et al. [32], consists of users, roles, permissions, and sessions. Users are assigned to roles that correspond to job functions, and roles are granted permissions that authorize specific operations on protected resources. This indirect association between users and permissions through roles simplifies security administration while maintaining the principle of least privilege.

Several extensions to the basic RBAC model have been developed to address more complex security requirements. Sandhu et al. [32] describe hierarchical RBAC, which introduces role hierarchies that enable permission inheritance between roles, supporting organizational structures with senior roles inheriting permissions from junior roles.

Sandhu and Bhamidipati [31] present administrative RBAC (ARBAC), which addresses the management of the RBAC system itself, defining who can assign users to roles and modify role permissions. This extension is particularly relevant for enterprise systems where role and permission management is distributed across different administrative domains.

Ferraiolo et al. [15] describe policy-enhanced RBAC, which combines role-based permissions with attribute-based policies to provide context-sensitive access control. This hybrid approach is particularly valuable for systems where access decisions



depend on both user roles and context-specific factors such as time, location, or resource attributes.

3.3.2 RBAC in Database Systems

Modern database management systems provide varying levels of support for RBAC principles. Elmasri and Navathe [14] describe the evolution of database security mechanisms from simple user-based privileges to more sophisticated role-based models. Most enterprise database systems now include native support for roles, user-role assignments, and permission management through SQL statements like GRANT and REVOKE.

Obe and Hsu [28] detail PostgreSQL's implementation of RBAC concepts, including role hierarchies through role inheritance, permission management through fine-grained privileges, and row-level security policies for content-based access control. These capabilities provide a foundation for implementing domain-specific access control models on top of the database system's native security features.

However, database-level RBAC implementations typically focus on controlling access to database objects like tables, views, and functions, rather than providing application-level access control that considers domain-specific entities and operations. For complex applications like automotive parameter management, database-level RBAC must be complemented with application-level access control logic that maps domain-specific concepts to database operations [15].

3.3.3 Access Control for Parameter Management

Access control for automotive parameter management presents specific requirements that extend beyond basic RBAC models. Drawing from the literature, several key access control requirements can be identified:

Sandhu et al. [32] provide the theoretical foundation for role-based access control, which aligns with the organizational structure of automotive development teams. Different roles such as parameter engineers, module developers, and system integrators require different access rights to parameter data.

Ferraiolo et al. [15] describe policy-enhanced RBAC, which combines role-based permissions with attribute-based policies. This hybrid approach is particularly relevant for parameter management, where access rights may depend on both user roles and



attributes of the parameters being accessed, such as their development phase or module assignment.

Hu et al. [19] discuss practical aspects of implementing and managing policy rules in attribute-based access control, providing insights into the challenges of combining role-based and attribute-based approaches. Their work highlights the importance of balancing security requirements with usability considerations, which is particularly relevant for parameter management systems used by diverse stakeholder groups.

Sandhu and Bhamidipati [31] address the administrative aspects of RBAC, which are important for parameter management systems where access control administration may be distributed across different organizational units. Their ARBAC97 model provides a framework for delegating administrative responsibilities while maintaining central governance.

3.4 Database Integration with Enterprise Systems

Integration between database systems and enterprise applications presents significant challenges in automotive development environments, where parameter management must interact with numerous other systems across the development lifecycle. Effective integration strategies must address both technical interoperability and semantic consistency while maintaining performance and security [18].

3.4.1 Enterprise Integration Patterns

Enterprise integration patterns, as documented by Hohpe and Woolf [18], provide a catalog of solutions for common integration challenges. These patterns address various aspects of system integration, including messaging styles, messaging channels, message construction, and message transformation.

For database-centric applications like parameter management systems, several integration patterns are particularly relevant. Fowler [16] describes the Repository pattern, which provides a structured approach to data access, abstracting the database implementation details behind a domain-focused interface. This abstraction simplifies integration by providing a stable API for other systems to interact with the parameter repository.

Fowler [16] also documents the Data Transfer Object (DTO) pattern, which addresses the challenge of transferring data between systems with different data models. By



defining specialized objects for inter-system communication, this pattern enables consistent data exchange while isolating each system's internal representation.

The Canonical Data Model pattern, as described by Hohpe and Woolf [18], establishes a common data representation across multiple systems, simplifying data transformation and ensuring consistent interpretation. This pattern is particularly valuable for parameter management, where the same parameter concepts may be represented differently in various systems across the development lifecycle.

3.4.2 Database Synchronization Approaches

Database synchronization presents specific challenges when integrating parameter management systems with other enterprise data sources. Mueller and Müller [26] describe approaches to database versioning and synchronization between research institutes, highlighting the challenges of maintaining consistency across systems with different update cycles.

Bhattacherjee et al. [4] discuss the principles of dataset versioning, which are relevant for synchronization between parameter management systems and other enterprise databases. Their analysis of the trade-offs between storage and recreation costs provides insights into designing efficient synchronization mechanisms that minimize both data transfer volumes and processing overhead.

Seenivasan and Vaithianathan [34] examine change data capture (CDC) techniques, which enable incremental synchronization by identifying and propagating only changed data between systems. These techniques reduce synchronization overhead compared to full dataset transfers but require reliable change detection mechanisms and careful handling of interdependent changes.

Kleppmann and Beresford [23] address the challenges of conflict resolution in distributed data systems, which are relevant for parameter management systems that must synchronize with multiple enterprise data sources. Their work on conflict-free replicated data types provides theoretical foundations for designing synchronization mechanisms that maintain consistency across distributed systems.

3.5 Summary and Research Gaps

The review of existing literature reveals several research gaps in the domain of database systems for automotive parameter management:



Current database versioning approaches, as described by Bhattacherjee et al. [4] and Snodgrass [35], provide general frameworks for managing time-varying data but do not specifically address the phase-based development processes common in automotive parameter management. There is a need for specialized versioning approaches that align directly with automotive development workflows while providing the traceability and auditability required for regulatory compliance.

The RBAC models described by Sandhu et al. [32] and Ferraiolo et al. [15] provide a foundation for access control but require extensions to address the specific requirements of parameter management, where access rights depend on both organizational roles and parameter-specific attributes such as module assignment and development phase.

Integration approaches documented by Hohpe and Woolf [18] and Mueller and Müller [26] provide general patterns for system integration but do not specifically address the challenges of integrating parameter management systems with automotive-specific enterprise systems such as parameter definition databases and vehicle configuration databases.

These research gaps highlight the need for domain-specific solutions that combine insights from database version control, access control models, and enterprise integration patterns with specialized knowledge of automotive development processes. TheVMAP system addresses these gaps by developing a database architecture tailored to the specific requirements of automotive parameter management, as will be detailed in subsequent chapters.

4 Methodology and Concept Development

This chapter presents the systematic approach taken in designing the Variant Management and Parametrization (VMAP) system. The methodology follows established software engineering principles to address the complex requirements of automotive parameter management. Beginning with a requirements analysis, the chapter proceeds to detail the conceptual architecture design, data model, validation mechanisms, and integration approaches developed to ensure system robustness and compatibility with existing enterprise infrastructure.

4.1 Requirements Analysis

The foundation of the VMAP system design was a comprehensive requirements analysis conducted through a series of structured interviews with stakeholders, detailed examination of the existing Excel-based process, and workshops with domain experts. This multi-faceted approach, following Sommerville's framework for requirements engineering, ensured that both functional and non-functional requirements would be thoroughly identified and prioritized [36].

4.1.1 Functional Requirements

The primary functional requirements were derived from direct observation of engineers' current Excel-based workflow combined with semi-structured interviews conducted with module developers and documentation specialists. Through this process, several critical requirements emerged for the VMAP system.

The system must support the hierarchical organization of parameters within ECUs, Modules, and Parameter IDs (PIDs), mirroring the domain-specific structure of automotive electronic systems as described by Staron [37]. This hierarchical organization is essential for maintaining the logical structure of vehicle parameters and aligning with established engineering practices.

Users must be able to create variants for parameters with specific code rules determining their applicability, and define segments representing modified parameter values. If no segment exists, the system must default to Parameter Definition Database values—an approach that allows efficient storage by tracking only modifications rather



than duplicating unchanged parameters, aligning with Bhattacherjee's principles of dataset versioning [4].

The system must track parameter values across four distinct release phases: Phase1, Phase2, Phase3, and Phase4, with changes in earlier phases propagating to later phases unless explicitly overridden. This phase-based approach represents a domain-specific adaptation particularly suited to automotive software development cycles as identified in Broy's research on automotive software engineering challenges [8].

All modifications require comprehensive logging with user information, timestamp, and detailed change data, supporting regulatory compliance and enabling parameter evolution tracking. Through the stakeholder interviews, it was determined that the system must also provide functionality to create parameter configuration snapshots at specific points, particularly at phase transitions, for documentation purposes—a capability identified as essential for quality assurance and regulatory compliance in automotive software development by Staron [37].

4.1.2 Integration with External Systems

The stakeholder interviews and process analysis revealed thatVMAP must integrate with two critical external enterprise systems: the Parameter Definition Database (PDD) and the Vehicle Configuration Database (VCD).

The Parameter Definition Database (PDD) serves as the authoritative source for the hierarchical structure of automotive electronic systems, containing definitions of ECUs, Modules, PIDs, and baseline parameter configurations. As noted by Pretschner et al. [29], maintaining this hierarchical structure is essential for automotive software development. The workshops with domain experts confirmed that while VMAP will manage parameter variants and customizations, it must rely on PDD for the underlying parameter definitions and structural relationships, requiring a robust synchronization mechanism to maintain consistency between the systems.

The Vehicle Configuration Database (VCD) contains comprehensive vehicle specifications and configuration codes that determine which parameter variants apply to specific vehicle configurations. Integration with this system is necessary for two critical functions: validating the boolean code rules associated with parameter variants to ensure they reference valid vehicle codes, and supporting parameter file generation for specific vehicle configurations by resolving the applicable parameter variants based on vehicle codes. This integration requirement aligns with Staron's analysis of automotive software architectures, which emphasizes the importance of configuration management in supporting variant-rich vehicle platforms [37].



These integration requirements necessitated careful consideration of data synchronization approaches, leading to an exploration of different strategies for maintaining consistency between VMAP and these external systems while minimizing performance impact and complexity.

4.1.3 User Role Requirements

A systematic analysis of the current Excel-based workflow, coupled with contextual inquiries with engineering teams, identified four distinct user roles with specific access requirements. This analysis included shadowing users in their daily work, documenting their tasks and access patterns, and conducting structured interviews to validate the observed patterns.

Module developers require write access to parameters within their assigned modules, with the ability to create and modify variants and segments. Documentation specialists need access to frozen data for documentation, comparison capabilities between phases, and comprehensive change history access. System administrators require comprehensive control over user management, release phases, and special operations like variant deletion and phase freezing. Read-only users need view access to all parameter data with parameter file generation capabilities but no modification rights.

These roles were defined based on the principle of least privilege as described by Sandhu [32], ensuring users have access only to functionality required for their specific responsibilities. This enhances system security while simplifying the user experience by presenting only relevant options.

4.1.4 Data Management Requirements

The system must maintain distinct parameter versions across different release phases, allowing simultaneous work on multiple phases while enabling access to parameter values from any point in the development lifecycle. As highlighted by Elmasri and Navathe [14], data integrity requires maintaining referential integrity across all related entities, particularly ensuring variants and segments associate with valid parameters.

Multi-dimensional parameter support is essential for complex automotive parameters such as mapping tables. Operations modifying multiple related entities must function as atomic transactions to maintain data consistency—particularly important for phase



transitions where numerous parameters, variants, and segments may change simultaneously, a requirement that aligns with Bhattacherjee's research on dataset versioning approaches [4].

Query performance analysis, based on projected usage patterns from the current Excelbased process, identified critical query paths including parameter retrieval by ECU, module, PID, release phase, and parameter name. These requirements influenced schema design decisions regarding normalization and indexing strategies to optimize common query patterns.

4.2 Use Case Modeling

Following the requirements gathering process, use case modeling was employed to formalize the system's functional requirements from a user perspective. This approach, as described by Jacobson [20], provides a structured way to represent the system's capabilities and the interactions between users and the system.

The use case diagram in Figure 4.1 illustrates the primary actors and their interactions with theVMAP system. Four primary actor types are identified, corresponding to the user roles established during requirements analysis: Module Developers, who create and modify parameter variants; Documentation Team members, who access parameter data for documentation purposes; Administrators, who manage system settings and user access; and Read-Only Users, who view parameter data without making modifications.

The diagram demonstrates how these actors interact with key system functionalities. Module Developers primarily interact with variant creation and modification use cases, while having limited access to parameter viewing and comparison features. Documentation Team members focus on viewing frozen configurations, comparing parameters across phases, and accessing change history. Administrators have access to all system functions, including user management, release configuration, and system settings. Read-Only Users are limited to viewing parameters and generating parameter files.

This use case model provides a clear visual representation of the system's scope and functionality, serving as a bridge between user requirements and technical implementation. By mapping user roles to specific system functions, the model ensures that the database design will support all required user interactions while maintaining appropriate access controls.



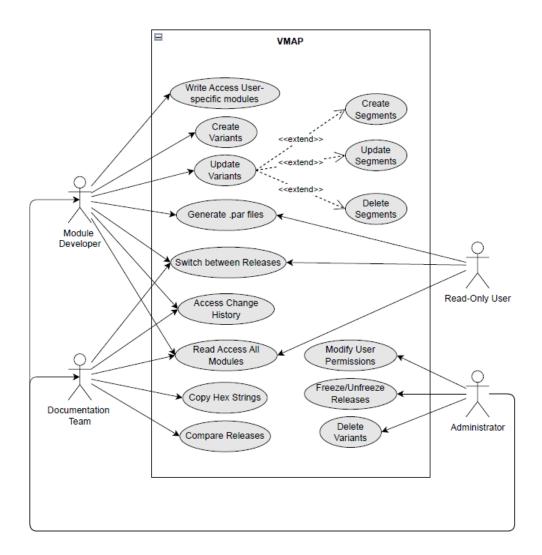


Figure 4.1: VMAP System Use Case Diagram

4.3 User Management Approaches

Based on the identified user role requirements, two distinct approaches to user management were considered for the VMAP system: a traditional role-based approach and a hybrid role-permission approach. Each approach offers different advantages in terms of flexibility, administrative complexity, and alignment with organizational needs.

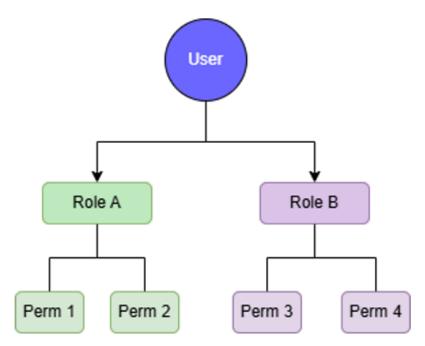


Figure 4.2: Traditional Role-Based Access Control Approach

4.3.1 Traditional Role-Based Approach

The traditional role-based approach, illustrated in Figure 4.2, assigns users to predefined roles that contain fixed sets of permissions, following the classic Role-Based Access Control (RBAC) model described by Sandhu [32]. In this approach, each user is assigned one or more roles (Administrator, Module Developer, Documentation Team, Read-Only User), and all permissions are granted through these role assignments without individual permission adjustments.

This approach offers administrative simplicity, as user management involves only assigning appropriate roles rather than configuring individual permissions. The role structure also provides clear organizational alignment, with roles directly corresponding to job functions within the development process. From an implementation perspective, this approach simplifies permission checking, typically requiring only verification of role membership rather than individual permission verification.

A key limitation of this approach is its reduced flexibility for accommodating exceptions or specialized access requirements. If a user requires a subset of permissions that doesn't align with existing roles, administrators must either create a new role specifically for that user or grant a role with more permissions than strictly necessary, potentially compromising the principle of least privilege [32].



4.3.2 Hybrid Role-Permission Approach

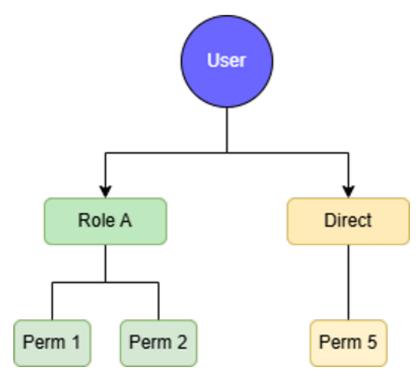


Figure 4.3: Hybrid Role-Permission Access Control Approach

The hybrid approach, illustrated in Figure 4.3, combines role-based permissions with direct permission assignments, similar to the model described by Ferraiolo et al. [15]. In this approach, users are assigned to primary roles defining their core permissions, but additional permissions can be granted on a per-user basis to address exceptional cases or specialized responsibilities.

This approach offers greater flexibility for accommodating exceptions without creating specialized roles, essential in environments where organizational structures evolve over time. It provides more granular permission control, allowing precise tailoring of access rights to individual responsibilities. However, this flexibility comes at the cost of increased administrative complexity, as both roles and individual permissions must be managed.

The hybrid approach is particularly valuable in the automotive parameter management context, where development responsibilities can vary between projects and temporary access adjustments may be needed for specific tasks or during transition periods. The approach balances structured role assignments with the flexibility to accommodate evolving access requirements, a common scenario in complex engineering environments like automotive development.



4.4 Parameter Synchronization Approaches

Integration with the Parameter Definition Database (PDD) represents a critical aspect of the VMAP system, requiring careful consideration of synchronization approaches. Two different conceptual approaches were explored for maintaining parameter data across the release phases: the change-based approach and the phase-based approach.

4.4.1 Change-Based Synchronization Approach

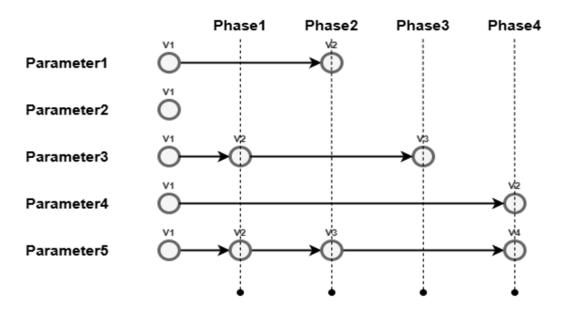


Figure 4.4: Change-Based Parameter Synchronization Approach

The change-based approach, illustrated in Figure 4.4, maintains parameter values by recording changes between phases rather than storing complete parameter sets for each phase. In this model, parameters are initially created before the first phase (Phase1), and subsequent modifications are recorded as change entries associated with specific phases.

When a parameter changes in a later phase (like Phase2), only the specific change is recorded rather than creating a new complete copy of the parameter. As shown in the figure, Parameter1 is created before Phase1 with version V1, then modified in Phase2 (creating version V2), but remains unchanged in Phase3 and Phase4. Similarly, Parameter3 changes in Phase2 and Phase3, while Parameter5 changes in every phase except Phase3.



This approach is conceptually aligned with traditional version control systems as described by Bhattacherjee et al. [4], where efficiency is achieved by storing only the differences between versions rather than complete copies. The approach potentially offers storage efficiency advantages by minimizing data duplication across phases, which could be significant for parameter sets containing thousands of entries.

However, this approach introduces conceptual complexity for retrieving parameter values in a specific phase. To determine a parameter's value for a given phase, the system must identify the most recent version of that parameter up to and including the target phase. This reconstruction process introduces additional processing steps compared to direct parameter retrieval.

4.4.2 Phase-Based Synchronization Approach

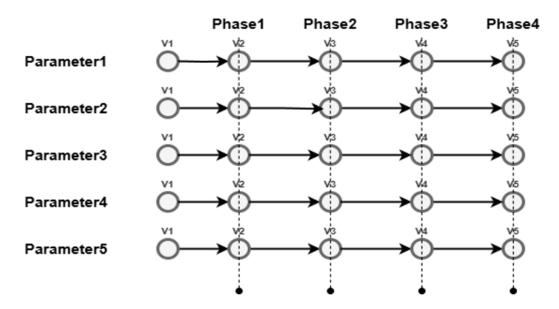


Figure 4.5: Phase-Based Parameter Synchronization Approach

The phase-based approach, illustrated in Figure 4.5, maintains complete parameter sets for each phase independently. When parameters are initially created before Phase1, each parameter has its specific version (V1). When transitioning to a new phase, all parameters are copied forward, even if they haven't changed. If a parameter is subsequently modified in the new phase, it receives a new version specific to that phase.

As shown in the figure, each parameter exists in every phase with a phase-specific version, regardless of whether the parameter value actually changed between phases.



This creates a clear separation between phases, with each phase maintaining its complete parameter configuration independently.

This approach aligns more directly with the phase-oriented structure of automotive development described by Broy [8], where distinct development milestones form the primary organizational principle. The approach simplifies conceptual understanding and parameter retrieval, as parameters for a specific phase can be accessed directly without reconstructing their values from change history.

The phase-based approach also simplifies phase inheritance by copying parameter configurations forward during phase transitions, allowing subsequent modifications in each phase without affecting previous phases. This copying mechanism preserves the integrity of phase data while supporting the automotive development process, where configurations stabilize progressively through successive phases.

The primary consideration with this approach is the increased storage requirements, as parameters are duplicated across phases even when they haven't changed. However, this trade-off potentially provides benefits in terms of conceptual clarity, query simplicity, and alignment with the automotive development workflow.

4.5 Database System Considerations

Selecting an appropriate database management system for VMAP required consideration of different options against the specific requirements of automotive parameter management. Four major relational database systems were considered as potential platforms for the VMAP implementation: PostgreSQL, Oracle, Microsoft SQL Server, and MySQL.

4.5.1 Database System Requirements

The requirements analysis identified several critical database capabilities needed for effective parameter management:

- 1. Support for complex data types, including arrays for multi-dimensional parameters and structured types for variant definitions.
- 2. Robust transaction support for maintaining data consistency during operations affecting multiple related entities.



- 3. Advanced indexing capabilities to optimize the performance of common query patterns, particularly parameter retrieval across different dimensions.
- 4. Extensibility for implementing domain-specific operations and validation rules.
- 5. Comprehensive access control mechanisms supporting the role-based security model.
- 6. Efficient storage and retrieval of historical data for audit and traceability purposes.

These requirements guided the evaluation of different database systems, focusing on their respective strengths and limitations in addressing the specific needs of automotive parameter management.

4.5.2 Comparative Analysis of Database Systems

Table 4.1 presents a comparative analysis of the four database systems considered for the VMAP implementation, evaluating each against criteria relevant to automotive parameter management.

The comparative analysis revealed different strengths among the database systems. PostgreSQL offers excellent support for complex data types and extensibility, particularly valuable for representing multi-dimensional parameters and implementing domain-specific operations. Oracle provides robust enterprise features with sophisticated optimization capabilities but introduces licensing complexity. SQL Server offers strong integration with Microsoft technologies, while MySQL provides simplicity but has limitations for complex data management requirements.

This analysis provides a foundation for database system selection, considering both technical capabilities and practical factors such as licensing and total cost of ownership. The implementation chapter will detail the specific database system selected and how its capabilities are leveraged in the VMAP implementation.

4.6 Entity-Relationship Model

Based on the requirements analysis and architectural considerations, a comprehensive entity-relationship (ER) model was developed to capture the complex relationships between system entities. This model follows the approach described by Chen [9], providing a conceptual foundation for the database implementation.



Table 4.1: Comparison of Database Systems for Automotive Parameter Management

| Feature | PostgreSQL | Oracle | SQL Server | MySQL |
|---------------|----------------------------|----------------------------|--------------------|-------------------------------|
| Complex | Excellent sup- | Good support, | Limited built- | Limited sup- |
| Data Types | port for arrays, | additional | in support, | port, improved |
| | JSON, custom | licensing for | extensions | in recent |
| | types [28] | advanced features [1] | required [3] | versions [42] |
| Transaction | Comprehensive | | Robust sup- | Limited in |
| Support | with serializ- | advanced op- | port with | some storage |
| | able isolation | tions [1] | multiple isola- | engines [33] |
| | [28] | | tion levels [3] | |
| Indexing Ca- | Diverse index | Advanced | Solid capa- | Basic indexing |
| pabilities | types includ- | indexing with | bilities with | with some lim- |
| | ing GIN for | optimizer | columnstore | itations [33] |
| | text search [28] | hints [1] | indexes [3] | |
| Extensibility | Highly extensi- | Extensible | Extensible | Limited exten- |
| | ble with cus- | with pro- | through .NET | sibility [42] |
| | tom types and | prietary | integration [3] | |
| | functions [28] | mechanisms | | |
| | | [1] | | |
| Access Con- | Fine-grained | • | Strong integra- | Basic capa- |
| trol | with role- | with advanced | tion with Ac- | bilities with |
| | based mecha- nisms [28] | security fea- tures [1] | tive Directory [3] | plugin archi- tecture [42] |
| Licensing | Open source, | Commercial, | Commercial | Dual licensing: |
| | PostgreSQL | complex | with edition- | GPL and com- |
| | License [28] | licensing | based pricing | mercial [42] |
| | _ | model [1] | [3] | |



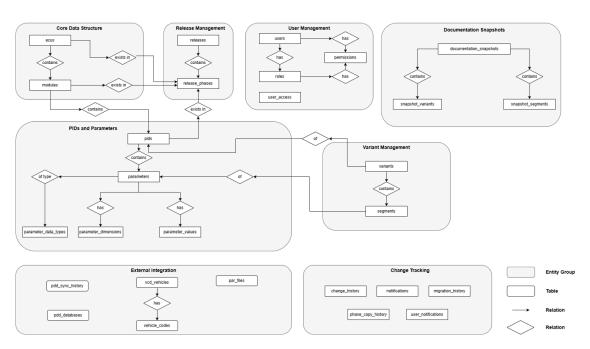


Figure 4.6: Comprehensive Entity-Relationship Diagram for the VMAP System

Figure 4.6 illustrates the complete entity-relationship model for theVMAP system, showing the logical organization of entities into functional groups and the relationships between them. The diagram captures the hierarchical structure of automotive parameter data, the phase-based versioning approach, the role-permission access control model, and the mechanisms for external system integration and change tracking.

4.6.1 Core Data Entities

The ER model includes several categories of entities representing different aspects of the system:

User management entities include Users, Roles, Permissions, and their relationships, implementing a role-permission model for access control. These entities support the authentication and authorization requirements identified during the requirements analysis.

Release management entities encompass Releases, Release Phases, and ECU Phase mappings, providing the foundation for the phase-based version control approach. These entities support the automotive development cycle with its distinct phases and milestones.

Parameter structure entities include ECUs, Modules, PIDs, Parameters, and Parameter Dimensions, representing the hierarchical organization of vehicle electronic systems as



described by Staron [37]. These entities capture both the structure and characteristics of parameters, including multi-dimensional parameters like maps and tables.

Variant management entities comprise Variants, Segments, and their relationships to parameters, implementing the core parameter customization functionality. These entities support the creation of parameter variations for different vehicle configurations and operating conditions.

Documentation entities include Documentation Snapshots, Snapshot Variants, and Snapshot Segments, supporting the preservation of historical parameter states for documentation and regulatory compliance. These entities enable the creation of complete parameter configuration snapshots at specific development milestones.

Integration entities consist of Synchronization Records, Vehicle Configurations, and Parameter File Records, supporting connectivity with the Parameter Definition Database and Vehicle Configuration Database. These entities track synchronization operations and store the necessary data for variant resolution and parameter file generation.

Audit entities encompass Change History, Transaction Records, and Phase Copy History, providing comprehensive traceability for all significant operations within the system. These entities support both regulatory compliance and diagnostic capabilities for investigating parameter evolution.

4.6.2 Relationship Structure

The relationships between entities in the ER model reflect the complex interactions between different aspects of automotive parameter management. Key relationships include:

Hierarchical relationships between ECUs, Modules, PIDs, and Parameters, representing the structural organization of automotive electronic systems. These relationships enforce the domain-specific hierarchy while supporting navigation from higher-level entities to their components.

Many-to-many relationships between parameters and phases, implemented through direct association rather than temporal versioning. This structure supports the phase-based versioning approach, allowing efficient retrieval of parameters for specific phases.

Complex relationships between variants, parameters, and segments, capturing the parameter customization process. These relationships ensure that segments are



associated with valid variants and parameters while supporting efficient resolution of effective parameter values based on vehicle configuration.

Temporal relationships for audit and history entities, capturing the evolution of parameter configurations over time. These relationships support both regulatory compliance and diagnostic capabilities, allowing reconstruction of parameter states at specific points in time.

4.6.3 Normalization and Optimization

The ER model was developed using data normalization principles to minimize redundancy while maintaining data integrity, following the approach described by Codd [10]. The model generally adheres to Third Normal Form (3NF), ensuring that non-key attributes depend on the primary key rather than on other non-key attributes.

Strategic denormalization was considered in specific areas to optimize performance for common operations, following the principles described by Molinaro [25]. For example, certain frequently accessed attributes from parent entities might be duplicated in child entities to reduce join operations in common queries, providing performance benefits that outweigh the controlled redundancy.

4.7 Validation Mechanisms

To ensure data integrity and consistency, multiple validation mechanism layers were conceptualized for the VMAP system, from basic constraints to sophisticated business rule validation. These mechanisms work together to maintain parameter data quality and reliability throughout the system lifecycle.

4.7.1 Data Integrity Constraints

Database-level constraints were identified as the foundation for enforcing basic integrity rules, following the principles described by Elmasri and Navathe [14]. These constraints include primary key constraints ensuring unique entity identifiers, foreign key constraints maintaining referential integrity between related entities, not-null constraints ensuring required fields contain values, unique constraints preventing duplicate values in specified columns, and check constraints enforcing domain-specific rules such as valid date ranges and parameter value ranges.



These constraints are designed into the database schema as integral parts of entity definitions, ensuring consistent enforcement throughout the system regardless of access path. By implementing constraints at the database level rather than in application code,VMAP ensures that all data modifications adhere to fundamental integrity rules regardless of the source of those modifications.

4.7.2 Business Rule Validation

Domain-specific business rules were conceptualized to be implemented through database triggers and stored procedures, providing a second layer of validation beyond basic constraints. These rules include parameter range validation automatically checking modified values against defined minimum and maximum bounds, phase status validation preventing modifications to frozen phases, segment validation ensuring segments reference valid parameters and variants, and user access validation ensuring users can only modify parameters, variants, and segments for modules to which they have been granted access.

4.7.3 Conflict Resolution Strategies

In a multi-user environment, concurrent modifications can create conflicts. Several strategies were conceptualized to detect and resolve these conflicts, maintaining data consistency while supporting collaborative parameter management.

For web-based interactions, optimistic concurrency control using version timestamps allows multiple users to view the same data concurrently, detecting conflicts only when updates collide. This approach aligns with the principles described by Bhattacherjee et al. [4], optimizing for the common case where conflicts are rare while still ensuring data integrity.

When changes propagate from one phase to the next, conflicts can arise if the target phase has already been modified. Explicit conflict resolution mechanisms compare the source variant or segment with existing target configurations during phase propagation operations. When conflicts are detected, resolution options allow users to make informed decisions: override the target with source values, preserve target values, or merge values based on specified rules.



4.7.4 Audit and Traceability Mechanisms

Comprehensive audit and traceability mechanisms were identified as essential for regulatory compliance and quality assurance in automotive parameter management. The core of this capability is the change history tracking mechanism, which automatically captures both before and after states for entity modifications. For each change, the system records the entity being modified, the type of change, the user making the change, the timestamp, and detailed before/after values.

To optimize performance, selective filtering of change data was conceptualized, excluding non-essential fields such as timestamps and large binary data. Additionally, asynchronous audit recording for bulk operations was considered, to reduce the performance impact on high-volume operations while ensuring that all changes are eventually recorded.

Beyond change tracking, specialized audit mechanisms were conceptualized for specific scenarios: phase transition logging to record all phase propagation operations, freeze operation logging to record phase freeze and unfreeze operations, user access logging to capture authentication and authorization events, and integration operation logging to record all external system interactions.

5 Implementation

This chapter presents the technical implementation of theVMAP system conceptual architecture described in Chapter 4. The discussion focuses on key architectural components and technical decisions that enable efficient parameter versioning in automotive software development, with concrete implementation details demonstrating how theoretical concepts translate to practical solutions.

5.1 Database Structure Implementation

Following the comparative analysis of database systems presented in Chapter 4, PostgreSQL was selected as the implementation platform due to its superior support for complex data types, extensibility features, and advanced indexing capabilities [28]. The database implementation transforms abstract entities and relationships into concrete database structures, implementing the hierarchical organization of automotive electronic systems through physical tables and relationships.

5.1.1 Core Data Entities

The hierarchical structure of automotive electronic systems is implemented through four primary entity types: ECUs, Modules, PIDs, and Parameters. The ECU and Module entities form the top levels of the hierarchy, with a many-to-many relationship reflecting the reality that modules can exist across multiple ECUs. This structure aligns with the domain model described by Staron [37], where logical groupings of software functions must be maintained across different hardware configurations.



```
CREATE TABLE modules (
      module id INTEGER PRIMARY KEY,
      shortcut VARCHAR (50) NOT NULL,
      name VARCHAR (255) NOT NULL,
      kind VARCHAR (255),
      created at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      external id INTEGER UNIQUE -- \ac{PDD}
16
        breference ID
  );
18
  CREATE TABLE ecu modules (
      ecu_id INTEGER REFERENCES ecus(ecu_id)
20
        □DELETE CASCADE,
      module id INTEGER REFERENCES modules (
        by module id) ON DELETE CASCADE,
      created at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      PRIMARY KEY (ecu id, module id)
  );
```

Listing 5.1: ECU and Module Table Implementation

The junction table ecu_modules implements the many-to-many relationship using foreign key constraints as described by Elmasri and Navathe [14]. Each entity includes an external_id attribute to maintain mapping with the Parameter Definition Database (PDD), facilitating integration through persistent identifier correlation as recommended by Hohpe and Woolf [18].

PIDs (Parameter IDs) and parameters constitute the lower levels of the hierarchy, with PIDs grouping related parameters for specific modules. The parameter table implements additional attributes for handling the complexity of automotive parameter specifications:

```
create table parameters (
parameter_id BIGINT PRIMARY KEY,
pid_id BIGINT REFERENCES pids(pid_id) ON
LDELETE CASCADE,
ecu_id INTEGER,
```



```
phase id INTEGER,
      name VARCHAR (255) NOT NULL,
      parameter_name VARCHAR (255),
      type id INTEGER REFERENCES
        parameter_data_types(data_type_id),
      array_definition VARCHAR(50),
      position INTEGER,
      factor DECIMAL,
      unit VARCHAR (50),
      bias_offset DECIMAL,
13
      is active BOOLEAN DEFAULT true,
14
      external_id INTEGER, -- \ac{PDD}reference
        ↓ I D
      FOREIGN KEY (ecu id, phase id) REFERENCES
16
        \ecu_phases(ecu_id, phase_id)
  );
```

Listing 5.2: Parameter Table Implementation

The parameter table incorporates a form of strategic denormalization by including direct references to ecu_id and phase_id alongside thePIDforeign key. While this introduces some redundancy, this approach aims to improve query execution for parameter queries filtered by release phase, a critical operation in the system. Bhattacherjee et al. [4] note that such denormalization can be justified when query performance benefits outweigh the overhead of maintaining consistency, especially for frequently accessed paths in the data model.

A particularly challenging aspect of the implementation was supporting multi-dimensional parameters, which are common in automotive applications for representing lookup tables and characteristic curves [22]. This was addressed through a specialized table structure:



```
UNIQUE (parameter_id, dimension_index)
);
```

Listing 5.3: Parameter Dimension Implementation

This implementation follows a modified entity-attribute-value (EAV) pattern while addressing the potential challenges typically associated with such models [13]. The introduction of a dimension index provides an ordered structure to parameter dimensions, enabling efficient representation of arrays and matrices while maintaining the relationship between parameters and their dimensional values.

5.1.2 Version Control Implementation

The version control implementation is a defining feature of the VMAP system, enabling parameter evolution management across different development phases. After evaluating multiple approaches described in Chapter 4, the phase-based versioning model was implemented, creating explicit relationships between parameters and development phases rather than using a generic temporal approach.

```
CREATE TABLE releases (
      release id INTEGER PRIMARY KEY,
      name VARCHAR (50) NOT NULL UNIQUE,

  □
  "24.1"

  □
  "24.3"

      description TEXT,
      is active BOOLEAN DEFAULT true,
      created_at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP ,

      created by BIGINT REFERENCES users (user id)
  );
  CREATE TABLE release_phases
      phase id INTEGER PRIMARY KEY,
      release id INTEGER REFERENCES releases(
         \release_id) ON DELETE CASCADE,
      name VARCHAR (50) NOT NULL, -- e.g.,
         Lanitial", "PreTest1", "PreTest2", "Final
         L II
      sequence_number INTEGER NOT NULL,
14
         Determines the order of phases
```



```
is active BOOLEAN DEFAULT true,
      created_at TIMESTAMP WITH TIME ZONE DEFAULT
16

↓ CURRENT_TIMESTAMP,

      created by BIGINT REFERENCES users(user_id)
17
      UNIQUE (release id, name),
18
      UNIQUE (release id, sequence number)
19
  );
20
21
  CREATE TABLE ecu_phases (
      ecu id INTEGER REFERENCES ecus(ecu id)
23
         DELETE CASCADE,
      phase id INTEGER REFERENCES release phases (
24

¬phase_id) ON DELETE CASCADE,

      is_active BOOLEAN DEFAULT true,
25
      is frozen BOOLEAN DEFAULT false,
      frozen at TIMESTAMP WITH TIME ZONE,
27
      frozen by BIGINT REFERENCES users (user id),
28
      created at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      created by BIGINT REFERENCES users (user id)
30
      PRIMARY KEY (ecu id, phase id)
31
  );
```

Listing 5.4: Release and Phase Management Implementation

This implementation supports the bi-annual release cycle with four sequential phases per release. The sequence_number field provides explicit ordering of phases within a release, while unique constraints ensure consistency of phase naming and sequencing. The ECU-phase mapping implements the association between ECUs and specific release phases, supporting independent progression of different ECUs through the development cycle as described by Broy [8].

The phase-based approach was selected over temporal versioning for several reasons supported by database design principles:

 Direct phase associations align with the process-based nature of automotive development, where parameters evolve through explicitly defined development stages.



- The explicit phase relationships simplify the implementation of phase transitions and comparison operations, which are fundamental to the automotive development process as described by Pretschner et al. [29].
- The approach aligns with the mental model of automotive development engineers, who conceptualize parameter evolution in terms of distinct development phases rather than continuous time.

Phase management functionality was implemented through a stored procedure for phase transitions and parameter propagation:

```
CREATE OR REPLACE FUNCTION copy phase data(
      source_ecu_id INTEGER,
      source_phase_id_INTEGER,
      target ecu id INTEGER,
      target_phase_id INTEGER,
      user id BIGINT
  )
  RETURNS TABLE (
      variants_copied INTEGER,
      segments copied INTEGER
   AS $$
  DECLARE
      variant count INTEGER
      segment_count INTEGER
      transaction id BIGINT;
  BEGIN
      -- Get a transaction ID for change tracking
      SELECT nextval('
18
        $\change history transaction id seq')
        transaction_id;
19
      -- Set the transaction ID for tracking in
        ↓logging trigger
      PERFORM set_config('app.transaction_id',
21
        \transaction_id::text, true);
      PERFORM set_config('app.user_id', user_id::
        btext, true);
```



```
-- Create a temporary table to map source
        4 variant IDs to target variant IDs
      CREATE TEMPORARY TABLE variant mapping (
25
          source variant id BIGINT,
26
          target_variant_id_BIGINT
27
      ) ON COMMIT DROP;
28
      -- Copy variants with modified attributes
        ↓for target phase
      WITH inserted_variants AS (
          INSERT INTO variants (
32
              pid_id, ecu_id, phase_id, name,
                 created at, created by, updated by
34
          )
35
          SELECT
              v.pid_id, target_ecu_id,
37
                 \target_phase_id, v.name, v.
                 CURRENT_TIMESTAMP, user id, user id
38
          FROM variants v
39
          WHERE v.ecu id = source ecu id AND v.
            bphase_id = source_phase_id
          RETURNING variant_id, pid_id, name,
41
            ↓code rule
42
      INSERT INTO variant mapping (
43
        \source_variant_id, target_variant_id)
      SELECT v.variant_id, iv.variant_id
44
      FROM variants v
45
      JOIN inserted variants iv ON v.pid id = iv.
        AND v.name = iv.
47
                                  ⊣name
                                AND v.code rule =
48
                                  ↓iv.code rule
      WHERE v.ecu_id = source_ecu_id AND v.
        phase id = source phase id;
50
      -- Get count of copied variants
```



Listing 5.5: Phase Transition Function

This implementation provides an atomic transaction for phase transition, copying parameter variants and segments from one phase to another while maintaining proper references and tracking the operation in the change history. The function returns counts of affected entities, providing immediate feedback on the operation's scope. The use of temporary tables for mapping entities between phases implements the identity map pattern described by Fowler [16], ensuring proper relationship preservation during complex operations.

5.1.3 Variant and Segment Management

The variant and segment management implementation realizes the core parameter customization capabilities of the VMAP system:



```
updated by BIGINT REFERENCES users (user id)
10
      FOREIGN KEY (ecu id, phase id) REFERENCES
         \ecu_phases(ecu_id, phase_id),
      UNIQUE (phase id, pid id, name),
12
      UNIQUE (phase_id, pid_id, code_rule)
13
  );
14
15
  CREATE TABLE segments
16
      segment id BIGINT PRIMARY KEY,
17
      variant id BIGINT REFERENCES variants (
18
         byvariant id) ON DELETE CASCADE,
      parameter id BIGINT REFERENCES parameters (
19
         bparameter_id) ON DELETE CASCADE,
      dimension index INTEGER NOT NULL,
20
      decimal NUMERIC NOT NULL,
      created at TIMESTAMP WITH TIME ZONE DEFAULT
22

↓ CURRENT_TIMESTAMP,

      created by BIGINT REFERENCES users (user id)
      updated by BIGINT REFERENCES users (user id)
24
  );
25
```

Listing 5.6: Variant and Segment Implementation

Each variant is associated with a specific PID, ECU, and phase, creating a three-way relationship that places the variant in both functional and temporal contexts. The uniqueness constraints prevent duplicate variant names or code rules within the samePIDand phase, enforcing a critical business rule identified during requirements analysis. The code_rule field stores boolean expressions determining when a variant applies based on vehicle configuration codes, implementing a domain-specific rule language for variant applicability.

Segments form the foundation of parameter customization, linking parameters to variants and storing modified values. The canonical decimal representation for all parameter values, regardless of native data type, implements the canonical model pattern described by Hohpe and Woolf [18], simplifying data manipulation while ensuring consistent value handling.

For parameter value resolution, a specialized database function was implemented:



```
CREATE OR REPLACE FUNCTION
    \resolve parameter value(
      p_parameter_id BIGINT,
      p_dimension_index INTEGER,
      p_variant_ids BIGINT[]
 RETURNS NUMERIC AS $$
  DECLARE
      v value NUMERIC;
      v_default_value NUMERIC;
 BEGIN
      -- First try to find a segment with the
        bgiven parameter and variant
      SELECT s.decimal INTO v value
      FROM segments s
13
      WHERE s.parameter id = p parameter id
        AND s.dimension_index = p_dimension_index
        AND s.variant_id = ANY(p_variant_ids)
16
      ORDER BY array_position(p_variant_ids, s.
        LIMIT 1;
18
      -- If no segment found, get the default
20

    value

      IF v value IS NULL THEN
          SELECT pd.default_value INTO
            FROM parameter_dimensions pd
          WHERE pd.parameter_id = p_parameter_id
24
            AND pd.dimension_index =
25
              bp_dimension_index;
26
          -- If no dimension record, try
            parameter default
          IF v_default_value IS NULL
                                      THEN
28
              SELECT p.default value

    v_default_value

              FROM parameters p
30
              WHERE p.parameter_id =
```



Listing 5.7: Parameter Resolution Function

This function implements the value resolution logic required for parameter file generation, handling the complex logic of finding applicable parameter values based on variant precedence. The array position-based ordering ensures that variants are applied in the correct priority sequence, a critical requirement for automotive parameter management as described by Staron [37].

To support documentation and compliance requirements, a snapshot mechanism was implemented to capture complete parameter configurations at specific points in time:

```
CREATE TABLE documentation_snapshots
      snapshot id INTEGER PRIMARY KEY,
      name VARCHAR (255) NOT NULL,
      description TEXT,
      ecu id INTEGER,
      phase_id INTEGER,
      variant_count INTEGER DEFAULT 0,
      segment_count INTEGER DEFAULT 0,
      created at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      created by BIGINT REFERENCES users (user id)
10
        Ļ,
      FOREIGN KEY (ecu id, phase id) REFERENCES
11
        \ecu_phases(ecu_id, phase_id)
  );
12
13
  CREATE TABLE snapshot_variants (
      snapshot variant id INTEGER PRIMARY KEY,
```



```
snapshot id INTEGER REFERENCES
        4 documentation snapshots (snapshot id)
        □DELETE CASCADE,
      original_variant_id BIGINT, -- Reference
        by the original variant
      pid_id BIGINT REFERENCES pids(pid_id)
18
        DELETE CASCADE,
      name VARCHAR (100) NOT NULL,
      code rule TEXT,
20
      created at TIMESTAMP WITH TIME ZONE,
      created by BIGINT REFERENCES users (user id)
  );
24
  CREATE TABLE snapshot segments (
      snapshot segment id INTEGER PRIMARY KEY,
26
      snapshot id INTEGER REFERENCES

\documentation_snapshots(snapshot_id) ON

        DELETE CASCADE,
      snapshot_variant_id INTEGER REFERENCES
28
        \snapshot_variants(snapshot_variant_id)
        ↓ON DELETE CASCADE,
      original segment id BIGINT,
                                   -- Reference to
        4 the original segment
      parameter id BIGINT REFERENCES parameters (
30
        parameter_id) ON DELETE CASCADE,
      dimension_index INTEGER NOT NULL,
      decimal NUMERIC NOT NULL,
      created at TIMESTAMP WITH TIME ZONE,
      created by BIGINT REFERENCES users (user id)
  );
35
```

Listing 5.8: Documentation Snapshot Implementation

This implementation follows the snapshot pattern described by Fowler [16], creating a complete copy of variant and segment data at specific points in time. Rather than using a temporal database approach with validity periods, the system explicitly materializes historical states, ensuring they remain accessible regardless of subsequent modifications to live data. References to original entities enable traceability between snapshot and live data, implementing the origin tracking pattern described by Tichy [40].



The snapshot creation function automates the process of capturing parameter configurations:

```
CREATE OR REPLACE FUNCTION

\create_documentation_snapshot(
      p_name VARCHAR(255),
      p_description TEXT,
      p_ecu_id INTEGER,
      p phase id INTEGER,
      p_user_id BIGINT
  )
  RETURNS INTEGER AS $$
 DECLARE
      v snapshot id INTEGER;
  BEGIN
11
      -- Create the snapshot record
      INSERT INTO documentation snapshots (
13
          name, description, ecu_id, phase_id,
14
          variant_count, segment_count,
             \created_at, created_by
      ) VALUES (
16
          p_name, p_description, p_ecu_id,
17
            0, 0, CURRENT_TIMESTAMP, p_user_id
18
      ) RETURNING snapshot id INTO v snapshot id;
20
      -- Copy all variants for the ECU and phase
21
      INSERT INTO snapshot variants (
          snapshot_id, original_variant_id,
23
             bpid_id, name,
          code_rule, created_at, created_by
      )
25
      SELECT
26
          v_snapshot_id, v.variant_id, v.pid_id,

    v.name,

          v.code_rule, v.created_at, v.created_by
28
      FROM variants v
      WHERE v.ecu_id = p_ecu_id AND v.phase_id =
30

¬p_phase_id;
```



Listing 5.9: Documentation Snapshot Function

In the automotive industry, these snapshots serve multiple purposes: providing immutable records of parameter configurations at significant development milestones, supporting quality assurance processes and regulatory compliance requirements, and facilitating comparative analysis between development phases.

5.2 Access Control Implementation

The access control implementation realizes the hybrid role-permission model described in Chapter 4, providing a flexible foundation for managing user permissions across the system.

5.2.1 Role-Based Permission Model

The core RBAC implementation follows the structure defined by Sandhu et al. [32], with tables for users, roles, permissions, and their relationships:

```
create table roles (
role_id INTEGER PRIMARY KEY,
name VARCHAR(255) UNIQUE NOT NULL,
description TEXT
);

CREATE TABLE permissions (
```



```
permission id INTEGER PRIMARY KEY,
      name VARCHAR (255) UNIQUE NOT NULL,
      description TEXT
  );
12
  CREATE TABLE role_permissions (
13
      role id INTEGER REFERENCES roles (role id)
         →ON DELETE CASCADE,
      permission id INTEGER REFERENCES
         permissions (permission id) ON DELETE

        ↓ CASCADE ,

      granted_at TIMESTAMP WITH TIME ZONE DEFAULT
16

↓ CURRENT_TIMESTAMP,

      granted by BIGINT REFERENCES users (user id)
      PRIMARY KEY (role id, permission id)
  );
19
20
  CREATE TABLE user_roles (
      user id BIGINT REFERENCES users (user id) ON
         → DELETE CASCADE,
      role id INTEGER REFERENCES roles (role id)
         ↓ON DELETE CASCADE,
      granted at TIMESTAMP WITH TIME ZONE DEFAULT
24

↓ CURRENT_TIMESTAMP,

      granted_by BIGINT REFERENCES users(user_id)
      PRIMARY KEY (user_id, role_id)
 );
27
```

Listing 5.10: Core RBAC Implementation

The model is extended with direct user permissions to implement the hybrid approach:

```
CREATE TABLE user_permissions (
user_permission_id INTEGER PRIMARY KEY,
user_id BIGINT REFERENCES users(user_id) ON

DELETE CASCADE,
permission_id INTEGER REFERENCES

permissions(permission_id) ON DELETE
```



Listing 5.11: User-Permission Implementation

This hybrid model allows permissions to be granted either through role assignments or directly to users, providing the flexibility to handle exceptional cases without creating specialized roles, as recommended by Ferraiolo et al. [15].

The permission checking function implements the core security verification logic:

```
CREATE OR REPLACE FUNCTION has_permission(
     p user id BIGINT,
     p_permission_name VARCHAR
 )
 RETURNS BOOLEAN AS $$
 BEGIN
     RETURN EXISTS (
         -- Check role-based permissions
8
         SELECT 1
         FROM user_roles ur
10
         JOIN role permissions rp ON ur.role id
           JOIN permissions p ON rp.permission_id
12
           WHERE ur.user id = p user id AND p.name
            = p_permission_name
14
         UNION
15
16
         -- Check direct user permissions
         SELECT 1
         FROM user_permissions up
19
         JOIN permissions p ON up.permission_id
20
```



Listing 5.12: Permission Check Implementation

The UNION-based approach combines role-based and direct permission checks in a single database query, implementing an efficient verification mechanism. This implementation supports the separation of permission checking from application logic, allowing security policies to be enforced consistently across different access paths.

5.2.2 Module-Based Access Control

To complement the role-based permission model, a module-based access control system was implemented to restrict write access based on module assignments:

```
CREATE TABLE user access (
      user id BIGINT REFERENCES users (user id)
        → DELETE CASCADE,
      ecu id INTEGER REFERENCES ecus(ecu id) ON
        □DELETE CASCADE,
      module id INTEGER REFERENCES modules (
        write_access BOOLEAN DEFAULT true,
      created at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      created by BIGINT REFERENCES users (user id)
        Ļ,
      PRIMARY KEY (user_id, ecu_id, module_id),
      CONSTRAINT user_access ecu module fk
        FOREIGN KEY (ecu id, module id)
      REFERENCES ecu modules (ecu id, module id)
10
  );
11
```

Listing 5.13: Module Access Control Implementation



This implementation establishes a three-way relationship between users, ECUs, and modules, with a boolean flag distinguishing between read and write access. The constraint ensures access is granted only for valid ECU-module combinations, enforcing structural integrity. Database functions verify both role-based permissions and module-based access rights during critical operations:

```
CREATE OR REPLACE FUNCTION
    has_module_write_access(
      p user id BIGINT,
      p_ecu_id INTEGER,
      p_module_id_INTEGER
  )
  RETURNS BOOLEAN AS
                     $$
  BEGIN
      -- First check if user has administrator
        4role (bypasses module restrictions)
      IF EXISTS (
9
          SELECT 1
          FROM user_roles ur
          JOIN roles r ON ur.role_id = r.role_id
          WHERE ur.user_id = p_user_id AND r.name
            THEN
14
          RETURN TRUE;
      END IF;
      -- Check module-specific write access
18
      RETURN EXISTS (
          SELECT 1
20
          FROM user_access ua
          WHERE ua.user_id = p_user_id
            AND ua.ecu_id = p_ecu_id
            AND ua.module id = p module id
24
            AND ua.write access = TRUE
      );
26
  END:
  $$ LANGUAGE plpgsql;
```

Listing 5.14: Module Access Check Implementation



This implementation combines role-based and attribute-based access control approaches, creating what Ferraiolo et al. describe as policy-enhanced RBAC [15]. The function first checks for administrator privileges, which bypass module-specific restrictions, then verifies specific module access rights if necessary.

5.3 Query Optimization Implementation

The complex relationships and large data volumes in automotive parameter management necessitated a comprehensive indexing strategy to support common query patterns.

5.3.1 Indexing Implementation

The indexing strategy focuses on accelerating hierarchical navigation and parameter retrieval operations:

```
-- Hierarchical path indexes
2 CREATE INDEX idx_pids_ecu_module ON pids(ecu_id
    ↳, module id);
 CREATE INDEX idx_parameters_pid_phase ON
    parameters(pid id, phase id);
 CREATE INDEX idx variants pid phase ON variants
    \(pid_id, phase_id);
5 CREATE INDEX idx segments variant ON segments (

¬variant id);
 CREATE INDEX idx_segments_parameter ON segments
    \( parameter_id);
 -- Phase-specific indexes
 CREATE INDEX idx parameters phase ON parameters
   \( phase_id);
 CREATE INDEX idx variants phase ON variants (
    -- Text search optimization
 CREATE EXTENSION IF NOT EXISTS pg trgm;
 CREATE INDEX idx parameters name trgm
```



```
ON parameters USING gin (name gin_trgm_ops);

CREATE INDEX idx_variants_name_trgm
ON variants USING gin (name gin_trgm_ops);
```

Listing 5.15: Core Index Implementation

These indexes match the natural navigation paths in the data model, implementing the access path optimization pattern described by Molinaro [25]. The hierarchical path indexes support efficient traversal from ECUs through modules and PIDs to parameters, while the phase-specific indexes accelerate queries filtered by release phase, a common operation in the system.

The trigram-based text search indexes utilize PostgreSQL's specialized text search capabilities as described by Obe and Hsu [28]. These indexes enable efficient pattern matching and similarity-based searches, accommodating the fuzzy search requirements identified during user interviews. According to Karwin [21], these specialized indexes can significantly improve the performance of text-based queries, which are common in parameter management systems where engineers need to locate parameters by name or partial name.

5.3.2 Function-Based Optimization

To optimize complex parameter retrieval operations, several specialized database functions were implemented:

```
CREATE OR REPLACE FUNCTION search_parameters(
search_term TEXT,
max_results INT DEFAULT 20,
ecu_id_filter INT DEFAULT NULL,
phase_id_filter INT DEFAULT NULL
)
RETURNS TABLE (
parameter_id BIGINT,
name VARCHAR,
parameter_name VARCHAR,
pid_id BIGINT,
pid_name VARCHAR,
module id INTEGER,
```



```
module name VARCHAR,
      ecu id INTEGER,
15
      ecu name VARCHAR,
16
      phase id INTEGER,
17
      relevance REAL
18
    AS $$
  BEGIN
    RETURN QUERY
21
    SELECT
22
      p.parameter_id,
23
      p.name,
24
      p.parameter_name,
25
      pid.pid id,
26
      pid.name AS pid name,
      m.module_id,
28
      m.name AS module name,
      e.ecu_id,
30
      e.name AS ecu_name,
31
      p.phase_id,
      -- Calculate relevance score based on
33
         4 different matching criteria
       (CASE
           -- Exact matches get highest score
35
           WHEN p.name ILIKE search term OR p.
36
             parameter_name ILIKE search_term
             → THEN 1.0
           -- Starts with gets high score
37
           WHEN p.name ILIKE (search_term || '%')
38
             \downarrow search_term || '\',') THEN 0.8
           -- Contains gets medium score
           ELSE GREATEST (
40
               similarity(p.name, search term),
41
               similarity(p.parameter_name,
42

\downarrow search_term)
43
      END) AS relevance
    FROM
45
      parameters p
46
    JOIN
```



```
pids\ac{PID}ON p.pid_id = pid.pid_id
           JOIN
49
                 modules m ON pid.module_id = m.module_id
            JOIN
                 ecus e ON pid.ecu_id = e.ecu_id
           WHERE
53
                 -- Apply filters if provided
                 (ecu_id_filter IS NULL OR e.ecu_id =
                       becu id filter) AND
                 (phase id filter IS NULL OR p.phase id =
                       bphase id filter) AND
                 -- Match parameter name or display name
                 (p.name ILIKE '%' || search term || '%'
                    p.parameter_name ILIKE '%' || search_term

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                 -- Only active parameters
                 p.is_active = true
           ORDER BY
62
                 relevance DESC,
                p.name ASC
           LIMIT max results;
65
     END;
     $$ LANGUAGE plpgsql;
```

Listing 5.16: Parameter Search Function

This function implements a sophisticated parameter search algorithm that combines exact matches, prefix matches, and similarity-based matching to provide relevant search results. The relevance scoring approach ensures that the most relevant parameters appear first in the results, following the information retrieval principles described by Obe and Hsu [28].

For phase comparison operations, a specialized function implements the parameter difference detection logic:



```
p target phase id INTEGER
  )
6
  RETURNS TABLE (
      pid_id BIGINT,
      name VARCHAR,
q
      has_param_changes BOOLEAN,
      has variant changes BOOLEAN,
      has segment changes BOOLEAN
  ) AS $
13
  BEGIN
      RETURN QUERY
15
      WITH common pids AS (
16
          -- Get PIDs available in at least one
17
             ↓of the phases
          SELECT DISTINCT p.pid id, p.name
18
          FROM pids p
19
          LEFT JOIN pid_phases pp1 ON p.pid_id =
20
             pp1.pid id AND pp1.phase id =
             bp_source_phase_id
          LEFT JOIN pid phases pp2 ON p.pid id =
21
             pp2.pid id AND pp2.phase id =
             WHERE p.ecu_id = p_ecu_id AND p.
22
             bmodule_id = p_module_id
            AND (pp1.phase_id IS NOT NULL OR pp2.
23
               phase_id IS NOT NULL)
      ),
24
         Parameter changes detection
      param_changes AS (
26
          SELECT cp.pid_id
27
          FROM common pids cp
          WHERE EXISTS (
               -- Parameters that exist in target
30
                 but not in source
               SELECT 1 FROM parameters pt
31
              WHERE pt.pid_id = cp.pid_id
32
                 AND pt.phase_id =
                   ↓p_target_phase_id
                 AND NOT EXISTS (
34
                     SELECT 1 FROM parameters ps
```



```
WHERE ps.pid id = cp.pid id
                        AND ps.phase id =
37
                          bp_source_phase_id
                        AND ps.external_id = pt.
38
                          bexternal_id
                 )
39
          ) OR EXISTS (
40
               -- Parameters that exist in source
                 but not in target
               SELECT 1 FROM parameters ps
42
               WHERE ps.pid id = cp.pid id
43
                 AND ps.phase_id =
44
                   bp_source_phase_id
                 AND NOT EXISTS (
45
                      SELECT 1 FROM parameters pt
46
                     WHERE pt.pid id = cp.pid id
                        AND pt.phase id =
48
                          bp_target_phase_id
                        AND pt.external_id = ps.
49
                          |external id
50
           ) OR EXISTS (
               -- Parameters that exist in both
52
                 but have differences
               SELECT 1 FROM parameters ps
               JOIN parameters pt ON ps.
                 \external_id = pt.external_id
               WHERE ps.pid_id = cp.pid_id
55
                 AND pt.pid_id = cp.pid_id
56
                 AND ps.phase_id =
57
                   ↓p source phase id
                 AND pt.phase_id =
58
                   ↓p target phase id
                 AND (
                     ps.name != pt.name OR
60
                     ps.parameter_name != pt.
61
                        bparameter_name OR
                     ps.type_id != pt.type_id OR
62
                     ps.factor != pt.factor OR
63
                     ps.bias_offset != pt.
```



```
bias offset OR
                      ps.unit != pt.unit
65
                  )
66
           )
67
      ),
68
       -- Variant changes detection
69
      variant changes AS (
           SELECT cp.pid_id
           FROM common pids cp
           WHERE EXISTS (
73
                -- Different variant counts between
74

↓ phases

               SELECT 1
75
               FROM (
76
                    SELECT pid id, count(*) as cnt
77
                    FROM variants
78
                    WHERE pid_id = cp.pid_id AND
79
                       phase_id = p_source_phase_id
                    GROUP BY pid_id
81
                    UNION ALL
82
                    SELECT pid_id, -count(*) as cnt
84
                    FROM variants
85
                    WHERE pid_id = cp.pid_id AND
                       bphase_id = p_target_phase_id
                    GROUP BY pid id
87
                ) v
               GROUP BY v.pid_id
89
               HAVING sum(v.cnt) != 0
90
           ) OR EXISTS (
                -- Variants with same name but
92
                  \different code rules
               SELECT 1
               FROM variants vs
94
                JOIN variants vt ON vs.name = vt.
95

\downarrow
 name
               WHERE vs.pid_id = cp.pid_id
96
                  AND vt.pid_id = cp.pid_id
97
                  AND vs.phase_id =
```



```
AND vt.phase id =
99
                   bp_target_phase_id
                 AND vs.code_rule != vt.code_rule
100
           )
101
       ),
102
          Segment changes detection
103
       segment_changes AS (
104
           SELECT cp.pid_id
105
           FROM common pids cp
106
           JOIN parameters ps ON ps.pid_id = cp.
             bpid_id AND ps.phase_id =
             bp_source_phase_id
           JOIN parameters pt ON pt.pid id = cp.
108
             bpid_id AND pt.phase_id =
             ↓p target phase id
                             AND pt.external_id =
109
                               □ps.external id
           JOIN variants vs ON vs.pid_id = cp.
110
             bpid_id AND vs.phase_id =

¬p_source_phase_id
           JOIN variants vt ON vt.pid id = cp.
111
             bpid_id AND vt.phase_id =
             bp_target_phase_id
                           AND vt.name = vs.name
           JOIN segments ss ON ss.parameter_id =
113
             \ps.parameter_id AND ss.variant_id =
             \lorvs.variant_id
           JOIN segments st ON st.parameter_id =
114
             bpt.parameter_id AND st.variant_id =
             AND st.dimension_index =
115

↓ ss.dimension index

           WHERE ss.decimal != st.decimal
116
117
           UNION
118
119
           -- Segments in source not in target
120
           SELECT cp.pid_id
121
           FROM common pids cp
```



```
JOIN parameters ps ON ps.pid_id = cp.
             ↓pid id AND ps.phase id =
             bp_source_phase_id
           JOIN parameters pt ON pt.pid_id = cp.
124
             bpid_id AND pt.phase_id =
             bp_target_phase_id
                            AND pt.external id =

bps.external_id
           JOIN variants vs ON vs.pid id = cp.
126
             bpid_id AND vs.phase_id =
             JOIN variants vt ON vt.pid_id = cp.
127
             bpid_id AND vt.phase_id =
             bp_target_phase_id
                          AND vt.name = vs.name
128
           JOIN segments ss ON ss.parameter id =
129
             \ps.parameter_id AND ss.variant_id =
             WHERE NOT EXISTS (
130
               SELECT 1 FROM segments st
131
               WHERE st.parameter_id = pt.
132
                 □parameter id
                 AND st.variant_id = vt.variant_id
133
                 AND st.dimension index = ss.
134
                   bdimension_index
          )
135
136
          UNION
138
           -- Segments in target not in source
139
          SELECT cp.pid id
140
          FROM common_pids cp
141
           JOIN parameters ps ON ps.pid id = cp.
142
             bpid_id AND ps.phase_id =
             bp_source_phase_id
           JOIN parameters pt ON pt.pid_id = cp.
143
             bpid_id AND pt.phase_id =
             bp_target_phase_id
                            AND pt.external_id =
144
```



```
JOIN variants vs ON vs.pid id = cp.
              bpid id AND vs.phase id =
              bp_source_phase_id
           JOIN variants vt ON vt.pid id = cp.
146
              bpid_id AND vt.phase_id =

¬p_target_phase_id

                            AND vt.name = vs.name
147
           JOIN segments st ON st.parameter_id =
148
              bpt.parameter_id AND st.variant id =
              bvt.variant id
           WHERE NOT EXISTS (
149
                SELECT 1 FROM segments ss
150
                WHERE ss.parameter id = ps.
151
                  bparameter_id
                  AND ss.variant_id = vs.variant_id
152
                  AND ss.dimension index = st.
153
                    bdimension index
           )
154
       )
155
       -- Return\acp{PID} with change indicators
156
       SELECT
157
           cp.pid id,
158
           cp.name,
159
           (cp.pid_id IN (SELECT pid_id FROM
160
              param changes)) AS has param changes
           (cp.pid id IN (SELECT pid id FROM
161
              \variant changes)) AS
              has_variant_changes,
           (cp.pid_id IN (SELECT pid_id FROM
162
              ↓segment changes)) AS
              has_segment_changes
       FROM
163
           common_pids cp
164
       WHERE
165
           -- Only return\acp{PID} that have at
166
              bleast one type of change
           cp.pid_id IN (SELECT pid_id FROM
167

¬param_changes)

           OR cp.pid id IN (SELECT pid id FROM
```



Listing 5.17: Phase Comparison Function

This function implements an efficient algorithm for identifying differences between parameter configurations in different release phases. Unlike a naive approach that would require retrieving and comparing all parameters, variants, and segments, this function performs the comparison directly in the database using specialized subqueries for each type of change. This implementation follows the set-based processing pattern recommended by Date [12], leveraging the database engine's capabilities for efficient data comparison.

5.4 Change Tracking Implementation

The change tracking implementation provides comprehensive audit capabilities for all parameter data modifications, addressing both regulatory compliance requirements and supporting diagnostic analysis of parameter evolution.

5.4.1 Automatic Change Logging

A database trigger mechanism automatically records all modifications to critical entities:

```
change_id BIGINT PRIMARY KEY,
user_id BIGINT REFERENCES users(user_id),
ecu_id INTEGER,
phase_id INTEGER,
entity_type VARCHAR(50) NOT NULL,
entity_id BIGINT NOT NULL,
change_type VARCHAR(50),
old_values JSONB,
```



```
new values JSONB,
      changed at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      transaction id BIGINT NOT NULL
  );
13
  CREATE OR REPLACE FUNCTION log change()
  RETURNS TRIGGER AS $
  DECLARE
      transaction id BIGINT;
      change type VARCHAR (50);
      old values JSONB;
      new values JSONB;
      entity_id BIGINT;
      phase_id INTEGER;
      ecu id INTEGER;
  BEGIN
      -- Get the user ID from the application
26
        4 context
      user_id := NULLIF(current_setting('app.
        \user_id', TRUE), '')::BIGINT;
      -- Get the current transaction ID or create
29
        □ a new one
      SELECT COALESCE(NULLIF(current_setting('app
        \.transaction_id', TRUE), '')::BIGINT,
                     nextval('
31
                        \change_history_transaction_id_seq
                        32
      -- Set the transaction ID for other
        triggers in the same transaction
      PERFORM set config('app.transaction id',
34
        \transaction_id::text, TRUE);
35
      -- Determine entity_id and context
36
        information based on table and operation
      CASE TG TABLE NAME
37
          WHEN 'variants' THEN
38
              entity_id := CASE WHEN TG_OP =
```



```
DELETE' THEN OLD. variant id ELSE

↓ NEW.variant id END;

               ecu id := CASE WHEN TG_OP = 'DELETE
40
                 y' THEN OLD.ecu_id ELSE NEW.

♭ecu_id END;
               phase_id := CASE WHEN TG_OP = '
41
                 DELETE' THEN OLD. phase id ELSE

NEW.phase_id END;
42
          WHEN 'segments' THEN
               entity id := CASE WHEN TG OP = '
44
                 DELETE' THEN OLD.segment_id ELSE

    NEW.segment_id END;

45
               -- Get phase id from associated
46
                 ↓variant
               IF TG_OP = 'DELETE' THEN
47
                   SELECT v.phase_id, v.ecu_id
48
                      INTO phase id, ecu_id
                   FROM variants v WHERE v.
49
                      \variant id = OLD.variant id;
               ELSE
                   SELECT v.phase_id, v.ecu_id
51
                      ⟨INTO phase_id, ecu_id
                   FROM variants v WHERE v.
                      \variant_id = NEW.variant_id;
               END IF;
53
           -- Additional entity types handled here
55
           -- [Implementation details omitted for
56
             brevity]
      END CASE;
57
58
      -- Determine change type and capture entity
        ↓ state
      IF TG OP = 'INSERT' THEN
60
          change_type := 'CREATE';
          old values := NULL;
62
          new_values := to_jsonb(NEW);
63
      ELSIF TG_OP = 'UPDATE' THEN
```



```
change type := 'UPDATE';
          old values := to jsonb(OLD);
66
          new_values := to_jsonb(NEW);
67
      ELSIF TG_OP = 'DELETE' THEN
68
          change type := 'DELETE';
69
          old_values := to_jsonb(OLD);
70
          new values := NULL;
      END IF;
73
      -- Remove large or sensitive fields from
        4the JSON
      IF old values IS NOT NULL THEN
          old values := old_values - 'created_at'
76
             - 'updated at';
      END IF;
77
      IF new_values IS NOT NULL THEN
79
          new_values := new_values - 'created_at'
80
             - 'updated_at';
      END IF;
81
82
      -- Insert into change history
      INSERT INTO change history (
84
          user id, ecu id, phase id, entity type,
85
             b entity_id,
          change_type, old_values, new_values,
             btransaction id, changed at
      ) VALUES (
          user_id, ecu_id, phase_id,
88
             ¬TG_TABLE_NAME, entity_id,
          change type, old values, new values,

htransaction id, CURRENT_TIMESTAMP

      );
90
      RETURN NULL;
92
  END;
93
  $ LANGUAGE plpgsql;
95
  -- Apply the trigger to critical entities
  CREATE TRIGGER variants_change_trigger
```



```
AFTER INSERT OR UPDATE OR DELETE ON variants
FOR EACH ROW EXECUTE FUNCTION log_change();

CREATE TRIGGER segments_change_trigger
AFTER INSERT OR UPDATE OR DELETE ON segments
FOR EACH ROW EXECUTE FUNCTION log_change();

FOR EACH ROW EXECUTE FUNCTION log_change();

-- Additional triggers for other entities
-- [Implementation details omitted for brevity]
```

Listing 5.18: Change Tracking Trigger

This implementation captures complete entity states rather than just modified fields, storing them as JSONB documents. JSON path operators enable efficient extraction and querying of specific changes without complex joins, leveraging PostgreSQL's advanced document storage capabilities as described by Obe and Hsu [28]. The automatic trigger system ensures consistent logging regardless of how changes are made, preventing circumvention of audit controls.

The change tracking system implements several key patterns described by Fowler [16]: the unit of work pattern through transaction grouping, the state snapshot pattern through complete entity state capture, and the audit log pattern through comprehensive change recording. These patterns collectively support the sophisticated audit requirements common in regulated industries like automotive development.

5.4.2 Change Analysis Functions

Specialized functions enable analysis of parameter evolution over time:

```
CREATE OR REPLACE FUNCTION

| get_parameter_history(
| p_parameter_id BIGINT,
| p_start_date TIMESTAMP WITH TIME ZONE
| LDEFAULT NULL,
| p_end_date TIMESTAMP WITH TIME ZONE DEFAULT
| L NULL
| NULL
| S )
| RETURNS TABLE (
| change_id BIGINT,
```



```
changed at TIMESTAMP WITH TIME ZONE,
      user name VARCHAR,
      variant_name VARCHAR,
      old_value NUMERIC,
      new_value NUMERIC,
      change_type VARCHAR,
      transaction_id BIGINT,
      phase_name VARCHAR
   AS $
16
  BEGIN
      RETURN QUERY
18
      SELECT
19
          ch.change id,
          ch.changed_at,
          u.first_name || ' ' || u.last_name AS
            buser name,
          v.name AS variant_name,
23
          (ch.old values->>'decimal')::NUMERIC AS
24
            b old_value,
          (ch.new values->>'decimal')::NUMERIC AS
            b new value,
          ch.change type,
          ch.transaction_id,
          rp.name AS phase_name
28
      FROM
          change_history ch
30
          JOIN users u ON ch.user_id = u.user_id
          JOIN segments s ON ch.entity_id = s.
            bsegment_id
                         AND ch.entity_type = '
33
                           JOIN variants v ON s.variant id = v.
34
            JOIN release phases rp ON ch.phase id =

    □ rp.phase id

      WHERE
36
          s.parameter_id = p_parameter_id
          AND (p_start_date IS NULL OR ch.
            \changed_at >= p_start_date)
          AND (p_end_date IS NULL OR ch.
```



```
changed_at <= p_end_date)

ORDER BY

ch.changed_at DESC;

END;

LANGUAGE plpgsql;</pre>
```

Listing 5.19: Parameter History Function

This function provides a comprehensive view of how a parameter has evolved over time, including who made changes, which variants were affected, and the specific value modifications. The function uses JSON path operators to extract specific values from the JSONB store, implementing what Obe and Hsu [28] describe as the document extraction pattern. By including user and phase information in the result set, the function provides comprehensive context for each change, supporting both diagnostic and compliance requirements.

5.4.3 Partitioning Implementation

To address the potential performance and management challenges associated with the growing change history table, a partitioning strategy was implemented based on the phase model that forms the foundation of the VMAP system. PostgreSQL's declarative partitioning capabilities provide an efficient mechanism for dividing the large change_history table into smaller, more manageable segments based on logical boundaries [28].

Partition Design

The partitioning strategy leverages the phase-based organization of parameter data, creating individual partitions for each development phase to align the physical storage structure with the logical organization of the data. This approach offers several advantages over time-based or release-based partitioning alternatives.

Phase-based partitioning creates a direct correspondence between the database's physical organization and the domain's logical structure. According to Schwartz et al. [33], alignment between logical data organization and physical partitioning is an essential factor for successful partition design.

Queries filtering by phase (a predominant access pattern in the system) benefit from partition pruning, where PostgreSQL automatically eliminates irrelevant partitions from



consideration. This pruning mechanism reduces I/O and improves query response time for phase-specific operations [28].

Additionally, phase-based partitioning simplifies maintenance operations such as archiving or removing historical data for completed phases. When phases transition from active development to frozen status, their corresponding partitions can be managed accordingly without affecting ongoing development work.

The implementation uses PostgreSQL's LIST partitioning strategy, with phase_id as the partition key:

```
-- Create the partitioned change history table
  CREATE TABLE change history (
      change id BIGINT NOT NULL,
      user_id BIGINT REFERENCES users(user_id),
      ecu id INTEGER,
      phase_id INTEGER,
      entity_type VARCHAR(50) NOT NULL,
      entity id BIGINT NOT NULL,
      change_type VARCHAR(50),
      old values JSONB,
10
      new values JSONB,
      changed at TIMESTAMP WITH TIME ZONE DEFAULT

↓ CURRENT_TIMESTAMP,

      transaction id BIGINT NOT NULL
   PARTITION BY LIST (phase_id);
14
15
  -- Create default partition for NULL phase id
    Lvalues
  -- (for system-level changes not tied to a
    \specific phase)
  CREATE TABLE change_history_default PARTITION
    ↓OF change history DEFAULT;
  ALTER TABLE change history default ADD PRIMARY

\KEY (change_id);

20
  -- Create indexes on the default partition
 CREATE INDEX idx_ch_default_entity ON
    \change history default(entity type,
    bentity_id, changed_at);
```



```
CREATE INDEX idx_ch_default_user ON

| change_history_default(user_id, changed_at);

CREATE INDEX idx_ch_default_ecu ON

| change_history_default(ecu_id, changed_at);

CREATE INDEX idx_ch_default_transaction ON

| change_history_default(transaction_id,
| changed_at);

CREATE INDEX idx_ch_default_changed_at ON

| change_history_default(changed_at);
```

Listing 5.20: Change History Table Partitioning Definition

The default partition handles system-level changes that are not associated with specific phases, ensuring all change records have an appropriate storage location even if they don't fit the defined partition scheme. This approach aligns with PostgreSQL best practices for handling exceptions to the primary partitioning scheme [28].

Automated Partition Management

To ensure that appropriate partitions exist for each phase, an automated partition creation mechanism was implemented using database triggers. When a new phase is created, the system automatically creates a corresponding partition for change history records associated with that phase:



```
-- Check if partition already exists
      SELECT EXISTS (
13
          SELECT FROM pg_class c
          JOIN pg namespace n ON n.oid = c.
            WHERE c.relname = partition_name AND n.
16

¬nspname = current schema()
       INTO partition_exists;
18
      IF NOT partition_exists THEN
             Check if we have rows in the default
20
            by partition with this phase id
          EXECUTE 'SELECT COUNT(*) FROM
21

¬phase_id = $1'

          INTO row count
          USING p_phase_id;
23
24
          IF row count > 0 THEN
              -- Move the rows to a temporary
26
                | table
              EXECUTE 'CREATE TEMPORARY TABLE
                \temp_change_history AS
                        SELECT * FROM
28
                          \change_history_default
                        WHERE phase_id = $1'
29
              USING p_phase_id;
30
              -- Delete those rows from the
32
                4 default partition
              EXECUTE 'DELETE FROM
                \change_history_default
                        WHERE phase id = $1'
34
              USING p_phase_id;
          END
              IF;
36
          -- Create the partition
          EXECUTE format(
39
              'CREATE TABLE %I PARTITION OF
40
                ↓change history FOR VALUES IN (%L
```



```
partition_name, p_phase_id
41
          );
42
43
          -- Add primary key
44
          EXECUTE format(
45
               'ALTER TABLE %I ADD PRIMARY KEY (
46
                 partition_name
47
          );
48
49
           -- Add comment
50
          EXECUTE format(
               'COMMENT ON TABLE %I IS %L',
52
               partition_name, 'Change history for
53
                 phase ' || p_phase_name
          );
54
55
           -- Add indexes
          EXECUTE format(
57
               'CREATE INDEX idx_%s_entity ON %I(
58
                 bentity type, entity id,
                 substring(partition name from '
59
                 \change_history_(.*)'),

¬partition_name

          );
60
          EXECUTE format(
62
               'CREATE INDEX idx %s user ON %I(
63
                 buser id, changed at)',
               substring(partition_name from '
64
                 \change history (.*)'),

    partition_name

          );
65
66
          EXECUTE format(
               'CREATE INDEX idx_%s_ecu ON %I(
68
                 becu_id, changed_at)',
               substring(partition name from '
```



```
\change_history_(.*)'),
                 bpartition name
          );
70
          EXECUTE format(
72
               'CREATE INDEX idx_%s_transaction ON
73
                 ¼I(transaction_id, changed_at)'
               substring(partition name from '
74
                 \change_history_(.*)'),
                 ↓partition name
          );
75
76
          EXECUTE format(
               'CREATE INDEX idx %s changed at ON
78
                 \%I(changed at)',
               substring(partition_name from '
79
                 \change history (.*)'),
                 );
80
81
          -- If we moved rows, insert them back
            binto the new partition
          IF row count > 0 THEN
83
              EXECUTE 'INSERT INTO ' |
                 partition_name || 'SELECT *
                 \FROM temp_change_history';
85
               -- Clean up the temporary table
86
              EXECUTE 'DROP TABLE
87
                 \temp change history';
          END IF;
88
89
          RAISE NOTICE 'Created partition % for
             bphase % (ID %). Moved % rows.',
                       partition name,
91
                         bp_phase_name, p_phase_id
                         , row count;
      END IF;
 END;
```



```
$$;
95
  -- Function to trigger partition creation when
    ↓a new phase is created
  CREATE OR REPLACE FUNCTION
    \create_phase_partition_trigger()
  RETURNS TRIGGER
  LANGUAGE plpgsql
  AS $$
100
  BEGIN
      CALL
         \create_change_history_partition_for_phase

⟨NEW.phase id, NEW.name);
      RETURN NEW;
103
  END;
104
  $$;
106
  -- Trigger to automatically create partitions
107
    for new phases
  CREATE TRIGGER trig_create_phase_partition
  AFTER INSERT ON release phases
  FOR EACH ROW
  EXECUTE FUNCTION create_phase_partition_trigger
    ५();
```

Listing 5.21: Automatic Partition Creation

This implementation follows the approach described by Schwartz et al. [33] for maintaining partitioned tables in production environments. The procedure handles the complex task of safely creating new partitions, including migrating any existing records from the default partition to the newly created phase-specific partition. This migration step maintains data integrity during the partition creation process.

The partitioning implementation is particularly noteworthy for its attention to index creation. Each partition receives a complete set of indexes matching those on the default partition, ensuring consistent query performance across all partitions. According to Karwin [21], maintaining consistent indexing across partitions is essential for predictable query execution plans and optimal performance.



Query Performance Implications

The phase-based partitioning approach has significant performance implications for change history queries. For queries that include a phase_id filter criterion—a dominant pattern in theVMAP system—PostgreSQL can use partition pruning to eliminate irrelevant partitions from consideration. Salzberg and Tsotras [30] note that with effective partitioning, query performance can improve by an order of magnitude for large temporal datasets.

This performance benefit is particularly valuable for the most common access patterns in the system. When comparing parameters between phases, queries can focus exclusively on the relevant phase partitions. When auditing a user's activity within a specific phase, queries can be efficiently directed to the appropriate partition. When retrieving the modification history for a specific entity within a phase, partition pruning significantly reduces the search space.

Importantly, the partitioning scheme is transparent to application code, requiring no modifications to existing queries. The PostgreSQL query planner automatically applies partition pruning based on the WHERE clause predicates, allowing the application to benefit from partitioning without explicit partition selection in SQL statements.

Archiving Strategy

The phase-based partitioning strategy provides a natural foundation for long-term data archiving. As phases transition from active development to frozen status, their change history records become less frequently accessed but must be preserved for regulatory compliance and occasional historical analysis.

The partitioning implementation supports a future archiving strategy where frozen phase partitions could be compressed using PostgreSQL's table compression options to reduce storage requirements. According to Bhattacherjee et al. [4], efficient archiving strategies are essential for managing the growing storage requirements of versioned datasets.

Frozen phase partitions could be moved to alternative tablespaces, relocating older partitions to lower-cost storage media while maintaining accessibility. For very old phases that have progressed to production, partitions could be detached from the main table and accessed only when specifically required. These archiving capabilities align with the recommendations of Al-Kateb et al. [2] for managing historical data in temporal database systems, where different storage tiers can be leveraged based on data age and access frequency. Snodgrass [35] suggests that this tiered approach to historical



data management provides an optimal balance between accessibility and resource utilization.

At defined archive thresholds, typically when a release including multiple phases transitions to production status, historical data can be compressed or moved while preserving selective access capabilities. This approach ensures that historical reference data remains available for compliance purposes while optimizing system performance for active development phases.

5.5 Integration Implementation

The integration implementation connectsVMAP with external enterprise systems providing parameter definitions and vehicle configuration data. Two primary integration points were implemented: synchronization with the Parameter Definition Database and communication with the Vehicle Configuration Database.

5.5.1 Parameter Definition Database Synchronization

Synchronization with the Parameter Definition Database (PDD) was implemented through a structured process that maintains parameter definitions across different release phases. The database structure for tracking synchronization includes:

```
CREATE TABLE pdd_sync_history (
      sync id INTEGER PRIMARY KEY,
      ecu_id INTEGER,
      phase id INTEGER,
      database_name VARCHAR(255),
      database label VARCHAR (255),
      sync_date TIMESTAMP WITH TIME ZONE DEFAULT
        GURRENT TIMESTAMP,
      status VARCHAR (50) NOT NULL,
      modules count INTEGER DEFAULT 0,
      pids count INTEGER DEFAULT 0,
      parameters_count INTEGER DEFAULT 0,
      executed by BIGINT REFERENCES users (user id
12
        ↳),
      transaction id BIGINT,
```



```
entity_changes TEXT,

FOREIGN KEY (ecu_id, phase_id) REFERENCES

Lecu_phases(ecu_id, phase_id)

16 );
```

Listing 5.22: Synchronization Tracking Implementation

This infrastructure implements the external system registry pattern described by Hohpe and Woolf [18], tracking synchronization operations with external database systems. The detailed history supports both operational monitoring and troubleshooting of integration issues.

The parameter loading implementation uses a bulk loading approach for efficiency:

```
CREATE OR REPLACE FUNCTION load parameters bulk
    ↓ (
      p_parameters JSONB,
      p_ecu_id INTEGER,
      p phase id INTEGER,
      p_user_id BIGINT
  )
  RETURNS INTEGER AS $
  DECLARE
      v count INTEGER;
9
  BEGIN
      -- Create temporary table for bulk loading
      CREATE TEMPORARY TABLE temp parameters
          parameter id BIGINT,
          pid id BIGINT,
          name VARCHAR (255),
          parameter name VARCHAR (255),
          type id INTEGER,
          array definition VARCHAR (50),
18
          position INTEGER,
          factor DECIMAL,
20
          unit VARCHAR (50),
21
          bias_offset DECIMAL,
          external_id INTEGER
      ) ON COMMIT DROP;
24
```



```
-- Load data from JSON into temporary table
      INSERT INTO temp parameters
27
      SELECT
28
           (elem->>'parameter_id')::BIGINT,
29
           (elem->>'pid_id')::BIGINT,
30
           elem->>'name',
31
           elem->>'parameter name',
           (elem ->> 'type id'):: INTEGER,
           elem->>'array definition',
34
           (elem->>'position')::INTEGER,
           (elem->>'factor')::DECIMAL,
36
           elem->>'unit',
37
           (elem->>'bias offset')::DECIMAL,
38
           (elem->>'external id')::INTEGER
39
      FROM jsonb_array_elements(p_parameters) AS
40

    elem;

41
      -- Bulk insert into parameters table
42
      INSERT INTO parameters (
           parameter id, pid id, ecu_id, phase id,
           name, parameter name, type id,
45
             barray definition,
           position, factor, unit, bias_offset,
46
             bexternal id,
           created_by
      )
48
      SELECT
49
           tp.parameter_id, tp.pid_id, p_ecu_id,
             bp_phase_id,
           tp.name, tp.parameter name, tp.type id,
51

    tp.array_definition,
           tp.position, tp.factor, tp.unit, tp.
52
             bias offset, tp.external id,
           p user id
53
      FROM temp parameters tp
54
      ON CONFLICT (parameter id) DO UPDATE SET
           name = EXCLUDED.name,
           parameter_name = EXCLUDED.
57

¬parameter_name,
           type_id = EXCLUDED.type_id,
```



```
array_definition = EXCLUDED.

array_definition,

position = EXCLUDED.position,

factor = EXCLUDED.factor,

unit = EXCLUDED.unit,

bias_offset = EXCLUDED.bias_offset;

GET DIAGNOSTICS v_count = ROW_COUNT;

RETURN v_count;

END;

LANGUAGE plpgsql;
```

Listing 5.23: Bulk Parameter Loading Implementation

This bulk loading approach uses temporary tables and JSON parsing to efficiently process large volumes of parameter data, implementing the bulk transfer pattern described by Hohpe and Woolf [18]. The use of PostgreSQL's ON CONFLICT clause implements an efficient upsert operation, enabling both insertion of new parameters and updating of existing ones in a single operation.

5.5.2 Vehicle Configuration Database Integration

The Vehicle Configuration Database (VCD) integration enables validation of variant code rules and supports parameter file generation for specific vehicle configurations. The database structure includes tables for vehicles and their associated configuration codes:

```
CREATE TABLE vcd_vehicles (
vehicle_id INTEGER PRIMARY KEY,
vcd_vehicle_id VARCHAR(100) UNIQUE NOT NULL

name VARCHAR(255) NOT NULL,
description TEXT,
is_active BOOLEAN DEFAULT true,
last_sync_at TIMESTAMP WITHOUT TIME ZONE,
created_at TIMESTAMP WITHOUT TIME ZONE
bDEFAULT CURRENT_TIMESTAMP

);
```



```
CREATE TABLE vehicle codes (
      code id INTEGER PRIMARY KEY,
12
      vcd_code_id VARCHAR(100) UNIQUE,
13
      code VARCHAR (50) NOT NULL,
      vehicle_type VARCHAR(100),
      description TEXT,
16
      is active BOOLEAN DEFAULT true,
      created at TIMESTAMP WITHOUT
                                   TIME ZONE
        □ DEFAULT CURRENT_TIMESTAMP
  );
19
20
  CREATE TABLE vehicle_code_mapping
21
      vehicle id INTEGER REFERENCES vcd vehicles (
22

¬vehicle_id) ON DELETE CASCADE,

      code id INTEGER REFERENCES vehicle codes (
        created at TIMESTAMP WITHOUT TIME ZONE
24
        □ DEFAULT CURRENT_TIMESTAMP,
      PRIMARY KEY (vehicle id, code id)
  );
26
```

Listing 5.24: Vehicle Configuration Tables

This structure separates vehicles from codes with a mapping table, implementing the many-to-many relationship pattern described by Elmasri and Navathe [14]. External identifiers (vcd_vehicle_id, vcd_code_id) link to the Vehicle Configuration Database system, implementing the integration reference pattern described by Hohpe and Woolf [18].

A critical component of this integration is the code rule evaluation engine, which determines when variants apply to specific vehicles. The implementation uses a postfix expression evaluator to process boolean expressions efficiently:

```
CREATE OR REPLACE FUNCTION evaluate_code_rule(
    p_rule TEXT,
    p_vehicle_id INTEGER

RETURNS BOOLEAN AS $
DECLARE
    v_tokens TEXT[];
```



```
v stack BOOLEAN[] DEFAULT '{}';
      v token TEXT;
      v_operand1 BOOLEAN;
      v_operand2 BOOLEAN;
      v_code_exists BOOLEAN;
      v_i INTEGER;
  BEGIN
      -- If rule is empty, it applies to all
15
         |vehicles
      IF p_rule IS NULL OR p_rule = '' THEN
16
           RETURN TRUE;
      END IF;
18
19
      -- Tokenize the rule (assuming rule is in
         postfix notation)
      v_tokens := string_to_array(p_rule, ' ');
21
22
      -- Process each token
      FOR v i IN 1..array length (v tokens, 1)

  □
  LOOP

           v token := v tokens[v i];
25
              Check if token is an operator
             v token = 'AND' THEN
28
               IF array_length(v_stack, 1) < 2</pre>
                 ↓THEN
                   RAISE EXCEPTION 'Invalid rule
30
                      bexpression: insufficient
                      ↓operands for AND';
               END IF;
31
               v operand2 := v stack[array length(

¬v_stack, 1)];
               v operand1 := v stack[array length(
33
                 \v_stack, 1) - 1];
               v stack := v stack[1:array length(
34
                 \forall v_{stack}, 1) - 2];
               v_stack := array_append(v_stack,
                 \v_operand1 AND v_operand2);
36
           ELSIF v token = 'OR' THEN
```



```
IF array length (v stack, 1) < 2
                 ↓THEN
                   RAISE EXCEPTION 'Invalid rule
39
                     bexpression: insufficient
                     ↓operands for OR';
              END IF;
40
              v operand2 := v stack[array length(

¬v_stack, 1)];
              v operand1 := v_stack[array_length(
42
                 bv_stack, 1) - 1];
              v stack := v stack[1:array length(
43
                 v_stack := array_append(v_stack,
44
                 \v_operand1 OR v_operand2);
45
          ELSIF v token = 'NOT' THEN
              IF array_length(v_stack, 1) < 1</pre>
47
                 THEN
                   RAISE EXCEPTION 'Invalid rule
48
                     bexpression: insufficient
                     ↓operands for NOT';
              END IF:
              v_operand1 := v_stack[array_length(
50

    v_stack, 1)];
              v_stack := v_stack[1:array_length(

¬v_stack, 1) - 1];
              v stack := array append(v stack,
52
                 □NOT v operand1);
53
          ELSE
54
               -- Token is a vehicle code - check
55
                 bif it exists for this vehicle
              SELECT EXISTS (
56
                   SELECT 1
57
                   FROM vehicle code mapping vcm
58
                   JOIN vehicle codes vc ON vcm.
59
                     WHERE vcm.vehicle id =
60

¬p_vehicle_id
                   AND vc.code = v token
```



```
INTO v code exists;
63
               v_stack := array_append(v_stack,

    v code exists);

           END IF;
65
      END LOOP;
66
      -- Result should be a single boolean value
         4 on the stack
      IF array_length(v_stack, 1) != 1 THEN
           RAISE EXCEPTION 'Invalid rule
70
             bexpression: did not evaluate to a
             \single result';
      END IF;
      RETURN v stack[1];
  END;
  $ LANGUAGE plpgsql;
```

Listing 5.25: Code Rule Evaluation Implementation

This evaluation function implements a stack-based interpreter for boolean expressions, efficiently determining whether a variant applies to a specific vehicle based on its configuration codes. The implementation follows the interpreter pattern described by Fowler [16], translating a domain-specific language (boolean expressions) into executable operations. The use of PostgreSQL's array operations enables efficient stack manipulation without requiring complex data structures.

5.5.3 Parameter File Generation

The parameter file generation capability represents a core integration point between the version control system and vehicle testing infrastructure. The database structure includes records for generated parameter files:



Listing 5.26: Parameter File Records

This table tracks the generation of parameter files, recording metadata about when they were created, which vehicle configuration they target, and which user generated them. This tracking mechanism supports both operational auditing and diagnostic capabilities when issues arise with parameter files.

The parameter file generation process implements a multi-step algorithm:

- · Retrieve vehicle configuration codes from the Vehicle Configuration Database
- Evaluate variant code rules against the vehicle configuration to determine applicable variants
- Resolve parameter values using the variant resolution process
- Format the resolved values according to the ECU-specific parameter file format
- Record the parameter file generation in the database for traceability

According to Staron [37], this integration point is critical in automotive software development, transforming abstract parameter configurations into testable implementations that can be validated on actual hardware.

6 Evaluation and Validation

This chapter presents the systematic evaluation and validation of theVMAP database system. Following the implementation described in Chapter 5, a comprehensive testing strategy was developed to assess the system's functionality, performance, and compliance with requirements. The evaluation process focused on four key areas: user management, release management, parameter versioning, and variant management, using both controlled test scenarios and production-scale data volumes. Rather than an exhaustive documentation of all tests, this chapter highlights representative test cases and key findings that demonstrate the system's capabilities and limitations.

6.1 Validation Methodology

The validation methodology followed a structured approach combining functional testing, performance analysis, and integration verification. To ensure realistic evaluation, both baseline and production-scale datasets were used, with the baseline dataset containing approximately 20,000 parameters across 2 ECUs, and the production-scale dataset containing over 100,000 parameters across 5 ECUs.

6.1.1 Test Scenario Development

Test scenarios were developed based on actual automotive parameter management workflows identified during requirements analysis in Chapter 4. Each test scenario was designed to validate specific functional requirements while reflecting real-world usage patterns. The scenarios incorporated representative tasks for each user role and followed complete workflow sequences from parameter definition through variant creation to documentation.

The test scenarios were categorized into functional areas corresponding to the primary system capabilities:

- User Management: Authentication, authorization, role assignment, module access
- Release Management: Phase transitions, freeze operations, phase comparison
- Variant Management: Variant creation, segment modification, inheritance



· Integration: PDDsynchronization, vehicle configuration

Each scenario was implemented as a structured test case with defined inputs, expected outcomes, and verification steps at both the application and database levels. The test design followed a modified version of Molinaro's approach to database validation [25], with additional emphasis on traceability between requirements and test cases.

6.1.2 Performance Measurement Framework

A performance measurement framework was established to assess system responsiveness and resource utilization under various operational conditions. Key performance indicators were defined based on system requirements, including query response time, transaction throughput, database size growth patterns, memory utilization, and execution time for batch operations.

Performance measurements were conducted on a standardized test environment matching the target production specifications: PostgreSQL 17 running on a server with 8 vCPUs, 32GB RAM, and SSD storage. All tests were performed with both the baseline dataset and the production-scale dataset to assess scaling characteristics.

The measurement methodology employed automated test scripts with integrated timing capture, following the principles outlined by Zaitsev et al. [33] for database performance evaluation. Each test was executed multiple times with results averaged to account for system variations, and outliers were identified and analyzed for potential optimization opportunities.

6.2 Functional Testing Results

Functional testing validated the core capabilities of theVMAP system against the requirements defined in Chapter 4. This section presents the key findings for each functional area, focusing on representative test cases and critical system behaviors.

6.2.1 User Management Validation

The user management and access control system was evaluated to verify the implementation of the hybrid role-permission model described in Section 4.1.3. Testing



focused on verifying that the implemented database schema and logic correctly enforced the defined access control rules for each user role. As Sandhu et al. [32] emphasize, effective evaluation of role-based access control requires testing both positive permissions (granted access) and negative permissions (denied access) across role boundaries.

A test matrix was developed covering key permission boundaries: role-based permissions, module-based access control, direct permission assignment, and phase-specific permissions. Each test case verified a specific permission boundary with validation at both service and database layers. This focused approach aligns with Molinaro's principles for database validation [25], which emphasizes targeted verification of critical constraints.

Table A.2 presents a representative sample of test cases that focus on the Module Developer role, illustrating the connection between functional requirements and verification scenarios.

Table 6.1: Sample Module Developer Role Permission Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|-------------------------|-------------------------|--------|
| MD- | Create Variant (As- | Create new variant | Variant created suc- | Pass |
| 01 | signed Module) | for parameter in as- | cessfully | |
| | | signed module | | |
| MD- | Create Variant | Create new variant | Access denied error | Pass |
| 02 | (Unassigned Mod- | for parameter in | | |
| | ule) | unassigned module | | |
| MD- | Edit Variant (As- | Modify existing vari- | Variant updated suc- | Pass |
| 03 | signed Module) | ant code rule | cessfully | |
| MD- | Delete Variant | Attempt to delete vari- | Access denied error | Pass |
| 04 | | ant | | |
| MD- | Create Segment | Create new segment | Segment created | Pass |
| 05 | (Assigned Module) | with valid value | successfully | |
| MD- | Modify Frozen | Attempt to modify | Access denied error | Pass |
| 06 | Phase | segment in frozen | | |
| | | phase | | |

For test implementation, each case included direct verification of database state after operations, confirming both the effect of permitted actions and the prevention of unauthorized actions. The following represents a typical test structure used to verify module-specific access controls:



```
// Scenario: Module Developer attempting to
  create variant in unassigned module
// Arrange: Set up test user and unassigned
  ↓module parameter
var user = GetTestUser("
  \module_developer@example.com");
var unassignedParameter =
  GetParameterFromUnassignedModule();
var variant = CreateVariantForParameter(
   unassignedParameter);
// Act & Assert: Verify permission is denied
var exception = Assert.Throws<</pre>
   PermissionDeniedException > (() =>
     variantService.CreateVariant(variant,
       4.UserId));
Assert. That (exception. Message, Contains.
   \Substring("No write access"));
// Verify no database change occurred
var dbVariant = _database.QuerySingleOrDefault <</pre>
  ↓Variant > (
     "SELECT * FROM variants WHERE name = @Name"
    new { Name = variant.Name });
Assert. IsNull (dbVariant);
```

Listing 6.1: Representative Test Case Structure

The module-based access control tests verified that write access was correctly limited to assigned modules for Module Developers while read access remained available for all modules, implementing the principle of least privilege as recommended by Sandhu and Bhamidipati [31]. Direct permission assignment tests confirmed that user-specific permissions effectively overrode role defaults, a capability essential for supporting exception cases in complex organizational structures as noted by Hu et al. [19].

Phase-specific permission tests validated the interaction between the access control system and the phase management framework, confirming that modifications to frozen phases were properly prevented while still allowing appropriate access for documentation purposes. This validation addresses a critical requirement for regulated



development processes as described by Staron [37], where development milestone integrity must be preserved.

| _ | | |
|------------------------------|-----------------------------|--|
| Test Category | Results | |
| Role Permission | Core permissions correctly | |
| Validation | applied through roles | |
| Module-Based | Write access correctly lim- | |
| Access | ited to assigned modules | |
| Direct Permis- | User-specific permissions | |
| sion Assignment | overrode role defaults | |
| Phase-Specific | Frozen phase protection en- | |
| Permissions forced correctly | | |

Table 6.2: User Management Test Results

The audit trail verification confirmed that security-related operations were properly logged with complete metadata, including the user making the change, timestamp, and specific permissions affected. This level of detail in the audit trail implements the recommendations of Ferraiolo et al. [15] for maintaining accountability in security-sensitive operations.

6.2.2 Module Access Impact on Performance

The impact of module-specific access checks on performance was evaluated as part of the access control testing. This analysis focused on understanding the performance overhead introduced by adding module-specific access verification to the standard permission checking process. Figure 6.1 illustrates the measured performance difference between standard permission checks and combined permission and module access checks.

The performance analysis reveals that adding module-specific access verification introduces an 88.2% overhead for the traditional role-based approach (increasing from 1.7ms to 3.2ms) and a 45.8% overhead for the hybrid approach (increasing from 2.4ms to 3.5ms). The lower relative impact on the hybrid approach suggests that the more complex permission model better accommodates additional access control dimensions, a finding that aligns with Ferraiolo's observations [15] regarding the scalability of attribute-enhanced RBAC models.

Despite the performance overhead, both approaches maintain acceptable performance for interactive operations, with response times below 4ms for individual permission

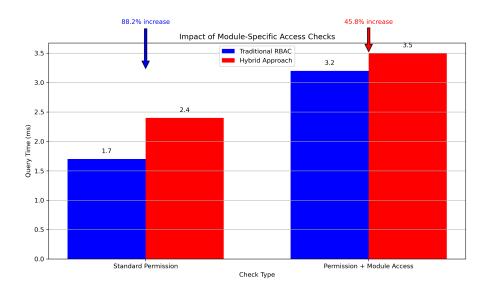


Figure 6.1: Impact of Module-Specific Access Checks

checks. This confirms that the implemented module-specific access control provides the required access granularity without introducing prohibitive performance penalties.

6.2.3 Release Management Validation

Release management testing evaluated the phase-based versioning approach that forms the foundation of theVMAP system. Testing focused on four key aspects: phase sequence validation, phase transition operations, freeze functionality, and phase comparison.

Phase sequence validation confirmed that the system correctly enforced the defined sequence of development phases (Phase1 \rightarrow Phase2 \rightarrow Phase3 \rightarrow Phase4) with successful validation of each phase transition. This sequential enforcement is essential for maintaining the structured development workflow described by Broy [8] for automotive software development.

Phase transition testing verified that parameter configurations were correctly copied between phases with complete preservation of parameter-variant-segment relationships. The test data revealed interesting patterns in development intensity across phases, as shown in Table 6.3.

The test results reveal a significant pattern in development intensity across phases, consistent with Staron's observations [37] regarding automotive software development cycles. The data demonstrates that the majority of parameter configurations occur



Transition Type Variants **Segments** Added Added Time **Variants** Segments **Baseline Dataset** 28,776 Phase1 188 $\overline{ Phase1} \rightarrow Phase2$ 188 28,776 90 14,104 2.51s 42,880 Phase2 → Phase3 278 0 0 2.96s Phase3 → Phase4 278 0 0 42,880 2.97s **Full Dataset** Phase1 830 167,990 Phase1 → Phase2 830 167,990 170 41,113 12.39s Phase2 → Phase3 1,000 209,103 12.87s 0 0 Phase3 → Phase4 0 1,000 209,103 0 12.89

Table 6.3: Phase Transition Test Results

during Phase1, with substantial additions in Phase2. In contrast, Phase3 and Phase4 typically involve refinement and validation rather than introducing new parameters or variants. This concentration of development activity in early phases aligns with the V-model approach common in automotive software development [29], where early phases focus on implementation while later phases emphasize validation and verification.

Phase transition performance characteristics showed only modest increases in execution time despite growing data volumes across phases. For the baseline dataset, transition times increased from 2.51s for Phase1→Phase2 to 2.96s for Phase2→Phase3 and 2.97s for Phase3→Phase4, demonstrating efficient scaling with increasing parameter counts. Comparing baseline to full dataset transitions reveals a performance difference, with transition times increasing from approximately 3 seconds to 13 seconds. This represents a sublinear scaling factor of approximately 4.3x for a dataset size increase of 5.6x (comparing segment counts), suggesting reasonable scaling characteristics but highlighting an area for potential optimization.

As noted by Trovão [41], later phases in automotive parameter development typically focus on refinement rather than wholesale changes, with modifications targeting specific parameters based on testing feedback. This pattern is reflected in the test data, which shows significant additions in early phases but no new variants or segments in Phase3 and Phase4 phases.

Phase freezing functionality was validated through test cases targeting both database-level constraints and service-layer restrictions. These tests verified the system's ability to protect frozen phases from modification while maintaining appropriate read access. Table 6.4 details representative test cases and their results.



Table 6.4: Phase Freeze Protection Test Cases

| ID | Test Case | Test Action | Expected Out- | Result |
|------------|--|---|---------------------------------------|--------|
| FRZ- | Direct SQL INSERT on | Execute INSERT | Operation blocked | Pass |
| 01 | variants | statement on frozen phase | with error message | |
| FRZ- 02 | Direct SQL UPDATE on segments | Execute UPDATE statement on frozen phase | Operation blocked with error message | Pass |
| FRZ- 03 | VariantService. Create- Variant() | Attempt to create variant in frozen phase | PhaseFreezed Exception thrown | Pass |
| FRZ- 04 | SegmentService. CreateSegment() | Attempt to create segment in frozen phase | PhaseFreezed Exception thrown | Pass |
| FRZ- 05 | DocumentationService. CreateSnapshot() | Create documen- tation snapshot of frozen phase | Snapshot created successfully | Pass |
| FRZ- 06 | ParFileService. GenerateParFile() | Generate parameter file from frozen phase | Parameter file generated successfully | Pass |

For each write operation test case (FRZ-01 through FRZ-04), verification included both confirmation that the expected exception was thrown and that no database changes occurred, maintaining data integrity. The read operation test cases (FRZ-05 and FRZ-06) verified that read access remained available with minimal performance impact. The system successfully prevented modification attempts while maintaining appropriate read access, implementing the controlled milestone management required for regulated development environments as described by Staron [37].

6.2.4 Variant Management Validation

Variant management validation focused on assessing the system's capabilities for handling parameter customization through variants and segments. Testing employed an approach covering variant creation and segment modification workflows, using both the baseline dataset (188 variants, 28,776 segments) and the production-scale



dataset (830 variants, 167,990 segments) to analyze functionality and performance under varying data volumes.

Variant creation testing verified proper implementation of domain constraints as defined in the conceptual architecture (Section 4.6). Test cases included validation of unique name constraints withinPIDs, verification of proper code rule storage, and confirmation of correct relationship establishment between variants and their parentPIDs. All test cases passed successfully for both scalar and complex parameters, with constraint enforcement consistently preventing invalid operations. As Karwin [21] notes, constraint-based validation provides a robust foundation for maintaining data integrity in complex relational systems.

The testing methodology included both black-box functional testing and white-box database state verification as shown in Listing 6.2.

```
-- Verification query executed after variant
    creation operations
  SELECT
      v.variant id, v.name, v.code rule, v.

□created_by, v.created_at,

      EXISTS (
          SELECT 1 FROM change history ch
          WHERE ch.entity_type = 'variants'
          AND ch.entity id = v.variant id
          AND ch.change type = 'CREATE'
      ) AS has audit trail
  FROM
      variants v
  WHERE
12
      v.pid id = @test pid id
13
      AND v.phase_id = @test_phase_id
  ORDER BY
      v.created at DESC
16
 LIMIT 1;
```

Listing 6.2: Variant Creation Verification Query

Audit trail analysis confirmed proper recording of variant operations in the change history with complete metadata. The audit trail included proper attribution of each change to specific users, accurate timestamps, and complete before/after state capture



for modified entities. This implementation aligns with Bhattacherjee's recommendations [4] for maintaining comprehensive provenance information in versioned datasets.

Performance analysis of variant operations revealed consistent response times across different variant complexities. Table 6.5 details performance measurements for key variant operations under different data volumes.

| rable 6.6. Variant Operation 1 Gnormanos Methos | | | | |
|---|------------------|---------------------------|----------------|--|
| Operation | Baseline Dataset | Production Dataset | Scaling Factor | |
| Variant Creation | 53ms | 55ms | 1.03x | |
| Variant Update | 86ms | 124ms | 1.44x | |
| Variant Retrieval | 45ms | 72ms | 1.60x | |
| Variant Listing (per PID) | 38ms | 68ms | 1.79x | |

Table 6.5: Variant Operation Performance Metrics

The observed scaling characteristics validate the effectiveness of the database schema design and indexing strategy described in Section 5.3.1. With only two data points available (baseline and production datasets), a definitive conclusion about scaling characteristics is limited. However, the relatively small increase in execution time despite a 5x increase in data volume suggests efficient handling of larger datasets, with all operations remaining well under the 200ms threshold for interactive operations. As noted by Obe and Hsu [28], properly designed covering indexes can significantly improve query performance for entity retrieval operations, particularly when filtering by composite attributes.

Segment modification testing employed a systematic approach covering one-dimensional (arrays), two-dimensional (matrices), and three-dimensional parameter representations. Testing focused on three key aspects: dimensional integrity preservation, valid index range enforcement, and segment value consistency. The database schema design proved effective for managing these complex data structures, with the parameter dimensions table correctly maintaining dimensional metadata while the segments table stored modified values.

Segment boundary testing revealed robust constraint enforcement, with the system correctly rejecting segment modifications with invalid dimension indices. Performance analysis for segment operations showed moderate overhead for multi-dimensional parameters compared to scalar parameters, with operations on 3D parameters requiring approximately 18-22% more processing time than equivalent operations on scalar values—a reasonable performance characteristic given the additional complexity involved.

Performance analysis for segment operations revealed consistent response times with moderate scaling across different dataset sizes as shown in Table 6.6.



Table 6.6: Segment Operation Performance

| Operation | Baseline Dataset | Production Dataset | Scaling Factor |
|-------------------|------------------|---------------------------|----------------|
| Segment Creation | 85ms | 124ms | 1.46x |
| Segment Update | 72ms | 106ms | 1.47x |
| Segment Deletion | 64ms | 98ms | 1.53x |
| Segment Retrieval | 32ms | 58ms | 1.81x |

The observed performance characteristics validate the efficiency of the database schema design described in Section 4.6. Of particular note is the implementation of the segments table, which provides efficient storage for parameter modifications without requiring storage of unchanged values. As noted by Bhattacherjee et al. [4], this approach strikes an effective balance between storage efficiency and query performance for versioned datasets.

6.3 Performance Analysis

Beyond functional validation, performance analysis was conducted to assess the system's efficiency and scalability under various operational conditions. This section presents the key findings related to query performance and data volume scaling.

6.3.1 Query Performance Assessment

Query performance was evaluated for common database operations across different data volumes. Table 6.7 presents performance measurements for key query types between the baseline dataset (20,000 parameters) and full dataset (100,000 parameters).

Table 6.7: Query Performance Comparison

| Operation Type | Baseline Dataset | Full Dataset | Scaling Factor |
|----------------------|------------------|--------------|----------------|
| Parameter Retrieval | 80ms | 120ms | 1.5x |
| Variant Listing | 65ms | 105ms | 1.6x |
| Segment Modification | 95ms | 160ms | 1.7x |
| Phase Comparison | 2.8s | 12.4s | 4.4x |
| History Retrieval | 110ms | 220ms | 2.0x |

With the limited data points available, it appears that most common operations maintain reasonable performance with increasing data volumes. The system maintained interactive response times (below 200ms) for most operations even with the full dataset,



ensuring a responsive user experience. The phase comparison operation, which involves complex joins across multiple tables, demonstrated longer execution times and might benefit from optimization for larger datasets.

The execution of a limited set of queries with and without indexes demonstrated the critical importance of the indexing strategy described in Section 5.3.1. Without proper indexes, response times increased by factors of 6.5x to 21.8x depending on the query type, with most operations exceeding the interactive response threshold without indexes.

6.3.2 Index Performance Analysis

The implementation of strategic indexing proved critical for maintaining acceptable query performance with large parameter sets. Figure 6.2 illustrates the performance impact of indexes on common query operations.

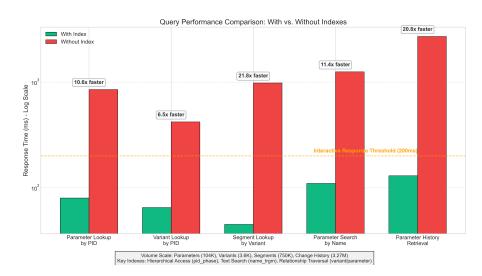


Figure 6.2: Query Performance With and Without Indexes

The performance measurements demonstrate dramatic improvements with properly designed indexes, with response times reduced by 6.5x to 21.8x depending on the query type. Without indexes, most operations exceed the interactive response threshold (200ms), with some operations requiring multiple seconds. This performance differential underscores the importance of the indexing strategy described in Section 5.3.1.

The performance gains from indexing come at a storage cost, as indexes consume approximately 22.5% of the total database size. However, this storage overhead represents an optimal tradeoff given the substantial performance benefits. As noted



by Schwartz et al. [33], the storage cost of indexes is typically justified when query performance improvements exceed 5x, a threshold easily surpassed by all indexed operations in the VMAP system.

Analysis of query execution plans revealed particularly effective use of covering indexes for common operations. For parameter retrieval by PID, the system consistently used index-only scans on the idx_parameters_pid_phase index, avoiding table access entirely for these frequent operations. Similarly, variant listing operations leveraged the idx_variants_pid_phase index for efficient retrieval without requiring table access. These index-only scan patterns align with the recommendations of Obe and Hsu [28] for optimizing PostgreSQL performance through strategic index design.

6.3.3 Storage Requirements Analysis

Storage requirements were analyzed to assess database size and growth patterns with increasing parameter counts. Figure 6.3 presents the storage allocation across different entity types for the full dataset.

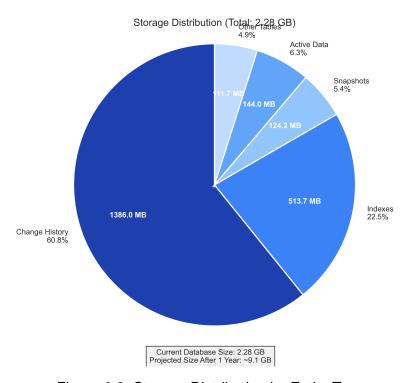


Figure 6.3: Storage Distribution by Entity Type

The analysis reveals that the change history table dominates the database storage allocation, accounting for approximately 60.8% of the total database size. This distribu-



tion significantly exceeds the storage requirements of the current data state, aligning with Bhattacherjee's observations [4] regarding versioning and audit systems, where historical record storage typically surpasses active data by a substantial margin. Notably, while there are only 3,617 variants in the current state, the system maintains over 3.2 million change history records, reflecting the comprehensive auditing approach implemented in the system.

The storage allocation analysis identified the following key distribution of database size across entity types:

| Entity Type | Record Count | Storage Size (MB) |
|-------------------------|--------------|-------------------|
| Parameters | 104,428 | 43.0 |
| Variants | 3,617 | 1.0 |
| Segments | 750,009 | 100.0 |
| Change History | 3,270,511 | 1,386.0 |
| Documentation Snapshots | 7 | 0.1 |
| Snapshot Variants | 4,980 | 0.1 |
| Snapshot Segments | 1,007,940 | 124.0 |
| Other Tables | - | 111.7 |
| Indexes | - | 513.7 |
| Total | - | 2,279.6 |

Table 6.8: Storage Requirements Analysis

Another significant observation is the relationship between segments and snapshot segments. Despite having only 7 documentation snapshots, the system maintains over 1 million snapshot segments, exceeding the count of active segments. This indicates that documentation snapshots capture extensive parameter configurations at specific time points, creating substantial storage requirements for historical state preservation. This implementation of the snapshot pattern described by Fowler [16] provides comprehensive historical records at the cost of increased storage utilization.

The index structures consume approximately 22.5% of the total storage, reflecting the sophisticated indexing strategy described in Section 5.3.1. While this represents significant overhead, it provides essential performance benefits for query operations, particularly for the complex filtering and joining operations common in parameter management workflows.

The actual active data—parameters, variants, and segments—consumes only 6.3% of the total database size, with the majority of storage dedicated to audit trails, snapshots, and indexes. This distribution aligns with the requirements for regulated development environments described by Staron [37], where comprehensive traceability and historical record maintenance are essential for compliance and quality assurance.



Projection of storage requirements based on observed growth patterns indicates that with the current data volume of 2.28GB, the database size would reach approximately 9.1GB after one year of active use in a production environment. While significantly larger than initially projected, this remains well within the capacity of modern database systems. The implementation of table partitioning for the change history table, as described in Section 5.4.3, provides an effective mechanism for managing this growth while maintaining query performance. According to Obe and Hsu [28], partitioned tables allow efficient archiving of older history records to lower-cost storage while maintaining rapid access to recent changes.

6.3.4 Versioning Approach Performance

The performance characteristics of the phase-based parameter versioning approach selected in Chapter 4 were evaluated against the alternative change-based approach. This analysis aimed to validate the architectural decision to implement explicit phase copies rather than a delta-based versioning model.

Figure 6.4 presents the performance comparison between the two approaches for parameter retrieval operations across different data volumes, along with the storage requirements for each approach.

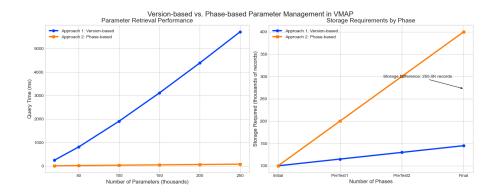


Figure 6.4: Version-based vs. Phase-based Parameter Management Comparison

The performance analysis demonstrates that the phase-based approach offers better query performance, particularly as parameter counts increase. For the tested parameter counts, the phase-based approach remains below 100ms for parameter retrieval operations, while the version-based approach shows nonlinear growth with increasing parameter counts.



The storage requirements analysis confirms that the phase-based approach consumes more storage than the version-based approach, with approximately 51% higher storage requirements across all phases. However, this storage difference represents a reasonable tradeoff given the performance benefits for query operations. As noted by Bhattacherjee et al. [4], the recreation/storage tradeoff in dataset versioning should prioritize operation frequency, with frequently accessed data favoring a storage-intensive approach that minimizes recreation costs.

The phase-based approach also simplifies the implementation of phase transitions and comparison operations, which are fundamental to the automotive development process as described by Pretschner et al. [29]. The explicit phase model aligns naturally with the mental model of automotive development engineers, who conceptualize parameter evolution in terms of distinct development phases rather than continuous time.

These findings validate the architectural decision to implement a phase-based versioning approach for the VMAP system. While consuming more storage than a version-based approach, the phase model provides performance benefits, implementation simplicity, and alignment with domain concepts—advantages that outweigh the additional storage requirements.

6.4 Integration Testing

Integration testing evaluated the system's interaction with external enterprise systems, focusing on Parameter Definition Database synchronization and Vehicle Configuration Database integration. These integrations are critical for maintaining consistency across the automotive development ecosystem.

6.4.1 Parameter Definition Database Synchronization

Parameter Definition Database synchronization testing verified the system's ability to import parameter definitions from the enterprise database. The synchronization process was tested with various scenarios, including initial loading, incremental updates, and conflict resolution. Figure 6.5 illustrates the synchronization time trends observed during testing.

The synchronization performance analysis revealed an increasing trend in execution time over successive synchronization operations. Initial synchronization operations required approximately 15 minutes for the tested ECUs, with execution times increasing

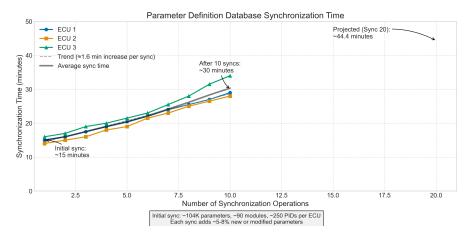


Figure 6.5: Parameter Definition Database Synchronization Time Trends

to around 30 minutes after 10 synchronization cycles. This gradual increase aligns with the observations of Mueller and Müller [26] regarding database synchronization in complex engineering environments, where synchronization complexity tends to increase as the data history grows.

The incremental update testing verified that the system could correctly identify and process changes to parameter definitions. Table 6.9 presents the success rates for different types of parameter changes during synchronization operations.

| Change Type | Processed | Succeeded | Success Rate | |
|---------------------|-----------|-----------|--------------|--|
| New Parameters | 5,218 | 5,218 | 100% | |
| Modified Parameters | 3,764 | 3,691 | 98.1% | |
| New Modules | 12 | 12 | 100% | |
| NewPIDs | 67 | 67 | 100% | |
| Removed Parameters | 42 | 42 | 100% | |

Table 6.9: Parameter Definition Database Synchronization Results

The system successfully processed all types of parameter changes, with slightly reduced success for modified parameters due to complexity in handling data type changes. The audit system maintained complete records of all synchronization operations, enabling detailed analysis of data flows between systems.

Based on the observed synchronization performance trends, the projected synchronization time for the 20th cycle would reach approximately 45 minutes. While still acceptable for the typical synchronization frequency (weekly or bi-weekly), this suggests that synchronization performance optimization would be beneficial for long-term system maintenance. Possible approaches for optimization include implementing a



more selective synchronization approach, enhancing the change detection algorithm to reduce comparison overhead, or implementing parallel processing for independent ECUs.

6.4.2 Vehicle Configuration Integration

Vehicle Configuration Database integration testing verified the system's ability to use vehicle configuration data for code rule evaluation and parameter file generation. Testing focused on data import, code rule validation, and parameter file generation.

Vehicle configuration data import testing confirmed that the system could correctly import and store vehicle configuration codes, with proper mapping between codes and vehicles. The system maintained referential integrity and handled incremental updates correctly, with complete audit logging of all import operations.

Code rule validation testing verified that the system could evaluate boolean expressions against vehicle configurations. Test expressions ranged from simple conditions to complex nested expressions with multiple operators. The evaluation engine correctly interpreted both simple logical operators and complex nested expressions with precedence rules.

Parameter file generation testing confirmed that the system could produce valid parameter files for vehicle testing, with correct application of variant selection logic based on vehicle configuration codes. The generated files included all required parameters with appropriate values, providing a complete configuration for ECU testing and validation.

6.5 Feature Comparison with Excel-Based Approach

To assess the improvements provided by the VMAP system, a feature comparison was conducted against the Excel-based approach currently used for parameter management. This comparison focused on capability coverage rather than performance metrics, providing a qualitative assessment of system improvements.

Table 6.10 presents a comparison of key features between the VMAP system and the Excel-based approach.

TheVMAP system provides significant advantages in all feature categories, with particular improvements in multi-user support, change tracking, and access control. These



| Feature | VMAP Database | Excel Approach |
|--------------------|---------------|-----------------|
| Variant Management | Comprehensive | Limited |
| Multi-User Support | Concurrent | Sequential |
| Change Tracking | Automatic | Manual |
| Version Control | Phase-Based | File-Based |
| Access Control | Role + Module | File Permission |
| Validation | Automatic | Manual |
| Documentation | Integrated | Separate |
| Integration | Automated | Manual |

Table 6.10: Feature Comparison with Excel-Based Approach

improvements address the limitations identified in the requirements analysis phase, providing a more robust and scalable solution for automotive parameter management.

6.5.1 Data Integrity Improvements

Data integrity was evaluated through a set of controlled test cases where both systems were subjected to various validation scenarios including invalid values, constraint violations, and relationship consistency. TheVMAP system demonstrated superior data integrity protection, correctly preventing invalid operations through database constraints and business rule validation.

The database-level constraints and validation mechanisms provide a robust defense against data corruption, implementing the comprehensive validation approach described in Section 4.7. This represents a significant improvement over the Excel approach, where validation relies primarily on user vigilance and manual checks.

6.5.2 Development Process Impact

Beyond technical improvements, theVMAP system introduces significant enhancements to the automotive parameter development process. The centralized database approach enables concurrent work by multiple engineers, eliminating the file sharing bottlenecks common in the Excel-based approach. The role-based access control ensures that engineers can modify only their assigned modules, preventing accidental changes to other areas.

The phase-based versioning approach aligns naturally with the automotive development lifecycle, supporting the structured progression from initial development through



testing to final release. The explicit phase transitions provide clear development milestones, while the freezing mechanism ensures configuration stability at critical points.

The comprehensive change tracking and documentation features address regulatory compliance requirements, providing complete traceability for all parameter modifications. This capability is increasingly important in the context of functional safety standards like ISO 26262, which require rigorous configuration management and change documentation.

7 Conclusion and Future Work

This chapter presents the conclusions drawn from the design, implementation, and evaluation of the Variant Management and Parametrization (VMAP) database system. The research has addressed critical challenges in automotive parameter management through a comprehensive database solution that supports version control, role-based access, and integration with enterprise systems. This chapter summarizes the key contributions and findings, discusses limitations of the current implementation, and outlines directions for future research and development.

7.1 Summary of Contributions

TheVMAP database system represents a significant advancement in automotive parameter management, replacing error-prone spreadsheet-based approaches with a structured, database-driven solution. The research has produced several key contributions to both academic knowledge and industry practice:

First, the implementation of a phase-based versioning model provides an effective approach to managing parameter evolution across automotive development cycles. Unlike generic temporal database approaches, this domain-specific versioning strategy aligns naturally with the structured development phases common in automotive engineering, supporting a clear separation between development stages while maintaining traceability [8]. The phase model enables simultaneous work on different development stages while preserving milestone integrity, addressing a fundamental challenge in automotive software development.

Second, the hybrid role-permission access control model delivers a flexible framework for managing user access across complex organizational structures. By combining role-based permissions with module-specific access controls, the system achieves a balance between administrative simplicity and access granularity [15]. This approach accommodates the specialized access patterns common in automotive development, where engineers typically have responsibility for specific vehicle subsystems rather than entire parameter sets.

Third, the integration architecture establishes robust connections with enterprise systems, enabling consistent data flow across the automotive development environment. The synchronization mechanisms maintain parameter definitions across different development phases while supporting variant customization and parameter file generation.



This integration addresses a critical requirement in automotive software development, where parameter management must operate within a complex ecosystem of engineering tools and databases [18].

Fourth, the comprehensive audit and documentation capabilities provide complete traceability for all parameter modifications, supporting both diagnostic analysis and regulatory compliance. The snapshot-based documentation approach creates immutable records of parameter configurations at significant development milestones, addressing the increasing regulatory requirements for configuration management in automotive software development [37].

Finally, the research provides valuable insights into database design patterns for complex domain-specific applications. The strategic denormalization approach, specialized indexing strategies, and query optimization techniques demonstrate effective database engineering practices for managing complex data relationships and supporting domain-specific workflows [28].

7.2 Key Findings

The evaluation of the VMAP system revealed several significant findings that validate the design decisions made during system development and provide insights for future database system implementations in automotive contexts.

7.2.1 Versioning Approach Effectiveness

The comparison between the implemented phase-based versioning approach and alternative change-based approach demonstrated significant performance advantages for common operations at the cost of moderate storage overhead. As shown in Section 6.3.4, the phase-based approach maintained query response times below 100ms even with large parameter sets, while the change-based approach exhibited nonlinear growth with increasing parameter counts.

This finding validates the architectural decision to implement explicit phase copies rather than a delta-based versioning model. While consuming approximately 51% more storage across all phases, the performance benefits for query operations justify this tradeoff, particularly for interactive operations where response time directly impacts user productivity. As noted by Bhattacherjee et al. [4], the recreation/storage tradeoff



in dataset versioning should prioritize operation frequency, with frequently accessed data favoring a storage-intensive approach that minimizes recreation costs.

The phase-based approach also demonstrated significant advantages for implementation simplicity and alignment with domain concepts. The explicit phase model aligned naturally with the mental model of automotive development engineers, reducing conceptual complexity and providing clear development milestones. These advantages extend beyond performance metrics to affect usability and adoption, particularly important factors for systems that must be integrated into established engineering workflows [27].

7.2.2 Access Control Performance

The evaluation of the hybrid role-permission model revealed that the approach effectively balanced flexibility with performance. As shown in Section 6.2.2, the addition of module-specific access checks introduced a performance overhead of 45.8% for permission verification in the hybrid model, compared to 88.2% in a traditional role-based approach.

This finding suggests that the more complex permission model accommodates additional access control dimensions with relatively lower overhead, validating Ferraiolo's observations [15] regarding the scalability of attribute-enhanced RBAC models. Despite the performance overhead, both approaches maintained acceptable performance for interactive operations, with response times below 4ms for individual permission checks.

The hybrid approach also demonstrated the flexibility required for complex organizational structures, supporting both role-based permission inheritance and direct permission assignments for exception cases. This flexibility addresses a common challenge in engineering organizations, where standard role definitions often require customization for specific projects or temporary access requirements [31].

7.2.3 Storage Distribution Insights

The analysis of storage requirements revealed significant insights about data distribution in parameter management systems with comprehensive audit capabilities. As shown in Section 6.3.3, the change history dominated storage allocation, consuming approximately 60.8% of the total database size, while active parameter data accounted for only 6.3%.



This distribution aligns with Bhattacherjee's observations [4] regarding versioning and audit systems, where historical record storage typically surpasses active data by a substantial margin. The finding has important implications for storage planning and database administration in regulated industries where comprehensive audit trails are mandatory.

Another significant finding was the storage impact of documentation snapshots, which consumed approximately 124MB despite representing only 7 distinct snapshots. This substantial storage footprint validates Fowler's observations [16] about the storage implications of the snapshot pattern, particularly when applied to complex data structures with many relationships.

7.2.4 External System Integration Challenges

The evaluation of integration with external enterprise systems revealed increasing synchronization time over successive operations, with execution times rising from approximately 15 minutes initially to around 30 minutes after 10 synchronization cycles. This gradual increase aligns with the observations of Mueller and Müller [26] regarding database synchronization in complex engineering environments.

This finding highlights a significant challenge for long-term system maintenance, with projected synchronization times potentially reaching 45 minutes after 20 synchronization cycles. While still acceptable for typical synchronization frequencies, this trend suggests that synchronization performance optimization should be a priority for future development.

The integration testing also validated the effectiveness of the vehicle configuration integration, with the system successfully evaluating complex boolean expressions against vehicle configurations. This capability is essential for supporting the variant-rich development approach common in the automotive industry, where parameters must be customized for numerous vehicle configurations [38].

7.3 Limitations

While the VMAP system successfully addresses the core requirements for automotive parameter management, several limitations were identified during implementation and evaluation that could affect its application in specific contexts.



7.3.1 Performance Limitations

The phase comparison operation demonstrated suboptimal scaling with increasing data volumes, with execution times increasing from 2.8 seconds for the baseline dataset to 12.4 seconds for the full dataset—a scaling factor of 4.4x. This operation, which involves complex joins across multiple tables to identify differences between parameter configurations, exceeds interactive response thresholds for larger datasets, potentially affecting usability for certain workflows.

The database change history mechanism, while providing comprehensive audit capabilities, contributes significantly to storage growth, with change records accounting for approximately 60.8% of the total database size. This growth pattern could present challenges for long-term database management, potentially requiring archiving strategies for older change records to maintain acceptable storage utilization and backup performance.

The integration with the PDD showed increasing synchronization times over successive operations, with execution times rising to approximately 30 minutes after 10 synchronization cycles. This performance degradation could impact the practicality of frequent synchronization operations in active development environments, where timely parameter updates are essential for effective engineering workflows.

7.3.2 Architectural Limitations

The phase-based versioning approach, while effective for structured development cycles, offers limited support for non-sequential development models where parameters might evolve through different paths. For organizations employing alternative development methodologies, this structured approach might introduce unnecessary rigidity in parameter management workflows.

The current implementation provides limited support for parameter dependency tracking, where changes to one parameter might necessitate adjustments to related parameters. Without explicit dependency management, ensuring parameter consistency across complex interrelationships relies primarily on user knowledge rather than system enforcement, potentially allowing inconsistent configurations.

The database schema, while well-optimized for the identified requirements, follows a relatively traditional relational approach rather than incorporating newer capabilities like native JSON support or graph database structures. This design choice prioritizes



reliability and compatibility over potential performance or modeling advantages offered by more specialized database technologies.

7.3.3 Integration Limitations

The current integration with external systems relies heavily on scheduled synchronization operations rather than event-driven updates. This approach introduces potential synchronization delays, where parameter changes in source systems might not be immediately reflected in the VMAP database. For time-sensitive engineering processes, these delays could impact workflow efficiency.

The vehicle configuration code rule evaluation mechanism, while effective for static rules, provides limited support for dynamic rule evaluation based on calculated values or complex relationships. This limitation could restrict the expressiveness of variant selection logic in advanced engineering scenarios where dynamic configuration determination is required.

7.4 Future Work

Based on the findings and limitations identified during system implementation and evaluation, several directions for future research and development have been identified. These opportunities range from technical optimizations to architectural extensions that could enhance the VMAP system's capabilities and address emerging requirements in automotive parameter management.

7.4.1 Performance Optimizations

Change History Partitioning Enhancements

The current implementation includes basic partitioning for the change history table based on release IDs, as described in Section 5.4. This approach could be extended with time-based subpartitioning to further improve performance for historical queries. By implementing a hierarchical partitioning scheme that combines release-based partitioning with time-based subpartitions, the system could achieve more granular data distribution while maintaining logical organization by release [28].



Additionally, implementing automated partition management for the change history table would enhance long-term maintainability. This extension would include partition rotation procedures that archive older partitions to lower-cost storage while maintaining rapid access to recent changes. According to Schwartz et al. [33], such automated partition management is essential for sustaining performance in rapidly growing audit systems.

Query Optimization for Phase Comparison

The phase comparison operation demonstrated suboptimal scaling with increasing data volumes, suggesting an opportunity for query optimization. Future development could explore materialized view approaches for phase comparison, where difference data is pre-computed and incrementally maintained rather than calculated on demand. This approach could significantly reduce response times for comparison operations, potentially bringing even complex comparisons within interactive response thresholds.

Another promising avenue is the implementation of parallel query execution for phase comparison operations. By partitioning the comparison workload across multiple execution threads, the system could leverage modern multi-core processors more effectively, reducing execution time for this resource-intensive operation [28].

Synchronization Performance Improvements

To address the observed synchronization performance degradation, future development could implement an incremental synchronization approach that focuses on changed entities rather than performing comprehensive comparisons. By leveraging change data capture techniques described by Seenivasan and Vaithianathan [34], the system could achieve more efficient synchronization with minimal overhead growth over time.

Additionally, the implementation of parallel synchronization for independent ECUs could significantly reduce overall synchronization time. By processing multiple ECUs concurrently, the system could better utilize available resources and reduce the total time required for synchronization operations, particularly in environments with numerous ECUs.



7.4.2 Architectural Enhancements

Advanced Versioning Capabilities

The current phase-based versioning model could be extended with branching capabilities to support parallel development streams. This enhancement would allow engineers to create specialized branches for experimental parameter configurations while maintaining the stability of the main development sequence. By incorporating branching concepts from software version control systems like Git, the database could support more flexible development workflows while preserving traceability [4].

The implementation of parametric inheritance mechanisms would enable more efficient management of variant similarities and differences. Rather than treating each variant as an independent entity, the system could implement inheritance hierarchies where specialized variants inherit parameter values from base configurations, reducing redundancy and simplifying maintenance of related variants [37].

Data Archiving and Retention

A comprehensive data archiving framework would enhance long-term database maintainability by providing structured processes for managing historical data. This framework would include configurable retention policies for different data types, automated archiving procedures for historical data, and seamless access mechanisms for archived information.

The implementation could leverage PostgreSQL's tablespace capabilities to physically separate active and archived data while maintaining logical accessibility through the database schema. By transitioning older data to compressed tablespaces or lower-cost storage, the system could maintain performance for active operations while preserving historical records for compliance purposes [28].

For regulatory compliance scenarios, the framework would include legal hold mechanisms that override standard retention policies for data subject to specific preservation requirements. This capability is particularly important in the automotive industry, where product liability considerations may require extended retention of development records [37].



Advanced Parameter Dependency Management

The implementation of explicit parameter dependency tracking would enhance system capabilities for maintaining consistency across related parameters. By modeling the relationships between parameters, the system could automatically identify potential inconsistencies when parameters are modified, prompting engineers to review and update related parameters as needed.

This capability could be extended with rule-based validation for parameter relationships, implementing domain-specific constraints that enforce engineering rules across parameter sets. Such validation would reduce the risk of inconsistent configurations and improve overall parameter quality by capturing domain expertise in executable rules [29].

7.4.3 Integration Enhancements

Event-Driven Integration

To address the limitations of scheduled synchronization, future development could implement event-driven integration with external systems. By leveraging database change data capture capabilities, the system could detect and react to changes in source systems in near-real-time, significantly reducing synchronization delays.

This approach would involve implementing a message-based integration architecture as described by Hohpe and Woolf [18], where change events from source systems trigger specific synchronization actions in the VMAP database. By processing changes incrementally as they occur rather than in batch operations, the system could maintain more consistent data across systems while reducing synchronization overhead.

Enhanced Vehicle Configuration Integration

The vehicle configuration integration could be enhanced with support for dynamic rule evaluation, allowing more sophisticated variant selection logic based on calculated values or complex relationships. This extension would enable more flexible variant applicability definitions, addressing advanced engineering scenarios where static rule evaluation is insufficient.

Additionally, implementing a simulation framework for vehicle configurations would enhance parameter validation capabilities, allowing engineers to verify parameter



behavior across diverse vehicle scenarios. By combining configuration simulation with parameter resolution, the system could provide valuable insights into parameter behavior before physical testing, potentially reducing development cycles [37].

7.5 Broader Implications

The development and evaluation of the VMAP database system offers several broader implications for database research and automotive software development beyond the specific implementation described in this thesis.

7.5.1 Implications for Database Research

TheVMAP system demonstrates the effectiveness of domain-specific versioning approaches over generic temporal database techniques for specialized applications. By aligning database versioning with the natural structure of the application domain—in this case, automotive development phases—the system achieves both conceptual clarity and performance advantages. This finding suggests that domain-specific adaptations of established database patterns may offer significant benefits in specialized contexts, a perspective that could inform future research in applied database design [4].

The implementation also highlights the evolving relationship between relational database systems and document-oriented approaches. While maintaining a fundamentally relational structure, the VMAP system leverages Postgre SQL's JSONB capabilities for change tracking, combining structured schema enforcement with flexible document storage. This hybrid approach suggests promising directions for database research that bridges traditional relational models with document-oriented flexibility [28].

The performance analysis provides empirical evidence of the storage-performance tradeoffs inherent in different versioning strategies, offering valuable data points for researchers exploring versioning approaches in other domains. These findings contribute to the ongoing discussion of recreation/storage tradeoffs in versioned datasets, providing concrete measurements of these tradeoffs in a real-world implementation [4].



7.5.2 Implications for Automotive Software Development

For automotive software development, the VMAP system demonstrates the feasibility and advantages of transitioning from document-based parameter management to structured database approaches. The evaluation results provide empirical evidence of the performance, consistency, and traceability improvements possible through this transition, potentially encouraging similar transformations in other aspects of automotive software development [8].

The integration architecture establishes patterns for connecting specialized engineering databases with enterprise systems, addressing a significant challenge in automotive development environments. By demonstrating effective synchronization between parameter management and vehicle configuration systems, the implementation provides a template for similar integrations across the automotive development ecosystem [18].

The comprehensive audit capabilities implemented in theVMAP system highlight the increasing importance of traceability in automotive software development, particularly as vehicles become more software-defined and subject to regulatory oversight. The storage implications of these audit requirements, as revealed in the evaluation, provide valuable planning insights for automotive organizations implementing similar systems [37].

More broadly, the VMAP system represents a step toward the more rigorous software engineering practices increasingly required in automotive development. By applying established database principles to the specialized domain of parameter management, the system contributes to the evolution of automotive software development from document-centric approaches to structured, database-driven methodologies [29].

7.6 Conclusion

TheVMAP database system represents a significant advancement in automotive parameter management, addressing fundamental challenges in version control, user access, and system integration. Through a carefully designed database architecture implemented in PostgreSQL, the system provides a robust foundation for managing parameter configurations across development phases, supporting variant customization for diverse vehicle configurations.

The phase-based versioning approach, hybrid role-permission model, and comprehensive audit capabilities collectively enable a more structured and reliable parameter management process. These capabilities address critical limitations of spreadsheet-based



approaches, providing improved data integrity, enhanced traceability, and systematic validation of parameter configurations.

While the implementation has demonstrated clear advantages over existing approaches, several opportunities for enhancement remain in areas such as performance optimization, architectural extensions, and integration refinements. Future development in these areas could further improve system capabilities and address emerging requirements in automotive parameter management.

Beyond its specific technical contributions, theVMAP system illustrates the value of applying proven database engineering principles to specialized domains like automotive development. By combining domain-specific knowledge with established database patterns, the system achieves both technical excellence and practical usability, demonstrating the potential for database-driven approaches to transform complex engineering workflows.

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A User Management Test Cases

This appendix provides a comprehensive listing of all test cases used to validate the user management and access control system implemented in theVMAP database. The test cases are organized by category and include detailed information about test actions, expected outcomes, and test results.

A.1 Role-Based Permission Test Cases

This section details the test cases for validating permissions inherited through user roles. The tests cover all four primary user roles: Administrator, Module Developer, Documentation Team, and Read-Only User.

Table A.1: Administrator Role Permission Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|------------------|-----------------------|-------------------------|--------|
| AD- | Create User | Add new user with | User created suc- | Pass |
| 01 | | valid details | cessfully | |
| AD- | Modify User Role | Change user's as- | Role updated suc- | Pass |
| 02 | | signed role | cessfully | |
| AD- | Delete User | Remove existing user | User deleted suc- | Pass |
| 03 | | | cessfully | |
| AD- | Create Role | Create new role with | Role created suc- | Pass |
| 04 | | permissions | cessfully | |
| AD- | Delete Variant | Delete existing vari- | Variant deleted suc- | Pass |
| 05 | | ant | cessfully | |
| AD- | Freeze Phase | Set phase status to | Phase frozen suc- | Pass |
| 06 | | frozen | cessfully | |

Table A.2: Module Developer Role Permission Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|----------------------------------|-------------------------|--------|
| MD- | Create Variant (As- | Create new variant | Variant created suc- | Pass |
| 01 | signed Module) | for parameter in assigned module | cessfully | |



Table A.2 – Continued from previous page

| ID | Description | Test Action | Expected Outcome | Status |
|-----|-------------------|-------------------------|----------------------|--------|
| MD- | Create Variant | Create new variant | Access denied error | Pass |
| 02 | (Unassigned Mod- | for parameter in | | |
| | ule) | unassigned module | | |
| MD- | Edit Variant (As- | Modify existing vari- | Variant updated suc- | Pass |
| 03 | signed Module) | ant code rule | cessfully | |
| MD- | Delete Variant | Attempt to delete vari- | Access denied error | Pass |
| 04 | | ant | | |
| MD- | Create Segment | Create new segment | Segment created | Pass |
| 05 | (Assigned Module) | with valid value | successfully | |
| MD- | Modify Frozen | Attempt to modify | Access denied error | Pass |
| 06 | Phase | segment in frozen | | |
| | | phase | | |
| MD- | Generate Parame- | Create parameter file | File generated suc- | Pass |
| 07 | ter File | for testing | cessfully | |
| MD- | Read Parameters | View parameters | Parameters dis- | Pass |
| 80 | (Any Module) | from any module | played successfully | |

Table A.3: Documentation Team Role Permission Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|-----------------------|-------------------------|--------|
| DT- | Create Documenta- | Create snapshot of | Snapshot created | Pass |
| 01 | tion Snapshot | frozen phase | successfully | |
| DT- | Compare Phases | Compare parameters | Comparison results | Pass |
| 02 | | between two phases | displayed | |
| DT- | View Parameter His- | View change history | History displayed | Pass |
| 03 | tory | for parameter | successfully | |
| DT- | Export Hex String | Copy parameter hex | Hex string copied | Pass |
| 04 | | string | successfully | |
| DT- | Modify Parameter | Attempt to modify pa- | Access denied error | Pass |
| 05 | | rameter | | |
| DT- | Access All Phases | View parameters | Parameters dis- | Pass |
| 06 | | across all phases | played successfully | |
| DT- | Generate Parame- | Create parameter file | File generated suc- | Pass |
| 07 | ter File | for reference | cessfully | |



Table A.4: Read-Only User Role Permission Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|------------------|-----------------------|-------------------------|--------|
| RO- | View Parameters | Access parameter | Parameters dis- | Pass |
| 01 | | details | played successfully | |
| RO- | View Variants | Access variant de- | Variants displayed | Pass |
| 02 | | tails | successfully | |
| RO- | Modify Parameter | Attempt to modify pa- | Access denied error | Pass |
| 03 | | rameter | | |
| RO- | Modify Variant | Attempt to modify | Access denied error | Pass |
| 04 | | variant | | |
| RO- | Generate Parame- | Create parameter file | File generated suc- | Pass |
| 05 | ter File | for reference | cessfully | |

A.2 Module-Based Access Control Test Cases

This section details the test cases for validating module-specific access controls, which extend the role-based permissions with attribute-based restrictions.

Table A.5: Module-Based Access Control Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|-----------------------|-------------------------|--------|
| MA- | Assign Module Ac- | Grant write access to | Access granted suc- | Pass |
| 01 | cess | specific module | cessfully | |
| MA- | Revoke Module Ac- | Remove write access | Access revoked suc- | Pass |
| 02 | cess | to specific module | cessfully | |
| MA- | Read Access Cross- | Access parameters | Read access suc- | Pass |
| 03 | Module | from unassigned | cessful | |
| | | module | | |
| MA- | Write Access As- | Create variant in as- | Variant created suc- | Pass |
| 04 | signed Module | signed module | cessfully | |
| MA- | Write Access Unas- | Create variant in | Access denied error | Pass |
| 05 | signed Module | unassigned module | | |
| MA- | Multiple Module As- | Create variants in | All variants created | Pass |
| 06 | signment | multiple assigned | successfully | |
| | | modules | | |



Table A.5 – Continued from previous page

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|--------------------|-----------------------|--------|
| MA- | Edit Segment As- | Modify segment in | Segment updated | Pass |
| 07 | signed Module | assigned module | successfully | |
| MA- | Edit Segment Unas- | Modify segment in | Access denied error | Pass |
| 80 | signed Module | unassigned module | | |
| MA- | Administrator Over- | Admin modifies any | Modification success- | Pass |
| 09 | ride | module | ful | |
| MA- | Module Permission | User with role | Access updated suc- | Pass |
| 10 | Inheritance | change inherits | cessfully | |
| | | proper module | | |
| | | access | | |

A.3 Direct Permission Assignment Test Cases

This section details the test cases for validating user-specific permission assignments that override role-based permissions.

Table A.6: Direct Permission Assignment Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|----------------------|-----------------------|--------|
| DP- | Grant Additional | Assign permission | Permission applied | Pass |
| 01 | Permission | not in user's role | successfully | |
| DP- | Revoke Role Per- | Remove permission | Permission restric- | Pass |
| 02 | mission | normally granted by | tion applied | |
| | | role | | |
| DP- | Grant Delete Per- | Give read-only user | Deletion operation | Pass |
| 03 | mission | delete permission | successful | |
| DP- | Permission Conflict | Conflicting role and | Direct permission | Pass |
| 04 | Resolution | direct permissions | takes precedence | |
| DP- | Role Change with | Change user's role | Custom permissions | Pass |
| 05 | Custom Permission | with custom permis- | preserved | |
| | | sions | | |
| DP- | Permission Audit | Track changes to | Audit trail correctly | Pass |
| 06 | Trail | user permissions | recorded | |



A.4 Phase-Specific Permission Test Cases

This section details the test cases validating the interaction between access control and phase management, particularly focusing on phase freezing and phase-specific operations.

Table A.7: Phase-Specific Permission Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|--------------------|-----------------------|-------------------------|--------|
| PP- | Frozen Phase Modi- | Attempt to modify | Access denied error | Pass |
| 01 | fication | variant in frozen | | |
| | | phase | | |
| PP- | Documentation | Documentation team | Access granted suc- | Pass |
| 02 | Access to Frozen | accesses frozen | cessfully | |
| | Phase | phase | | |
| PP- | Administrator | Administrator un- | Phase unfrozen suc- | Pass |
| 03 | Unfreeze | freezes a phase | cessfully | |
| PP- | Non-Administrator | Module developer | Access denied error | Pass |
| 04 | Freeze Attempt | attempts to freeze | | |
| | | phase | | |
| PP- | Read Access to | Read-only user ac- | Access granted suc- | Pass |
| 05 | Frozen Phase | cesses frozen phase | cessfully | |
| PP- | Phase Transition | Module developer ini- | Transition completed | Pass |
| 06 | Permission | tiates phase transi- | successfully | |
| | | tion | | |

A.5 Boundary Case Test Cases

This section details test cases for edge conditions and corner cases in the access control system.

| ID | Description | Test Action | Expected Outcome | Status |
|-----|-------------------|-----------------------|-------------------------|--------|
| BC- | No Role Assign- | User with no as- | Access limited to | Pass |
| 01 | ment | signed role attempts | public content | |
| | | access | | |
| BC- | Multiple Role As- | User with multiple | Most permissive role | Pass |
| 02 | signment | roles attempts action | takes effect | |
| BC- | Role With No Per- | Assign user to empty | No permissions | Pass |
| 03 | missions | role | granted | |
| BC- | Session Timeout | Session expires dur- | User properly redi- | Pass |
| 04 | Handling | ing operation | rected to login | |

Table A.8: Boundary Case Test Cases

A.6 Test Implementation Details

Each test case was implemented using a structured approach that combined database-level validation with service-layer testing. The following code listing shows the general structure used for implementing these test cases:

```
[Test]
2 public void
    \TestCaseID_Description_ExpectedOutcome()
  {
      // Arrange: Set up test environment
      var testUser = CreateTestUser("[UserRole]")
        ↓;
      var testEntity = CreateTestEntity();
      // Configure specific test conditions
      ConfigureTestConditions();
      // Act: Perform the operation being tested
      if (ShouldSucceed)
      {
13
          var result = _service.PerformOperation(
            \testEntity, testUser.UserId);
15
          // Assert: Verify operation succeeded
```



```
Assert. IsNotNull (result);
17
           Assert.That(result.Status, Is.EqualTo(
18
             ↓OperationStatus.Success));
19
           // Verify database state reflects the
20
             var dbEntity = _database.
             □QuerySingleOrDefault <Entity > (
               "SELECT * FROM entities WHERE id =
                 □@Id",
               new { Id = testEntity.Id });
23
           Assert.IsNotNull(dbEntity);
24
           Assert. That (dbEntity. Property, Is.
25
             \EqualTo(testEntity.Property));
      }
26
      else
      {
28
           // Assert: Verify operation is denied
29
             with appropriate error
           var exception = Assert.Throws<</pre>
30
             PermissionDeniedException > (() =>
               service.PerformOperation(
                 \testEntity, testUser.UserId));
           Assert. That (exception. Message, Contains
32
             \( .Substring("expected error message")
             ↳);
33
           // Verify database state was not
             \modified
           var dbEntity = _database.
35
             └QuerySingleOrDefault <Entity > (
               "SELECT * FROM entities WHERE id =
36
                 □@Id",
               new { Id = testEntity.Id });
           Assert.That(dbEntity, Is.Null().Or.
38

Property("Property")

                                    .Not.EqualTo(
39

    Property));
```



41 }

Listing A.1: Test Case Implementation Template

This standardized approach ensured consistent validation across all test cases while providing clear evidence of both successful permission grants and appropriate permission denials. Each test verified both the immediate operation result and the resulting database state, ensuring comprehensive validation of the access control system.

A.7 Role Permission Matrix

Table A.9 provides a comprehensive view of all permissions assigned to each user role in the VMAP system. This matrix formed the basis for the permission validation test cases.

Permission Admin Module Dev Doc Team Read-Only manage_users \checkmark \checkmark manage roles × × × delete_variants \checkmark × × × create_variants \checkmark \checkmark × × \checkmark edit_variants × × **√** create segments \checkmark × × edit segments delete_segments \checkmark \checkmark × × \checkmark \checkmark create_snapshots × view history \checkmark \checkmark \checkmark \checkmark generate_par_files freeze_phases \checkmark view_all \checkmark \checkmark \checkmark \checkmark

Table A.9: Role Permission Matrix

Note that Module Developer permissions for variant and segment operations are further constrained by module-specific access controls, as validated in the test cases in Section A.2.

B Variant Management Test Cases

This appendix provides a comprehensive listing of all test cases used to validate the variant management functionality implemented in the VMAP database. The test cases are organized by category and include detailed information about test actions, expected outcomes, and test results.

B.1 Variant Creation Test Cases

This section details the test cases for validating variant creation functionality across different parameter types and constraints.

Table B.1: Variant Creation Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|--------------------|-----------------------|-------------------------|--------|
| VC- | Basic Variant Cre- | Create variant with | Variant created suc- | Pass |
| 01 | ation | valid name and code | cessfully | |
| | | rule | | |
| VC- | Duplicate Variant | Create variant with | Name uniqueness er- | Pass |
| 02 | Name | name that already ex- | ror | |
| | | ists inPID | | |
| VC- | Empty Variant | Create variant with | Validation error | Pass |
| 03 | Name | empty name | | |
| VC- | Special Characters | Create variant with | Variant created suc- | Pass |
| 04 | in Name | special characters in | cessfully | |
| | | name | | |
| VC- | Maximum Name | Create variant with | Variant created suc- | Pass |
| 05 | Length | 100-character name | cessfully | |
| | | (maximum length) | | |
| VC- | Exceed Name | Create variant with | Validation error | Pass |
| 06 | Length | name exceeding 100 | | |
| | | characters | | |
| VC- | Valid Code Rule | Create variant with | Variant created suc- | Pass |
| 07 | | syntactically valid | cessfully | |
| | | code rule | | |



Table B.1 – Continued from previous page

| ID | Description | Test Action | Expected Outcome | Status |
|-----|----------------------|------------------------|-------------------------|--------|
| VC- | Complex Code Rule | Create variant with | Variant created suc- | Pass |
| 08 | | complex rule contain- | cessfully | |
| | | ing multiple operators | | |
| VC- | InvalidPIDReference | Create variant with | Foreign key con- | Pass |
| 09 | | non-existentPID | straint error | |
| VC- | Creation in Frozen | Create variant in a | Phase frozen error | Pass |
| 10 | Phase | frozen phase | | |
| VC- | Variant in Inac- | Create variant for | Validation error | Pass |
| 11 | tivePID | parameter in inac- | | |
| | | tivePID | | |
| VC- | Null Code Rule | Create variant with | Variant created suc- | Pass |
| 12 | | null code rule | cessfully | |
| VC- | Variant Audit Trail | Create variant and | Audit record created | Pass |
| 13 | | verify audit trail | correctly | |
| VC- | Variant for Boolean | Create variant for pa- | Variant created suc- | Pass |
| 14 | Parameter | rameter with boolean | cessfully | |
| | | type | | |
| VC- | Variant for Enum Pa- | Create variant for pa- | Variant created suc- | Pass |
| 15 | rameter | rameter with enumer- | cessfully | |
| | | ation type | | |
| VC- | Concurrent Variant | Create variants con- | All variants created | Pass |
| 16 | Creation | currently from multi- | successfully | |
| | | ple sessions | | |
| VC- | Transaction Roll- | Begin transaction, | No variant created | Pass |
| 17 | back | create variant, then | | |
| | | force rollback | | |
| VC- | Permission Verifica- | Create variant with | Permission denied er- | Pass |
| 18 | tion | insufficient permis- | ror | |
| | | sions | | |

B.2 Segment Modification Test Cases

This section details the test cases for validating segment modification functionality across different parameter dimensions and value types.



Table B.2: Segment Creation Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|--------------------|---------------------------|--------------------|--------|
| SC- | Create Scalar Seg- | Create segment for | Segment created | Pass |
| 01 | ment | scalar parameter | successfully | |
| SC- | Create Array Seg- | Create segment for | Segment created | Pass |
| 02 | ment (1D) | 1D array parameter | successfully | |
| SC- | Create Matrix Seg- | Create segment for | Segment created | Pass |
| 03 | ment (2D) | 2D matrix parameter | successfully | |
| SC- | Create 3D Array | Create segment for | Segment created | Pass |
| 04 | Segment | 3D array parameter | successfully | |
| SC- | Invalid Dimension | Create segment with | Validation error | Pass |
| 05 | Index | out-of-bounds dimen- | | |
| | | sion index | | |
| SC- | Invalid Parameter | Create segment with | Foreign key con- | Pass |
| 06 | Reference | non-existent parame- | straint error | |
| | | ter ID | | |
| SC- | Integer Parameter | Create segment with | Segment created | Pass |
| 07 | Value | integer parameter | successfully | |
| | | type | | |
| SC- | Float Parameter | Create segment with | Segment created | Pass |
| 08 | Value | float parameter type | successfully | |
| SC- | Boolean Parameter | Create segment with | Segment created | Pass |
| 09 | Value | boolean parameter | successfully | |
| 00 | NA'-' A | type | 0 | Desir |
| SC- | Minimum Value | Create segment with | Segment created | Pass |
| 10 | Boundary | minimum allowed | successfully | |
| SC- | Maximum Value | value Create segment with | Segment created | Pass |
| 11 | Boundary | maximum allowed | | rass |
| 11 | Bouridary | value | successfully | |
| SC- | Below Minimum | Create segment with | Validation error | Pass |
| 12 | Value | value below minimum | | |
| SC- | Above Maximum | Create segment with | Validation error | Pass |
| 13 | Value | value above maxi- | | |
| | | mum | | |
| SC- | Creation in Frozen | Create segment in a | Phase frozen error | Pass |
| 14 | Phase | frozen phase | | |



Table B.2 – Continued from previous page

| ID | Description | on | Test Action | Expected Outcome | Status |
|-----|-------------|-----------|----------------------|-------------------------|--------|
| SC- | Duplicate | | Create segment for | Unique constraint er- | Pass |
| 15 | Parameter | ſ- | already modified pa- | ror | |
| | Dimension | 1 | rameter dimension | | |
| SC- | High | Precision | Create segment with | Segment created | Pass |
| 16 | Value | | high precision deci- | successfully | |
| | | | mal value | | |

Table B.3: Segment Update Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|------------------------|-------------------------|--------|
| SU- | Update Scalar Seg- | Modify existing scalar | Segment updated | Pass |
| 01 | ment | segment value | successfully | |
| SU- | Update 1D Array El- | Modify element in 1D | Segment updated | Pass |
| 02 | ement | array segment | successfully | |
| SU- | Update 2D Matrix | Modify element in 2D | Segment updated | Pass |
| 03 | Element | matrix segment | successfully | |
| SU- | Value Range Verifi- | Update segment with | Validation error | Pass |
| 04 | cation | value outside valid | | |
| | | range | | |
| SU- | Update in Frozen | Modify segment in a | Phase frozen error | Pass |
| 05 | Phase | frozen phase | | |
| SU- | Concurrent Updates | Update same seg- | Last update pre- | Pass |
| 06 | | ment from multiple | served with proper | |
| | | sessions | locking | |
| SU- | Update Non- | Update segment that | Not found error | Pass |
| 07 | Existent Segment | doesn't exist | | |
| SU- | Change to Default | Update segment to | Segment updated | Pass |
| 08 | Value | match default param- | successfully | |
| | | eter value | | |

Table B.4: Segment Deletion Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|--------------------|----------------------|-------------------------|--------|
| SD- | Delete Single Seg- | Remove existing seg- | Segment deleted | Pass |
| 01 | ment | ment | successfully | |



Table B.4 – Continued from previous page

| ID | Description | Test Action | Expected Outcome | Status |
|-----|----------------------|----------------------|-------------------------|--------|
| SD- | Delete Non-Existent | Delete segment that | Not found error | Pass |
| 02 | Segment | doesn't exist | | |
| SD- | Delete in Frozen | Delete segment in a | Phase frozen error | Pass |
| 03 | Phase | frozen phase | | |
| SD- | Cascade Delete via | Delete variant and | All segments deleted | Pass |
| 04 | Variant | verify segments cas- | | |
| | | cade | | |
| SD- | Cascade Delete via | Delete parameter | All segments deleted | Pass |
| 05 | Parameter | and verify segments | | |
| | | cascade | | |
| SD- | Segment Deletion | Delete segment and | Audit record created | Pass |
| 06 | Audit | verify audit trail | correctly | |
| SD- | Permission Verifica- | Delete segment | Permission denied er- | Pass |
| 07 | tion | with insufficient | ror | |
| | | permissions | | |
| SD- | Transaction Roll- | Begin transaction, | Segment not deleted | Pass |
| 08 | back | delete segment, then | | |
| | | force rollback | | |

B.3 Performance Test Cases

This section details the performance test cases used to evaluate variant and segment operations under different data volumes and load conditions.

Table B.5: Variant and Segment Performance Test Cases

| ID | Description | Test Action | Expected Outcome | Status |
|-----|---------------------|---------------------|-------------------------|--------|
| VP- | Baseline Variant | Create 10 variants | < 2 seconds total | Pass |
| 01 | Creation | and measure time | time | |
| VP- | Baseline Segment | Create 100 seg- | < 10 seconds total | Pass |
| 02 | Creation | ments and measure | time | |
| | | time | | |
| VP- | High Volume Variant | Create 100 variants | < 20 seconds total | Pass |
| 03 | Creation | for singlePID | time | |



| Table b | Table 6.5 – Continued Ironi previous page | | | |
|---------|---|----------------------|--|--|
| n | Test Action | Expected Outc | | |
| ne Seg- | Create 1000 seg- | < 2 minutes total | | |
| ion | ments across multi- | | | |

Table P. F. Continued from provious page

ID **Description** come **Status** VP-High Volum ıl time Pass 04 ment Creation ple variants VP-SinglePIDLoad Test Create 500 variants System remains re-**Pass** 05 for singlePID sponsive VP-Multi-dimensional < 3 minutes total time **Pass** Create segments for Parameter Load 06 3D parameter with 1000 elements VP-Concurrent No deadlocks or er-**Pass** User 10 concurrent users Simulation 07 creating variants rors VP-Variant Retrieval Response Retrieve variants time < **Pass** 80 Scaling fromPIDs with 10, 250ms 100, and 500 variants

B.4 Test Implementation Details

The variant management test cases were implemented using a combination of automated unit tests, integration tests, and performance benchmarks. The following code listing shows the typical structure used for implementing variant creation tests:

```
[Test]
 public void VC01 BasicVariantCreation Success()
  {
3
      // Arrange
      var testUser = userRepository.GetTestUser(
        \"module developer@example.com");
      var testPid = _pidRepository.GetTestPid();
      var variant = new VariantCreationPayload
      {
          PidId = testPid.PidId,
          EcuId = testPid.EcuId,
          PhaseId = _activePhaseId,
          Name = "Test Variant " + Guid.NewGuid()
13
            4. ToString(). Substring(0, 8),
```



```
CodeRule = "A AND (B OR C)"
      };
15
16
      // Act
17
      var result = variantService.CreateVariant(
18
         \variant, testUser.UserId);
      // Assert
20
      Assert. IsNotNull (result);
21
      Assert.That(result.VariantId, Is.
         GreaterThan(0));
23
      // Verify database state
      var dbVariant = _database.
         □QuerySingleOrDefault < Variant > (
           "SELECT * FROM variants WHERE
26
             \variant_id = @VariantId",
           new { VariantId = result.VariantId });
27
      Assert. IsNotNull(dbVariant);
      Assert.That(dbVariant.Name, Is.EqualTo(
30

¬variant.Name));
      Assert.That(dbVariant.CodeRule, Is.EqualTo(
31
         \variant.CodeRule));
      Assert.That(dbVariant.CreatedBy, Is.EqualTo
         \( (testUser.UserId));
33
      // Verify audit trail
      var auditRecord = _database.
35
         └QuerySingleOrDefault < ChangeRecord > (
           "SELECT * FROM change history WHERE
             \entity_type = 'variants' " +
           "AND entity_id = @VariantId AND
37
             \change type = 'CREATE'",
           new { VariantId = result.VariantId });
38
39
      Assert.IsNotNull(auditRecord);
      Assert.That(auditRecord.UserId, Is.EqualTo(
41
         \testUser.UserId));
  }
42
```



Listing B.1: Variant Creation Test Implementation Example

Similarly, segment modification tests followed this structure but with appropriate adaptations for the specific operations:

```
[Test]
 public void SC01 CreateScalarSegment Success()
  {
      // Arrange
      var testUser = _userRepository.GetTestUser(
        \"module_developer@example.com");
      var testVariant = _variantRepository.
        □GetTestVariant();
      var testParameter = parameterRepository.
        GetScalarParameter(testVariant.PidId);
      var segment = new SegmentCreationPayload
      {
          VariantId = testVariant.VariantId,
          ParameterId = testParameter.ParameterId
          DimensionIndex = 0,
13
          Decimal = 42.5m
      };
      // Act
      var result = segmentService.CreateSegment(
18
        \segment, testUser.UserId);
      // Assert
      Assert. IsNotNull (result);
      Assert.That(result.SegmentId, Is.
        GreaterThan(0));
23
      // Verify database state
      var dbSegment = _database.
        QuerySingleOrDefault < Segment > (
          "SELECT * FROM segments WHERE
```



```
\segment_id = @SegmentId",
          new { SegmentId = result.SegmentId });
27
28
      Assert. IsNotNull (dbSegment);
      Assert.That(dbSegment.VariantId, Is.EqualTo
30
         └(segment.VariantId));
      Assert. That (dbSegment. ParameterId, Is.
         \EqualTo(segment.ParameterId));
      Assert.That(dbSegment.DimensionIndex, Is.
32
         \EqualTo(segment.DimensionIndex));
      Assert.That(dbSegment.Decimal, Is.EqualTo(
33

¬segment.Decimal));
      Assert.That(dbSegment.CreatedBy, Is.EqualTo
         \((testUser.UserId));
35
      // Verify parameter value is within valid
         ¬range
      var parameterRange = _database.
37
         \QuerySingleOrDefault < ParameterRange > (
           "SELECT * FROM parameter values WHERE
38
             bparameter_id = @ParameterId",
          new { ParameterId = testParameter.
             ParameterId });
40
         (parameterRange != null)
42
          Assert. That (segment. Decimal, Is.
43
             GreaterThanOrEqualTo(parameterRange.
             ↓ValueRangeBegin));
           Assert.That(segment.Decimal, Is.
44
             LessThanOrEqualTo(parameterRange.
             ↓ValueRangeEnd));
      }
45
  }
46
```

Listing B.2: Segment Modification Test Implementation Example

Performance tests were implemented using a benchmarking approach that measured execution time across multiple iterations:



```
[Test]
  public void
    \VP01_BaselineVariantCreation_Performance()
  {
      // Arrange
      var testUser = _userRepository.GetTestUser(
         \"module developer@example.com");
      var testPid = _pidRepository.GetTestPid();
      var variants = new List<</pre>
         \VariantCreationPayload > ();
      for (int i = 0; i < 10; i++)</pre>
      {
           variants.Add(new VariantCreationPayload
           {
               PidId = testPid.PidId,
               EcuId = testPid.EcuId,
14
               PhaseId = _activePhaseId,
               Name = $"Perf Test Variant {i} {
16
                 □ Guid. NewGuid(). ToString().
                 \Substring(0, 8)}",
               CodeRule = "A AND B"
           });
18
      }
19
      // Act
      var stopwatch = new Stopwatch();
      stopwatch.Start();
24
      foreach (var variant in variants)
      {
           _variantService.CreateVariant(variant,
             ↓testUser.UserId);
      }
29
      stopwatch.Stop();
30
      // Assert
      Assert.That(stopwatch.ElapsedMilliseconds,
```



Listing B.3: Performance Test Implementation Example

This standardized approach ensured comprehensive validation of the variant management functionality while providing detailed performance metrics for system evaluation.

B.5 Test Environment Configuration

All variant management tests were conducted in a controlled test environment with the following specifications:

- PostgreSQL 17 running on Windows Server 2022
- Database server: 8 vCPUs, 32GB RAM, SSD storage
- · Application server: 4 vCPUs, 16GB RAM
- Database containing baseline dataset (20,000 parameters, 188 variants, 28,776 segments)
- Testing conducted with both the baseline dataset and scaled dataset (100,000 parameters, 830 variants, 167,990 segments)
- Network latency between application and database servers < 1ms
- PostgreSQL configuration optimized for test environment with appropriate memory allocation for shared buffers, work memory, and maintenance work memory

The test environment was reset to a known state between test runs using database snapshots, ensuring consistent starting conditions for each test execution.