

# Parsing Arithmetic Expressions



<https://courses.missouristate.edu/anthonyclark/333/>

# Outline

## Topics and Learning Objectives

- Learn about parsing arithmetic expressions
- Learn how to handle associativity with a grammar
- Learn how to handle precedence with a grammar

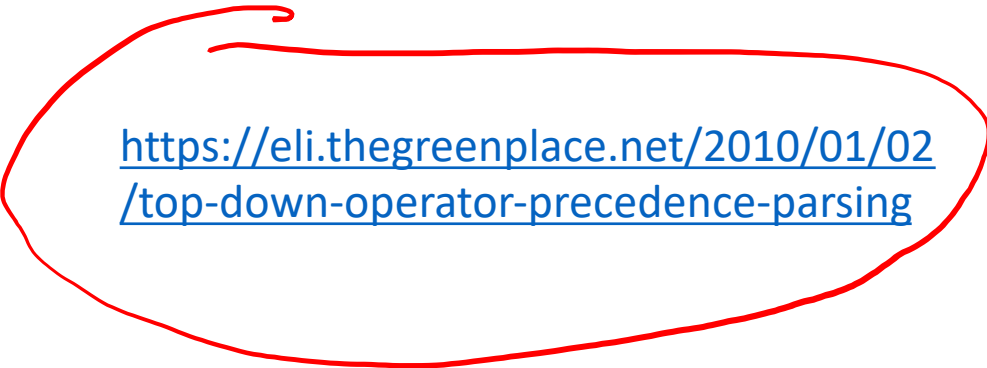
## Assessments

- ANTLR grammar for math

# Parsing Expressions

There are a variety of special purpose algorithms to make this task more efficient:

- The shunting yard algorithm
- Precedence climbing
- Pratt parsing



<https://eli.thegreenplace.net/2010/01/02/top-down-operator-precedence-parsing>

For this class we are just going to use recursive descent

- 
- Simpler
  - Same as the rest of our parser

# Grammar for Expressions

Needs to account for operator **associativity**

- Also known as fixity
- Determines how you apply operators of the same precedence
- Operators can be **left-associative** or **right-associative** (Binary Operator)

Needs to account for operator **precedence**

- Precedence is a concept that you know from mathematics
- Think PEMDAS
- Apply higher precedence operators first

# Associativity

By convention

$7 + 3 + 1$  is equivalent to  $(7 + 3) + 1$ ,

$7 - 3 - 1$  is equivalent to  $(7 - 3) - 1$  and

$(12 / 3) * 4$  is equivalent to  $(12 / (3 * 4))$

- If we treated  $7 - 3 - 1$  as  $7 - (3 - 1)$   
the result would be 5 instead of the 3.

- Another way to state this convention is *associativity*

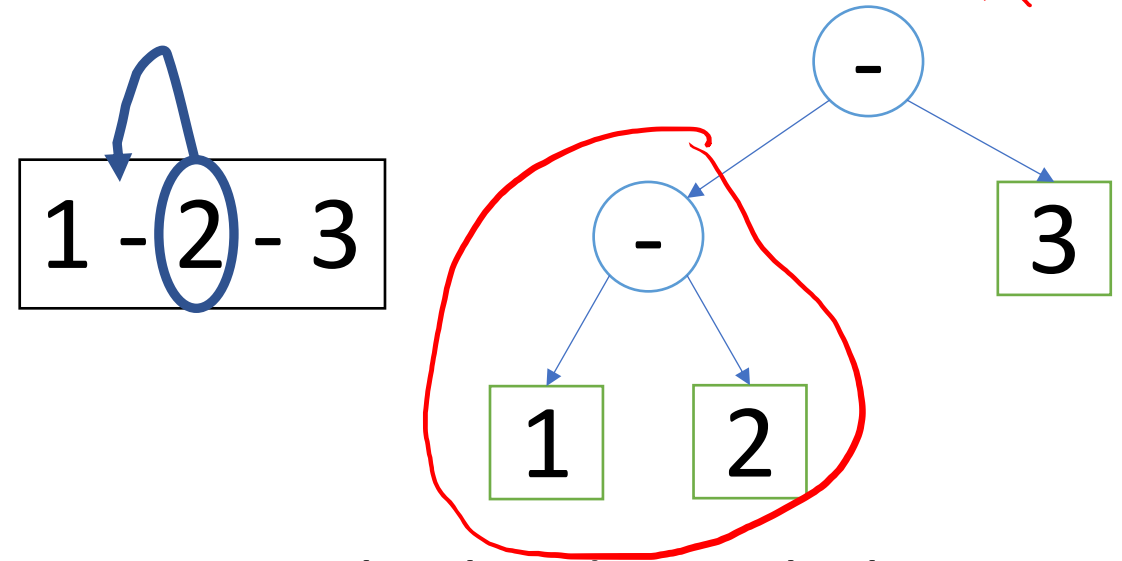
# Associativity

Addition, subtraction, multiplication, and division are **left-associative** *AST*

What does this mean?

You have:

- **operators** (+, -, \*, /, etc.) and
- **operands** (numbers, ids, etc.)
- Left-associativity: if an **operand** has **operators** on both sides with the same precedence, it associates with the operator on its **left**



# Associativity in the Grammar

For **left**-associativity use **left**-recursion

For **right**-associativity use **right**-recursion

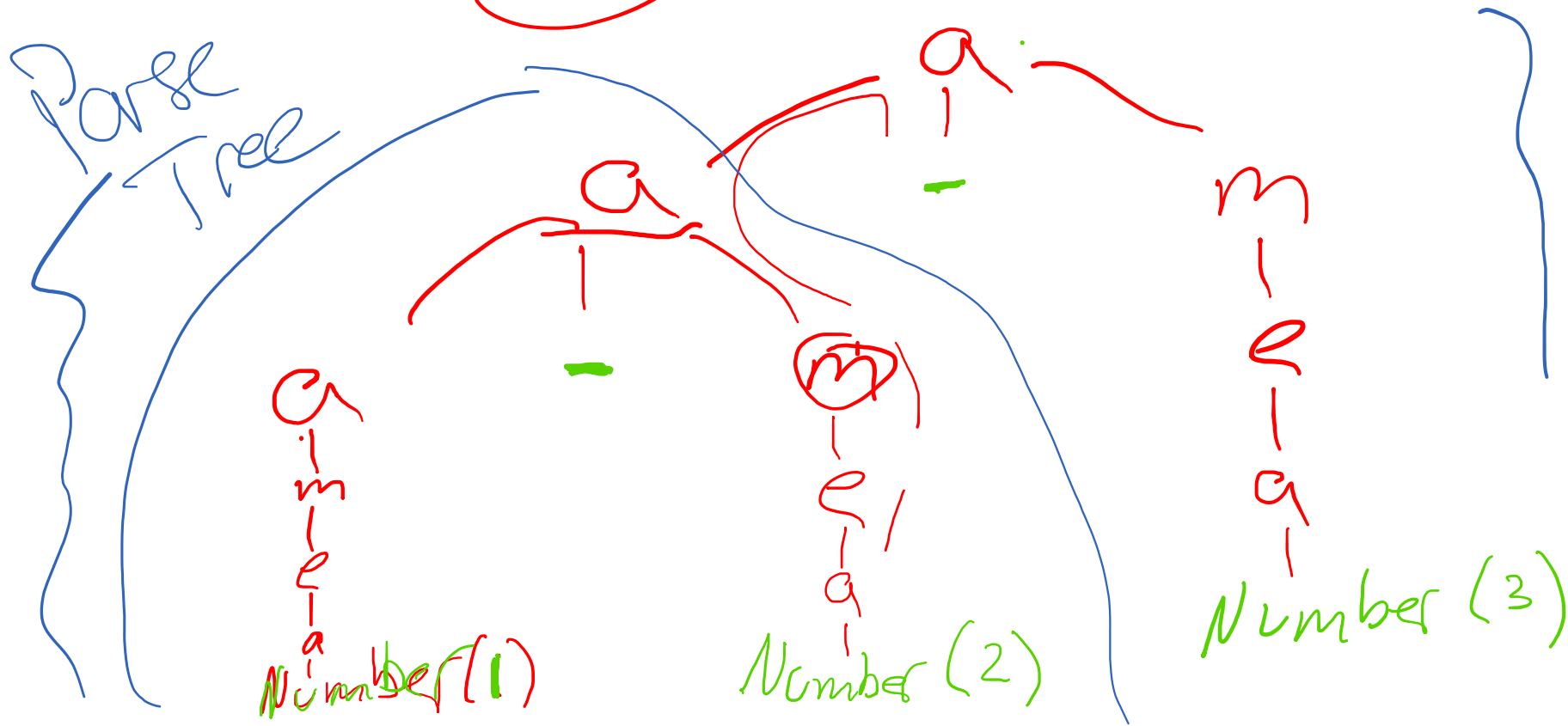
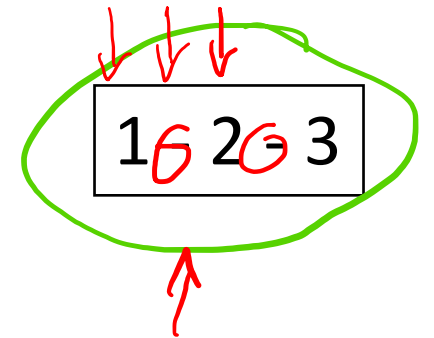
```
add_sub      : add_sub ('+' | '-') mul_div | mul_div;
```

```
mul_div      : mul_div ('*' | '/') exponentiate | exponentiate;
```

```
exponentiate : atom_unit '^' exponentiate | atom_unit;
```

```
atom_unit    : Number | (' ' add_sub '');
```

add\_sub : add\_sub ('+' | '-') mul\_div | mul\_div;  
 mul\_div : mul\_div ('\*' | '/') exponentiate | exponentiate;  
 exponentiate : atom\_unit '^' exponentiate | atom\_unit;  
 atom\_unit : Number | '(' add\_sub ')';





add\_sub : add\_sub ('+' | '-') mul\_div | mul\_div;

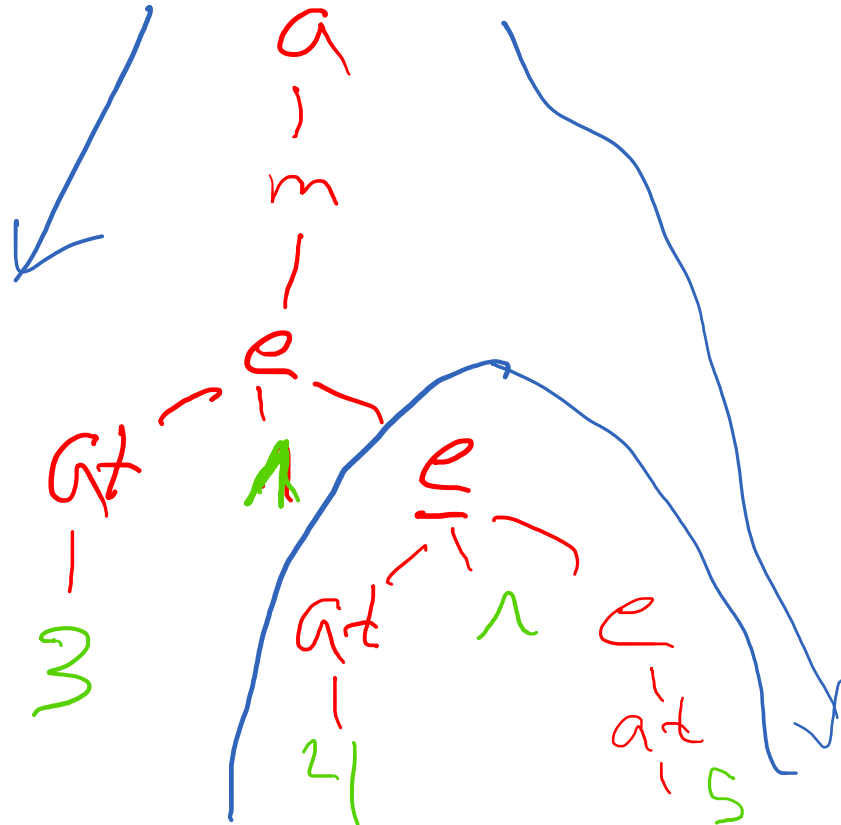
mul\_div : mul\_div ('\*' | '/') exponentiate | exponentiate;

exponentiate : atom\_unit '^' exponentiate | atom\_unit;

atom\_unit : Number | '(' add\_sub ')';

3 ^ 4 ^ 5

3<sup>4<sup>5</sup></sup>



Parse Tree



# Associativity in the Grammar

*We cannot implement  
this grammar.  
Recursive Descent  
Parser*

For **left**-associativity use **left**-recursion

For **right**-associativity use **right**-recursion

Any issues with implementing this grammar?

**add\_sub** : **add\_sub** ( '+' | '-' ) **mul\_div** | **mul\_div**;

**mul\_div** : **mul\_div** ( '\*' | '/' ) **exponentiate** | **exponentiate**;

**exponentiate** : **atom\_unit** '^' **exponentiate** | **atom\_unit**;

**atom\_unit** : Number | '(' **add\_sub** ')';

# Precedence

What about the following expression:  $7 + 5 * 2$

$$\begin{array}{l} 7 + (5 * 2) \quad \swarrow \\ (7 + 5) * 2 \quad \swarrow \end{array}$$

- Associativity is not helpful here...
- Now you just fall back to PEMDAS

P	E	M	D	A	S
---	---	---	---	---	---

# Precedence in the Grammar

Higher-precedence should appear in “lower” (called later) grammar rules

```
add_sub      : add_sub ('+' | '-') mul_div | mul_div;
mul_div      : mul_div ('*' | '/') exponentiate | exponentiate;
exponentiate : atom_unit '^' exponentiate | atom_unit;
atom_unit    : Number | '(' add_sub ')';
```

E	M	D	A	S
---	---	---	---	---

# Precedence in the Grammar

Higher-precedence should appear in “lower” (called later) grammar rules

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add_sub      : add_sub ('+' | '-') mul_div | mul_div;
mul_div      : mul_div ('*' | '/') exponentiate | exponentiate;
exponentiate : atom_unit '^' exponentiate | atom_unit;
atom_unit    : Number | '(' add_sub ')'; P
```

M	D	A	S
---	---	---	---

# Precedence in the Grammar

Higher-precedence should appear in “lower” (called later) grammar rules

```

add_sub      : add_sub ('+' | '-') mul_div | mul_div;
mul_div      : mul_div ('*' | '/') exponentiate | exponentiate;
exponentiate : atom_unit '^' E exponentiate | atom_unit;
atom_unit    : Number | '(' add_sub ')'; P
  
```

# Precedence in the Grammar

Higher-precedence should appear in “lower” (called later) grammar rules

```

add_sub      : add_sub ('+' | '-') mul_div | mul_div;
mul_div      : mul_div M ('*' | '/') D exponentiate | exponentiate;
exponentiate : atom_unit '^' E exponentiate | atom_unit;
atom_unit    : Number | '(' add_sub ')'; P
  
```

# Precedence in the Grammar

Higher-precedence should appear in “lower” (called later) grammar rules

**add\_sub** : **add\_sub** A ( '+' | '-' ) S **mul\_div** | **mul\_div**;

**mul\_div** : **mul\_div** M ( '\*' | '/' ) D **exponentiate** | **exponentiate**;

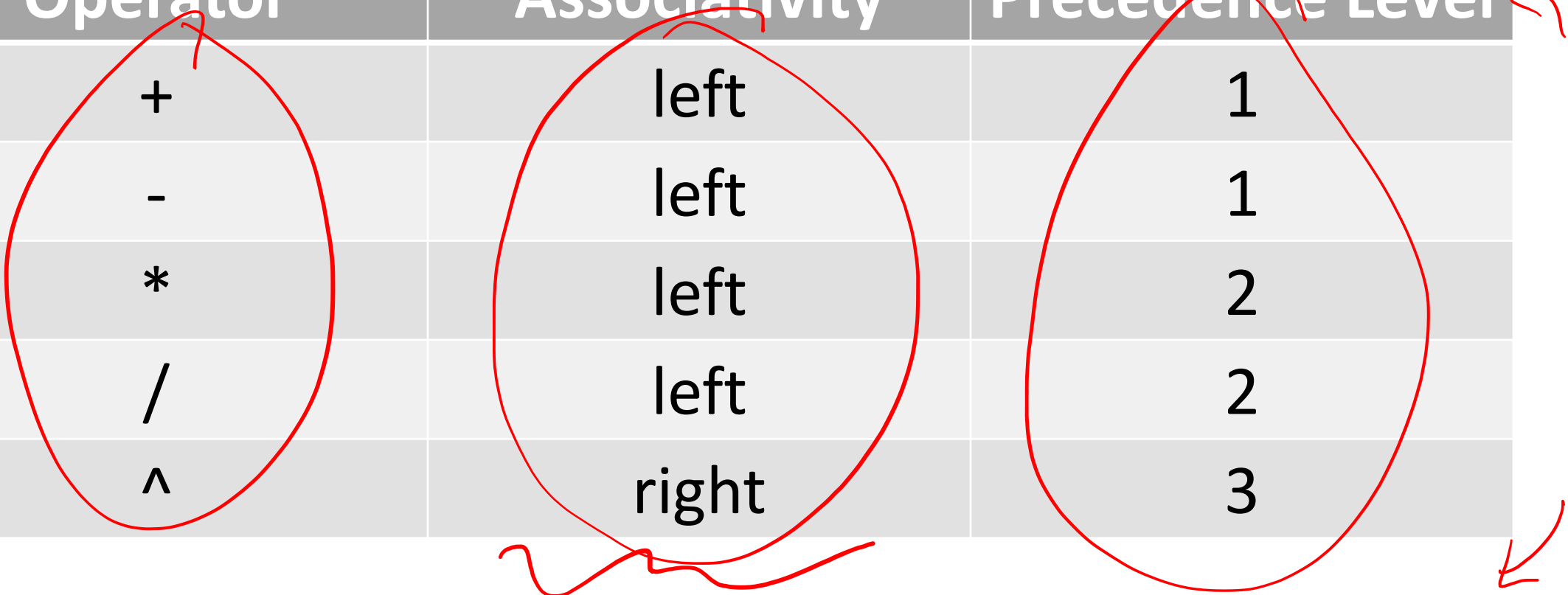
**exponentiate** : **atom\_unit** '^' E **exponentiate** | **atom\_unit**;

**atom\_unit** : Number | '(' **add\_sub** ')'; P



# Table of Operators

Operator	Associativity	Precedence Level
+	left	1
-	left	1
*	left	2
/	left	2
^	right	3



The table is annotated with hand-drawn red circles and arrows. A red circle encircles the entire 'Operator' column. Another red circle encircles the 'Associativity' column, with a red arrow pointing to the 'right' entry for the '^' operator. A third red circle encircles the 'Precedence Level' column, with a red arrow pointing to the '3' entry for the '^' operator. Additionally, a red arrow points from the top right towards the 'Precedence Level' header, and another red arrow points from the bottom right towards the '3' entry.

# A Grammar from the Table

Operator	Associativity	Level
+	left	1
-	left	1
*	left	2
/	left	2
^	right	3

For each **precedence level** in the table, define a non-terminal

- The RHS of the rule should match operators for that level
- The RHS should include non-terminals for the next higher precedence
- For **left**-associativity use **left**-recursion
- For **right**-associativity use **right**-recursion

If you have **n** precedence levels, you will need **n + 1** rules

( ) Numbers, Ids, etc.

# A Grammar from the Table

Operator	Associativity	Level
+	left	1
-	left	1
*	left	2
/	left	2
^	right	3

*Parse*

For each **precedence level** in the table, define a non-terminal

- The RHS of the rule should match operators for that level
- The RHS should include **non-terminals** for the next higher precedence
- For **left**-associativity use **left**-recursion
- For **right**-associativity use **right**-recursion

**Activity**

If you have **n** precedence levels, you will need **n + 1** rules

hashat : hashat '#' dollar

| dollar '@' hashat

| dollar ;

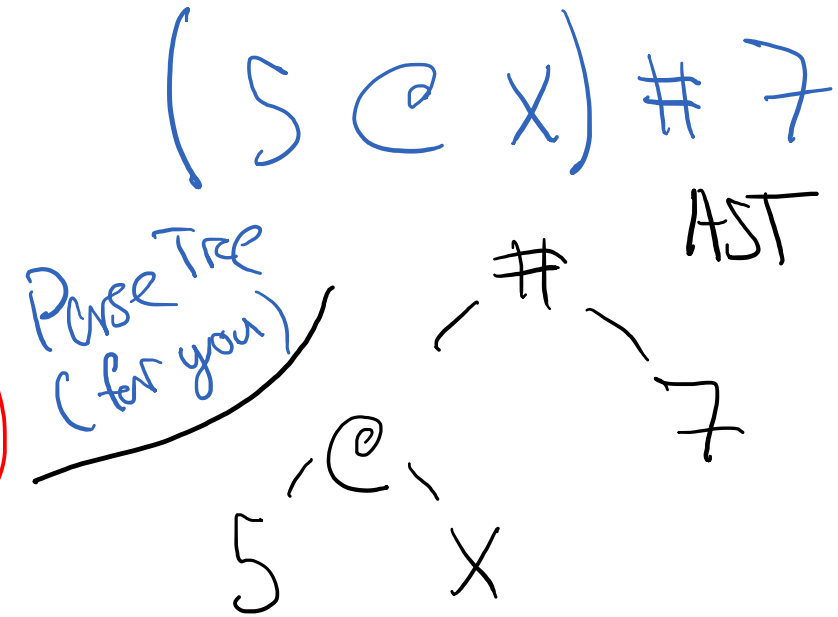
dollar : dollar '\$' atom

| atom ;

atom : Numbers | Identifiers | '(' hashat ')';

Operator	Associativity	Level
#	Left	1
@	Right	1
\$	Left	2

The basic units are Numbers, Identifiers, and parenthetical expressions.



```

hashat : hashat '#' dollar | dollar '@' hashat | dollar ;
dollar : dollar '$' unit | unit ;
unit   : Number | Identifier | '(' hashat ')' ;

```

hashat : dollar @ hashat hashat' | dollar ;  
 hashat' : # dollar hashat' | ε ;

dollar : unit dollar' ;  
 dollar' : \$ unit dollar' | ε ;

Left-Recursion Transformation	
Input Pattern:	
$A$	$\rightarrow A\alpha_1 \mid A\alpha_2 \mid \dots \mid A\alpha_n \mid \beta_1 \mid \beta_2 \mid \dots \mid \beta_m$
Output Pattern:	
$A$	$\rightarrow \beta_1 A' \mid \beta_2 A' \mid \dots \mid \beta_m A'$
$A'$	$\rightarrow \alpha_1 A' \mid \alpha_2 A' \mid \dots \mid \alpha_n A' \mid \varepsilon$

expr : ID '+' expr  
| ID '\*' expr  
| '(' expr ')'  
| ID  
| NUMBER

Try to parse the following:

x + y + z  
(x + y) \* z  
z \* (x + y)  
x + y \* z  
x \* y + z

Is the grammar expressive?

No

Is precedence correct?

No

Is associativity correct?

No

Can we implement the grammar as is?

Yes

expr : ID '+' expr  
     | ID '\*' expr  
     | '(' expr ')'  
     | ID  
     | NUMBER

Try to parse the following:

x + y + z

(x + y) \* z

z \* (x + y)

x + y \* z

x \* y + z

Is the grammar expressive?

No

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Can we implement the grammar as is?

Yes

expr : expr '+' expr  
| expr '\*' expr  
| '(' expr ')'  
| ID  
| NUMBER

Try to parse the following:

(x + y) \* z  
x + y + z

Is the grammar expressive?

Yes

Is precedence correct?

No

Is associativity correct?

No (ambiguous)

Can we implement the grammar as is?

No (left recursions)



expr : term '+' expr  
| term '\*' expr  
| term  
term : '(' expr ')'  
| ID  
| NUMBER

Try to parse the following:

(x + y) \* z  
x + y \* z  
x \* y + z

Pause  
⌞

Is the grammar expressive?

No

Is precedence correct?

No

Is associativity correct?

No

Can we implement the grammar as is?

Yes



Pause

expr : expr '+' term  
     | expr '\*' term  
     | term  
term : '(' expr ')'  
     | ID  
     | NUMBER

Try to parse the following:

{ (x + y) \* z  
  x + y \* z  
  x \* y + z  
  x + y + z

Is the grammar expressive?

Yes

Is precedence correct?

No

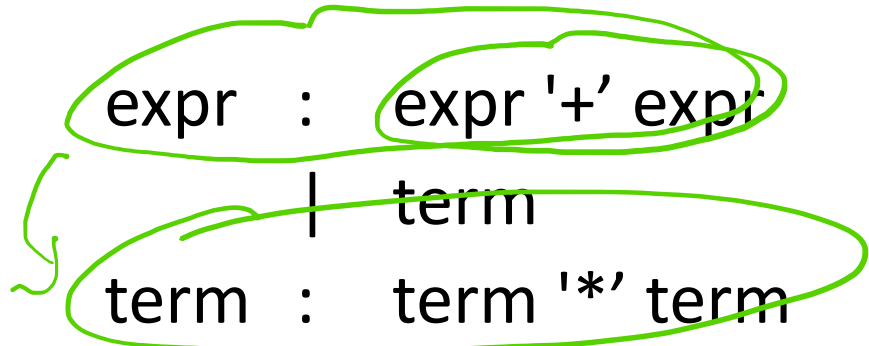
Is associativity correct?

Yes

Can we implement the grammar as is?

No (left recursions)

Pause



expr : expr '+' expr  
| term  
term : term '\*' term  
| factor  
factor : '(' expr ')'  
| ID  
| NUMBER

Try to parse the following:

(x + y) \* z  
x + y \* z  
x \* y + z  
x + y + z

Is the grammar expressive?

Yes

Is precedence correct?

Yes

Is associativity correct?

No (ambiguous)

Can we implement the grammar as is?

No (left recursions)

expr : expr '+' term  
     | term  
term : term '\*' factor  
     | factor  
factor: '(' expr ')'  
     | ID  
     | NUMBER

Try to parse the following:

(x + y) \* z  
x + y \* z  
x \* y + z  
x + y + z

Is the grammar expressive?

Yes

Is precedence correct?

Yes

Is associativity correct?

Yes

Can we implement the grammar as is?

No (left recursions)

Pause  
Do it  
Do it now



# Expression Grammar

Start with this:

$$\begin{aligned} \text{expr} &\rightarrow \text{expr } ('+' \mid '-') \text{term} \mid \text{term} \\ \text{term} &\rightarrow \text{term } ('*' \mid '/') \text{factor} \mid \text{factor} \\ \text{factor} &\rightarrow \text{NUMBER} \mid \text{ID} \mid '(' \text{expr} ') \end{aligned}$$

*Equivalent*

Turn it in to this (after removing left recursion):

$$\begin{aligned} \text{expr} &\rightarrow \text{term expr}' \\ \text{expr}' &\rightarrow ('+' \mid '-') \text{term expr}' \mid \text{ep} \\ \text{term} &\rightarrow \text{factor term}' \\ \text{term}' &\rightarrow ('*' \mid '/') \text{factor term}' \mid \text{ep} \\ \text{factor} &\rightarrow \text{NUMBER} \mid \text{ID} \mid '(' \text{expr} ') \end{aligned}$$

*Associativity ✓  
Precedence ✓  
• Expression ✓  
Not left recursive ✓*

Left recursive grammar (cannot be implemented using recursive descent)

# Expression Grammar

```
expr  -> expr ('+' | '-') term | term
term  -> term ('*' | '/') factor | factor
factor -> NUMBER | ID | '(' expr ')'
```

BNF (Backus–Naur form):

```
expr  -> term expr'
expr' -> ('+' | '-') term expr' | ep
term  -> factor term'
term' -> ('*' | '/') factor term' | ep
factor -> NUMBER | ID | '(' expr ')'
```

Becomes while-loop

**EBNF** (Extended Backus–Naur form):

```
expr  -> term [ ('+' | '-') term ]*
term  -> factor [ ('*' | '/') factor ]*
factor -> '(' expr ')' | ID | NUMBER
```

Much easier to  
implement!