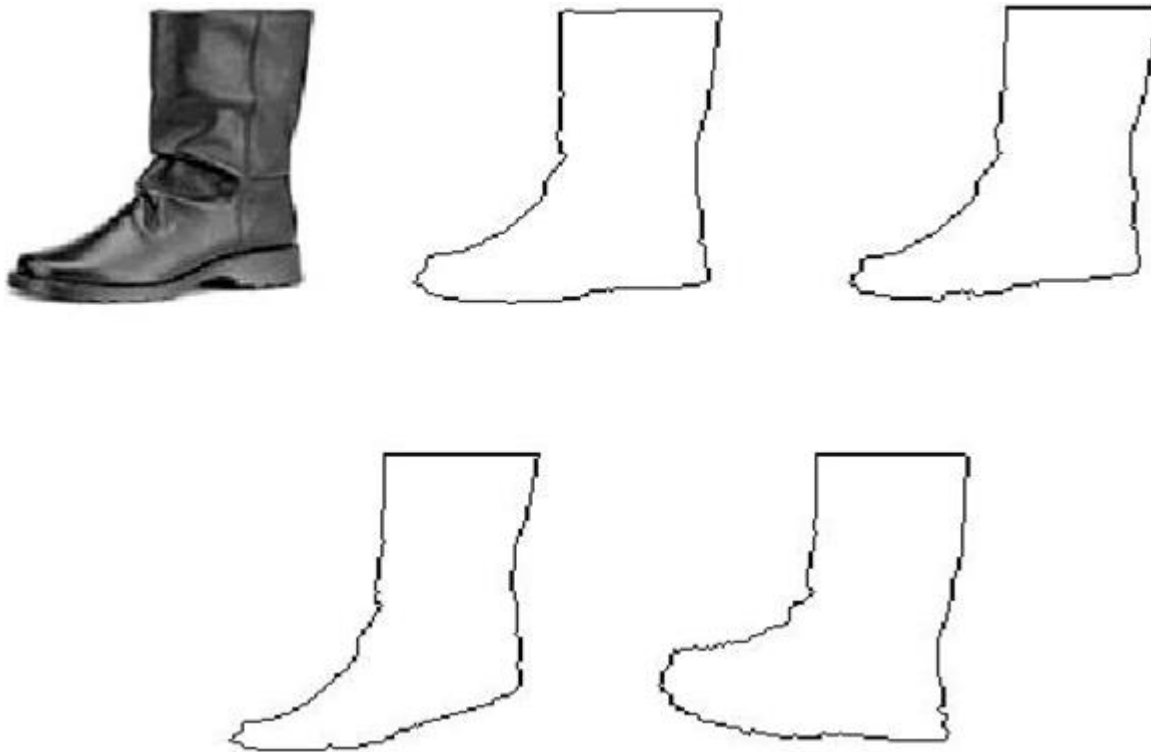


CSE 473
Pattern Recognition

Template Matching



Template Matching

- Typical Applications
 - Speech Recognition
 - Motion Estimation in Video Coding
 - Data Base Image Retrieval
 - Written Word Recognition
 - Bioinformatics

Template Matching

- The Goal:
 - Given a set of reference patterns known as **TEMPLATES**,
 - find the best match for unknown pattern
 - each class represented by a single typical pattern.
- requires an appropriate “measure” to quantify similarity or matching.

Template Matching

- The cost “measure”:
 - deviations between the **template** and the **test pattern**.

Template Matching

- The cost “measure”:
 - deviations between the **template** and the **test pattern**.
 - For example:
 - The word **beauty** may have been read as **beeauty** or **beuty**, etc., due to errors.
 - The **same person** may speak the **same word** **differently**.

Template Matching Methods

- Optimal path searching techniques
- Correlation
- Deformable models

TM using Optimal Path Searching

- Representation: Represent the template by a **sequence** of **measurement vectors** or **string patterns**

Template: $\underline{r}(1), \underline{r}(2), \dots, \underline{r}(I)$

Test pattern: $\underline{t}(1), \underline{t}(2), \dots, \underline{t}(J)$

TM using Optimal Path Searching

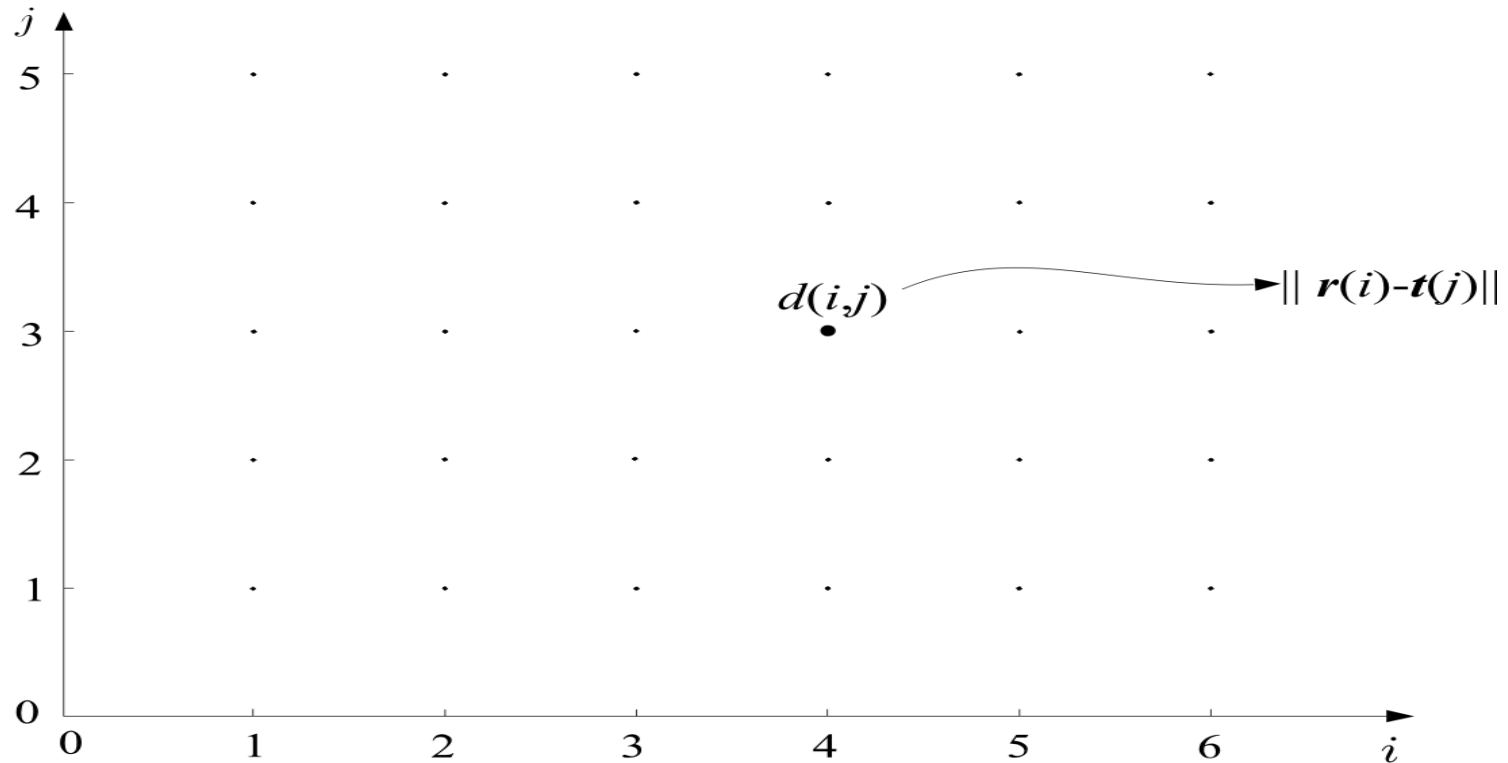
Template: $\underline{r}(1), \underline{r}(2), \dots, \underline{r}(I)$

Test pattern: $\underline{t}(1), \underline{t}(2), \dots, \underline{t}(J)$

- In general $I \neq J$
- We need to find an appropriate distance measure between test and reference patterns.

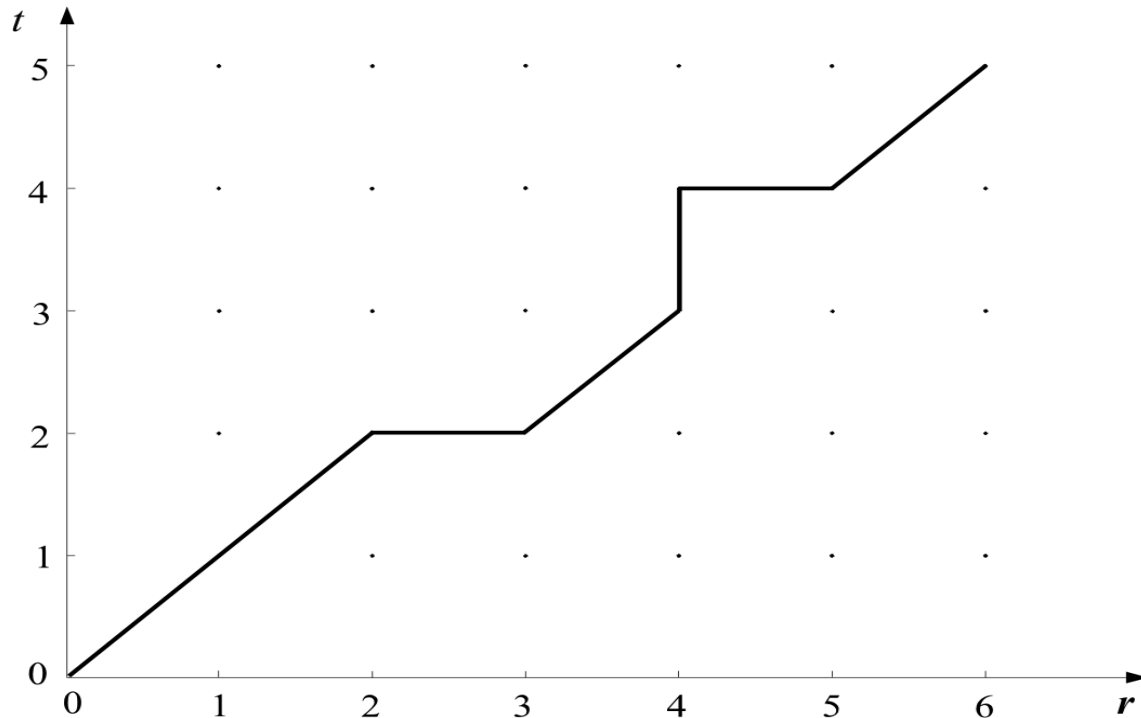
TM using Optimal Path Searching

- Form a grid with I points (template) in horizontal and J points (test) in vertical
- Each point (i,j) of the grid measures the **distance** between $\underline{r}(i)$ and $\underline{t}(j)$



TM using Optimal Path Searching

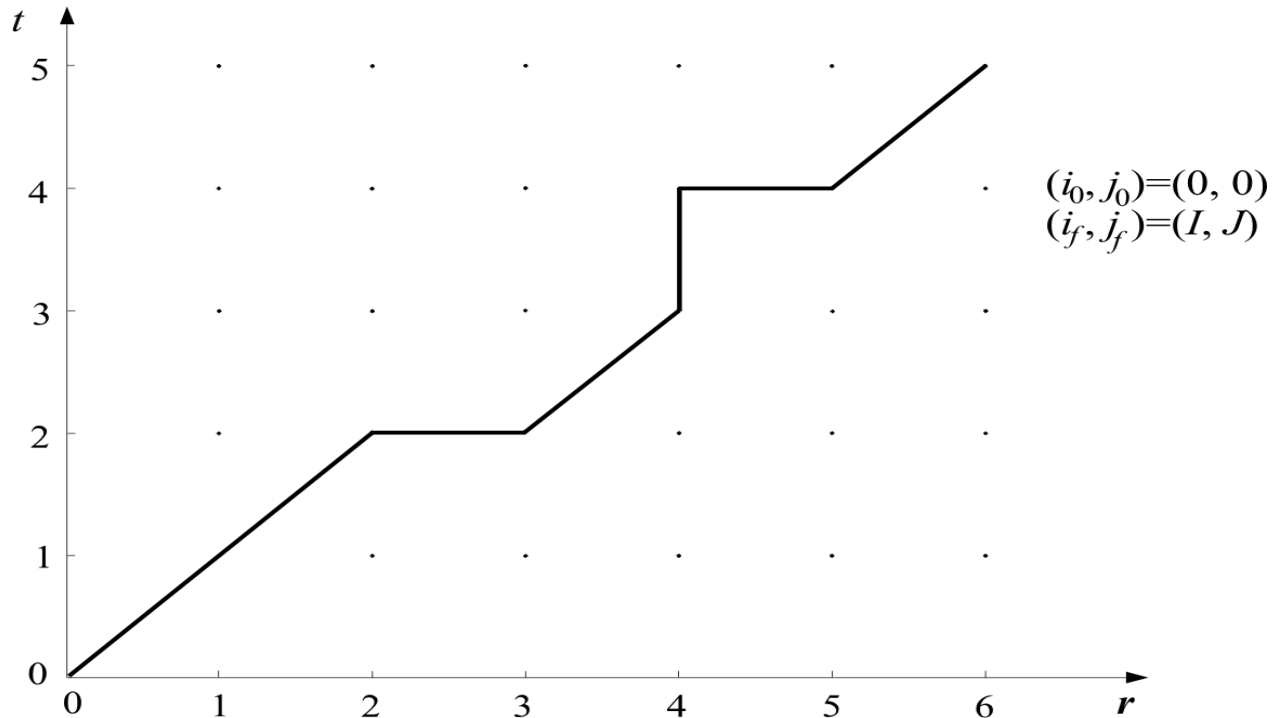
- **Path:** A path through the grid, from an **initial node** (i_0, j_0) to a **final one** (i_f, j_f) , is an **ordered set** of nodes $(i_0, j_0), (i_1, j_1), (i_2, j_2) \dots (i_k, j_k) \dots (i_f, j_f)$



TM using Optimal Path Searching

– **Path**: A path is complete path if:

$$(i_0, j_0) = (0, 0), (i_1, j_1), (i_2, j_2), \dots, (i_f, j_f) = (I, J)$$

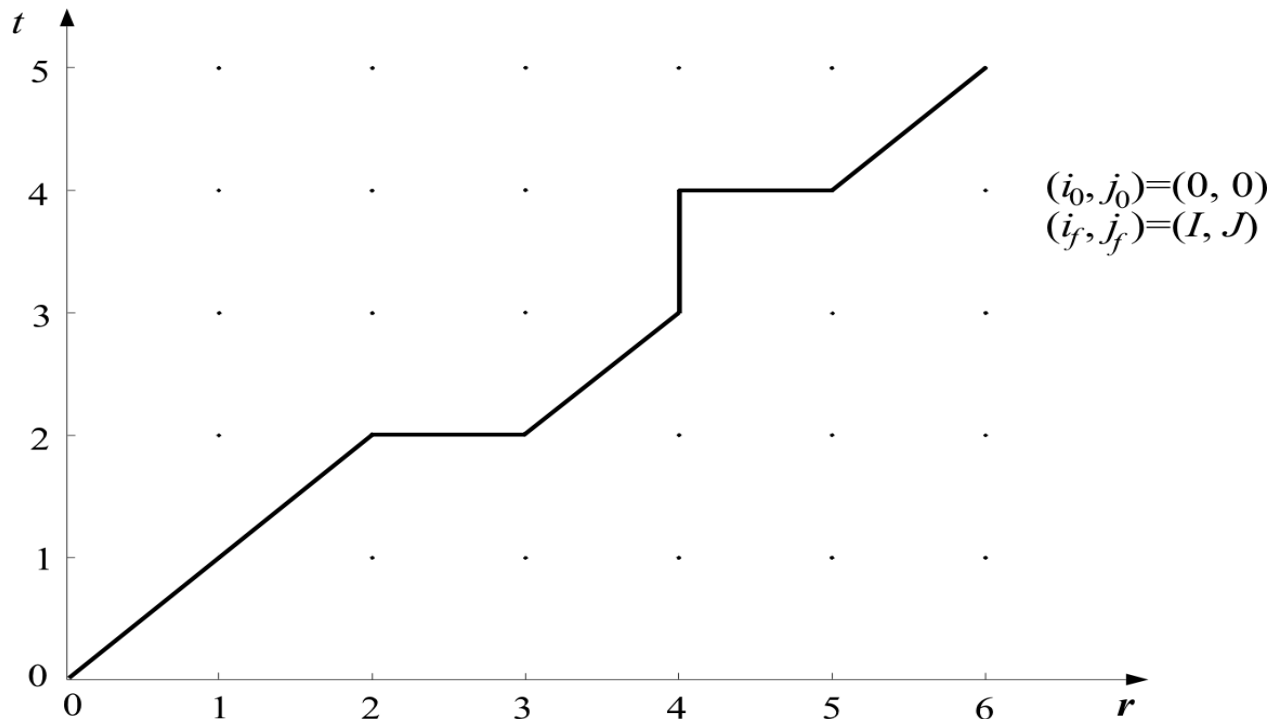


TM using Optimal Path Searching

- Each path is associated with a cost

$$D = \sum_{k=0}^{K-1} d(i_k, j_k)$$

where K is the number of nodes across the path



TM using Optimal Path Searching

- Let the cost up to node (i_k, j_k) be $D(i_k, j_k)$
- By convention
 - $D(0, 0)=0$
 - $d(0,0)=0$

TM using Optimal Path Searching

- The equation

$$D = \sum_{k=0}^{K-1} d(i_k, j_k)$$

assumes that each node has been associated with some cost

TM using Optimal Path Searching

- The equation

$$D = \sum_{k=0}^{K-1} d(i_k, j_k)$$

assumes that each node has been associated with some cost

- However, each transition (i_{k-1}, j_{k-1}) to (i_k, j_k) may also associate with a cost
- The new equation is:

$$D = \sum_k d(i_k, j_k | i_{k-1}, j_{k-1})$$

TM using Optimal Path Searching

$$D = \sum_k d(i_k, j_k | i_{k-1}, j_{k-1})$$

- Search for the path with the optimal cost D_{opt} .
- The matching cost between template \underline{r} and test pattern \underline{t} is D_{opt} .
- Costly operation
- Needs efficient computation

Bellman's Optimality Principle

- Optimal path:

$$(i_0, j_0) \xrightarrow{opt} (i_f, j_f)$$

Bellman's Optimality Principle

- Optimal path:

$$(i_0, j_0) \xrightarrow{opt} (i_f, j_f)$$

- Let (i, j) be an intermediate node, i.e.

$$(i_0, j_0) \rightarrow \dots \rightarrow (i, j) \rightarrow \dots \rightarrow (i_f, j_f)$$

Bellman's Optimality Principle

- Optimal path:

$$(i_0, j_0) \xrightarrow{opt} (i_f, j_f)$$

- Let (i, j) be an intermediate node, i.e.

$$(i_0, j_0) \rightarrow \dots \rightarrow (i, j) \rightarrow \dots \rightarrow (i_f, j_f)$$

Then, write the optimal path **through** (i, j)

$$(i_0, j_0) \xrightarrow[(i, j)]{opt} (i_f, j_f)$$

Bellman's Optimality Principle

- Bellman's Principle:

$(i_0, j_0) \xrightarrow{opt} (i_f, j_f)$ can be obtained as

$$(i_0, j_0) \xrightarrow{opt} (i, j) \oplus (i, j) \xrightarrow{opt} (i_f, j_f)$$

- meaning: The overall optimal path from (i_0, j_0) to (i_f, j_f) through (i, j) is the concatenation of the optimal paths from (i_0, j_0) to (i, j) and from (i, j) to (i_f, j_f)

Bellman's Optimality Principle

- Bellman's Principle:

$$(i_0, j_0) \xrightarrow{opt} (i_f, j_f) \Leftrightarrow (i_0, j_0) \xrightarrow{opt} (i, j) \oplus (i, j) \xrightarrow{opt} (i_f, j_f)$$

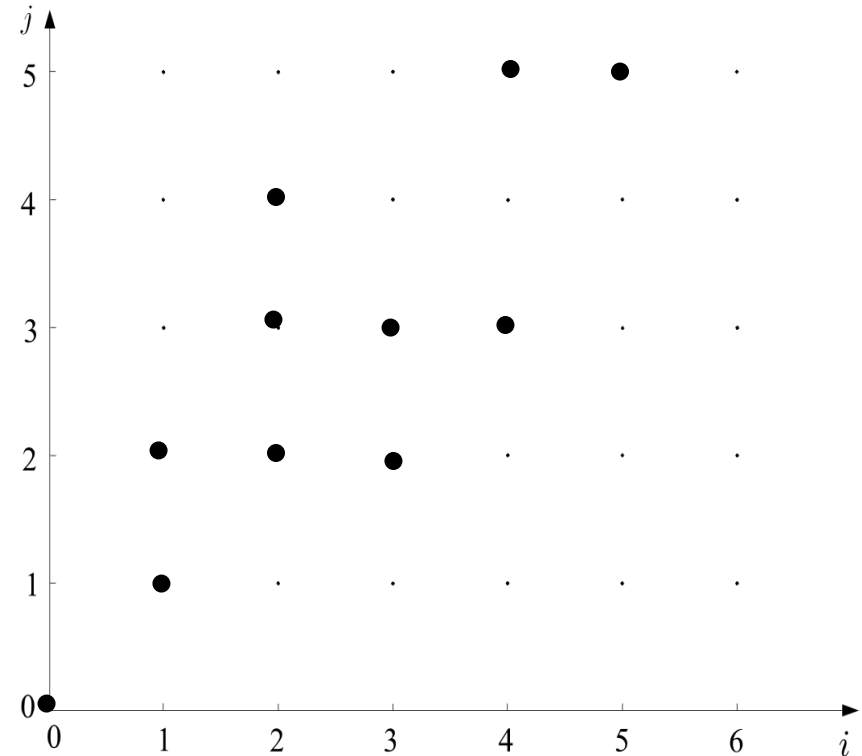
- Let $D_{opt.}(i_{k-1}, j_{k-1})$ is the optimal path to reach (i_{k-1}, j_{k-1}) from (i_0, j_0) , then Bellman's principle is stated as:

$$D_{opt}(i_k, j_k) = opt\{D_{opt}(i_{k-1}, j_{k-1}) + d(i_k, j_k | i_{k-1}, j_{k-1})\}$$

Bellman's Optimality Principle

$$D_{opt}(i_k, j_k) = \text{opt}\{D_{opt}(i_{k-1}, j_{k-1}) + d(i_k, j_k | i_{k-1}, j_{k-1})\}$$

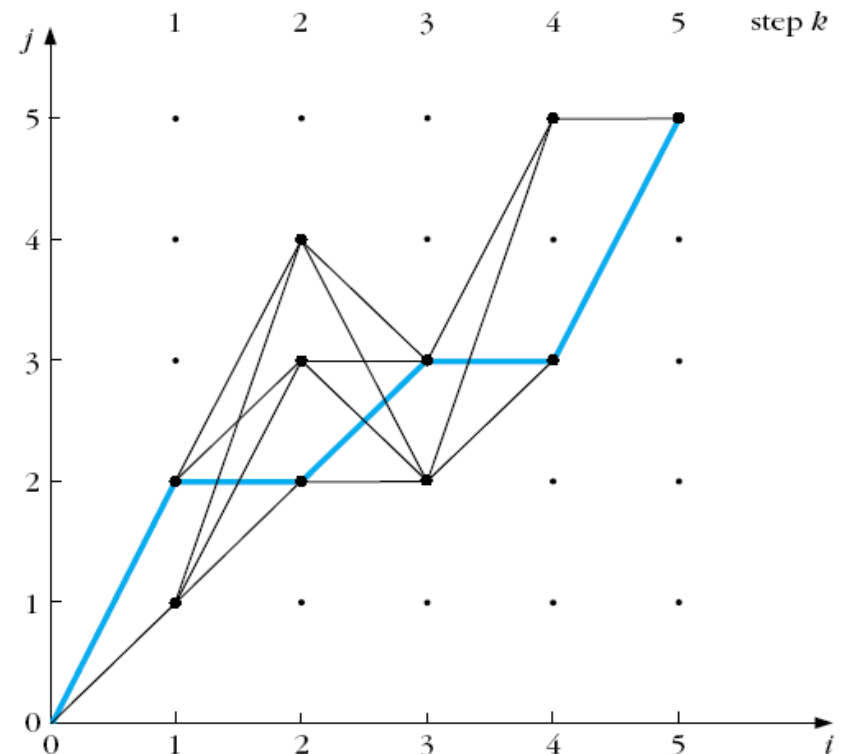
- We don't need to search the whole space to find the optimal path
- Global and local constraints may be imposed to reduce the search space



Bellman's Optimality Principle

$$D_{opt}(i_k, j_k) = \text{opt}\{D_{opt}(i_{k-1}, j_{k-1}) + d(i_k, j_k | i_{k-1}, j_{k-1})\}$$

- We don't need to search the whole space to find the optimal path
- Global and local constraints may be imposed to reduce the search space



Application of TM in Text Matching: The Edit Distance

- The Edit distance
 - It is used for matching written words.
- Applications:
- Automatic Editing
 - Text Retrieval

Application of TM in Text Matching: The Edit Distance

- The Edit distance
 - It is used for matching written words.
Applications:
 - Automatic Editing
 - Text Retrieval
 - The measure to be adopted for matching, must take into account:
 - **Wrongly identified** symbols
e.g. “befuty” instead of “beauty”
 - **Insertion errors**, e.g. “bearuty”
 - **Deletion errors**, e.g. “beuty”

The Edit Distance

- Edit distance: **Minimal** total number of **changes**, ***C***, **insertions** ***I*** and **deletions** ***R***, required to change pattern A into pattern B ,

$$D(A, B) = \min_j [C(j) + I(j) + R(j)]$$

where j runs over **All** possible variations of symbols, in order to convert $A \longrightarrow B$

The Edit Distance

- Edit distance: **Minimal** total number of **changes**, **C**, **insertions** **I** and **deletions** **R**, required to change pattern A into pattern B ,

$$D(A, B) = \min_j [C(j) + I(j) + R(j)]$$

where j runs over **All** possible variations of symbols, in order to convert $A \longrightarrow B$

- *Example*: many ways to change **beuty** to **beauty**

The Edit Distance

- The optimal path search algorithm can be used, provided we know
 - Initial conditions
 - Search space
 - Allowable transitions
 - Distance measure

The Edit Distance

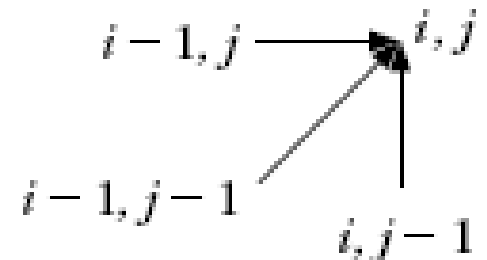
- Cost $D(0,0) = 0$,
- Complete path is searched
- Allowable predecessors and costs:

- $(i-1, j-1) \rightarrow (i, j)$

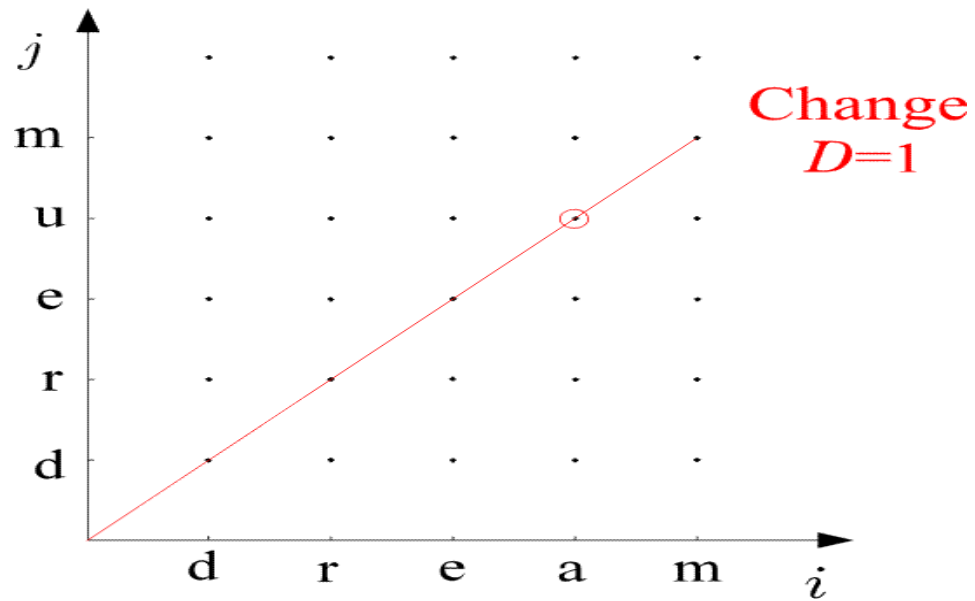
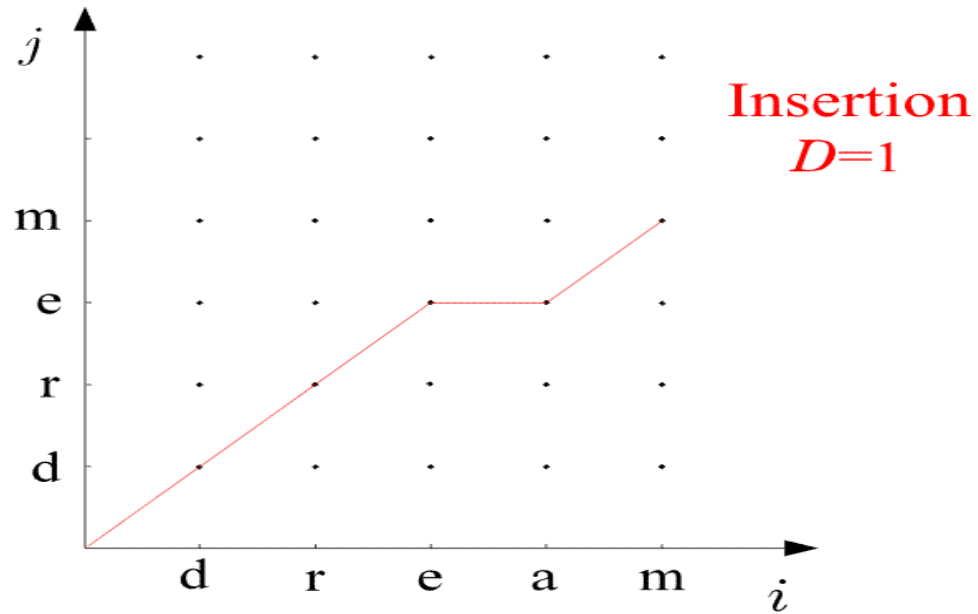
$$d(i, j | i-1, j-1) = \begin{cases} 0, & \text{if } t(i) = r(j) \\ 1, & t(i) \neq r(j) \end{cases}$$

- Horizontal $d(i, j | i-1, j) = 1$

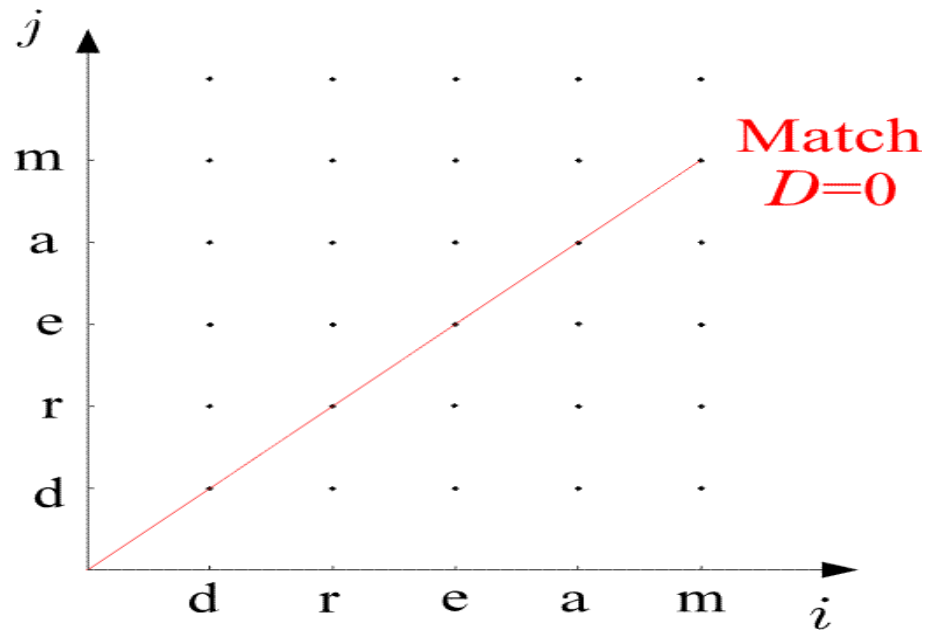
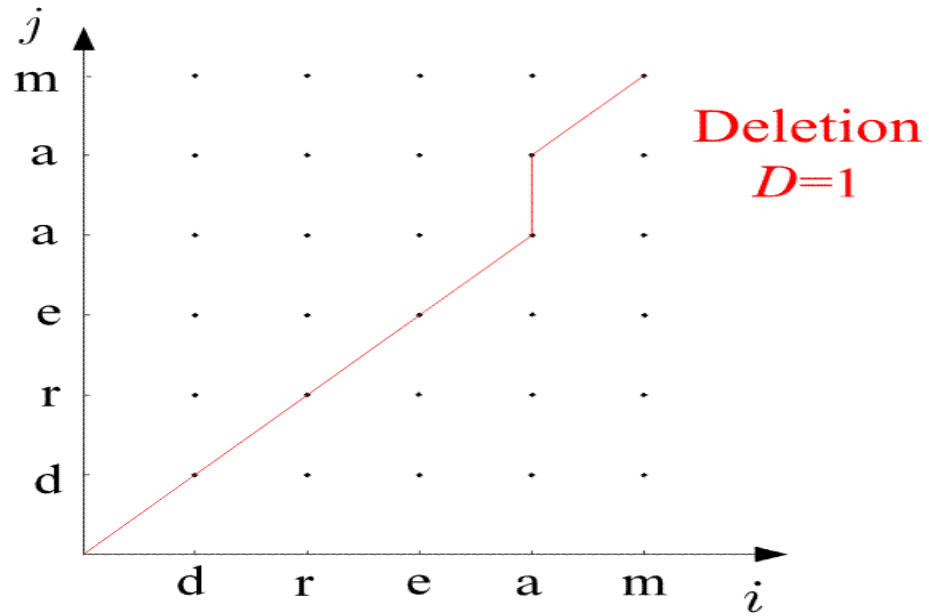
- Vertical $d(i, j | i, j-1) = 1$



- Examples:

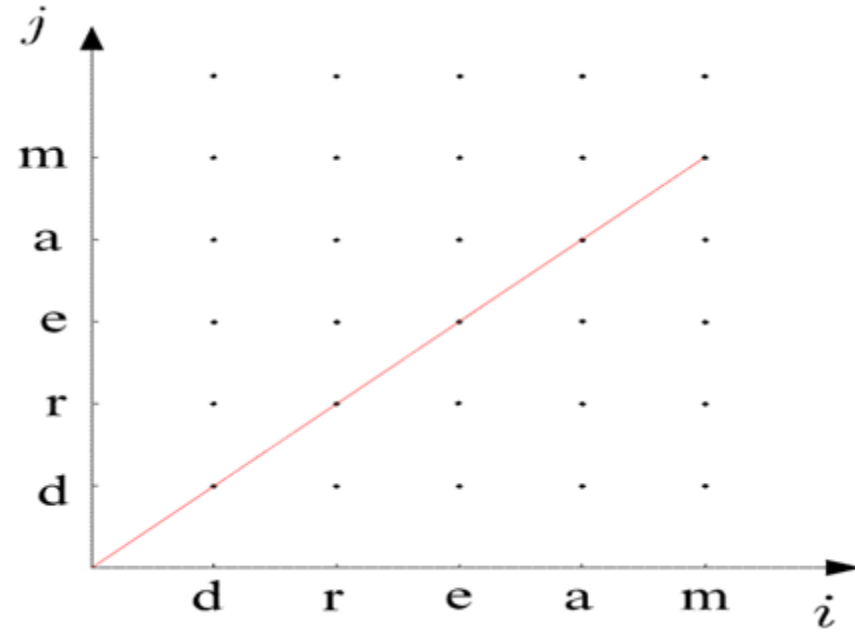


- Examples:



The Edit Distance

- The Algorithm
 - $D(0,0)=0$
 - For $i=1$, to I
 - $D(i,0)=D(i-1,0)+1$
 - END {FOR}
 - For $j=1$ to J
 - $D(0,j)=D(0,j-1)+1$
 - END{FOR}
 - For $i=1$ to I
 - For $j=1$, to J
 - $C_1=D(i-1,j-1)+d(i,j \mid i-1,j-1)$
 - $C_2=D(i-1,j)+1$
 - $C_3=D(i,j-1)+1$
 - $D(i,j)=\min (C_1,C_2,C_3)$
 - END {FOR}
 - END {FOR}
 - $D(A,B)=D(I,J)$



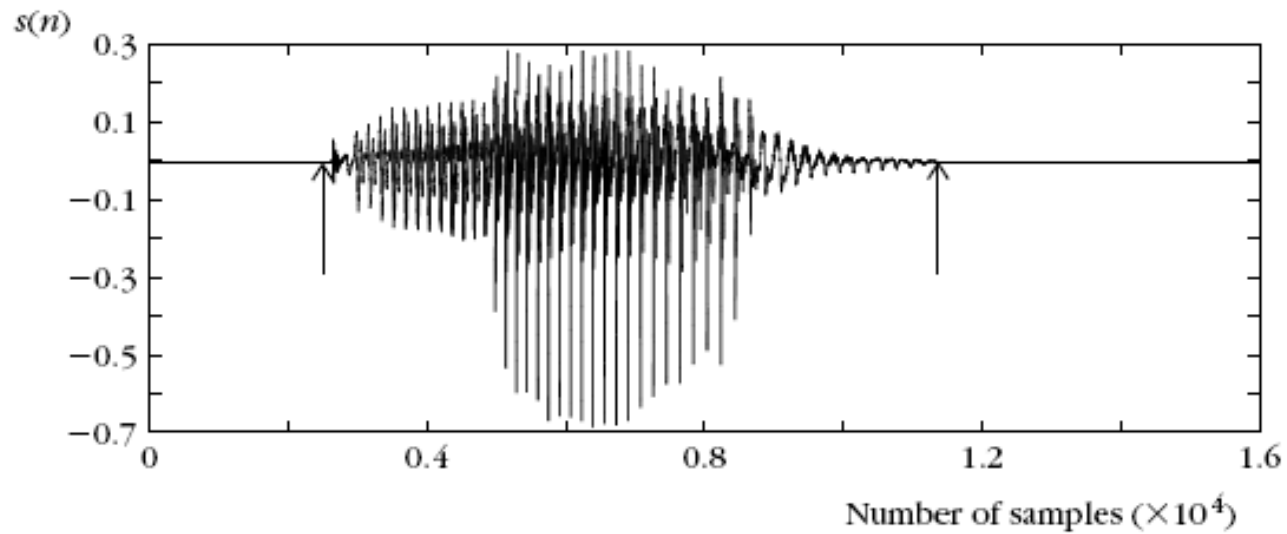
Application of TM in Speech Recognition

- A number of variations
 - Speaker Independent Speech Recognition
 - Speaker Dependent Speech Recognition
 - Continuous Speech Recognition
 - Isolated word recognition (IWR)

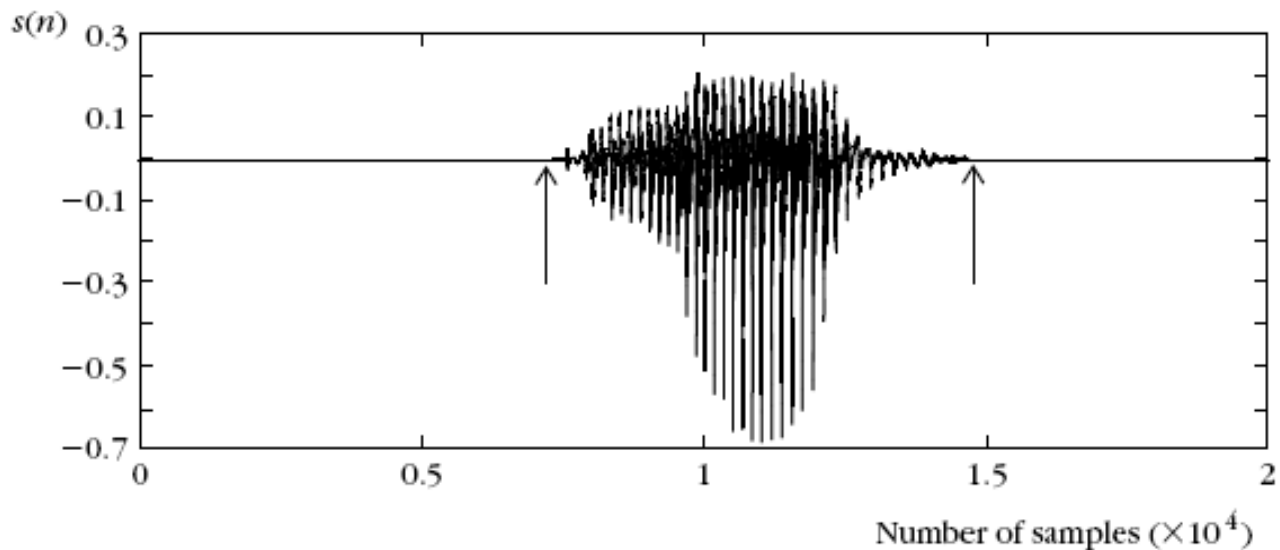
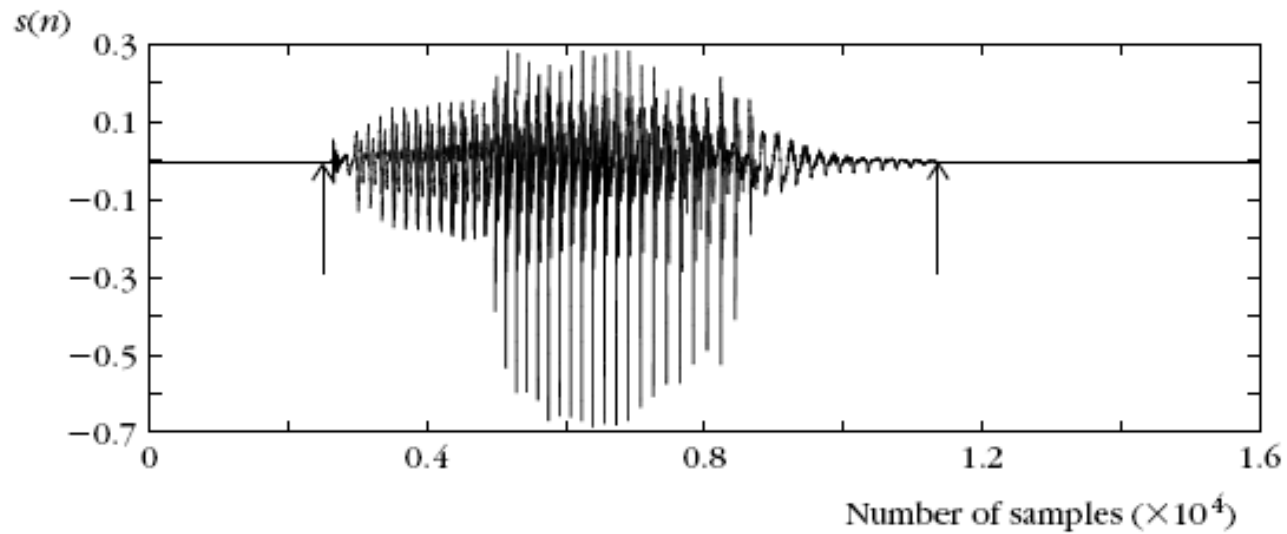
Application of TM in IWR

- The goal:
 - Given a number of known spoken words in a data base (reference patterns)
 - find the best match of an unknown spoken word (test pattern).
- Procedure:
 - compare the test word against reference words

Application of TM in IWR

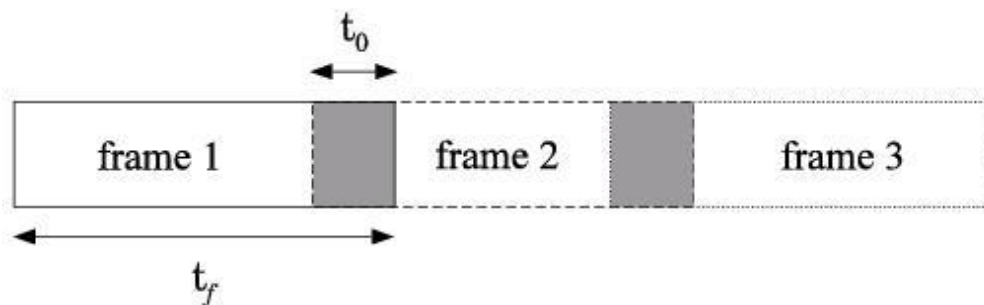


Application of TM in IWR



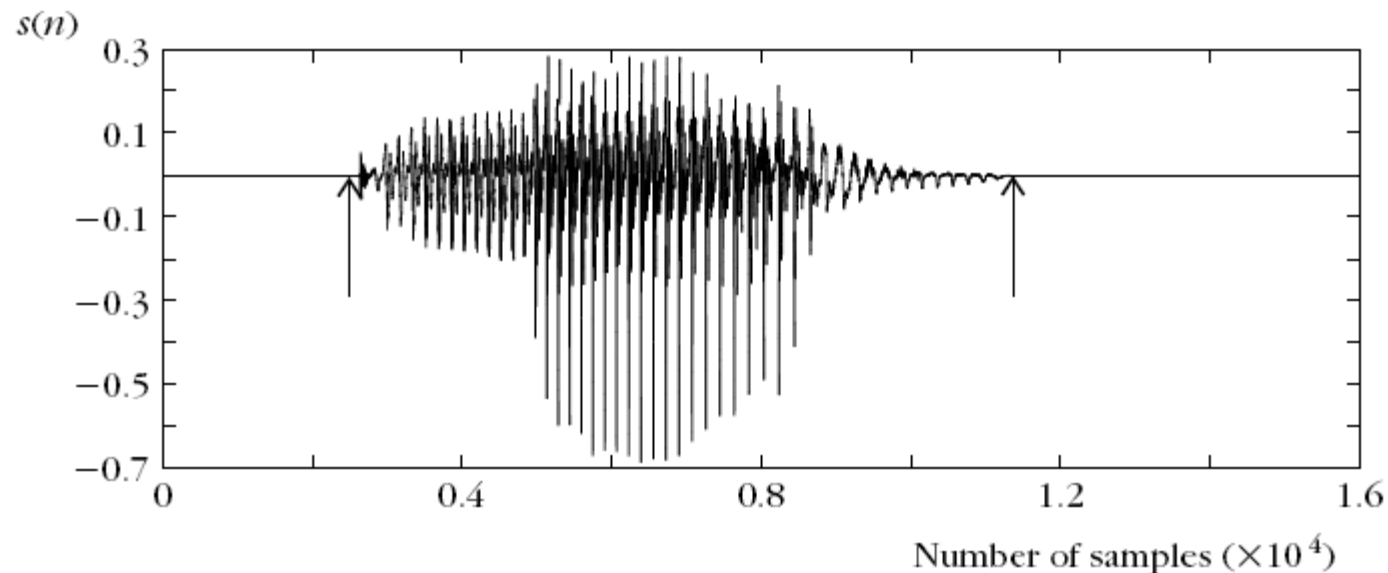
Application of TM in IWR

- The procedure:
 - Express the test and each of the reference patterns as sequences of feature vectors $\underline{r}(i)$, $\underline{t}(j)$.
 - To this end, divide each of the speech segments in a number of successive frames.



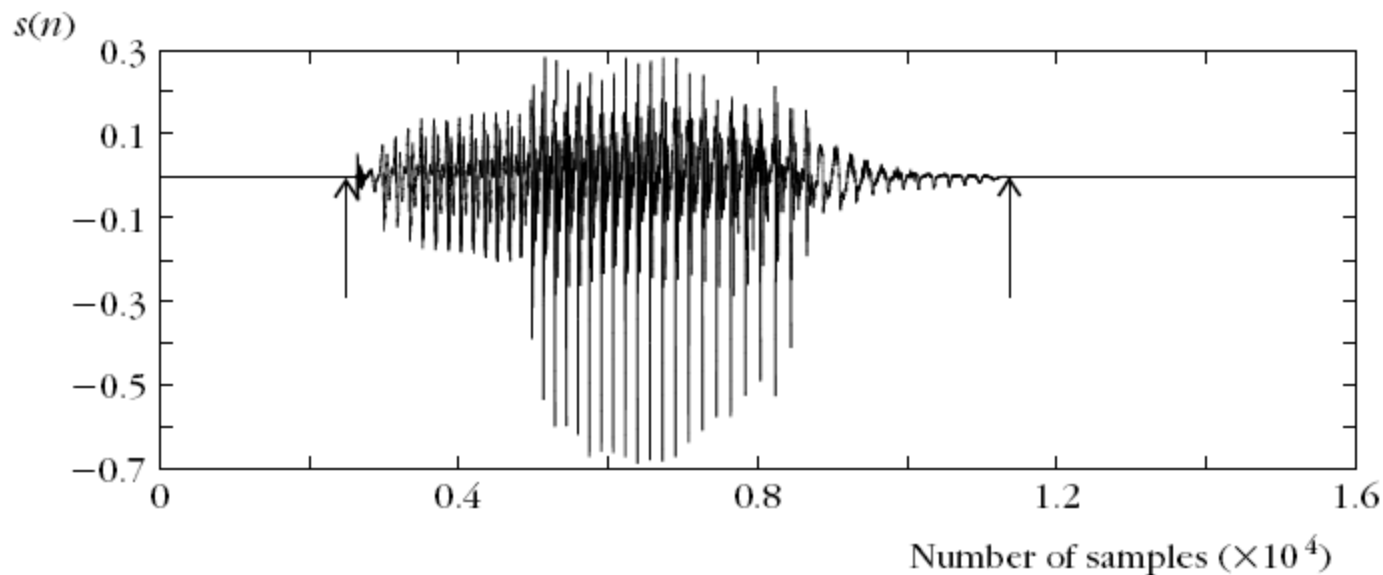
Application of TM in IWR

- The procedure:
 - Sample a speech segment from a microphone:



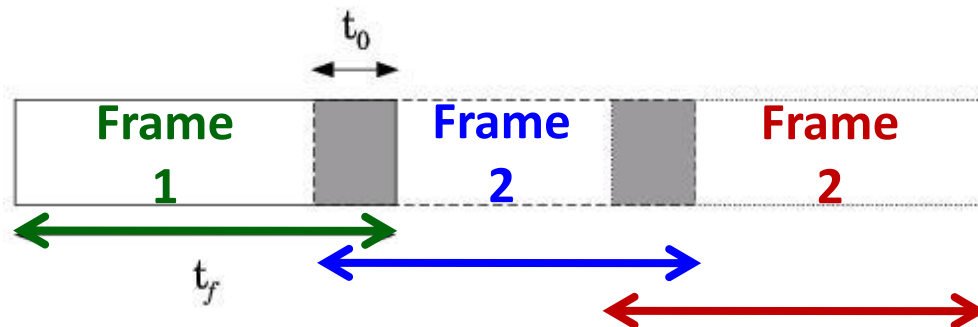
Application of TM in IWR

- The procedure:



$$t_f = 512$$

$$t_0 = 100$$



- each frame is represented by a vector of 512 samples

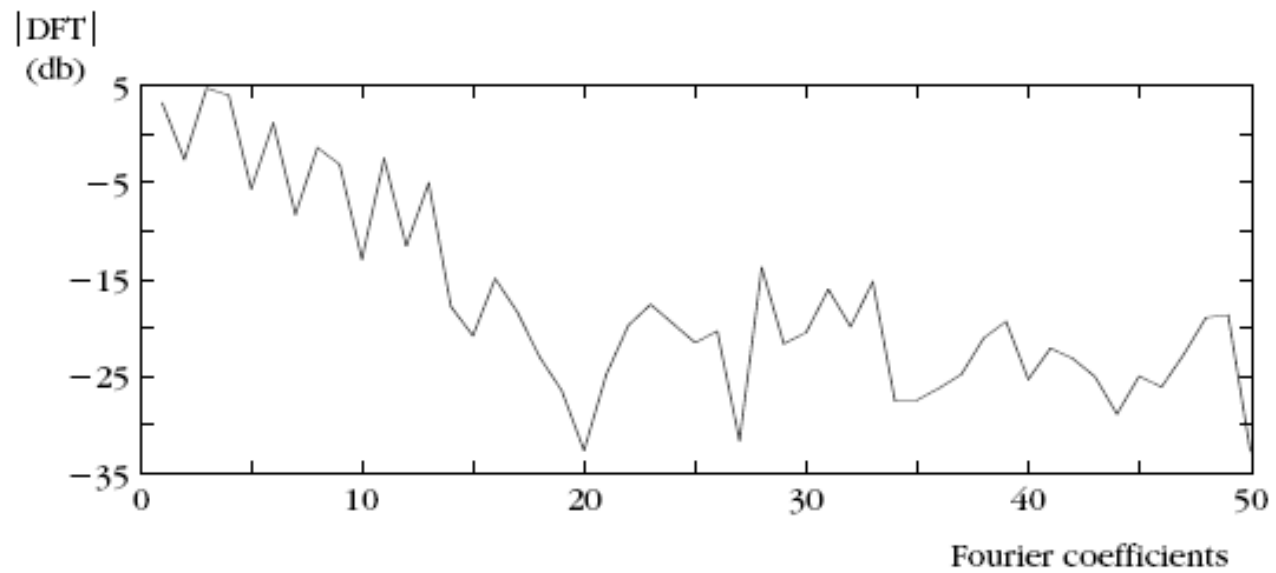
$$\underline{r}(i) = \begin{bmatrix} x_i(0) \\ x_i(1) \\ \dots \\ \dots \\ x_i(512) \end{bmatrix}, \quad i = 1, \dots, I \quad \underline{t}(j) = \begin{bmatrix} x_j(0) \\ x_j(1) \\ \dots \\ \dots \\ x_j(512) \end{bmatrix}, \quad j = 1, \dots, J$$

- convert them to DFT

$$DFT(\underline{r}(i)) = DFT\left(\begin{bmatrix} x_i(0) \\ x_i(1) \\ \dots \\ \dots \\ x_i(512) \end{bmatrix}\right) = \begin{bmatrix} X_i(0) \\ X_i(1) \\ \dots \\ \dots \\ X_i(512) \end{bmatrix}$$

$$DFT(\underline{t}(j)) = DFT\left(\begin{bmatrix} x_i(0) \\ x_i(1) \\ \dots \\ \dots \\ x_i(512) \end{bmatrix}\right) = \begin{bmatrix} X_i(0) \\ X_i(1) \\ \dots \\ \dots \\ X_i(512) \end{bmatrix}$$

- convert them to DFT



- For each frame compute a feature vector. For example, the DFT coefficients and use, say, ℓ of those:

$$\underline{r}(i) = \begin{bmatrix} X_i(0) \\ X_i(1) \\ \dots \\ \dots \\ X_i(\ell-1) \end{bmatrix}, \quad i = 1, \dots, I \quad \underline{t}(j) = \begin{bmatrix} X_j(0) \\ X_j(1) \\ \dots \\ \dots \\ X_j(\ell-1) \end{bmatrix}, \quad j = 1, \dots, J$$

- For each frame compute a feature vector. For example, the DFT coefficients and use, say, ℓ of those:

$$\underline{r}(i) = \begin{bmatrix} X_i(0) \\ X_i(1) \\ \dots \\ X_i(\ell-1) \end{bmatrix}, \quad i = 1, \dots, I \quad \underline{t}(j) = \begin{bmatrix} X_j(0) \\ X_j(1) \\ \dots \\ X_j(\ell-1) \end{bmatrix}, \quad j = 1, \dots, J$$

- Choose a cost function associated with each node across a path, e.g., the Euclidean distance

$$\|\underline{r}(i_k) - \underline{t}(j_k)\| = d(i_k, j_k)$$

- For each frame compute a feature vector. For example, the DFT coefficients and use, say, ℓ of those:

$$\underline{r}(i) = \begin{bmatrix} X_i(0) \\ X_i(1) \\ \dots \\ X_i(\ell-1) \end{bmatrix}, \quad i = 1, \dots, I \quad \underline{t}(j) = \begin{bmatrix} X_j(0) \\ X_j(1) \\ \dots \\ X_j(\ell-1) \end{bmatrix}, \quad j = 1, \dots, J$$

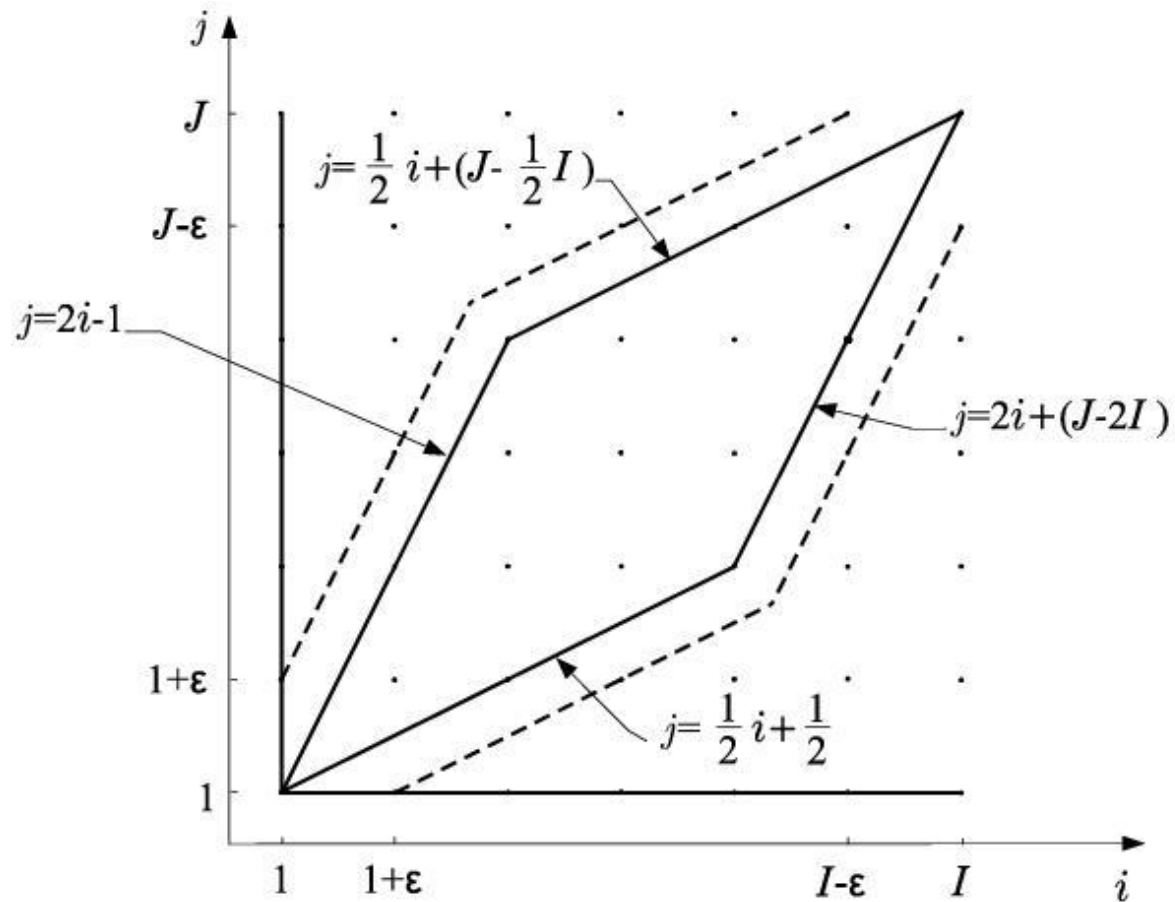
- Choose a cost function associated with each node across a path, e.g., the Euclidean distance

$$\|\underline{r}(i_k) - \underline{t}(j_k)\| = d(i_k, j_k)$$

- find the optimal path in the grid
- Match the test pattern to the reference pattern associated with the optimal path

- Prior to performing the math one has to choose:
 - end point constraints
 - global constraints
 - local constraints
 - distance

- Prior to performing the math one has to choose:
 - **The global constraints:** Defining the region of space within which the search for the optimal path will be performed.



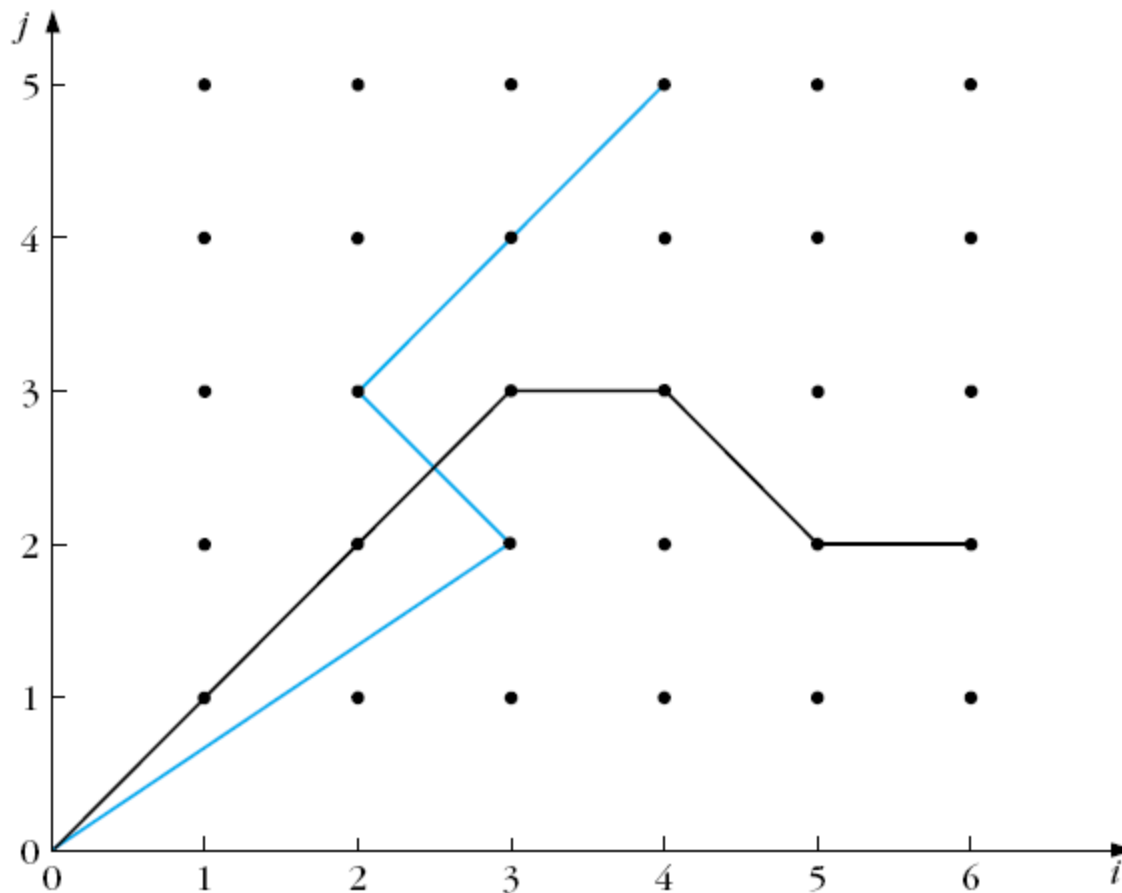
- The local constraints: monotonic path

$$i_{k-1} \leq i_k \quad \text{and} \quad j_{k-1} \leq j_k$$

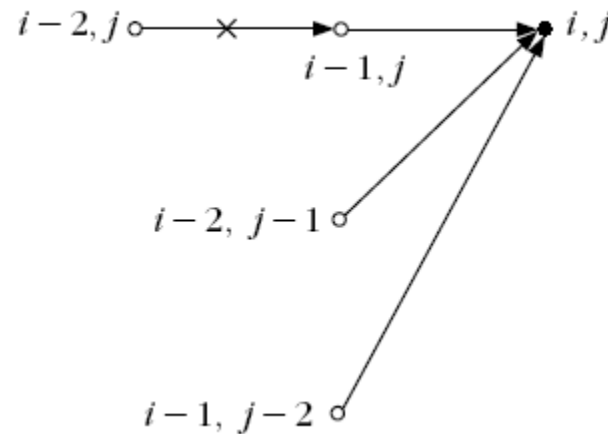
- The local constraints: monotonic path

$$i_{k-1} \leq i_k \quad \text{and} \quad j_{k-1} \leq j_k$$

- Non-monotonic path

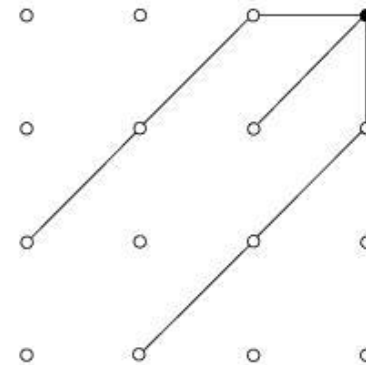
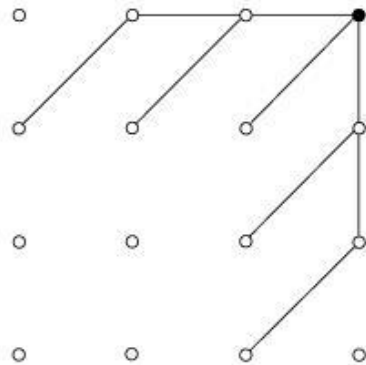
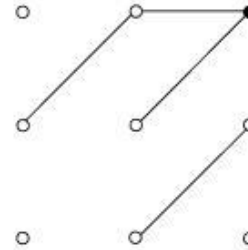
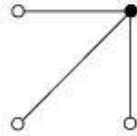


- **The local constraints:** Defining the type of transitions allowed between the nodes of the grid.



Itakura local constraints

- **The local constraints:** Defining the type of transitions allowed between the nodes of the grid.



Sakoe and Chiba local
constraints

- cost function:
 - Euclidean distance
 - only node distance

$$d(i_k, j_k \mid i_{k-1}, j_{k-1}) = d(i_k, j_k)$$

$$= \|\underline{r}(i_k) - \underline{t}(j_k)\|$$