

CSE 473: Pattern Recognition

Syntactic Pattern Recognition

- Objective:
 - Given a pattern description x,

determine:
$$x \in L(G_i)$$

SyntPR: A Simple Approach

- Training
 - For each class,
 - generate the grammar G_i
 - generate the Language $L(G_i)$
 - store the Language $L(G_i)$ in a database
- Classification
 - match the object pattern, x, with the patterns in each $L(G_i)$

Limitations in the Simple Approach

- $L(G_i)$ is often infinite
- Needs large storage
- Matching also computationally inefficient

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Alternate solution: Recognition by Parsing

Recognition by Parsing

- What we do in parsing
 - Training: Design a grammar for each class
 - Classification: Determine whether the object pattern, x, is syntactically well formed under a grammar

An Example of Parsing

```
• Grammar: G = (V_T, V_N, P, S)
• V_T = (program, computer, crashes, the)
• V_N = (SENTENCE, ADJ, VP, NP, NOUN, VERB)
• P = \{
       SENTENCE \rightarrow NP + VP
       NP \rightarrow ADJ + NOUN
       VP \rightarrow VERB + NP
      NOUN -> computer | program
       VERB -> crashes
      ADJ \rightarrow the
```

An Example of Parsing

Determine:

whether the sentence "the program crashes the computer" is valid under grammar G

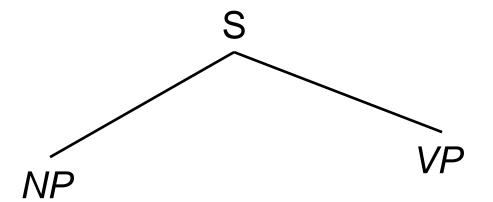
An Example of Parsing

Determine:

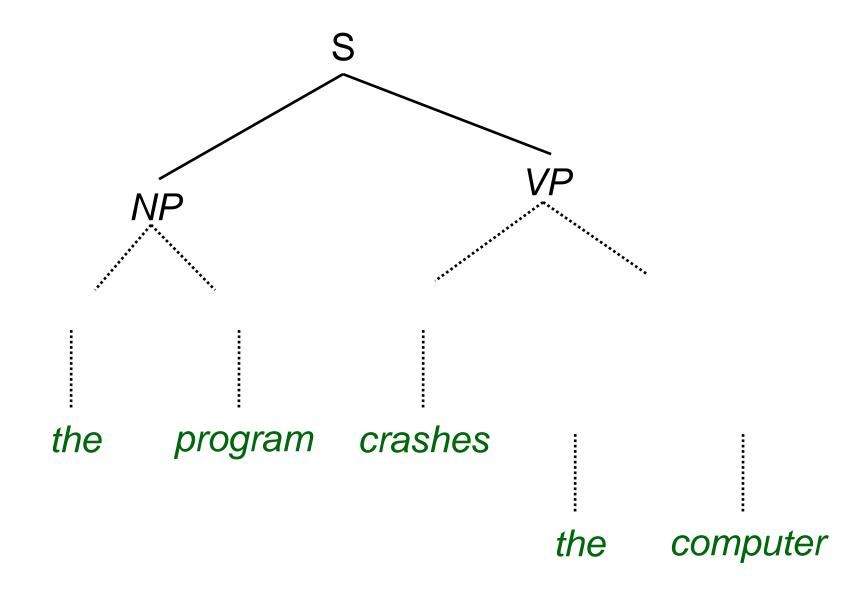
whether the sentence "the program crashes the computer" is valid under grammar G

We will use a derivation tree

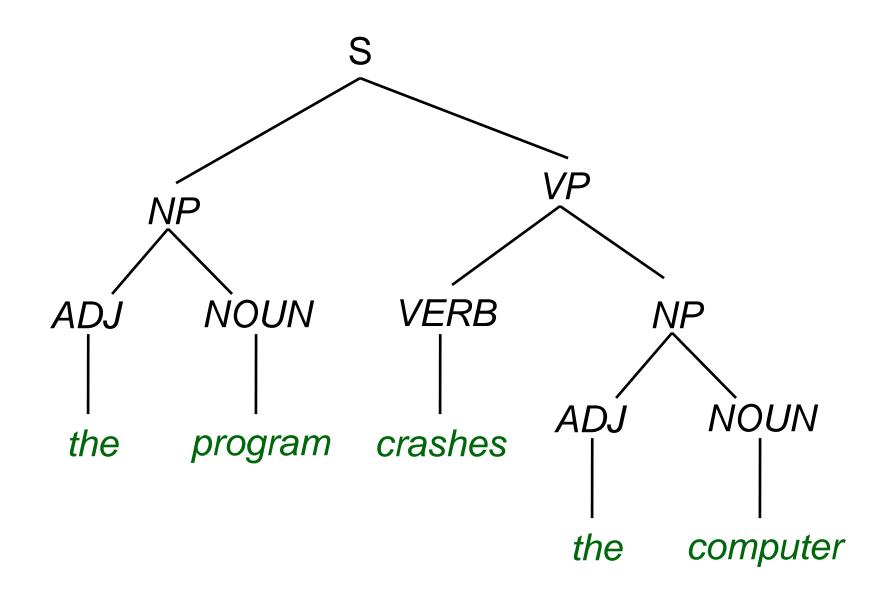
The Derivation Tree



The Derivation Tree

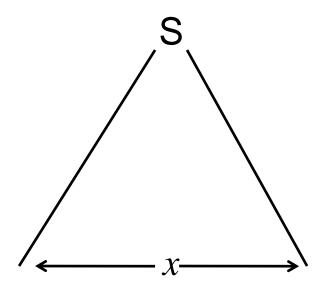


The Derivation Tree



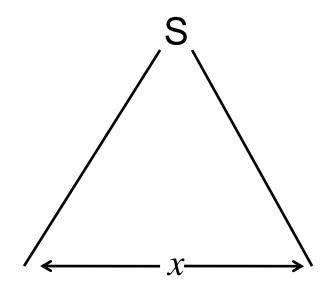
The Parsing Problem

• Given a grammar $G = (V_T, V_N, P, S)$ and a sentence x, form a triangle like this:



The Parsing Problem

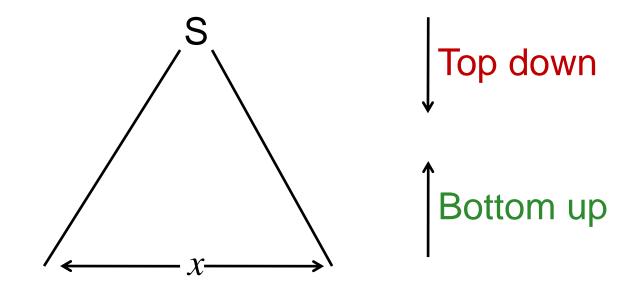
• Given a grammar $G = (V_T, V_N, P, S)$ and a sentence x, form a triangle like this:



 The parsing is the process of filling of interior of the triangle

The Parsing Problem

Use either top down or bottom up approach



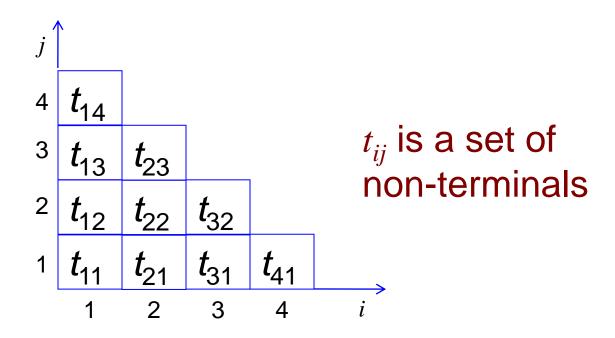
- Top down: proceed from S towards terminals
- Bottom up: proceed from terminals towards S

Cocke-Younger-Kasami (CYK) Parsing Algorithm

- A bottom up approach
- Complexity is $O(|x|^3)$

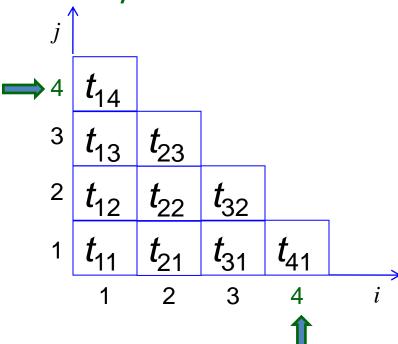
- Requirements:
 - The productions rules must be in CNF

• Given a grammar G and a string $x = x_1, x_2, \ldots, x_n$, form a table like this:

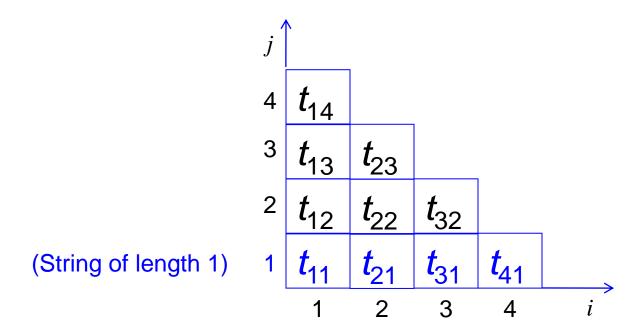


• Given a grammar G and a string $x = x_1, x_2, \ldots, x_n$, form a table like this:

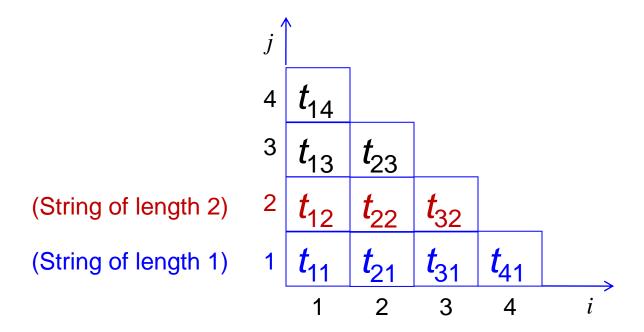
• The number of rows/columns = n



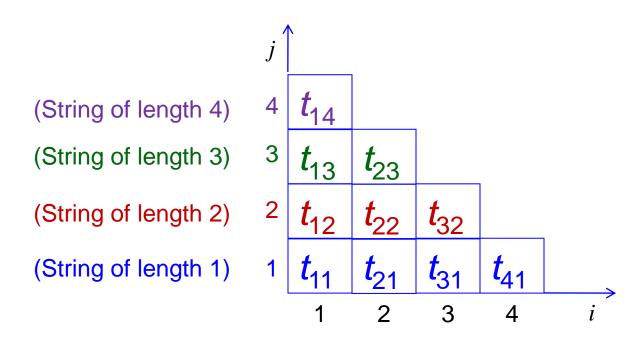
String covering



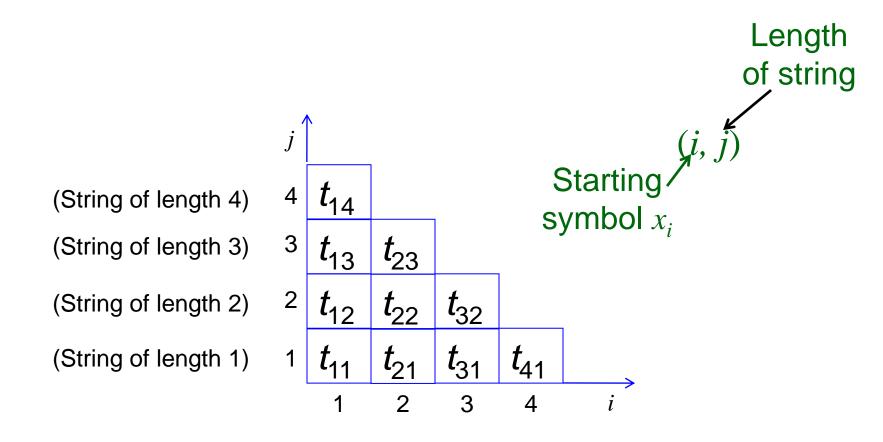
String covering



String covering

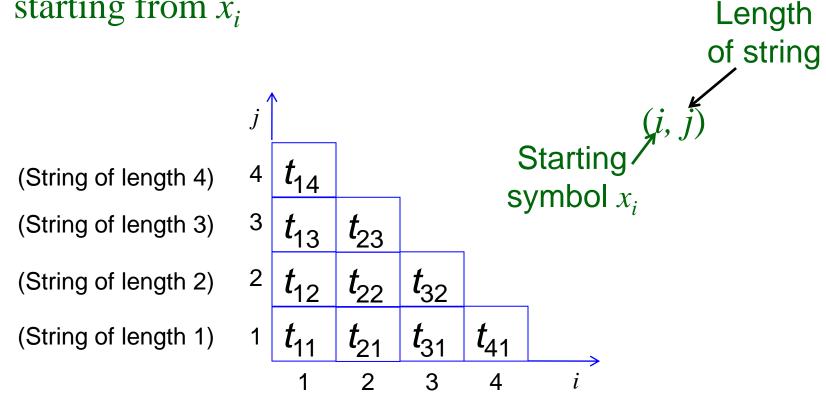


• String covering by entry t_{ij}

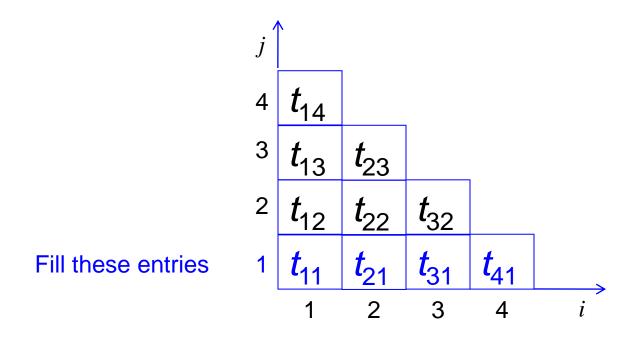


• String covering by entry t_{ij}

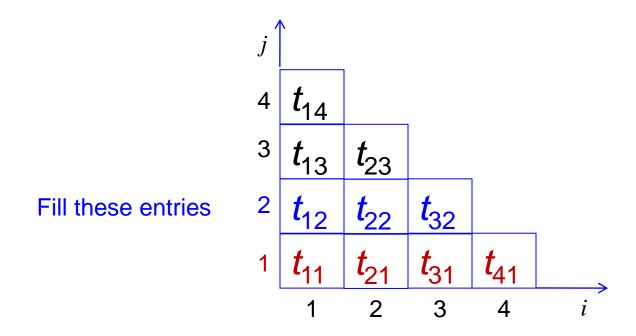
• A non-terminal in t_{ij} covers string of length j starting from x_i



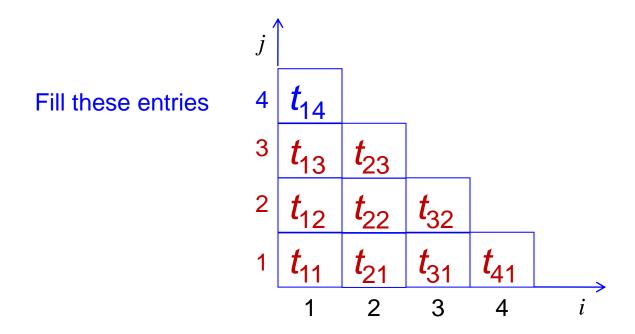
Filling CYK table



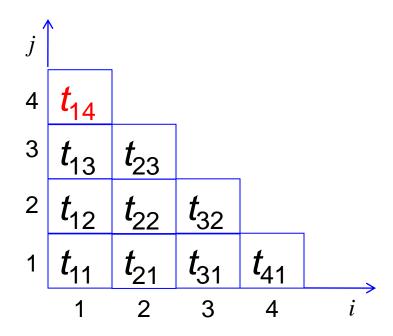
Filling CYK table



Filling CYK table



- With a grammar *G* and a string $x = x_1, x_2, ..., x_n$
- If t_{1n} (t_{14} , in this example) contains S, this means, x is defined in G



Example:

Let a grammar with following productions

$$S \rightarrow AB \mid BB$$
, $A \rightarrow CC \mid AB \mid a$

$$B \rightarrow BB \mid CA \mid b$$
, $C \rightarrow BA \mid AA \mid b$

Example:

Let a grammar with following productions

$$S \rightarrow AB \mid BB, \quad A \rightarrow CC \mid AB \mid a$$

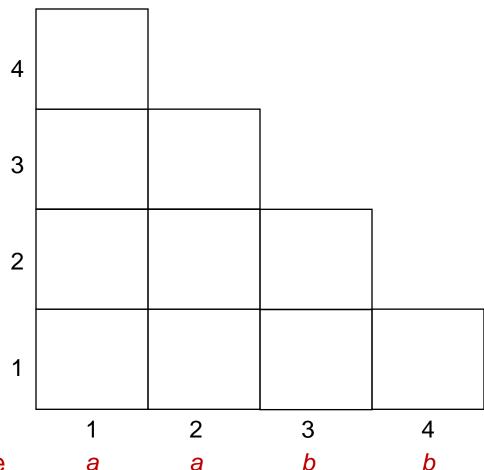
 $B \rightarrow BB \mid CA \mid b, \quad C \rightarrow BA \mid AA \mid b$

now parse, x = aabb

$$A \rightarrow CC \mid AB \mid a$$

$$B \rightarrow BB |CA| b$$
,

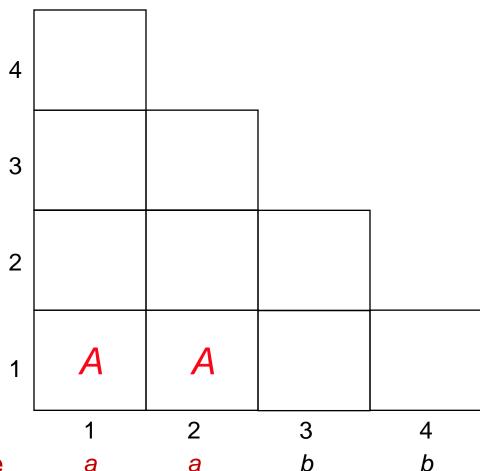
$$C \rightarrow BA \mid AA \mid b$$



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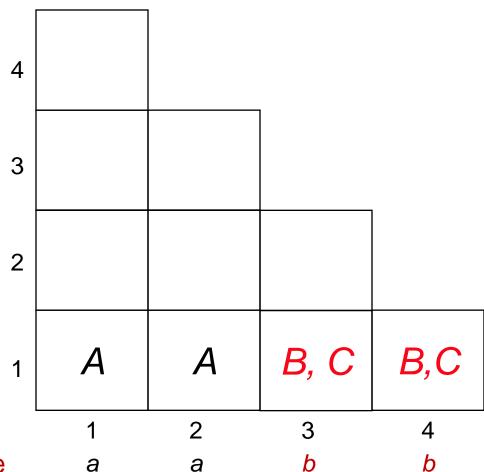
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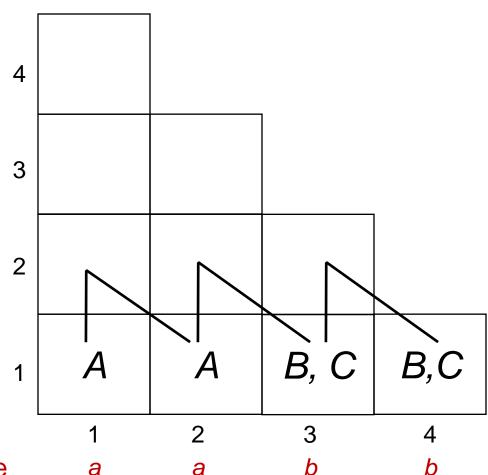
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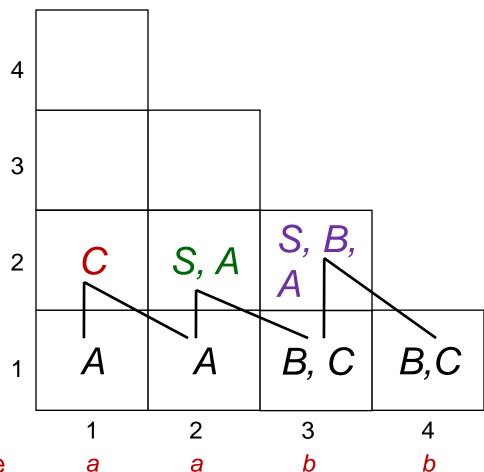
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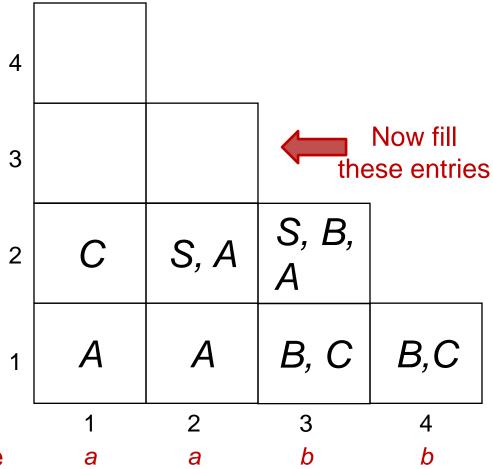
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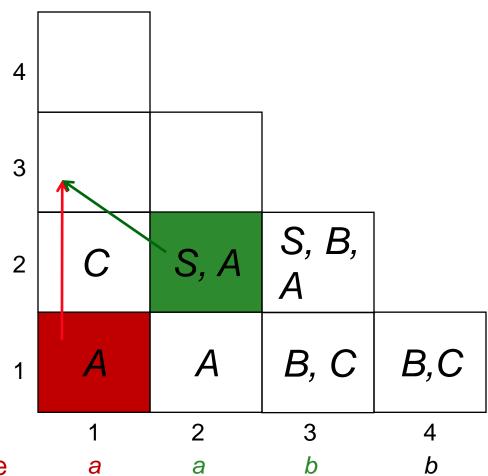
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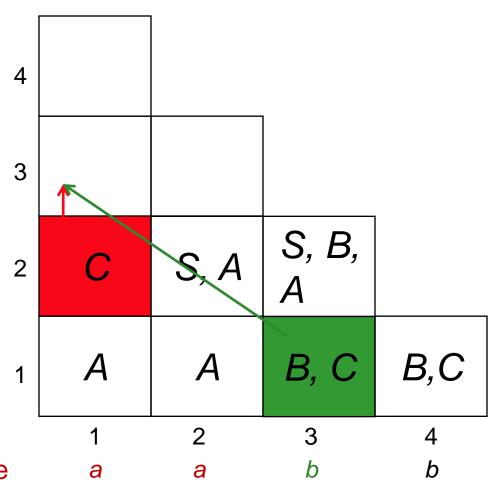
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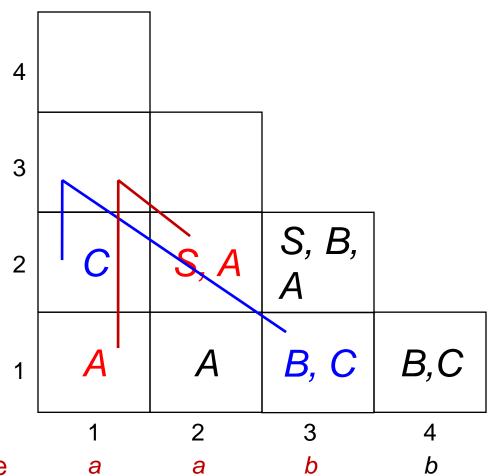
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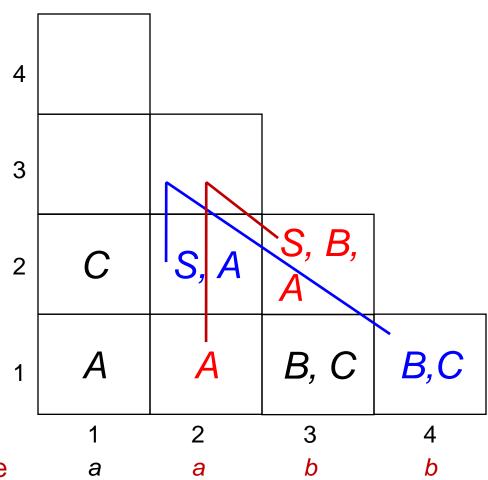
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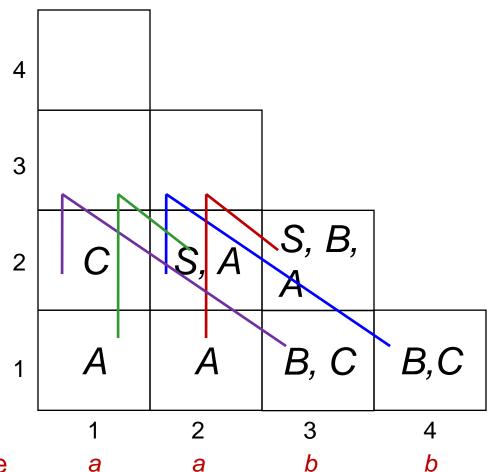
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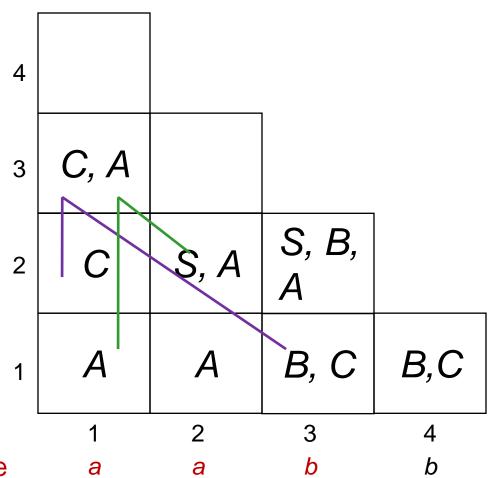
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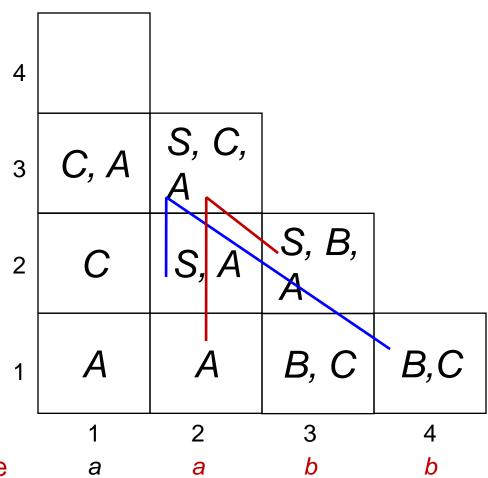
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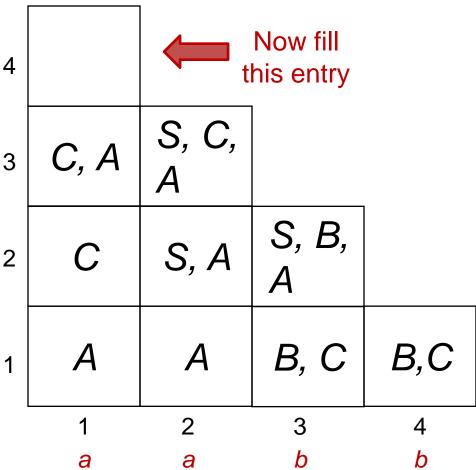
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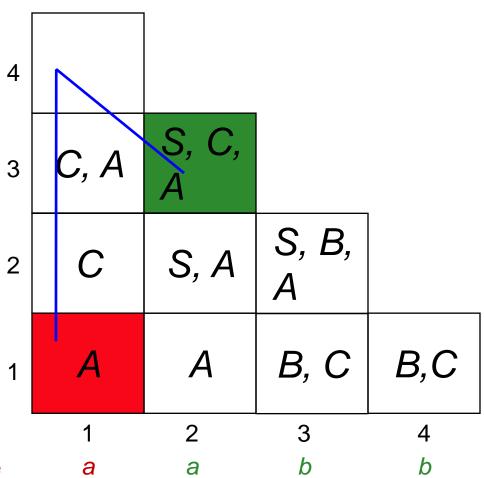
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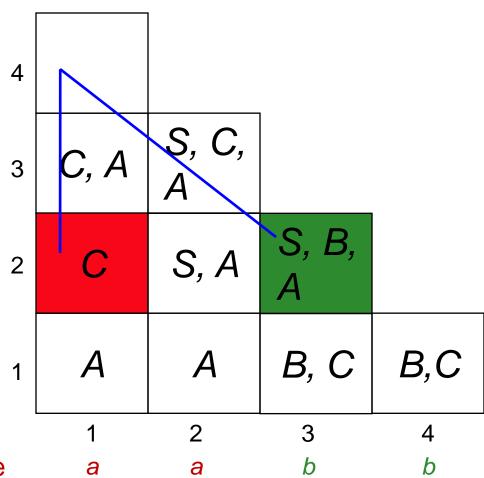
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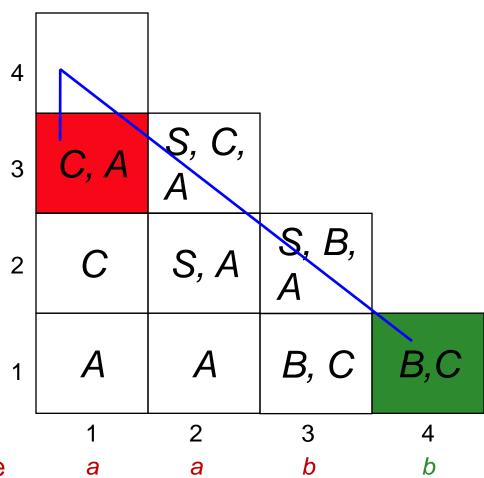
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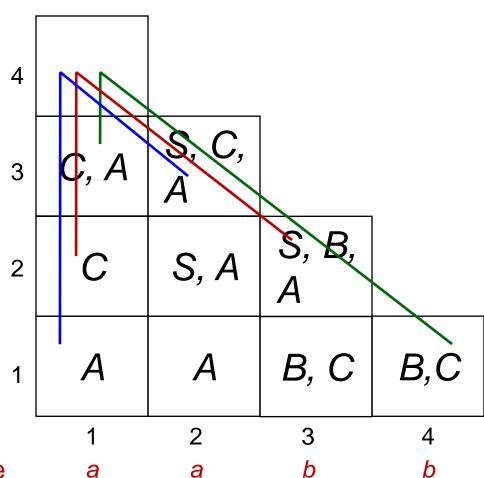
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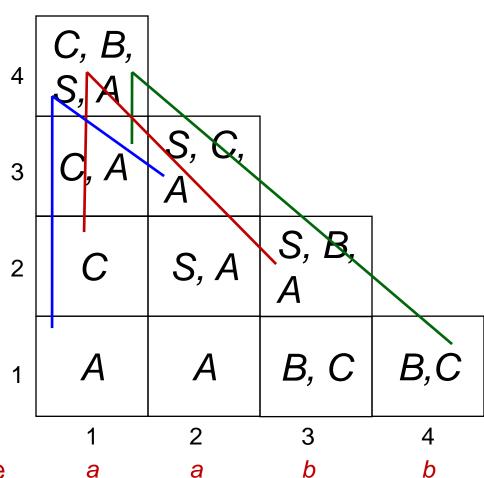
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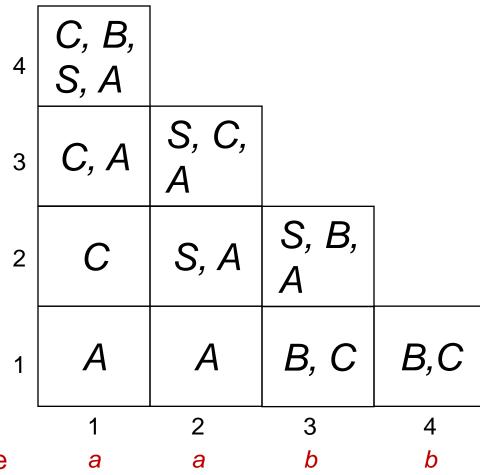
$$S \rightarrow AB \mid BB$$

$$A \rightarrow CC \mid AB \mid a$$

$$B \rightarrow BB | CA | b$$
,

$$C \rightarrow BA \mid AA \mid b$$

The string *aabb* is defined in *G*



What happens if the productions are not in CNF, e.g.,

$$S \rightarrow bA_1 \mid cA_2$$

 $A_1 \rightarrow bA_2$
 $A_2 \rightarrow b \mid aA_2$

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$$S \rightarrow bA_1 \mid cA_2$$

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We can convert them into CNF

What happens if the productions are not in CNF, e.g.,

$$S \to bA_1 \mid cA_2$$

 $A_1 \to bA_2$

$$A_2 \rightarrow b \mid aA_2$$

Equivalent CNF

$$S - > A_3 A_1$$

 $A_3 - > b$

What happens if the productions are not in CNF, e.g.,

$$S \rightarrow bA_1 \mid cA_2$$

 $A_1 \rightarrow bA_2$
 $A_2 \rightarrow b \mid aA_2$

Equivalent CNF

$$S - > A_3 A_1$$

 $A_3 - > b$
 $S - > A_4 A_2$
 $A_4 - > c$

What happens if the productions are not in CNF, e.g.,

$$S \rightarrow bA_1 \mid cA_2$$

 $A_1 \rightarrow bA_2$
 $A_2 \rightarrow b \mid aA_2$

Equivalent CNF

$$S - > A_3 A_1$$

 $A_3 - > b$
 $S - > A_4 A_2$
 $A_4 - > c$
 $A_1 - > A_3 A_2$
 $A_2 - > b$
 $A_2 - > b$
 $A_5 - > a$

Example 2: find whether bbaab is defined in G

$$S - > A_3 A_1$$

 $A_3 - > b$
 $S - > A_4 A_2$
 $A_4 - > c$
 $A_1 - > A_3 A_2$
 $A_2 - > b$
 $A_2 - > b$
 $A_5 - > a$

• Example 2: find whether *bbaab* is defined in *G*

$S - > A_3 A_1$
$A_3 - > b$
$S - > A_4 A_2$
$A_4 - > c$
$A_1 - > A_3 A_2$
$A_2 - > b$
$A_2 - > A_5 A_2$
$A_5 - > a$

5	S				
4	ı	A_1			
3	-	-	A_2		
2	A_1	-	-	A_2	
1	A_2,A_3	A_2,A_3	A_5	A_5	A_2,A_3
	b	b	а	а	b

Example 2: find whether cbab is defined in G

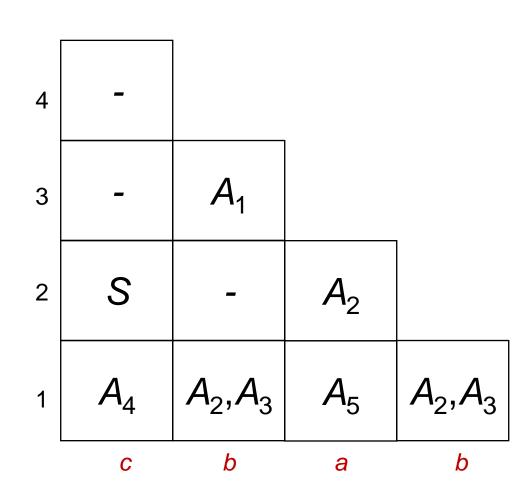
$$S - > A_3 A_1$$

 $A_3 - > b$
 $S - > A_4 A_2$
 $A_4 - > c$
 $A_1 - > A_3 A_2$
 $A_2 - > b$
 $A_2 - > b$
 $A_5 - > a$

Example 2: find whether cbab is defined in G

$$S - > A_3 A_1$$

 $A_3 - > b$
 $S - > A_4 A_2$
 $A_4 - > c$
 $A_1 - > A_3 A_2$
 $A_2 - > b$
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 $A_5 - > a$

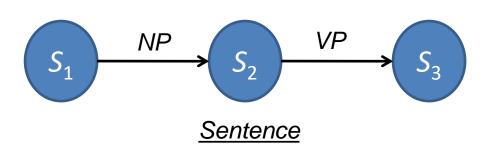


Parsing using Transition Network (TN)

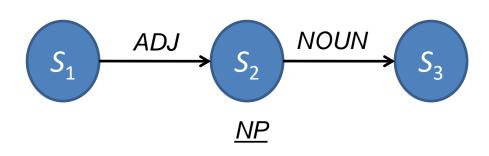
- Consists of nodes and arc
- Nodes represent states
- Arcs are labeled with terminal or non-terminal

An Example of Transition Network (TN)

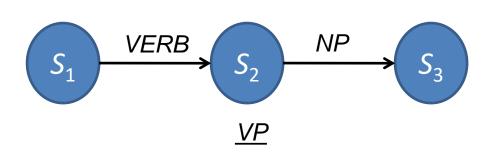
SENTENCE -> NP + VP



NP -> ADJ + NOUN

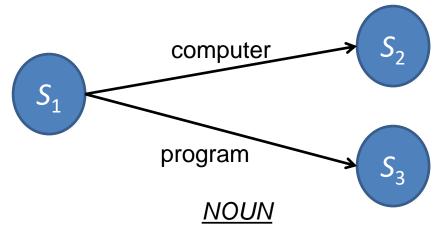


VP -> VERB + NP

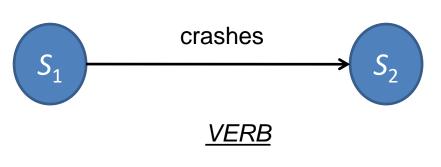


An Example of Transition Network (TN)

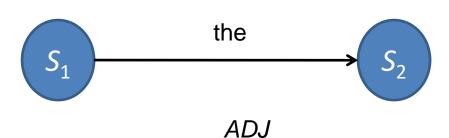
NOUN -> computer | program



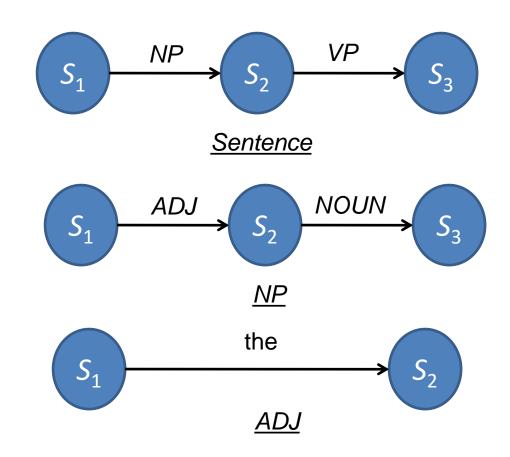
VERB -> crashes



ADJ -> the

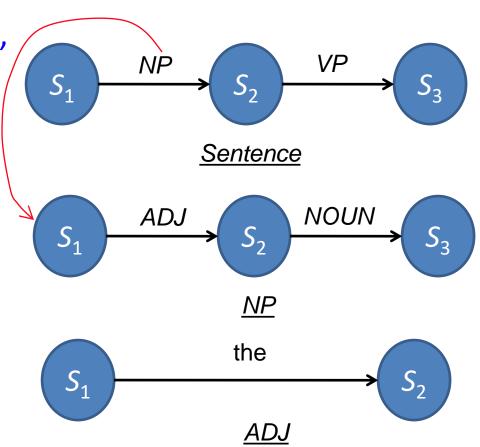


- Starts at an initial node
- Sequentially checks input string against the arc label



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- Sequentially checks input string against the arc label

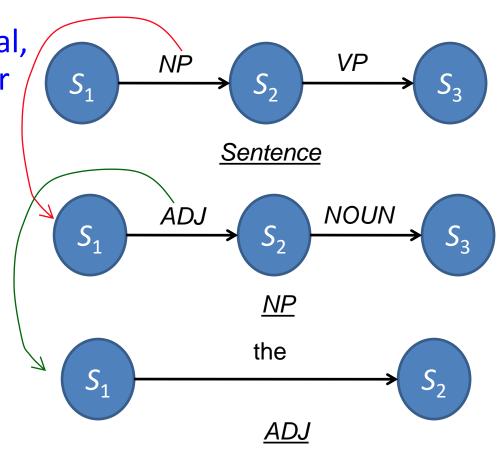
 If arc label is non-terminal, control passes to another TN



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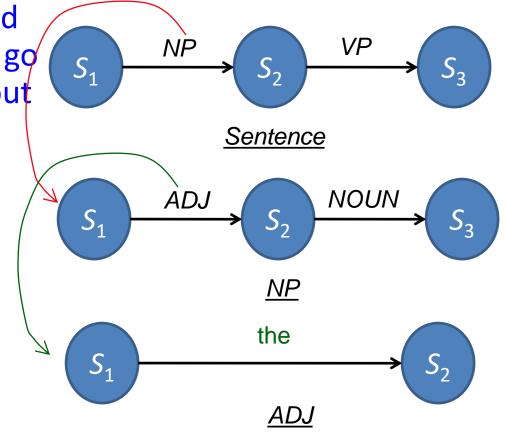
the program crashes the computer



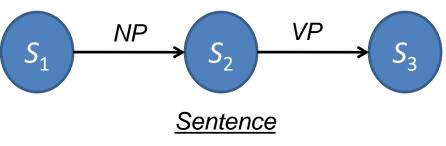
- Starts at an initial node
- Sequentially checks input string against the arc label

 If arc label is terminal and matches with the string, go to next node and the input is consumed

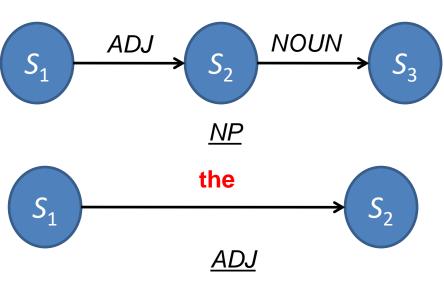
the program crashes the computer



- Starts at an initial node
- Sequentially checks input string against the arc label
- If arc label is terminal and doesn't match with the string, either backtrack or indicate a fail

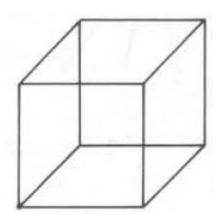


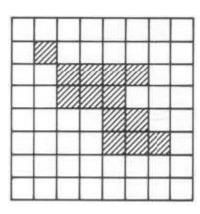
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Capability of String Grammar

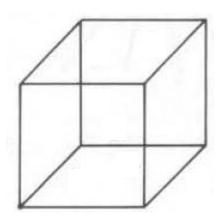
- able to represent 1D pattern
- difficult to classify patterns like these:

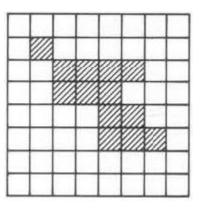




Capability of String Grammar

- These structure contains
 - information in both direction
 - hierarchical structure

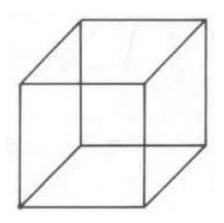


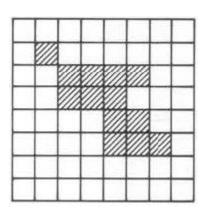


Capability of String Grammar

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Solution: Higher Order Grammar



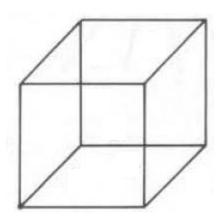


Higher Order Grammar: Tree Grammar

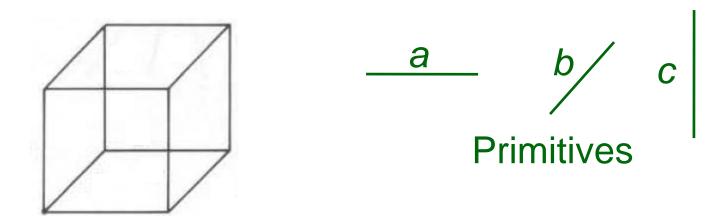
- Information is stored in
 - nodes: as primitives or sub-structures
 - edges: relation between primitives or sub-structures

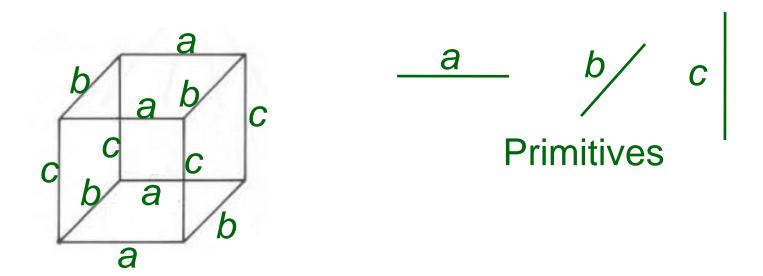
 Allows hierarchical decomposition of a complex structure

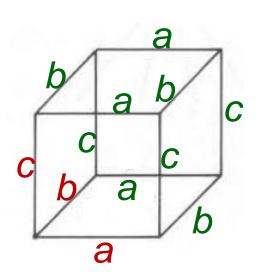
Representation using Tree (1)

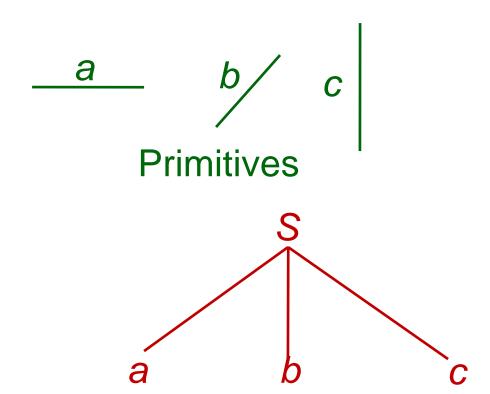


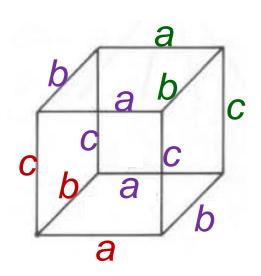
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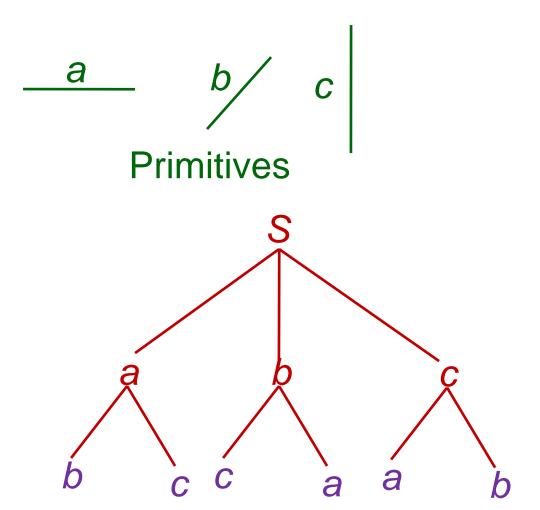


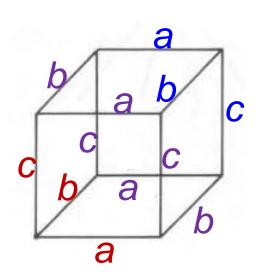


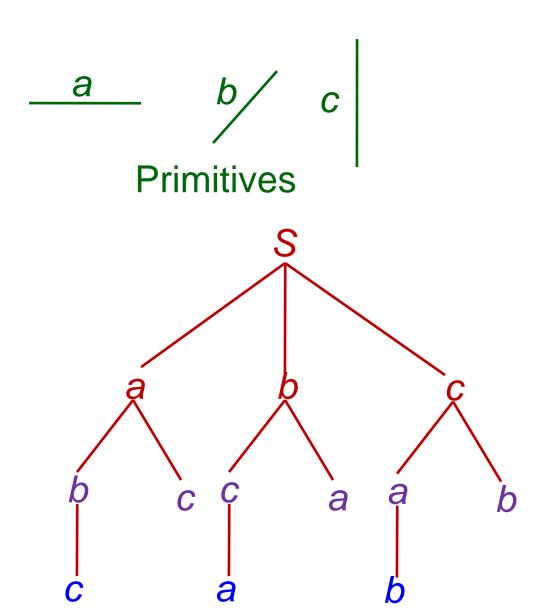


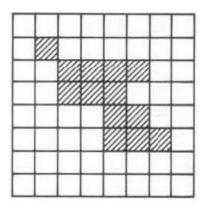


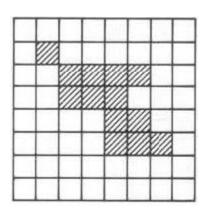














Primitives

Black region

White region

O Gray region

