



Assembly Language

Chapter 3 & 4

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Agenda

- CF and OF Flags ADD/SUB
- INC and DEC Instructions
- Data Operators
- Indirect Addressing

Flags

- Why flags are important:
 - To detect calculation errors
 - Check the outcome of an arithmetic operation
 - Activate conditional branching (loops and if/else)
- Let's see how each flag is affected by ADD/SUB instructions.

Carry Flag (CF)

- CF is set when the result of an **unsigned** operation is out of destination range
- ADD: CF is set when result is larger than the destination operand

```
mov al, FFh ; 255
```

```
add al, 1 ; CF = 1, SUM too large for AL
```

- SUB: CF is set when subtract a large operand from a smaller one (i.e., there's a borrow bit)

```
mov al, 3
```

```
sub al, 5 ; CF = 1, Res(-2)
```

cannot be placed in
unsigned number

00000011	00000011
- 00000101	+11111011
-----	-----
	11111110
	Carry out=0

Overflow Flag (OF)

- OF is set when the result of a **signed** operation is out of destination range

```
mov al, +127
```

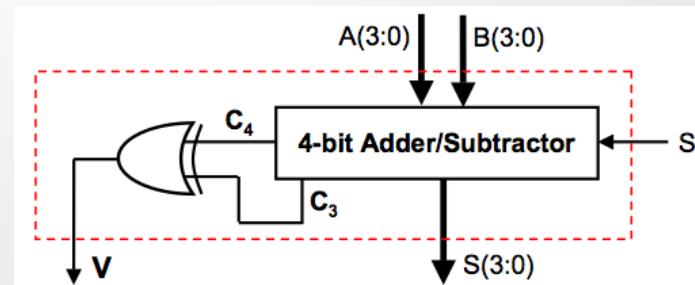
```
add al, 1 ; OF = 1, AL = 80h (-128)
```

- The overflow flag is set when ...
 - Two positive operands are added and their sum is negative → **Ex: +127 +1**
 - Two negative operands are added and their sum is positive → **EX: -128 -1**
 - Overflow cannot occur when adding operands of opposite signs

Hardware Perspective

- All CPU instructions operate exactly the same on signed and unsigned integers
- CPU does NOT distinguish signed from unsigned
- Hardware-wise, the CF and OF flags are set according to the following:
 - **CF** = Carry out for **ADD** instruction
 = INVERT(Carry out) for **SUB** instruction *
 - **OF** = (carry out of the MSB) XOR (carry in the MSB)

*[AF is inverted with SUB as well]



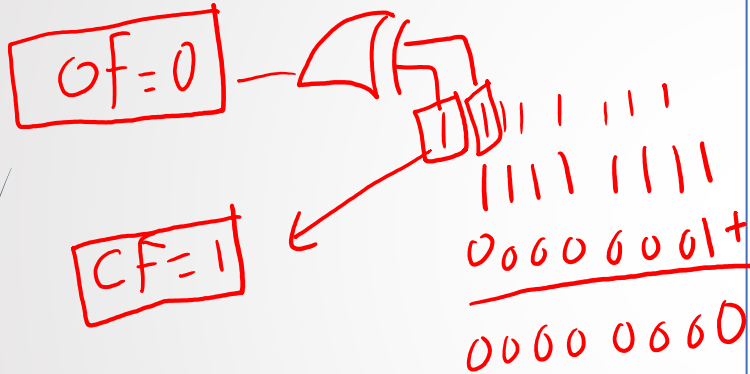
p	q	$p \oplus q$
1	1	0
1	0	1
0	1	1
0	0	0

(board)

7

Add

$$\begin{array}{r} 255 \\ + 1 \\ \hline \end{array}$$



U

S

CF
= 1

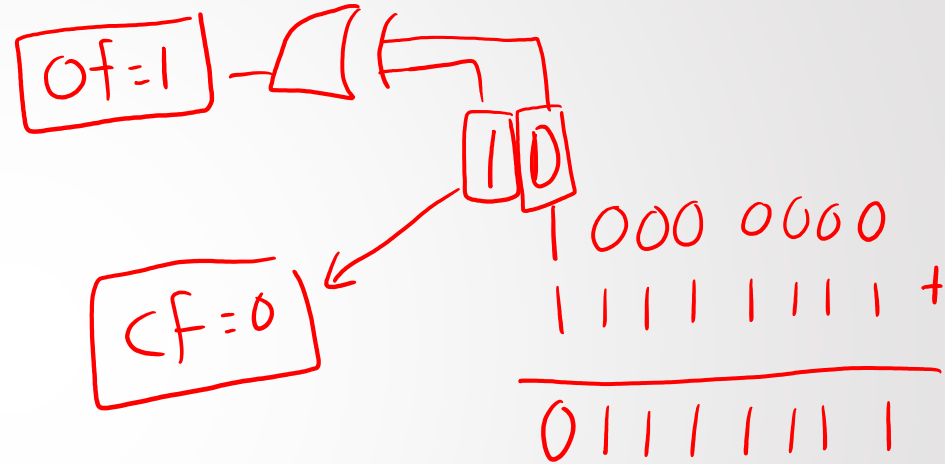
$$\begin{array}{r} 255 \\ + 1 \\ \hline 256 \end{array}$$

$$\begin{array}{r} -1 \\ + 1 \\ \hline 0 \end{array}$$

Of
= 0

Sub

$$\begin{array}{r} 128 \\ - 1 \\ \hline \end{array} \quad \begin{array}{r} 128 \\ + (-1) \\ \hline \end{array}$$



U

S

CF
= 0

$$\begin{array}{r} 128 \\ - 1 \\ \hline 127 \end{array}$$

$$\begin{array}{r} -128 \\ - 1 \\ \hline -129 \end{array}$$

Of
= 1

Overflow Flag (OF)

➤ What will be the values of the Overflow flag?

➤ `mov al, 80h`

-ve operands, +ve result

➤ `add al, 92h`

; al = 12h OF = 1

	1	0	0	0	0	0	0	
	1	0	0	0	0	0	0	80h (-128)
+	1	0	0	1	0	0	1	92h (-110)
	0	0	0	1	0	0	1	12h result should be -238 < -128

CF=1

- ➡ **add** al, 45h

$$; a_1 = 85h \quad OF = 1$$

0 1 0 0 0 0 0 0
 0 1 0 0 0 0 0 0 40h (64)
 + 0 1 0 0 0 1 0 1 45h (69)

 1 0 0 0 0 1 0 1 85h result
 should be 133 > 127

Overflow Flag (OF)

➤ What will be the values of the Overflow flag?

➤ `mov al, -128`

-ve operands, +ve result

➤ `sub al, 1`

; al = +127 OF = 1

	1	0	0	0	0	0	0	
	1	0	0	0	0	0	0	80h (-128)
+	1	1	1	1	1	1	1	FFh (-1)
CF=0	0	1	1	1	1	1	1	result should be -129 < -128

Addition and Subtraction

- ▶ Show the values of the destination operand and the six status flags
- ▶ Notice : AF is inverted in SUB like CF

<code>mov al,0FFh</code>	<code>; AL=-1</code>	MOV does NOT affect flags					
<code>add al,1</code>	<code>; AL=</code>	<code>CF=</code>	<code>OF=</code>	<code>SF=</code>	<code>ZF=</code>	<code>AF=</code>	<code>PF=</code>
<code>sub al,1</code>	<code>; AL=</code>	<code>CF=</code>	<code>OF=</code>	<code>SF=</code>	<code>ZF=</code>	<code>AF=</code>	<code>PF=</code>
<code>mov al,+127</code>	<code>; AL=7Fh</code>						
<code>add al,1</code>	<code>; AL=</code>	<code>CF=</code>	<code>OF=</code>	<code>SF=</code>	<code>ZF=</code>	<code>AF=</code>	<code>PF=</code>
<code>mov al,26h</code>							
<code>sub al,95h</code>	<code>; AL=</code>	<code>CF=</code>	<code>OF=</code>	<code>SF=</code>	<code>ZF=</code>	<code>AF=</code>	<code>PF=</code>

result = 0
 $\Rightarrow ZF = 1$

12

mov al, 0FFh	; AL= FFh	CF= X	OF= X	SF= X	ZF= X	AF= X	PF= X
add al, 1	; AL= 00h	CF= 1	OF= 0	SF= 0	ZF= 1	AF= 1	PF= 1
sub al, 1	; AL= ffh	CF= 1	OF= 0	SF= 1	ZF= 0	AF= 1	PF= 1
mov al, +127	; AL= 7Fh	CF= X	OF= X	SF= X	ZF= X	AF= X	PF= X
add al, 1	; AL= 80h	CF= 0	OF= 1	SF= 1	ZF= 0	AF= 1	PF= 0
mov al, 26h		X	X	X	X	X	X
sub al, 95h	; AL= 91h	CF= 1	OF= 1	SF= 1	ZF= 0	AF= 0	PF= 0

Add al, 1
 (FF + 1)

1	1	1	1	1	1	1	1
0	0	0	0	0	0	0	1
0	0	0	0	0	0	0	0

Add al, 1
 (127 + 1)

0	1	1	1	1	1	1	1
0	0	0	0	0	0	0	1
1	0	0	0	0	0	0	0

Sub al, 1
 (0 - 1)

$\begin{matrix} 0 \\ -1 \end{matrix} + \begin{matrix} 0 \\ +1 \end{matrix}$

0	0	0	0	0	0	0	0
1	1	1	1	1	1	1	1
1	1	1	1	1	1	1	1

Sub al, 95h
 (26h - 95h)

0	0	1	0	0	1	1	0
1	0	0	1	0	1	0	1

0	1	1	1	1	1	1	1
0	0	1	0	0	1	1	0
0	1	1	0	1	0	1	1
1	0	0	1	0	0	0	1

91h

Addition and Subtraction

- Show the values of the destination operand and the six status flags
- Notice : AF is inverted in SUB like CF

```

mov al,0FFh      ; AL=-1  MOV does NOT affect flags
add al,1         ; AL=00h  CF=1 OF=0 SF=0 ZF=1 AF=1 PF=1
sub al,1         ; AL=FFh  CF=1 OF=0 SF=1 ZF=0 AF=1 PF=1
mov al,+127      ; AL=7Fh
add al,1         ; AL=-128 CF=0 OF=1 SF=1 ZF=0 AF=1 PF=0
mov al,26h
sub al,95h       ; AL=91h  CF=1 OF=1 SF=1 ZF=0 AF=0 PF=0
  
```

1	0	0	1	0	0	0	1
0	0	1	0	0	1	1	0
26h (38)							
-							
1	0	0	1	0	1	0	1
95h (-107)							

0	1	1	0	1	1	1	0
0	0	1	0	0	1	1	0
26h (38)							
+							
0	1	1	0	1	0	1	1
6Bh (107)							
<hr/>							
1	0	0	1	0	0	0	1
91h (-111)							

Agenda

- CF and OF Flags ADD/SUB
- **INC and DEC Instructions**
- Data Operators
- Indirect Addressing

Addition and Subtraction

➤ INC, DEC, NEG instructions

➤ INC *dest*

- Like **dest ++** in C language
- More compact (uses less space) than: ADD dest, 1

➤ DEC *dest*

- Like **dest --** in C language
- More compact (uses less space) than: SUB dest, 1

➤ NEG *dest*

- **dest = 2's** complement of destination (0-dest)

Addition and Subtraction

- *dest* can be
 - 8-, 16-, or 32-bit operand
 - Memory or a register
 - **NO** immediate operand

Addition and Subtraction

- INC and DEC affect five status flags
 - Overflow, Sign, Zero, Auxiliary Carry, and Parity similar to ADD and SUB instruction
 - Carry flag is **NOT** modified
- NEG affects all the six status flags
 - Any **nonzero source operand** causes the **carry flag to be set**
 - **OF** is set if the result cannot be represented in the operand
 - Maximum negative number in range (e.g. +128 in 8-bit register)

(board)

$$\text{neg } x \equiv 0 - x$$

if x $\xrightarrow{0}$ $0-0 \Rightarrow$ no borrow $\Rightarrow \boxed{CF=0}$
 $\xrightarrow{\text{non zero}}$ borrow $\Rightarrow \boxed{CF=1}$

neg $\boxed{-128}$, in 8-bits range $-128 \rightarrow 127$

128 : 1000 0000 > 127
 -128 : 1000 0000 \Downarrow

can not be represented
 in 8-bits

INC/DEC/NEG

Trace the flags

.DATA

B SBYTE -1 ; 0FFh or 1111 1111

C SBYTE 127 ; 7Fh or 0111 1111

.CODE

inc B ; B=0 CF=X OF=0 SF=0 ZF=1 PF=1

dec B ; B=-1 CF=X OF=0 SF=1 ZF=0 PF=1

inc C ; C=-128=80h OF=1 SF=1 ZF=0 PF=0

neg C ; C=-128 CF=1 OF=1 SF=1 ZF=0 PF=0

0 1 1 1 1 1 1 1

0	1	1	1	1	1	1	1
---	---	---	---	---	---	---	---

 127

+

0	0	0	0	0	0	0	1
---	---	---	---	---	---	---	---

 1+

1	0	0	0	0	0	0	0
---	---	---	---	---	---	---	---

 128-

non-zero

max -ve number

← inc C

← Considered as -128

Agenda

- CF and OF Flags ADD/SUB
- INC and DEC Instructions
- **Data Operators**
- Indirect Addressing

Data-Related Operators and Directives

1. OFFSET Operator
2. PTR Operator
3. TYPE Operator
4. LENGTHOF Operator
5. SIZEOF Operator
6. LABEL Directive

OFFSET operator

- OFFSET returns the distance in bytes, of a label from the beginning of its enclosing segment
- Let's assume that the data segment begins at 00404000h

.data

bVal **BYTE** ?

wVal **WORD** ?

dVal **DWORD** ?

dVal2 **DWORD** ?

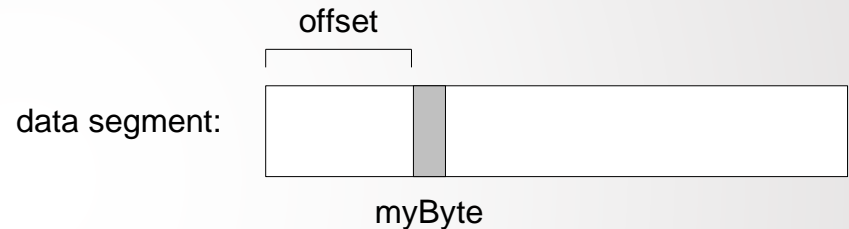
.code

mov esi,OFFSET bVal ; ESI = 00404000

mov esi,OFFSET wVal ; ESI = 00404001

mov esi,OFFSET dVal ; ESI = 00404003

mov esi,OFFSET dVal2 ; ESI = 00404007



bval		00404000
wval		00404001
dval		00404002
		00404003
		00404004
dval2		00404005
		00404006
		00404007
		00404007

OFFSET operator

- ▶ The value returned by OFFSET is a pointer

// C++ version:

```
char array[1000];  
char * p = array;
```

; Assembly language:

```
.data
```

```
array BYTE 1000 DUP(?)
```

```
.code
```

```
mov esi, OFFSET array
```

PTR operator

- ▶ PTR operator overrides the default type of a label (variable). Provides the flexibility to access part of a variable
- ▶ Similar to typecasting in C

.data

myDouble **DWORD** 12345678h

.code

mov ax, myDouble ; **error** - why?

mov ax, **WORD PTR** myDouble ; **AX = 5678h**

mov **WORD PTR** myDouble, 3412h ; **myDouble=12343412h**

myDouble	78	0x4020
	56	0x4021
	34	0x4022
	12	0x4023

PTR operator

.data

myDouble **DWORD** 12345678h

.code

mov al,**BYTE PTR** myDouble ;AL=78h

mov al,**BYTE PTR** [myDouble+1] ;AL=56h

mov al,**BYTE PTR** [myDouble+2] ;AL=34h

mov ax,**WORD PTR** myDouble ;AX=5678h

mov ax,**WORD PTR** [myDouble+2] ;AX=1234h

myDouble

78	0x4020
56	0x4021
34	0x4022
12	0x4023

PTR operator

- ▶ PTR can be used to **combine elements** of a smaller data type and move them into a larger operand
- ▶ The CPU will automatically reverse the bytes

.data

```
myBytes BYTE 12h, 34h, 56h, 78h
```

myBytes	12
	34
	56
	78

.code

```
mov ax, WORD PTR [myBytes] ;AX=3412h
```

```
mov ax, WORD PTR [myBytes+2];AX=7856h
```

```
mov eax, DWORD PTR myByte ;EAX=78563412h
```

PTR operator – Your turn

Ex: Write down the value of each destination

.data

varB **BYTE** 65h, 31h, 02h, 05h

varW **WORD** 6543h, 1202h

varD **DWORD** 12345678h

.code

```

mov ax, WORD PTR [varB+2]      ;0502h
mov bl, BYTE PTR varD          ;78h
mov bl, BYTE PTR [varW+2]      ;02h
mov ax, WORD PTR [varD+2]      ;1234h
mov eax,DWORD PTR varW         ;12026543h

```

varB	65
	31
	02
	05
varW	43
	65
	02
	12
varD	78
	56
	34
	12

TYPE operator

- ▶ The TYPE operator returns the size, in **bytes**, of a single element of a data declaration

.data

var1 BYTE ?

var2 WORD ?

var3 DWORD ?

var4 QWORD ?

.code

mov eax, TYPE var1 ; 1

mov eax, TYPE var2 ; 2

mov eax, TYPE var3 ; 4

mov eax, TYPE var4 ; 8

LENGTHOF operator

- ▶ The LENGTHOF operator counts the **number of elements** in a single data declaration

.data

```
byte1    BYTE 10, 20, 30                ; 3
array1    WORD 30 DUP(?), 0, 0           ; 32
array2    WORD 5 DUP(3 DUP(10))         ; 15
array3    DWORD 1, 2, 3, 4              ; 4
digitStr  BYTE "12345678", 0            ; 9
```

.code

```
mov ecx, LENGTHOF array1                ;ECX=32
mov eax, LENGTHOF digitStr              ;EAX=9
```

SIZEOF operator

- ▶ The SIZEOF operator returns number of **bytes** of a data declaration (i.e., **LENGTHOF * TYPE**)

.data

```
array1 WORD 30 DUP(?), 0, 0      ; 64 → 32*2
array2 WORD 5 DUP(3 DUP(10))    ; 30 → 15*2
array3 DWORD 10, 20,             } ; one array
                                30, 40, ; 24 → 6*4
                                50, 60
```

.code

```
mov ax, SIZEOF array3           ; 24
mov bx, LENGTHOF array3        ; 6
mov cx, TYPE array3            ; 4
```

SIZEOF operator

- ▶ Caution: Array spanning multiple lines

`.data`

`array WORD 10, 20`

`WORD 30`

`WORD 50, 60, 70`

This is NOT a single array
These are 3 different array
definitions

`.code`

`mov eax, LENGTHOF array ; 2`

`mov ebx, SIZEOF array ; 4`

LABEL directive

- ▶ Assigns an alternate label and type to the next storage location
- ▶ LABEL does **NOT** allocate any storage of its own

.data

dwList LABEL DWORD ;no declaration

wordList LABEL WORD ;no declaration

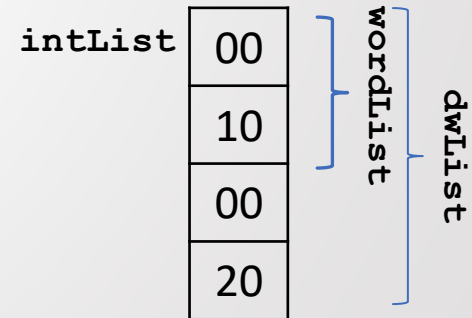
intList BYTE 00h, 10h, 00h, 20h

.code

mov eax, dwList ; 20001000h

mov cx, wordList ; 1000h

mov dl, intList ; 00h



LABEL directive

.data

dwList LABEL DWORD

arr1 byte 1, 2, 3, 4

byteList LABEL byte

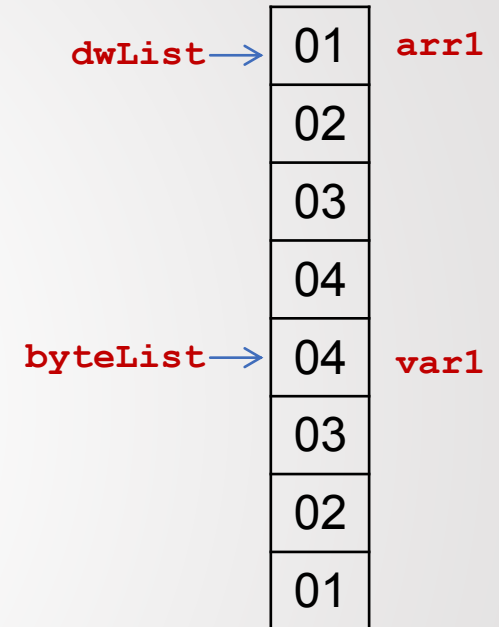
var1 DWORD 01020304h

.code

mov eax, dwList ; 04030201h

mov al, byteList ; 04h

mov dl, byteList+3 ; 01h



Agenda

- CF and OF Flags ADD/SUB
- INC and DEC Instructions
- NEG Instruction
- Data Operators
- Indirect Addressing

Remember Direct Offset Addressing?

- **Direct Offset Addressing** by adding a constant offset (displacement) to the label to produce the effective address

```
.data
```

```
    arrayB BYTE 10h,20h,30h,40h
```

```
.code
```

```
mov al, arrayB + 1 ; AL = 20h
```

```
mov al,[arrayB +1] ;alternative notation
```

```
mov al, arrayB[1] ;yet another notation
```

- **Note:** offset expresses on bytes not an item index

Register-indirect Addressing

- Direct offset mode is not so flexible since it is not practical to traverse an array using **constant offsets**
- Register-indirect addressing mode solves this problem by allowing storing the array address in a **register** (similar to pointers in C)

Register-indirect Addressing

- ▶ The operand offset is stored in a register
- ▶ Brackets [] are used to surround the register holding the address
- ▶ For 32-bit addressing, any 32-bit register can be used

```
mov ebx, OFFSET array1 ;ebx = array1 address
```

```
mov eax, [ebx]          ;[ebx] gets 1st item
```

Register-indirect Addressing

- ▶ In this mode, register holds the address of a variable, usually an **array** or **string**.

.data

```
var1 BYTE 10h, 20h, 30h
```

.code

```
mov esi, OFFSET var1
mov al, [esi]          ; AL = 10h
inc esi                ; esi ++
mov al, [esi]          ; AL = 20h
inc esi                ; esi ++
mov al, [esi]          ; AL = 30h
```

Register-indirect Addressing

- ▶ Register-indirect Addressing can point to any size of a memory location (byte, word, ..etc)
- ▶ Thus, you should use PTR to clarify the size of a memory operand

`.data`

`myCount WORD 0`

`.code`

`mov esi, OFFSET myCount`

`inc [esi] ; Error: operand must have size`

`inc WORD PTR [esi] ; OK`

*inc esi
inc [esi]
⇒ diff??*

Register-indirect Addressing

- This addressing mode is ideal for traversing an array.
- Note that the register in brackets must be incremented by a value that matches the **array type**
 - 1 for BYTE,
 - 2 for WORD,
 - 4 for DWORD

Register-indirect Addressing

- Write a program to **sum** the array elements

.data

```
arrayW WORD 1000h, 2000h, 3000h
```

.code

```
mov esi, OFFSET arrayW
```

```
mov ax, [esi]
```

```
add esi, 2      ; or: add esi, TYPE arrayW
```

```
add ax, [esi]
```

```
add esi, 2
```

```
add ax, [esi]   ; AX = sum of the array
```

Or we can write: **add esi, type arrayW**

Register-indexed Addressing

- ▶ Register-indexed operand treats the register value as an index added to the array offset to generate the effective address.
- ▶ There are two notational forms:

[label + reg] or label[reg]

.data

```
arrayW WORD 1000h, 2000h, 3000h
```

.code

```
mov esi, 0
```

```
mov ax, [arrayW + esi]    ; AX = 1000h
```

```
mov ax, arrayW[esi]       ; alternate format
```

```
add esi, 2
```

```
add ax, [arrayW + esi]
```

Register-indexed Addressing

- ▶ You can **scale** the indexed operand to the offset of an array element. This is done by multiplying the index by the array's TYPE

.data

```
arrB BYTE 0, 1, 2, 3, 4, 5
```

```
arrW WORD 0, 1, 2, 3, 4, 5
```

```
arrD DWORD 0, 1, 2, 3, 4, 5
```

.code

```
mov esi, 4 ;item index (0-based)
```

```
mov al, arrB[esi*TYPE arrB] ; arrW[esi*1] ;arrW[4]
```

```
mov ax, arrW[esi*TYPE arrW] ; arrW[esi*2] ;arrW[8]
```

```
mov eax, arrD[esi*TYPE arrD]; arrD[esi*4] ;arrW[16]
```

- ▶ All instructions access item of index 4

Reference

- Textbook:
 - Chapter 4: 4.1 – 4.4

Thank you

