

Competitive Programming

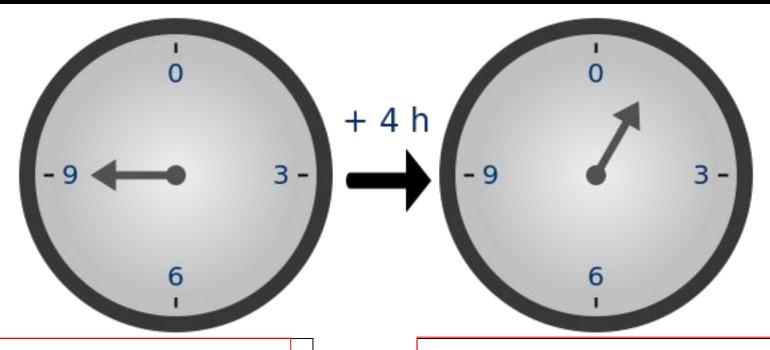
From Problem 2 Solution in O(1)

Number Theory Modular Arithmetic

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12-hour Clock Cycle



- If it is 9 now, what time:
- after 4 h? 1
- after 16 (4 + 12) h? 1
- \blacksquare after 17 (5 + 12) h? 2
- after 29 (5 + 2*12) h? 2
- before 24 (2*12) h? 9
- before 25 (1+2*12) h? 8

- Facts:
- N = x + m * 12
 - N is number, x < 12, m >= 0
- Every multiple of 12 is useless
- What is less than 12 affects us.
- We can go forward or backward

Modulo (modulus) operation

- a modulo n = finds the remainder of after division by n: In C++, operator is %
- let a = 27, n = 12, then r = a % n?
- 27 / 12 = (3+2*12)/12 = 3/12 + 2 = 2.25
 - q (quotient), the Integer division part is 2
 - r (remainder) of division is 3
 - $r = 27 \% 12 = 3 \Rightarrow$ Remainder from division
- a = nq + r (q multiple of n + r [< n])
- % operator is finally: $r = a n \left| \frac{a}{n} \right|$
- $|\mathbf{r}| < \mathbf{n}$

Back to the clock

- If it is 9 now, what time:
- after 4 h? => 9 + 4 = 13 h => 13 % 12 = 1
- after 16 (4 + 12) h? \Rightarrow 9 + 16 = 25 h \Rightarrow 25 % 12 = 1
- after 17 (5 + 12) h? \Rightarrow 9 + 17 = 26 h \Rightarrow 26 % 12 = 2
- after 29 (5 + 2*12) h? => 9 + 29 = 38 h => 38 % 12 = 2
- before 24 (2*12) h? $9 \Rightarrow 9 24 = -15h \Rightarrow hmm$
 - 15 % 12 = 3...hmm, we are sure results should be 9 too
 - +ve is not as same -ve
 - Fact: r = a % n = (a+qn)%n => I.e. adding multiplier on doesn't affect results
 - -15 + 12 = -3, still negative, add another 12
 - $-3 + 12 = 9 \dots$ Good! Done
 - In C++: -15h % 12 = -3, so you need to add 12 **only once**
- What time before 25 (1+2*12) h? 9 25 = -16 => -16 % 12 = -4 [in C++]
 - Add 1 cycle to make it positive: -4 + 12 = 8 hours
- In C++: for any r => (a % n + a) % n is always positive

modulus is expensive

- % and / are time expensive operations
- If you can avoid them, avoid them
- One scenario, when you are sure results can be fixed with little +/- of mod value
 - we can directly do: a = (a%n + n)%n
 - 1 addition and 2 mod operations
 - maybe we can fix results with e.g. 2 comparison/add
 - while $(a \ge n)$ a = n;
 - while (a < 0) a += n;

- To get modulus => add/remove cycles of n till
 - 0 <= r <= n-1
 - **27** % 12 => 15 % 12 => 3 % 12 = 3
 - -15 % 12 => -3 % 12 => 9 % 12
- |a%n| has n-1: 0, 1, ...n-1
- In C++:
 - a%3 = -2, -1, 0, 1, 2 [for a -ve or +ve]
 - \bullet a % n (for +ve) or (a % n + n) % n (generally)
- $a\%n = 0 \Rightarrow a \text{ divisible by } n$
- If a%n == b%n => (a-b)%n = 0
- largest n such that a%n = b%n is n = b-a

- (a % n) % n = a % n
- $(n \land x) \% n = 0 \text{ for any } x >= 0$
- -a\%n!= a\%n => (3 %12 = 3 vs -3 %12 = 9)
- ((-a%n)+(a%n))%n=0
- (a+b) % n = (a%n + b%n) % n
- a+b+c+d%n?
 - You can take mod of every one and sum
 - or ((((a%n+(b%n))%n+c%n)%n+d%n)%n
- x % (a+b) != x % a + x % b
- x%10 [the last digit]. x/10 [remove last digit]

- (a*b) % n = (a%n * b%n) % n
- (a^b) % $n = ((a\%n)^b) \% n$
- (a^b) % n => assume b even and x = b/2
 - $((a^x) \% n * (a^x) \% n)\% n$
- (1/a) % n? modular multiplicative inverse
- ((a*b) % n * (1/a)%n) %n = b % n
- a % (2^n) = a & (n-1) => E.g. a%4 = a&3
- a % 0 is undefined
- When -ve result => result = (result + n)%n

What is wrong here?

```
bool is_odd(int n) {
    return n % 2 == 1;
}
```

```
bool is_odd(int n) {
    return n % 2 == 1 || n % 2 == -1;
}
```

```
bool is_odd(int n) {
    return n % 2 != 0;
}
```

Cycling examples

- A machine keeps generating the sequence 5 2 7 1 for infinity...what is its value after 10^12 steps? 5 2 7 1 **5 2 7 1** 5 2 7 1
 - After 0 steps => 5 After 3 steps => 5
 - After 4 steps \Rightarrow 5 After 5 steps \Rightarrow 2
 - It keep cycling. Remove all cycling at once: 10¹² % 4
 - Rings, Cycles, ...should trigger the mod
- Given position X in array, iterate back M steps? We may cycle and back to array end

Why modulus?

- Either cycle (ring) is nature of the problem
 - 12-hour clock, week is 7 days, year is 356/366 days
- Encryption Algorithms, Pseudo-random Generators
- For fun, e.g. what is the last digit of 2^100?
- In competitions, final result is too big, but we want to avoid using big integers. Using mode, truncate results
- You are sure final results $\leq n$, but intermediate results overflow. Take intermediate % x (x > n)
 - \blacksquare 1001 1111 + 153 = **43** ...let x = 44
 - ((1001%44 + ((-1111%44)+44)%44 + 153%44)%44
 - (33+33+21)%44 = 87%44 = 43

UVA 408, 10006, CF447-A, CF284-A, 332A, 11155, 132A, 374, 128,

SRM 144-D2-1

CF476-D2-C

https://www.hackerrank.

com/domains/mathematics/fundamentals

تم بحمد الله

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