

# Trees 4

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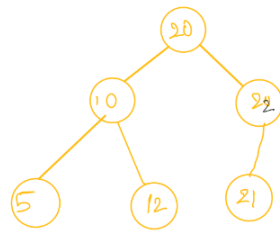
# Find Kth Smallest Element

Binary Tree : Hierarchical data structure where each node can have at max 2 children.

Binary Search Tree is Type of B tree with an additional property.  
for all nodes.

All ele in LST  $<=$  Node  $<=$  All ele in RST

Q Find  $K^{th}$  Smallest element in a BST.



$K=1 : 5$   
 $K=3 : 12$

→ Inorder traversal of BST is a sorted array.  
 $\{5, 10, 12, 20, 22, 21\}$

Approach 1 :-  
Do inorder traversal & store in array  
& return  $9^{th}$  index.  
 $T: O(N)$   
 $S: O(N)$

Approach 2 :-

Keep a tracker;

void inorder (Node root, int k) {

if (root == NULL)

return;

inorder (root->left, k);

cnt++;

if (cnt == k) {

ans = root->data;

return;

inorder (root->right, k);

}

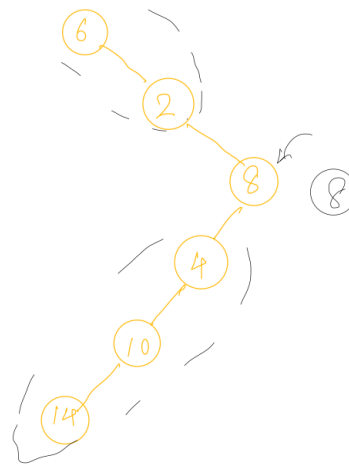
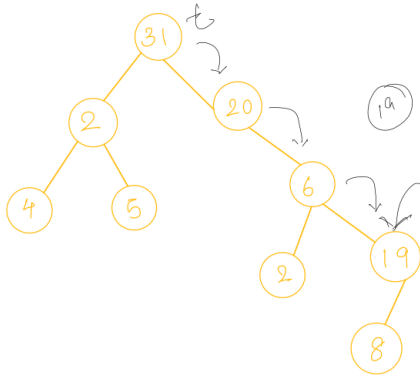
$T.C: O(N)$

$S.C: O(H)$

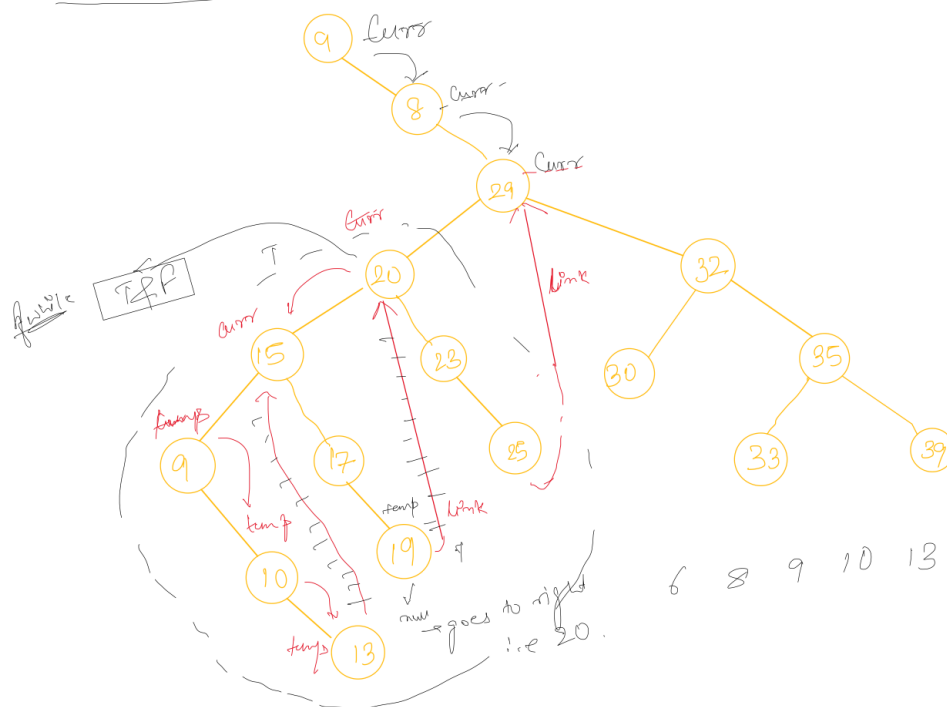
Approach 3 :- Inorder traversal using a stack.  
 $T.C: O(N)$   
 $S.C: O(H)$

# Morris Traversal

Q With inorder traversal on a tree, what's the last node we print?



Morris Traversal



6 8 9 10 13 15 17 19

Pseudo Code:-

```

if (curr.left == NULL)
{
    print (curr.data);
    curr = curr.right;
}
else
{

```

Node temp = curr.left;

while (temp.right != NULL && temp.right != curr)

```

{
    temp = temp.right
}

```

if (temp.right == NULL)

```

{
    temp.right = curr;
    curr = curr.left;
}

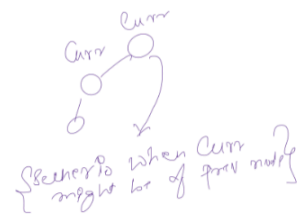
```

else // Visiting it for 2<sup>nd</sup> time;

```

{
    temp.right = NULL;
    print (curr.data);
    curr = curr.right;
}
}

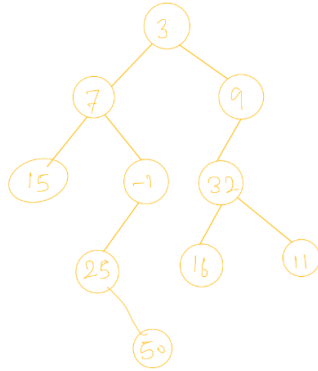
```



Visiting curr for second time

## Path from root to element's node

Q Search an ele in Binary tree; (All distinct) / Path from root to ele.



Search(25): True

Pseudo  
Code:-

list<int> ans

bool Search(Node root, int k)

{  
if (root == NULL)  
return false;

if (root->data == k)

{ ans.push\_back(root->data);  
return true;

}  
if (Search(root->left, k) || Search(root->right, k))

{ ans.push\_back(root->data);  
return true;

}

else

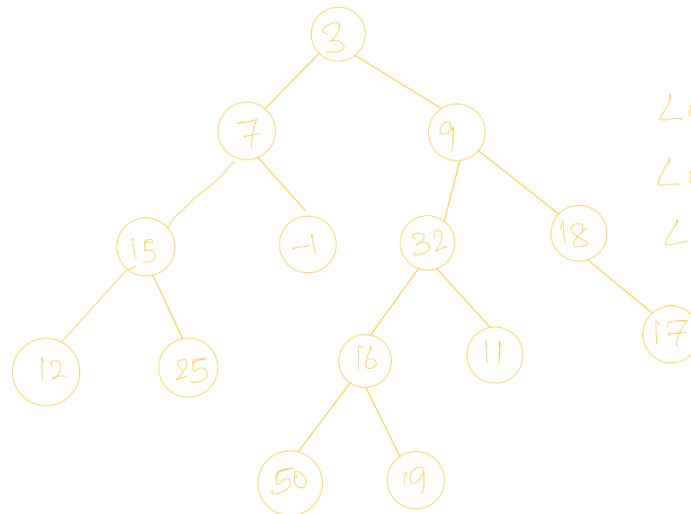
return false;

}

//reverse(ans) -> path from root to node.

## LCA in B tree

Q Find LCA (a, b) in a B-Tree.



$$\text{LCA}(16, 18) = 9$$

$$\text{LCA}(50, 19) = 16$$

$$\text{LCA}(11, 9) = 9$$

$$\text{LCA}(25, 9) = 3$$

LCA(a, b):

Path from Root to a: { 3 9 (32) 16 50 }

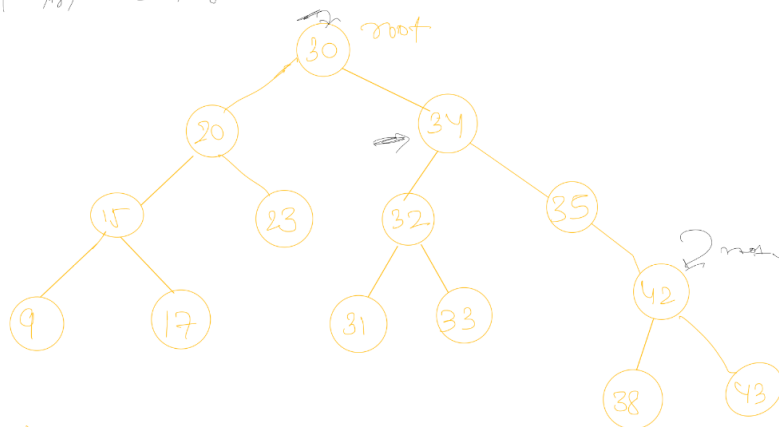
Path from Root to b: { 3 9 (32) 11 }

→ Traverse until you find first non common ancestor.

T.C:  $O(N)$

S.C:  $O(N)$

Q6 LCA in BST :-



$$LCA(31, 43) = 34$$

$$LCA(34, 33) = 34$$

$$LCA(23, 35) = 30$$

$$LCA(38, 43) = 42$$

Pseudo code

Node curr = root;

while (true)

{ if (a.data < curr.data && b.data < curr.data)

{ curr = curr.left;

} else if (a.data > curr.data && b.data > curr.data)

{ curr = curr.right;

}

else {

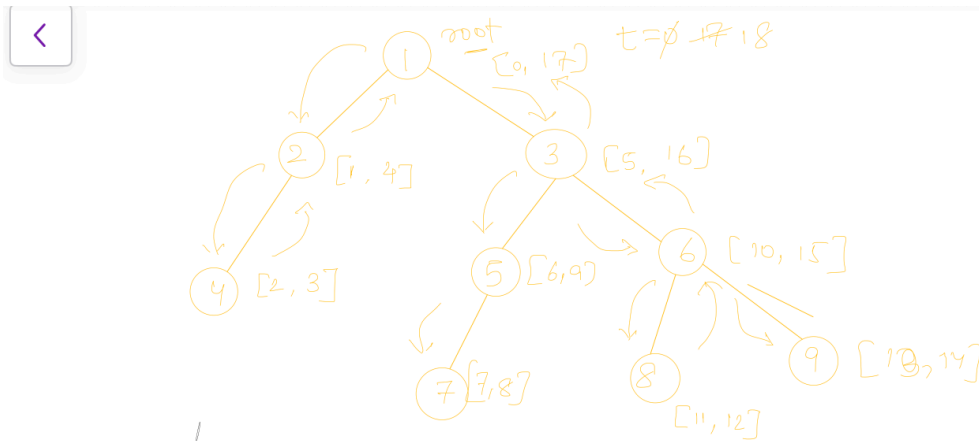
break;

}

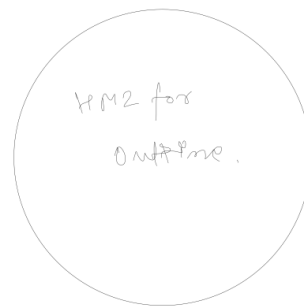
}

return curr.data;

# In Time and Out time in a B tree



for  $\forall$  nodes,  $\begin{cases} \rightarrow \text{in} \\ \rightarrow \text{out} \end{cases}$



pseudo code

```

int t = 0;
void traverse (Node root) {
    if (root == NULL)
        return;
    intime.insert (&root, time);
    time++;

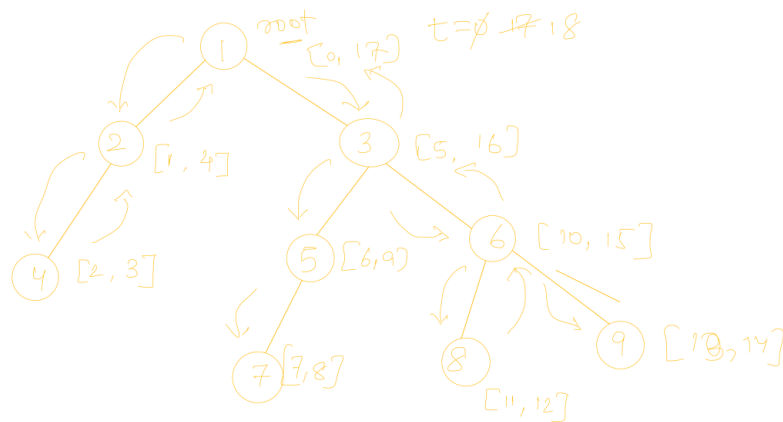
    traverse (root.left);
    traverse (root.right);
    Outtime.insert (&root, time);
    time++;
}
  
```

T.C :  $O(N)$   
S.C :  $O(N)$

// using pointers



# LCA(a, b) for multiple Queries



$LCA(5, 8) : 3$

$(6, 9)$   $(11, 12)$

$a, b$   
 $in(a) < in(b)$   
 $out(a) > out(b)$

} a is an ancestor.

Pseudo Code:-

```

curr = root;
while (curr != NULL)
{
    if (curr.left is an ancestor of a & b)
    {
        curr = curr.left;
    }
    else if (curr.right is an ancestor of a & b)
    {
        curr = curr.right;
    }
    else
    {
        return curr;
    }
}

```