

## Today's Agenda :-

Will Start at 9:05 PM

- 1) Intro to Prime Numbers
- 2) Get all primes from 1 to N (Sieve of Eratosthenes)
- 3) Print smallest prime factor for 2 to N
- 4) Prime factorization
- 5) Get the no of factors / divisors

Prime Number :- Nos having exactly 2 factors are Prime.



Eg  $N=1$  ✗

Eg  $N=2$  ✓

3, 5, 7, 11, ...

Q

```

bool checkPrime(int N) {
    int cnt = 0;
    for (int i = 1; i <= sqrt(N); i++) {
        if (N % i == 0) {
            if (i == N/i) cnt += 1;
            else cnt += 2;
        }
    }
    if (cnt == 2) return true;
    else return false;
}
  
```

TC:  $O(\sqrt{N})$

$$i * j = N$$

$$j = N/i$$

$N=12$

i	N/i	Count
1	12	+2
2	6	+2
3	4	+2
4	3	+6
6	2	
12	1	

$i <= \frac{N}{i}$   
 $i^2 <= N$   
 $i <= \sqrt{N}$

$N=36$

i	N/i	Count
1	36	+2
2	18	+2
3	12	+2
4	9	+2
6	6	+1
9	4	
12	3	
18	2	
36	1	

Q.) Given a no  $N$ . Find all primes from 1 to  $N$ .

$N = 10: [1 \ 10] : 2, 3, 5, 7$  ✓

$N=20$  :  $[1 \ 20]$  :  $2, 3, 5, 7, 11, 13, 17, 19$

BF: for every  $n_i$  from 1 to  $N$ , check prime or not.

```
for(int i=1; i<=N; i++) {
```

```

if ( checkPrime ( i ) ) {
    |
    | print ( i )
    |
}

```

$N = 50$

TC:  $O(N\sqrt{N})$   
SC:  $O(1)$

F F i

~~0~~ ~~1~~ 2 3 ~~4~~ ~~5~~ ~~6~~ ~~7~~

~~8~~ ~~9~~ ~~10~~ ~~11~~ ~~12~~ ~~13~~ ~~14~~

~~15~~ ~~16~~ ~~17~~ ~~18~~ ~~19~~ ~~20~~ ~~21~~

~~22~~ ~~23~~ ~~24~~ ~~25~~ ~~26~~ ~~27~~ ~~28~~

~~29~~ ~~30~~ ~~31~~ ~~32~~ ~~33~~ ~~34~~ ~~35~~

~~36~~ ~~37~~ ~~38~~ ~~39~~ ~~40~~ ~~41~~ ~~42~~

~~43~~ ~~44~~ ~~45~~ ~~46~~ ~~47~~ ~~48~~ ~~49~~

~~50~~

$$\begin{aligned} i &\neq i \\ i = 2 &: 4 \\ i = 3 &: 9 \\ i = 5 &: 25 \\ i = 7 &: 49 \end{aligned}$$
$$\begin{array}{l} i \quad \supset i \\ 2i \quad \supset i \\ 3i \quad \supset i \\ 4i \\ 5i \\ \vdots \\ N \end{array}$$
$$\dot{\lambda} = 7$$
$$\begin{array}{l} 2 \times 7 \quad (2) \\ 3 \times 7 \quad (3) \\ 4 \times 7 \quad (2) \\ 5 \times 7 \quad (5) \\ 6 \times 7 \quad (2) \\ 7 \times 7 \end{array}$$
$$\frac{i}{\perp} \\ i \neq i$$
$$N = 100$$

Pseudo Code :-

Given N

bool p[N+1] = {T}

p[0] = F, p[1] = F

for(int i=2; i <= N; i++) {

if(p[i] == T) {

// i is prime, all multiples of i

// are non-prime

for(int j=2\*i; j <= N; j+=i) {

p[j] = false

}

}

}

i

# iterations

2

N/2 ✓

3

N/3

5

N/5

7

N/7 ✓

11

N/11

...

P

N/P

$$\# \text{ iterations} = \frac{N}{2} + \frac{N}{3} + \frac{N}{5} + \frac{N}{7} + \dots + \frac{N}{P}$$

$$N \left[ \frac{1}{2} + \frac{1}{3} + \frac{1}{5} + \frac{1}{7} + \frac{1}{11} + \dots + \frac{1}{P} \right]$$

Sum of Reciprocals of prime  $\rightarrow$  last  $P < N$

$$N [\log(\log N)] \checkmark \approx \underline{\underline{N}}$$

$$N = 2^{30} \quad \log_2 N = \log_2 2^{30} = 30$$

$$\log_2 30 = 4.5$$

$$i = 2 \quad 4$$

$$i = 3 \quad 9$$

$$i = 5 \quad 25$$

$$i = 7 \quad 49$$

$$i = 11 \quad 121$$

$$\begin{array}{cc} N & \checkmark \\ i & j \\ \sqrt{N} & \sqrt{N} \end{array}$$

$$(\sqrt{N}+i) \quad \frac{(\sqrt{N}+1)^2}{N+1+2\sqrt{N}} > N$$

instead of starting from  $2 \times i$ , we can start from  $i \times i$

$$TC \approx O(N)$$

$$SC : O(N)$$

$$i - \sqrt{N}$$

$$i = \sqrt{N} + 1 \quad (\sqrt{N} + 1)^2 > N$$

Q) Given N. Find Smallest prime factor for all numbers from 2 to N.  
 $\text{SPF}$

Eg  $N=10$

$\begin{array}{c} 2 \\ \hline \downarrow \\ 2 \end{array}$	$\begin{array}{c} 3 \\ \hline \downarrow \\ 3 \end{array}$	$\begin{array}{c} 4 \\ \hline \downarrow \\ 2 \end{array}$	$\begin{array}{c} 5 \\ \hline \downarrow \\ 5 \end{array}$	$\begin{array}{c} 6 \\ \hline \downarrow \\ 2 \end{array}$	$\begin{array}{c} 7 \\ \hline \downarrow \\ 7 \end{array}$	$\begin{array}{c} 8 \\ \hline \downarrow \\ 2 \end{array}$	$\begin{array}{c} 9 \\ \hline \downarrow \\ 3 \end{array}$	$\begin{array}{c} 10 \\ \hline \downarrow \\ 2 \end{array}$
$\text{SPF}[ ]: 2$	$3$	$2$	$5$	$2$	$7$	$2$	$3$	$2$

$\text{SPF}[ ]$

2	3	4	5	6	7
2	3	<del>4</del>	5	<del>6</del>	7
		2		2	

  

8	9	10	11	12	13	14
<del>8</del>	<del>9</del>	<del>10</del>	11	<del>12</del>	13	<del>14</del>
2	3	2		2		2

  

15	16	17	18	19	20	21
<del>15</del>	<del>16</del>	17	<del>18</del>	19	<del>20</del>	<del>21</del>
3	2		2		2	3

  

22	23	24	25
<del>22</del>	23	<del>24</del>	<del>25</del>
2		2	5

$\text{SPF}[j] = \min(\text{SPF}[j], i)$

if ( $\text{SPF}[j] == j$ ) {  
 $\text{SPF}[j] = i$   
}

Pseudo Code :-

```

int SPF[N+1];

for (int i = 2; i <= N; i++) {
    SPF[i] = i;
}

for (int i = 2; i * i <= N; i++) {
    if (SPF[i] == i) { // i is prime
        for (int j = i * i; j <= N; j += i) {
            if (SPF[j] == j) {
                SPF[j] = i;
            }
        }
    }
}

return SPF;

```

$TC \simeq O(N)$   
 $SC \simeq O(1)$

Break for 5 min :- Prime factorization  
Count No of divisors

Prime Factorization :- Any number can be represented as product of power of prime nos.

2	72
2	36
2	18
3	9
3	3
	1

$$72 = 2^3 \times 3^2 \quad \checkmark$$

$$\text{No of divisors / factors} = (3+1)(2+1) = 12 \quad \checkmark$$

How ??

$$72 = 2^{\textcircled{3}} \cdot 3^{\textcircled{2}}$$

$$\downarrow \quad \hookrightarrow \quad [3^0, 3^1, 3^2]$$

$$[2^0, 2^1, 2^2, 2^3] \quad \hookrightarrow \quad {}^3C_1 = 12$$

$$2^0 \times 3^0 : 1 \quad \checkmark$$

$$2^0 \times 3^1 : 3 \quad \checkmark$$

$$2^0 \times 3^2 : 9 \quad \checkmark$$

$$2^1 \times 3^0 : 2 \quad \checkmark$$

$$2^1 \times 3^1 : 6 \quad \checkmark$$

$$2^1 \times 3^2 : 18 \quad \checkmark$$

$${}^4C_1 \quad \times$$

$$2^2 \times 3^0 : 4 \quad \checkmark$$

$$2^2 \times 3^1 : 12 \quad \checkmark$$

$$2^2 \times 3^2 : 36 \quad \checkmark$$

$$2^3 \times 3^0 : 8 \quad \checkmark$$

$$2^3 \times 3^1 : 24 \quad \checkmark$$

$$2^3 \times 3^2 : 72 \quad \checkmark$$

$\Rightarrow \textcircled{12}$

$$N = 600 \quad \checkmark$$

2	600
2	300
2	150
3	75
3	15
5	5
	1

$$\Rightarrow 2^3 \times 3^2 \times 5^1$$

$$\Rightarrow (3+1)(2+1)(1+1)$$

$$(4)(3)(2)$$

$$24 \quad \checkmark$$

Generalization :-

$$N = p_1^{x_1} \cdot p_2^{x_2} \cdot p_3^{x_3} \dots p_y^{x_y}$$

$$\text{No of factors} = (x_1 + 1)(x_2 + 1)(x_3 + 1) \dots (x_y + 1)$$

Q) Given N. Count no of divisors through prime factorization method.

$$1 * 4 * (3) * (2) = 24 \checkmark$$

$$\text{ans} = 1 * (\text{cnt} + 1)$$

$$N = 360$$

2	360	SPF(360) = 2
2	180	cnt = + 1 <u>2</u>
2	90	
3	45	x = 3, cnt = 0 + 2 <u>✓</u>
3	15	
5	5	x = 5, cnt = 0 + 1
1	1	

→ SPF[N+1] ✓

// assume that you have built the SPF array → N

while (N > 1) {

    x = SPF[N], cnt = 0

    while (N % x == 0) {

        N = N / x

        cnt ++

    }

    ans = ans \* (cnt + 1)

}

→ log N

$$TC : O(N + \log N)$$



Q) Given N. For every number from 1 to N, Get No. of factors.

N = 10 : 

1	2	3	4	5	6	7	8	9	10
1	2	2	3	2	4	2	4	3	4

 ✓

Build your SPF — N } N } (2nd Que) ✓

int cnt[N+1]; // O/P

cnt[1] = 1

```
for (int i = 2; i <= N; i++) {
    int tmp = i, ans = 1;
    while (tmp > 1) {
        int x = SPF[tmp], cnt = 0;
        while (tmp % x == 0) {
            tmp = tmp / x;
            cnt++;
        }
        ans = ans * (cnt + 1);
    }
    cnt[i] = ans;
}
return cnt;
```

log N

N log N

TC :  $O(N + N \log N) = N \log N$   
 SC :  $O(N)$  (because of SPF)

