Recursion

Recursion - Introduction
Function Call Tracing
Problems
Time Complexity For Recursive Function

Recursion - Introduction

- A function calling itself.
- Solving a problem using a smaller version of the same problems (subproblem).

3 Steps to implement Recursive code

```
    Assumption: Decide what your function does and assume that it does.
    Main Logic: Break down and solve problems using sub problems.
    Base Condition: Decide when your function should stop.
```

Function Call Tracing

```
Int add(N, M){
    return (N+M)
}

Int x = add(10, 20)
Int y = Square(x);
Int z = twice(y);
print(z);
}

Int twice(N){
    Return 2*N
}

main(){
    main(){
    main(){
    print(twice square(add(10, 20)))
    }
```

Add -> square -> twice 30 -> 900 -> 1800

Problems

Problem 1: Given N, find the factorial of N!

```
N! = 1 * & * & * 4 * . . . * N

5! = 1 * 2 * 3 * 4 * 5

= 120
```

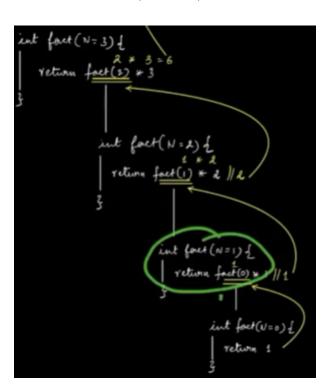
Assumption:

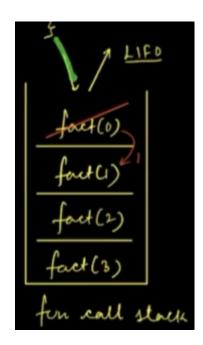
Int factorial(int N){
.....
}

Main Logic: Try to break the problem
N! = 1*2*3*4*5*6*7*.....*(N-1)*N

Base Condition:

When N==1, return 1;





Problem 2: Fibonacci Series / Sequence
Write a function to compute Nth Fibonacci

 Given N, the function will calculate Nth Fib and return int fib(int N){

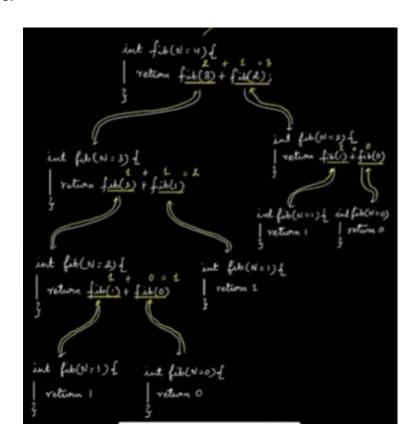
}

2) Main Logic,

return fib(N-1)+fib(N-2);

3) Base Condition

Function Call Trace:



Problem 3: Given 2 integers a and N. Find a^N using Recursion.

1) Assumption: given a and N, the function will calculate and return a^N. Int pow(int a, int N){

.....

2) Main Logic:

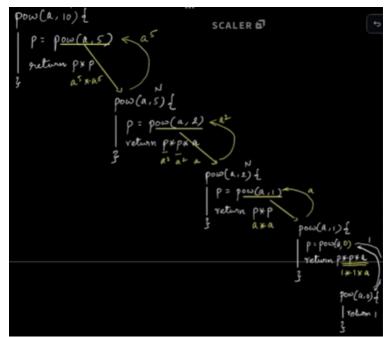
}

```
a^N = a*a*a*.....*a (N times)
Approach 1) a*pow(a, N-1);
Approach 2)
```

3) Base Condition:

```
if(N == 0){
     return 1;
}
```

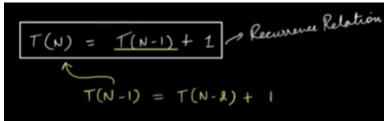
Function Call Stack:



Time Complexity For Recursive Function

Use the substitution method to find the recursive Functions.

1. For Factorial:



2. For Pow Function

$$T(N) = T(N|2) + T(N|2) + 1$$

$$T(N) = 2[T(N/2)] + 1;$$

$$T(N/2) = 2[2T(N/4) + 1] + 1 = 4T(N/4) + 3$$

$$= 8T(N/8) + 7$$

$$\dots$$
Generalized Eqn : $T(N) = 2^{K}[T(N/2^{K}) + (2^{K-1})]$

$$N/2^{K} = 1$$

$$N = 2^{K}$$

$$K = logN$$
T.c = $O(logN)$

Generalized Definition Of Time Complexity

Time Complexity = No of Function Calls * Time Taken in a Single Function Call.