

Verification Report

AHB-Lite Protocol

Introduction to the Device-Under-Test (DUT)

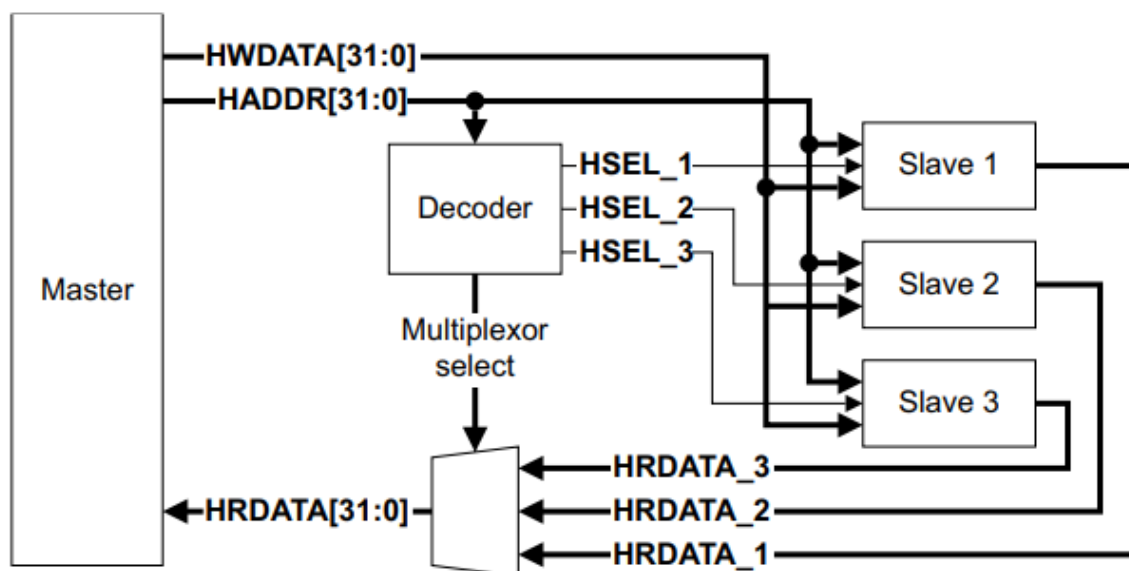
AMBA AHB-Lite addresses the requirements of high-performance synthesizable designs. It is a bus interface that supports a single bus master and provides high-bandwidth operation.

AHB-Lite implements the features required for high-performance, high clock frequency systems including:

- burst transfers
- Single-clock edge operation
- Non-tristate implementation
- Wide data bus configurations, 64, 128, 256, 512, and 1024 bits.

The most common AHB-Lite slaves are internal memory devices, external memory interfaces, and high bandwidth peripherals. Although low-bandwidth peripherals can be included as AHB-Lite slaves, for system performance reasons they typically reside on the AMBA Advanced Peripheral Bus (APB). Bridging between this higher level of bus and APB is done using an AHB-Lite slave, known as an APB bridge.

The figure shown below describes a single master AHB-Lite system design with one AHB-Lite master and three AHB-Lite slaves. The bus interconnect logic consists of one address decoder and a slave-to-master multiplexor. The decoder monitors the address from the master so that the appropriate slave is selected and the multiplexor routes the corresponding slave output data back to the master.



Verification Plan

| No. | Feature | Test Description | Ref. | Type | Expected Outcome | Result | Comments |
|-----|------------------------------------|--|-----------------------------|------|--|--------|--|
| 1. | Enable the value of HWRITE | HWRITE controls the direction of data transfer to or from the master | IHI0033.pdf/Sec3.1/pg3-2 | T | At positive clock edge, when HWRITE is HIGH, it indicates a write transfer and the master broadcasts data on the write data bus, HWDATA[31:0] | | The simplest transfer is one with no wait states, so the transfer consists of one address cycle and one data cycle |
| 2. | Disable the value of HWRITE | HWRITE controls the direction of data transfer to or from the master | IHI0033.pdf/Sec3.1/pg3-2 | T | At positive clock edge, when HWRITE is LOW, a read transfer is performed and the slave must generate the data on the read data bus, HRDATA[31:0] . | | The simplest transfer is one with no wait states, so the transfer consists of one address cycle and one data cycle |
| 3. | Enable HMASTLOCK signal | If the master requires locked accesses then it must assert the HMASTLOCK signal | IHI0033.pdf/Sec3.3/pg3-7 | A | At positive clock edge, when it is high, it indicates to any slave that the current transfer sequence is indivisible and must therefore be processed before any other transactions are processed | | Typically the locked transfer is used to maintain the integrity of a semaphore, by ensuring that the slave does not perform other operations between the read and write phases of a microprocessor SWP instruction |
| 4. | Bus termination | At the positive edge of the clock, after a | IHI0033.pdf/Sec3.5.1/pg3-10 | T | During an undefined length burst, INCR, the | | The protocol does not permit a master to end a |

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| | after a BUSY transfer for an Undefined Length Burst | burst has started, the master uses BUSY transfers if it requires more time before continuing with the next transfer in the burst. | | | master might insert BUSY transfers at the positive edge of the clock and then decide that no more data transfers are required. Under these circumstances, it is acceptable for the master to then perform a NONSEQ or IDLE transfer at the next positive edge that then effectively terminates the undefined length burst | | burst with a BUSY transfer for fixed length bursts of type incrementing or wrapping |
| 5. | Bus termination after a BUSY transfer for a SINGLE length | After a burst has started, the master uses BUSY transfers if it requires more time before continuing with the next transfer in the burst. | IHI0033.pdf/Sec3.5 .1/pg3-10 | T | At the positive edge of the clock, the fixed length burst types must terminate with a SEQ transfer | | The master is not permitted to perform a BUSY transfer immediately after a SINGLE burst. SINGLE bursts must be followed by an IDLE transfer or a NONSEQ transfer. |
| 6. | Change in transfer type to NONSEQ From IDLE Transfer | Slaves use HREADY to insert wait states if they require more time to provide or sample the data. During a waited transfer, the master is restricted to what changes it can make | IHI0033.pdf/Sec3.6 .1/pg3-16 | T | At positive edge of the clock, when the HTRANS transfer type changes to NONSEQ the master must keep HTRANS constant, until HREADY is HIGH at the second positive clock edge | | - |

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| | | to the transfer type and address. During a waited transfer, the master is permitted to change the transfer type from IDLE to NONSEQ | | | | | |
| 7. | Change in transfer type to SEQ from BUSY Transfer when Length of Burst is fixed | During a waited transfer for a fixed length burst, the master is permitted to change the transfer type from BUSY to SEQ | IHI0033.pdf/Sec3.6 .1/pg3-17 | T | At positive edge of the clock, when the HTRANS transfer type changes to SEQ the master must keep HTRANS constant, until HREADY is HIGH at the next positive edge | | Because BUSY transfers must only be inserted between successive beats of a burst, this does not apply to SINGLE bursts |
| 8. | Change in transfer type to any other type from BUSY Transfer when Length of Burst is Undefined | During a waited transfer for an undefined length burst, INCR, the master is permitted to change from BUSY to any other transfer type, when HREADY is LOW | IHI0033.pdf/Sec3.6 .1/pg3-18 | T | At the positive edge of the clock, the burst continues if a SEQ transfer is performed but terminates if an IDLE or NONSEQ transfer is performed | | - |
| 9. | Change in address during IDLE Transfer | When the slave is requesting wait states, the master can only change the address once except in case of an IDLE Transfer. During a waited transfer, the | IHI0033.pdf/Sec3.6 .2/pg3-19 | T | When the HTRANS transfer type changes to NONSEQ at the positive edge of clock, the master must keep the address constant, until HREADY is HIGH at the next clock edge | | - |

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| | | master is permitted to change the address for IDLE transfers | | | | | |
| 10. | Change in address after an ERROR response | When the slave is requesting wait states, the master can only change the address once except after an ERROR response | IHI0033.pdf/Sec3.6.2/pg3-20 | T | During a waited transfer, if the slave responds with an ERROR response then the master is permitted to change the address when HREADY is LOW at the positive edge of the clock | | - |
| 11. | Transfer done | After a master has started a transfer, the slave controls how the transfer progresses. A master cannot cancel a transfer after it has commenced. A slave must provide a response that indicates the status of the transfer when it is accessed. The transfer status is provided by the HRESP signal | IHI0033.pdf/Sec5.1.1/pg5-3 | A | At the positive edge of the clock, a successful completed transfer is signalled when HREADY is HIGH and HRESP is OKAY (0) | | - |
| 12. | Transfer pending | A slave uses HREADY to insert the appropriate number | IHI0033.pdf/Sec5.1.2/pg5-3 | A | In pending transfer, at the positive edge of the clock, HREADY is | | - |

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| | | of wait states into the data phase of the transfer. | | | LOW and HRESP is OKAY (0). A transfer then completes with HREADY HIGH and an OKAY response at the next edge | | |
| 13. | ERROR response | A slave uses the ERROR response to indicate some form of error condition with the associated transfer. Usually this denotes a protection error such as an attempt to write to a read-only memory location | IHI0033.pdf/Sec5.1.3/pg5-3 | A | The slave will give an ERROR response in the first cycle with HREADY LOW and HRESP ERROR (1). In the next cycle, HREADY is driven HIGH and HRESP remains HIGH to indicate error | | The two-cycle response is required because of the pipelined nature of the bus. By the time a slave starts to issue an ERROR response then the address for the following transfer has already been broadcast onto the bus. The two-cycle response provides sufficient time for the master to cancel this next access and drive HTRANS [1:0] to IDLE before the start of the next transfer |

Explanation of Different Fields

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| No. | The serial number of the test. |
| Feature | The feature which the current test is verifying in full or partially. The feature is usually on the abstraction level of a user. |
| Test Description | A detailed description of the test case being performed. You can be as verbose as you want. |
| Ref. | Reference to the section in the related standard document. The section number as well as page numbers should be described here. |
| Type | Type of the test. Whether the test is an assertion (A) or a transaction (T) type. |
| Result | Pass (P) or Fail (F). |

Comments Any other comments about the test or its results that you want to mention.