

# Intro

15 January 2018 20:33

$$SNR = \frac{\text{signal power}}{\text{noise power}}$$

External noise - interference

Internal noise - channel properties

Efficiency:

number of bits transmitted of unit: power, time, and bandwidth

Reliability:

expressed in terms of SNR or probability of error

## Information Theory:

- Shannon capacity formula
  - $C = W \log(1 + SNR) \text{ bps}$ 
    - $W(\text{Hz})$  – channel bandwidth
    - C - channel capacity

# Probability

## Normal Distribution

$$f_X(x) = \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-m)^2}{2\sigma^2}} \quad \text{for } -\infty < x < \infty$$

$$F_X(x) = \frac{1}{\sqrt{2\pi}\sigma} \int_{-\infty}^x e^{-\frac{(y-m)^2}{2\sigma^2}} dy$$

## Uniform Distribution

$$f_X(x) = \begin{cases} \frac{1}{b-a} & a \leq x \leq b \\ 0 & \text{elsewhere} \end{cases}$$
$$F_X(x) = \begin{cases} 0 & x < a \\ \frac{x-a}{b-a} & a \leq x \leq b \\ 1 & x > b \end{cases}$$

$$E[X] = \frac{a+b}{2}$$
$$\sigma_X^2 = \frac{(b-a)^2}{12}$$

## Joint Distribution

Joint distribution function for two random variables  $X$  and  $Y$

$$F_{XY}(x, y) = P(X \leq x, Y \leq y)$$

Joint probability density function

$$f_{XY}(x, y) = \frac{\partial^2 F_{XY}(x, y)}{\partial x \partial y}$$

Properties

$$1) \quad F_{XY}(\infty, \infty) = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} f_{XY}(u, v) du dv = 1$$

$$2) \quad f_X(x) = \int_{y=-\infty}^{\infty} f_{XY}(x, y) dy$$

$$3) \quad f_Y(y) = \int_{x=-\infty}^{\infty} f_{XY}(x, y) dx$$

$$4) \quad X, Y \text{ are independent} \Leftrightarrow f_{XY}(x, y) = f_X(x) f_Y(y)$$

$$5) \quad X, Y \text{ are uncorrelated} \Leftrightarrow E[XY] = E[X]E[Y]$$

If **normal** - uncorrelated  $\rightarrow$  independent

# Random Process

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Time-varying function which assigns to each outcome,  $s$ , a function of time

$$X(t, s_j)$$

Fixed  $s_j$  - sample function - signal

Fixed  $t$  - random variable

$$\begin{aligned} - \mu_X(t) &= E[X(t)] \\ - \text{mean varies with time} \end{aligned}$$

**Autocorrelation** - measures correlation between two samples

$$R_X(t_1, t_2) = E[X(t_1)X(t_2)] = \int_{-\infty}^{\infty} \int_{-\infty}^{\infty} xy f_X(x, y; t_1, t_2) dx dy$$

$R_X(\text{constant})$  - difference between  $t_1$  and  $t_2$  is a constant

$$\text{AutoCorr} = R_{ff}(t_1, t_2) - \mu_f(t_1)\mu_f(t_2)$$

For gaussian

- uncorrelated  $\rightarrow$  independent

## Stationary - homogenous

-check these as well for  $t_1-t_2$

$$R_X(0) = E[X^2(t)]$$

$$R_X(t) = R_X(-t)$$

$$|R_X(t)| \leq R_X(0)$$

### Strict-Sense Stationary

- all underlying  $n$  RVs time **shifted by a constant**
- previous joint PDF = new PDF
- implies WSS

1st order

$f_X(x; t)$  is invariant over time

2nd order

$$R_X(t_1, t_2) = R_X(t_2 - t_1)$$

### Wide-sense Stationary

both mean and autocorrelation are time invariant

$$\mu_X(t) = \mu_X$$

$$R_X(t, t + \tau) = R_X(\tau)$$

Asked about  $F(X(t_1), X(t_2))$

$$\begin{aligned} E \left[ F(X(t_1), X(t_2))^2 \right] \\ = G(R_X(0), R_X(\tau)), \text{ find } \tau \end{aligned}$$

$$\text{PSD} - S_X(f) = \int_{-\infty}^{\infty} R_X(\tau) e^{-j2\pi f\tau} d\tau \geq 0$$

Passing through an LTI filter

$y$  = convolution of  $x$  and impulse response  $h$

Mean

$$\mu_Y(t) = \mu_X(t) * h(t)$$

$$\mu_Y(t) = \mu_X * h(0) \quad (\text{WSS})$$

Autocorrelation

$$R_Y(t, u) = h(t) * [h(u) * R_X(t, u)]$$

$$R_Y(\tau) = h(\tau) * [h(-\tau) * R_X(\tau)] \quad (\text{WSS})$$

$$R_Y(\tau) = h(\tau) * [h^*(-\tau) * R_X(\tau)] \quad (\text{WSS} - \text{complex } X(t))$$

Autocorrelation

$$S_Y(f) = |H(f)|^2 S_X(f)$$

## Baseband and Passband

20 May 2018 18:09

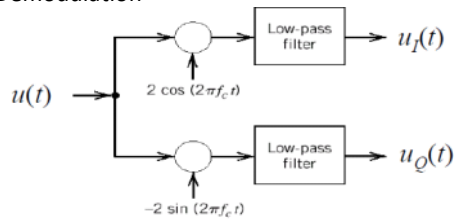
$$E = \int_{-\infty}^{+\infty} |s(t)|^2 dt = \int_{-\infty}^{+\infty} |S(f)|^2 df$$

$$P = \frac{1}{T} \int_{-T/2}^{+T/2} |s(t)|^2 dt$$

Baseband to passband - send two baseband in parallel

$$u(t) = u_I(t) \cos(2\pi f_c t) - u_Q(t) \sin(2\pi f_c t)$$

Demodulation



physical baseband channel - real baseband signals

physical passband channel - complex baseband signals

**Hilbert transform** - if FT of  $x(t)$  is  $X(f)$

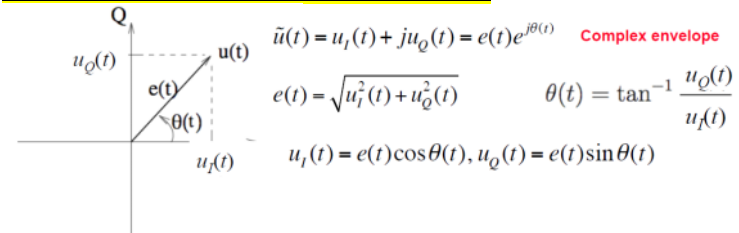
FT of  $\hat{x}(t) = -j \operatorname{sgn}(f) X(f)$

Pre-envelope

$$u_+(t) = u(t) + j\hat{u}(t)$$

$$u_-(t) = u(t) - j\hat{u}(t)$$

## Passband - complex envelope



$$u(t) = u_I(t) \cos(2\pi f_c t) - u_Q(t) \sin(2\pi f_c t) \quad \text{In terms of I and Q components}$$

$$u(t) = e(t) \cos(2\pi f_c t + \theta(t)) \quad \text{In terms of envelope and phase}$$

$$u(t) = \operatorname{Re} \{ \tilde{u}(t) e^{j2\pi f_c t} \} \quad \text{In terms of complex envelope}$$

# Noise in communications systems

20 May 2018 18:10

## White Noise

$$S_N(f) = \frac{N_0}{2} \quad -\infty < f < \infty \quad - \text{flat PSD} - \text{is a random process}$$

$$P_N = \int_{f_1}^{f_2} S_N(f) df$$

$$R_N(\tau) = \frac{N_0}{2} \delta(\tau)$$

Power of zero mean Gaussian:  
= Variance

White - samples are uncorrelated  
- since Gaussian -> independent

process is Gaussian if pdf of samples follows  
Gaussian distribution  
-CLT

N Power in LP Filter =  $N_0 B$

N Power in BP Filter =  $2N_0 B$

$n(t)$  in BP can be split into  $n_I(t)$  and  $n_Q(t)$  - (both base-band random processes)

- If noise  $n(t)$  has zero mean, then so do  $n_I(t)$  and  $n_Q(t)$ .
- If noise  $n(t)$  is Gaussian, then so are  $n_I(t)$  and  $n_Q(t)$ .
- If noise  $n(t)$  is stationary, then so are  $n_I(t)$  and  $n_Q(t)$ .
- If noise  $n(t)$  is Gaussian and its power spectral density  $S(f)$  is symmetric with respect to central frequency  $f_c$ , then  $n_I(t)$  and  $n_Q(t)$  are statistically independent.
- $n_I(t)$  and  $n_Q(t)$  have the same variance (i.e., same power) as  $n(t)$ .

$$S_{N_I}(f) = S_{N_Q}(f) = N_0 \quad |f| \leq B$$

$$P_{N_I} = P_{N_Q} = 2N_0 B$$

# Noise performance of DSB

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$$SNR_0 = \frac{\text{avg power of message signal at receiver}}{\text{avg power of noise}}$$

$$P = E\{m^2(t)\} = P \text{ of message}$$

$$SNR \text{ (dB)} = 10 \log_{10}(SNR)$$

increase  $P_T$  -  $P_S$  increases

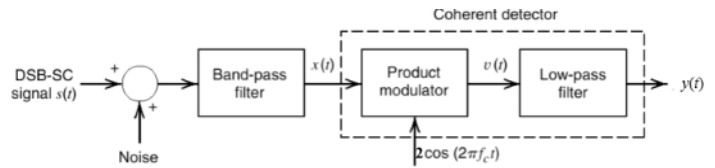
- uses same  $P_T$  to compare schemes
- use  $SNR_{\text{baseband}}$  as ref.

## Baseband SNR

- $P_S = P_T = P$
- $P_N = 2W * \frac{N_0}{2} = WN_0$
- $SNR_{\text{baseband}} = \frac{P_T}{N_0 W}$

## DSB-SC SNR

- received signal
  - $x(t) = Am(t) \cos(2\pi f_c t) + n_I(t) \cos(2\pi f_c t) - n_Q(t) \sin(2\pi f_c t)$



- multiply with  $2\cos(2\pi f_c t)$  and pass through LPF
  - $y(t) = Am(t) + n_I(t)$ 
    - $n_Q(t)$  removed!

- $P_S = E\{A^2 m^2(t)\} = A^2 P$
- $P_N = 2W * N_0 = 2N_0 W$

$$SNR = \frac{A^2 P}{2N_0 W}$$

$$P_T = E\{A^2 m^2(t) \cos^2(2\pi f_c t)\} = \frac{A^2 P}{2}$$

$$SNR_{DSB-SC} = \frac{P_T}{N_0 W} = SNR_{\text{baseband}}$$

- DSB-SC has same SNR, but less interference as at higher freq. transmission

# Noise performance of SSB and AM

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## SSB SNR

- modulated signal:  $s(t) = \frac{A}{2} m(t) \cos(2\pi f_c t) - \frac{A}{2} \hat{m}(t) \sin(2\pi f_c t)$ 
  - $m(t)$  and  $\hat{m}(t)$  have power  $P$
- $P_T = \frac{A^2 P}{4}$
- received signal:  $y(t) = \frac{A}{2} m(t) + n_I(t)$
- $P_S = \frac{A^2 P}{4}$
- $P_N = (W + W) * \frac{N_0}{2} = N_0 W$
- $SNR_{SSB} = \frac{A^2 P}{4N_0 W} = \frac{P_T}{N_0 W} = SNR_{baseband}$
- same SNR as baseband - less interference
- same SNR as DSB-SC - **req. half the bandwidth**

$$SSB = m(t) \cos \omega_c t \mp \hat{m}(t) \sin \omega_c t$$

**USB MINUS**

**LSB PLUS**

**USB - real M for real S, Im M for S as well**

## Full AM SNR

$$\mu = \frac{m_p}{A}$$

$$P_T = \frac{A^2 + P}{2}$$

$$\text{received signal after LPF: } y(t) = A + m(t) + n_I(t)$$

$$P_S = P$$

$$P_N = 4W * \frac{N_0}{2}$$

$$SNR_{AM} = \frac{P}{2N_0 W}$$

$$SNR_{baseband} = \frac{P_T}{N_0 W} = \frac{A^2 + P}{2N_0 W}$$

$$SNR_{AM} = \frac{P}{A^2 + P} SNR_{baseband}$$

$$\frac{P}{A^2 + P} < 1$$

□ lower SNR than baseband for coherent AM

## Envelope detection:

- RC circuit
- if small noise, same as above
- large noise - signal is lost

## Summary of AM SNR

24 May 2018 00:06

(De-) Modulation Format	Output SNR	Transmitted Power	Baseband Reference SNR	Figure of Merit (= Output SNR / Reference SNR)
AM Coherent Detection	$\frac{P}{2N_0W}$	$\frac{A^2 + P}{2}$	$\frac{A^2 + P}{2N_0W}$	$\frac{P}{A^2 + P} < 1$
DSB-SC Coherent Detection	$\frac{A^2P}{2N_0W}$	$\frac{A^2P}{2}$	$\frac{A^2P}{2N_0W}$	1
SSB Coherent Detection	$\frac{A^2P}{4N_0W}$	$\frac{A^2P}{4}$	$\frac{A^2P}{4N_0W}$	1
AM Envelope Detection (Small Noise)	$\frac{P}{2N_0W}$	$\frac{A^2 + P}{2}$	$\frac{A^2 + P}{2N_0W}$	$\frac{P}{A^2 + P} < 1$
AM Envelope Detection (Large Noise)	Poor	$\frac{A^2 + P}{2}$	$\frac{A^2 + P}{2N_0W}$	Poor

$A$ : carrier amplitude,  $P$ : power of message signal,  $N_0$ : single-sided PSD of noise,  $W$ : message bandwidth.



# Noise performance of FM

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$$f_i(t) = f_c + k_f m(t)$$

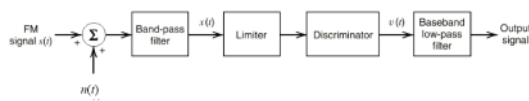
$$s(t) = A \cos \left[ 2\pi f_c t + 2\pi k_f \int_0^t m(\tau) d\tau \right]$$

$$\Delta f = k_f m_p$$

$$\beta = \frac{\Delta f}{W} \quad W = \text{message bandwidth}$$

Transmission Bandwidth

$$B_T = 2(\Delta f + W) = 2W(\beta + 1)$$



$$\text{BP Filter: } \pm f_c \pm \frac{B_T}{2}$$

Limiter - remove amplitude variations

## FM SNR

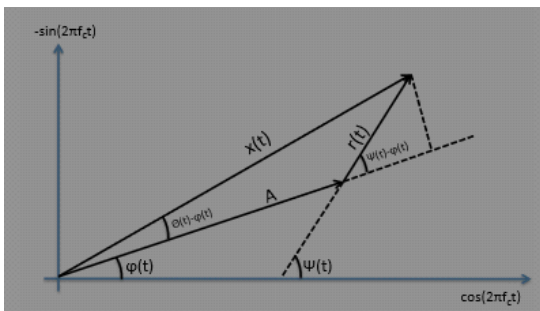
- $P_S = k_f^2 P$
- $P_N = \frac{2N_0 W^3}{3A^2}$
- $SNR_O = \frac{3A^2 k_f^2 P}{2N_0 W^3}$
- $P_T = \frac{A^2}{2}$
- $SNR_{baseband} = \frac{P_T}{N_0 W} = \frac{A^2}{2N_0 W}$

$$SNR_{FM} = \frac{3k_f^2 P}{W^2} SNR_{baseband}$$

$$\text{since } \frac{k_f}{W} = \frac{\beta}{m_p}$$

$$SNR_{FM} = 3\beta^2 \frac{P}{m_p^2} SNR_{baseband} \propto \beta^2 SNR_{baseband}$$

- doubling bandwidth,  $B_T$ , gives quadratic increase in SNR



$$\begin{aligned} x(t) &= s(t) + n(t), \\ &= A \cos [2\pi f_c t + \phi(t)] + r(t) \cos [2\pi f_c t + \psi(t)]. \end{aligned}$$

$$v(t) = \frac{1}{2\pi} \frac{d\theta(t)}{dt} \simeq k_f m(t) + n_d(t).$$

$$n_d(t) = \frac{1}{2\pi A} \frac{dn_\phi(t)}{dt}$$

$$S_{N_o}(f) = \frac{f^2}{A^2} N_0 \quad |f| \leq W$$

## Pre/de-emphasis for FM

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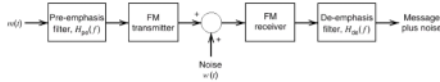
Threshold:

- carrier-to-noise ratio
  - $\rho = \frac{A^2}{2N_0 B_T}$
  - below 10, FM receiver breaks
  - small changes to phase can result in changes of  $2\pi$

### Improve Output SNR:

- noise is proportional to square of frequency at LPF output
- increase message amplitude at higher frequency
  - increase SNR at high freq.

### Pre-emphasis and De-emphasis



- $H_{pre}(f)$ : artificially emphasizes high frequency components of the message prior to modulation (before noise is introduced).
- $H_{de}(f)$ : de-emphasizes high frequency components at the receiver, and restore the original PSD of the message.
- In theory,  $H_{pre}(f) \propto f$ ,  $H_{de}(f) \propto 1/f$
- This can improve output SNR by around 13 dB.

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### Improvement Factor:

$$I = \frac{P_{Nold}}{P_{Nnew}} = \frac{2W^3}{3 \int_{-W}^W f^2 |H_{de}(f)|^2 df}$$

### Example

- (a) Pre-emphasis filter

$$H_{pre}(f) = 1 + j \frac{f}{f_0}$$

$$f_0 = 1 / (2\pi rC), \quad R \ll r, \quad 2\pi f rC \ll 1$$

- (b) De-emphasis filter

$$H_{de}(f) = \frac{1}{1 + jf / f_0}$$

- Improvement

$$I = \frac{2W^3}{3 \int_{-W}^W f^2 / (1 + f^2 / f_0^2) df}$$

$$= \frac{(W / f_0)^3}{3[(W / f_0) - \tan^{-1}(W / f_0)]}$$

- (Full) AM**: SNR performance is 4.8 dB worse than a baseband system, transmission bandwidth is  $B_T = 2W$
- DSB**: SNR performance is identical to a baseband system, transmission bandwidth is  $B_T = 2W$
- SSB**: SNR performance is again identical, transmission bandwidth is only  $B_T = W$
- FM**: SNR performance is 15.7 dB better than a baseband system, transmission bandwidth is  $B_T = 2(\beta + 1)W = 12W$  (with pre- and de-emphasis SNR performance is increased by about 13 dB with the same transmission bandwidth).

# Digital representation of signals

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sample at  $f_s = 2W$

## Quantisation: Uniform

- difference between actual and taken =  $q$

$$f_Q(q) = \begin{cases} \frac{1}{\Delta}, & -\frac{\Delta}{2} \leq q \leq \frac{\Delta}{2} \\ 0, & \text{otherwise} \end{cases}$$

has variance  $P_N = \frac{\Delta^2}{12}$

$n$  bits

$2^n$  levels

$2^n \Delta$  dynamic range

$$\text{max error} = \frac{\Delta}{2}$$

$$n = \log_2 \frac{1}{\text{max\_error}}$$

$$m_p = 2^{n-1} \Delta$$

$$SNR_0 = \frac{P_S}{P_N} = \frac{12P}{\Delta^2} = \frac{12P2^{(n-1)2}}{m_p^2} = \frac{3P}{m_p^2} 2^{2n}$$

$$SNR_0(dB) = 6n + 10 \log_{10} \left( \frac{3P}{m_p^2} \right)$$

## Companding:

- compress signal
  - tighter grouping
  - reduces peak to average power ratio
- quantise using **uniform quantizer**
- expand signal
  - ideally - inverse of compression
- can increase SNR
  - **make output SNR insensitive to peak:avg power ratio**

# Baseband digital transmission

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Analog Communication - quality via SNR

Digital Communication - **quality via BER**  
- less susceptible to noise

## Matched filter

- filter on receiver that makes instantaneous power at sampling time as large as possible compared to noise component

impulse response of filter:

$$h_{opt}(t) = kg(T_b - t)$$

$T_b$  - symbol period

$g(t)$  - transmitter pulse shape

$k$  - gain

if  $h_{opt}(t)$

$$\eta_{max} = \frac{2E}{N_0} = SNR$$

- proportional to signal energy

for Rectangular pulse:

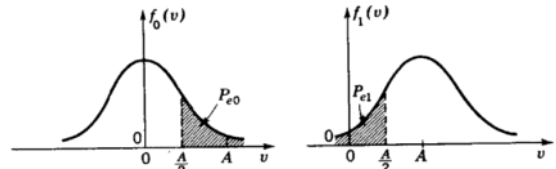
match with same pulse but scaled by  $k$   
equivalent: integrate

## n-ary PAM:

$$\frac{\log n \text{ bits per symbol}}{T \text{ secs per symbol}} = \frac{\log n}{T} \text{ bps}$$

## Error:

- binary PCM on-off
- noise is AWGN  $N \sim N(0, \sigma^2)$
- if 0, signal  $Y = N < \lambda$ 
  - $Y \sim N(0, \sigma^2)$
- if 1, signal  $Y = A + N > \lambda$ 
  - $Y \sim N(A, \sigma^2)$



$$P_e(\lambda) = p_1 P_{e1} + (1 - p_1) P_{e0}$$

Choose  $\lambda$  so that  $\frac{dP_e(\lambda)}{d\lambda} = 0$

$\rightarrow$  minimise  $P_e(\lambda) \rightarrow -\frac{\sigma^2}{A} \ln \frac{p_1}{1 - p_1} \rightarrow \lambda = \frac{A}{2}$  for equi probable

$$P_e = P_0 \int_{-\infty}^{\lambda} \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{x^2}{2\sigma^2}} dx + P_1 \int_{-\infty}^{\lambda} \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-A)^2}{2\sigma^2}} dx$$

$$\frac{dP_e}{d\lambda} = -P_0 \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(\lambda)^2}{2\sigma^2}} + P_1 \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(\lambda-A)^2}{2\sigma^2}} = 0$$

$$P_e = Q\left(\frac{A}{2} * \frac{1}{\sigma}\right)$$

$$E_{pulse} = \frac{A^2}{T_b} \rightarrow \text{avg energy per bit, } E_b = \frac{A^2}{2T_b}$$

$$P_N = \sigma^2 = \frac{N_0}{2T_b}$$

$$P_e = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$

- small increase in SNR - HUGE increase in reliability

# Digital modulation - coherent

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## ASK - A = 0, A

$$\underline{x(t) = A(t)\cos(2\pi f_c t) - A \in \{0, A\}}$$

$$\underline{y_o(t) = A(t) + n_i(t)}$$

$$E_{pulse} = \frac{A^2}{2T_b} \text{ (half from cosine)}$$

$$E_b = \frac{A^2}{4T_b}, \sigma^2 = \frac{N_0}{T_b}$$

$$P_{e,ASK} = Q\left(\frac{A}{2\sigma}\right) = Q\left(\sqrt{\frac{E_b}{N_0}}\right) - \text{same as baseband}$$

## PSK

BPSK - change phase of cosine by pi  
 - time by 1 or -1  
 - -1 - angle increases by pi

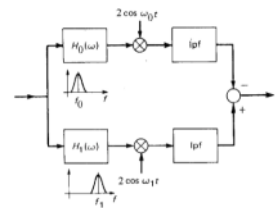
$$\underline{x(t) = A(t)\cos(2\pi f_c t) - A \in \{-A, A\}}$$

$$\underline{y_o(t) = A(t) + n_i(t)}$$

$$P_{e,PSK} = Q\left(\frac{A}{\sigma}\right) = Q\left(\sqrt{\frac{2E_b}{N_0}}\right) < \text{better than baseband}$$

## FSK

$$f = f_c + (0,1)$$



LPF output on each branch =  $\begin{cases} A + \text{noise} & \text{if symbol present} \\ \text{noise} & \text{if symbol not present} \end{cases}$

$$y(t) = \pm A + [n_1(t) - n_0(t)] \quad \{+ = 1, - = 0, n_i(t) = n_1 \text{ and } n_0\}$$

-> noise variance doubles

$$P_{e,FSK} = Q\left(\frac{A}{\sqrt{2}\sigma}\right) = Q\left(\sqrt{\frac{E_b}{N_0}}\right)$$

# non-coherent demodulation

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$$P_{e,\text{ASK},\text{noncoherent}} \approx \frac{1}{2} e^{-\frac{A^2}{8\sigma^2}} \geq Q\left(\frac{A}{2\sigma}\right) \text{ (coherent)}$$

- lose some performance, but for large SNR, lose less

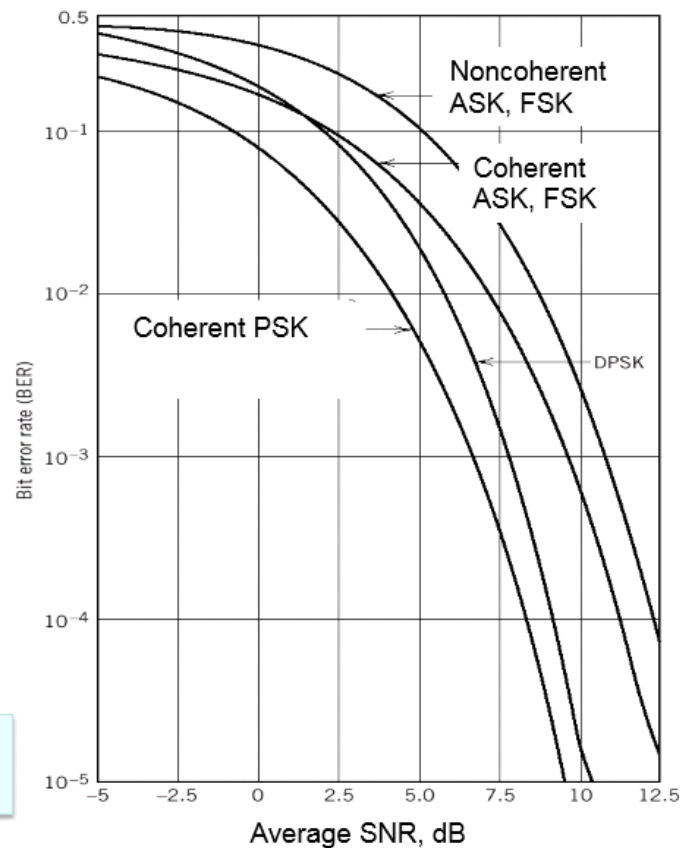
$$P_{e,\text{FSK},\text{noncoherent}} \approx \frac{1}{2} e^{-\frac{A^2}{4\sigma^2}} \geq Q\left(\frac{A}{\sqrt{2}\sigma}\right) \text{ (coherent)}$$

$$P_{e,\text{DPSK},\text{noncoherent}} \approx \frac{1}{2} e^{-\frac{A^2}{2\sigma^2}} \geq P_{e,\text{PSK}} = Q\left(\frac{A}{\sqrt{2}\sigma}\right) \text{ (coherent)}$$

- cannot do PSK using envelope
- DPSK:  $b_n = b_{n-1} * a_n$

Scheme	Bit Error Rate
Coherent ASK	$Q(A/2\sigma)$
Coherent FSK	$Q(A/\sqrt{2}\sigma)$
Coherent PSK	$Q(A/\sigma)$
Noncoherent ASK	$\frac{1}{2} \exp(-A^2/8\sigma^2)$
Noncoherent FSK	$\frac{1}{2} \exp(-A^2/4\sigma^2)$
DPSK	$\frac{1}{2} \exp(-A^2/2\sigma^2)$

**Caution:** ASK and FSK have the same bit error rate with respect to average SNR.



# Entropy and source Coding

20 May 2018 18:14

$$I(s) = \log \frac{1}{p}$$

Information in a symbol,  $s$   
a unique symbol  $s$ , has probability,  $p$

## Discrete Memoryless Source

- sequence of symbols from a finite alphabet
- successive symbols are **independent and identically distributed**

## Source Entropy

$$H(S) = E[I(s_k)] = \sum_{k=1}^K p_k I(s_k) = - \sum_{k=1}^K p_k \log_2 p_k$$

- units: bits/symbol

## Average Codeword Length

$l_k$  = number of bits to encode symbol  $k$

$$\bar{L} = \sum_{k=1}^K p_k l_k$$

Typical sequence of length  $N$

$$p(S_N) = p_1^{Np_1} \times p_2^{Np_2} \times \dots \times p_K^{Np_K}$$

$$L_N = \log_2 \frac{1}{p(S_N)} = -\log(p_1^{Np_1} \times p_2^{Np_2} \times \dots \times p_K^{Np_K})$$
$$\rightarrow L_N = -N \sum_{k=1}^K p_k \log_2 p_k = NH(S)$$

$L_N$  is average number of symbols to code a sequence of length  $N$ , so

$$\bar{L} = \frac{L_N}{N} = H(S) \rightarrow \bar{L} \geq H(S)$$

## Huffman Coding:

- assign each of the two smallest probabilities a bit, then merge
  - repeat until only one symbol left
- (read backwards for codeword of a symbol)
- not unique - can reorder probabilities
- the most efficient **prefix code** - any codeword does not contain, at its start, the code of any other codeword
- drawback - need to know probabilities

# Channel capacity

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## Conditional Entropy

$$H(X|Y = y_k) = \sum_{j=0}^{J-1} p(x_j | y_k) \log_2 \left[ \frac{1}{p(x_j | y_k)} \right]$$

## Mean Entropy:

$$H(X|Y) = \sum_{k=0}^{K-1} \sum_{j=0}^{J-1} p(x_j, y_k) \log_2 \left[ \frac{1}{p(x_j | y_k)} \right]$$

- amount of uncertainty after observing output

## Mutual Information:

$$I(X; Y) = H(X) - H(X|Y) = \sum_{k=0}^{K-1} \sum_{j=0}^{J-1} p(x_j, y_k) \log_2 \left[ \frac{p(x_j | y_k)}{p(y_k)} \right]$$

- $I(X; Y) = I(Y; X) \geq 0$
- at worst, Y is independent of X, and  $I(X; Y) = 0$

## Channel Capacity

$$C = \max_{p(x_j)} I(X; Y)$$

## AWGN Channel

$$C = B \log_2(1 + SNR) = B \log_2 \left( 1 + \frac{P}{N_0 B} \right) \text{ bps}$$

- B - channel bandwidth

## Simple Error Checks - offers no correction

Repetition, halves bit rate

Parity bit, detect odd amount of error

## Block Codes:

- encode source block into longer codeword
- detect - not valid code word
- correct - match corrupted block to a valid one via its proximity
- (n,k) - binary linear block code
  - k bits of source
  - n bit codeword
  - code rate,  $R = \frac{k}{n}$
- generator matrix, G
  - n by k matrix
- **Hamming weight**
  - $w_H(a) = \text{number of non-zero elements in the vector}$
- **Hamming Distance**
  - $d_H(a, b) = w_H(a + b)$
- error detection
  - $d_{min} = \min\{w_H(\text{codeword}), c \in C, c \neq 0\}$
  - number of errors that can be detected =  $d_{min} - 1$
- error correction
  - number of errors that can be corrected =  $\frac{d_{min}-1}{2}$