# DISCOVERY THROUGH GOSSIP

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Figure 1: Discovery Methods

#### 1 INTRODUCTION

We want to study Gossip based discovery processes in both directed and undirected graphs using push discovery (triangulation) and pull discovery (two-hop walk process).

We are interested in studying the time taken by process to converge to the transitive closure of the graph.

#### 1.1 Notation

Table 1: Table of Notations

Notation	Description
$\delta_{t}$	Minimum degree of G <sub>t</sub>
$N_t^i(u)$	Set of nodes at distance i from $u$ in $G_t$
$d_{t}(u)$	Degree of u in G <sub>t</sub>
$d_t(u, S)$	Degree induced on S

#### 1.2 Useful lemmas

 $\textit{Lemma} \ \text{1.} \ |\cup_{i=1}^4 N_t^i(u)| \geqslant \text{min}\{2\delta_t, n-1\} \ \text{for all} \ u \in G_t.$ 

*Proof.* If  $N_t^3 \neq \emptyset$ , then  $|\cup_{i=2}^4 N_t^i(u)| \geqslant \delta_t$  and  $|N_t^1(u)| \geqslant \delta_t$ . So  $|\cup_{i=1}^4 N_t^i(u)| \geqslant 2\delta_t$  since the two sets are disjoint.

If 
$$N_t^3 = \emptyset$$
,  $N_t^1(u) \cup N_t^2(u) = n - 1$  since  $G_t$  is connected.

**Lemma** 2. Consider k Bernoulli experiments in which the success probability of the ith experiment is at least i/m where  $m \ge k$ . If  $X_i$  denotes the number of trials needed for experiment i to output a success and  $X = \sum_{i=1}^k X_i$ , then  $\Pr[X > (c+1)m \ln m] < \frac{1}{m^c}$ .

*Proof.* w.l.o.g assume that k=m. The problem can be seen as *coupon collector problem* where  $X_{m+i-1}$  is the number of steps to collect ith coupon. Consider the probability of not obtaining the ith coupon after  $(c+1)m\ln m$  steps, we have:  $(1-\frac{1}{m})^{(c+1)m\ln m} < e^{-(c+1)\ln m} = \frac{1}{m^{c+1}}$  By union bound, the probability that some coupon has not been collected after  $(c+1)m\ln m$  steps is less than  $\frac{1}{m^c}$ .

### 2.1 Upper Bound

#### 2.2 Lower Bound

**Theorem 3** (Upper bound for triangulation process). For any connected undirected graph, the triangulation process converges to a complete graph in  $O(n \log^2 n)$  rounds with high probability.

In order to prove Theorem 3, we prove that the minimum degree of the graph increases by a constant factor (or equals to n-1) in  $O(n \log n)$  steps. We say that a node  $\nu$  is **weakly tied** to a set of nodes S if  $d_t(\nu, S) < \delta_0/2$ , and **strongly tied** to a set of nodes S if  $d_t(\nu, S) \ge \delta_0/2$ .

**Lemma** 4. If  $d_t(u) < \min\{n-1, (1+\frac{1}{4}\delta_0)\}$  and  $w \in N_t^1(u)$  has at least  $\frac{\delta_0}{4}$  edges to  $N_t^2(u)$ , then the probability that u connects to a node in  $N_t^2(u)$  through w in round t is at least  $\frac{1}{6n}$ .

*Proof.* The probability that u connects to a node in  $N_t^2(u)$  through w in round t is:

$$\frac{d_t(w,N_t^2(u))}{D_t(w)} \times \frac{1}{d_t(w)} \geqslant \frac{d_t(w,N_t^2(u))}{D_t(w)} \times \frac{1}{n} \geqslant \frac{d_t(w,N_t^2(u))}{|N_t^1(u)| + d_t(w,N_t^2(u))} \times \frac{1}{n} \geqslant \frac{d_t(w,N_t^2(u))}{\frac{\delta_0}{(1+\frac{1}{4})\delta_0 + d_t(w,N_t^2(u))}} \times \frac{1}{n} \geqslant \frac{\frac{\delta_0}{4}}{(1+\frac{1}{4})\delta_0 + \frac{\delta_0}{4}} \times \frac{1}{n} = \frac{1}{6n}$$

**Lemma** 5. If  $d_t(u) < \min\{n-1, (1+\frac{1}{4}\delta_0)\}$  and  $w \in N_t^1(u)$  is weakly tied to  $N_t^2(u)$ , and  $v \in N_0^1(u) \cap N_0^1(w)$ , then u connects to v through w in round t with probability at least  $\frac{1}{4\delta_2^2}$ .

*Proof.* Since w is weakly tied to  $N_t^2(u)$  and  $d_t(w)$  is at most  $|N_t^1(u)| + d_t(w, N_t^2(u))$ , we obtain that  $d_t(w)$  is at most  $(1 + \frac{1}{4})\delta_0 + \frac{\delta_0}{2}$ . Therefore, the probability that u connects to v through w in round t equals:

$$\frac{1}{\mathrm{d}_{\mathsf{t}}(w)^2} \geqslant \frac{1}{((1+\frac{1}{4})\delta_0 + \frac{\delta_0}{2})^2} \geqslant \frac{1}{\frac{7\delta_0}{4}^2} \geqslant \frac{1}{4\delta_0^2}.$$

To analyze the growth in the degree of a node u, we consider two overlapping cases. The first case is when more than  $\delta_0/4$  nodes of  $N_t^1(u)$  are strongly tied to  $N_t^2(u)$ , and the second is when less than  $\delta_0/3$  nodes of  $N_t^1(u)$  are strongly tied to  $N_t^2(u)$ .

**Lemma** 6 (Several nodes are strongly tied to two-hop neighbors). There exists a  $T = O(n \log n)$  such that if more than

### 2.3 Paragraphs

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#### 2.4 Math

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$$\cos^3 \theta = \frac{1}{4} \cos \theta + \frac{3}{4} \cos 3\theta \tag{1}$$

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**Definition 1** (Gauss). To a mathematician it is obvious that  $\int_{-\infty}^{+\infty} e^{-x^2} dx = \sqrt{\pi}$ .

**Theorem 7** (Pythagoras). *The square of the hypotenuse (the side opposite the right angle)* is equal to the sum of the squares of the other two sides.

*Proof.* We have that  $\log(1)^2 = 2\log(1)$ . But we also have that  $\log(-1)^2 = \log(1) = 0$ . Then  $2\log(-1) = 0$ , from which the proof.

#### 3 RESULTS AND DISCUSSION

Reference to Figure 2 on the following page.

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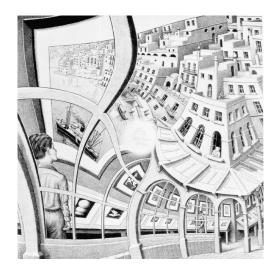


Figure 2: An example of a floating figure (a reproduction from the Gallery of prints, M. Escher, from http://www.mcescher.com/).

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#### Subsection

Subsubsection

word Definition

**CONCEPT** Explanation

#### IDEA Text

- First item in a list
- Second item in a list
- Third item in a list

Reference to Table 1 on page 2.

### 3.2 Figure Composed of Subfigures

Reference the figure composed of multiple subfigures as Figure 3 on the following page. Reference one of the subfigures as Figure 3b on the next page.

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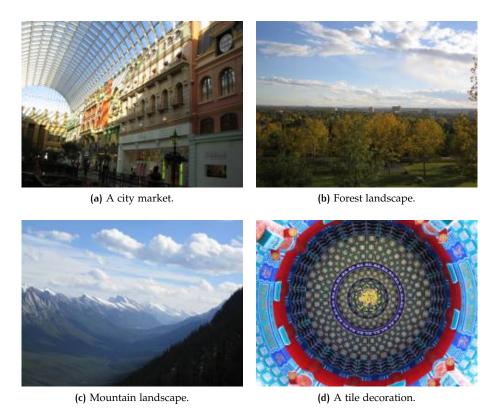


Figure 3: A number of pictures with no common theme.

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#### REFERENCES

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