

DIGITAL DESIGN

CS/ECE/EEE/INSTR F215

Lecture 3
Sarang Dhongdi

BOOLEAN ALGEBRA

Axiomatic definitions of Boolean Algebra

Huntington Postulates

Postulate 1 – Closure

- (a) Operator +
- (b) Operator .

Postulate 2 – Identity

- (a) $x+0 = 0+x = x$
- (b) $x.1 = 1.x = x$

Postulate 3 – Commutative

- a) $x+y = y+x$
- b) $x.y = y.x$

Postulate 4 – Distributive

- a) $x.(y+z) = (x.y) + (x.z)$
- b) $x+(y.z) = (x+y).(x+z)$

Postulate 5 – Complement

- (a) $x+x' = 1$
- (b) $x.x' = 0$

Postulate 6 –

- There exists at least two elements x and y in set B, such that $x \neq y$

- Associative law is not part of postulates, but it can be derived from other postulates.

- No inverses for addition and multiplication – i.e. no subtraction and division.

- Complement is not available for ordinary algebra.

- Postulate 4b is not valid for ordinary algebra.

Two-valued Boolean Algebra

- It is defined over the set of two elements $B=\{0,1\}$, with rules for the binary operators + and ., along with complement operator.

x	y	x.y
0	0	0
0	1	0
1	0	0
1	1	1

AND

x	y	x+y
0	0	0
0	1	1
1	0	1
1	1	1

OR

x	x'
0	1
1	0

NOT

Postulate 1 – Closure

$$a) (x+y)+z = x+(y+z)$$

$$b) (x.y).z = x.(y.z)$$

Postulate 4 – Distributive

- a) $x.(y+z) = (x.y) + (x.z)$
- b) $x+(y.z) = (x+y).(x+z)$

x	y	z	y+z	x.(y+z)	x.y	x.z	(x.y) + (x.z)
0	0	0	0	0	0	0	0
0	0	1	1	0	0	0	0
0	1	0	1	0	0	0	0
0	1	1	1	0	0	0	0
1	0	0	0	0	0	0	0
1	0	1	1	1	0	1	1
1	1	0	1	1	1	0	1
1	1	1	1	1	1	1	1

Duality

- For the dual of the algebraic expression,
 - Interchange AND and OR
 - Replace 1's with 0's and 0's with 1's.

Postulate	Expression (a)	Dual (b)
Postulate 2	$x+0 = x$	$x.1 = x$
Postulate 3	$x+y = y+x$	$x.y = y.x$
Postulate 4	$x.(y+z) = (x.y) + (x.z)$	$x+(y.z) = (x+y).(x+z)$
Postulate 5	$x+x' = 1$	$x.x' = 0$

Theorems of Boolean Algebra

Theorem	Expression (a)	Dual (b)
Theorem 1 - Idempotency (Sameness)	$x+x = x$	$x.x = x$
Theorem 2	$x+1 = 1$	$x.0 = 0$
Theorem 3 - Involution	$(x')' = x$	
Theorem 4 - Associative	$x+(y+z) = (x+y)+z$	$x(yz) = (xy)z$
Theorem 5 - De Morgan	$(x+y)' = x' y'$	$(xy)' = x' + y'$
Theorem 6 - Absorption	$x+xy = x$	$x(x+y) = x$

Theorem 1 (a) $x + x = x$

Statement	Rule/postulate used
$x+x = (x+x).1$	2 (b) $x.1 = x$
$= (x+x)(x+x')$	5 (a) $x+x' = 1$
$= x+xx'$	4 (b) $x+(y.z) = (x+y).(x+z)$
$= x+0$	5 (b) $x.x' = 0$
$= x$	2 (a) $x+0 = x$


Logic operations

- Total 2^{2n} functions for n binary variables.
- For two variables, total 16 different functions are possible.


Boolean Functions	Operator symbol	Name
$F_0 = 0$		Null
$F_1 = x.y$	$x.y$	AND
$F_2 = x$		Transfer
$F_3 = y$		Transfer
$F_4 = x'y$	y/x	Inhibition
$F_5 = x.y'$	x/y	Inhibition
$F_6 = x'y+xy'$	$x \oplus y$	Ex-OR
$F_7 = x+y$	$x+y$	OR

Boolean Functions	Operator symbol	Name
$F_8 = (x+y)'$	$x \downarrow y$	NOR
$F_9 = xy + x'y'$	$x \text{O} y$	Equivalence
$F_{10} = x+y'$	$x \subset y$	Implication
$F_{11} = x'+y$	$x \supset y$	Implication
$F_{12} = y'$	y'	Complement
$F_{13} = x'$	x'	Complement
$F_{14} = (xy)'$	$x \uparrow y$	NAND
$F_{15} = 1$		Identity


Digital Logic gates

AND



x	y	$F=x.y$
0	0	0
0	1	0
1	0	0
1	1	1

NOT


x	$F=x'$
0	1
1	0

OR



x	y	$F=x+y$
0	0	0
0	1	1
1	0	1
1	1	1

Buffer


x	$F=x$
0	0
1	1


Digital Logic gates

NAND




x	y	$F=(x \cdot y)'$
0	0	1
0	1	1
1	0	1
1	1	0

XOR




x	y	$F=x \cdot y' + x' \cdot y$ $F= X \oplus Y$
0	0	0
0	1	1
1	0	1
1	1	0

NOR



x	y	$F=(x+y)'$
0	0	1
0	1	0
1	0	0
1	1	0

XNOR

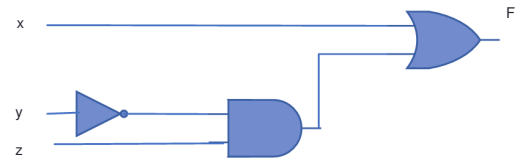


x	y	$F=xy+x'y'$ $F=(X \oplus Y)'$ $F=(X \odot Y)$
0	0	1
0	1	0
1	0	0
1	1	1

Boolean Function

- An algebraic expression consisting of binary variables, constants (0 and 1), and logic operation symbols.

$$F = x + y'z$$



Truth table for $F = x + y'z$

x	y	z	y'	y'z	x+y'z
0	0	0	1	0	0
0	0	1	1	1	1
0	1	0	0	0	0
0	1	1	0	0	0
1	0	0	1	0	1
1	0	1	1	1	1
1	1	0	0	0	1
1	1	1	0	0	1

Complement

- Complement of the Boolean expression can be obtained using De-Morgan's theorem.

- Or simply

- Take dual
- Invert the literals

- Ex. $F = x'yz' + x'y'z$
- Dual of F is $= (x' + y + z')(x' + y' + z)$
- Complement is $F' = (x + y' + z)(x + y + z')$

Definitions

- Literal –Ex. If x is the variable, both x and x' are literals.
- Product term –Ex. Terms such as x, xy, x'yz
- Sum term - Ex. Terms such as x, x'+y, x'+y+z
- Sum of products – SOP is the logical OR of multiple product terms.
- Ex. $xy' + x' + xy'z'$
- Product of sum – POS is the logical AND of the multiple OR terms.
- Ex. $(x+y')(x+y+z')(y'+z')$

Definitions

- Minterms – Special case product (AND) term. A minterm is a product term that contains all of the input variables that make up a Boolean expression. – Each literal no more than once.
- Maxterm – Special case sum (OR) term. A maxterm is a sum term that contains all of the input variables that make up a Boolean expression. – Each literal no more than once.

Definitions

- Canonical sum of products – It is a complete set of minterms that defines when an output variable is a logical 1.
- Each minterm corresponds to the row in the truth table where the output function is 1.
- Canonical product of sums – It is a complete set of maxterms that defines when an output variable is a logical 0.
- Each maxterm corresponds to the row in the truth table where the output function is 0.

Minterms and Maxterms

A	B	Minterms	Maxterms
0	0	$m_0 = A'B'$	$M_0 = A+B$
0	1	$m_1 = A'B$	$M_1 = A+B'$
1	0	$m_2 = AB'$	$M_2 = A'+B$
1	1	$m_3 = AB$	$M_3 = A'+B'$

Minterms and Maxterms -Properties

- Maxterm is logical complement of minterm (and vice-versa)
- For $m_0 = A'B' \rightarrow m_0' = (A'B')' = A+B = M_0$

- Logical OR of all 2^n minterms is equal to logical 1.

$$\sum_{i=0}^{2^n-1} m_i = 1$$

- Sum = $A'B' + A'B + AB' + AB$
 $= A'(B'+B) + A(B'+B)$
 $= A' + A$
 $= 1$

Minterms and Maxterms -Properties

- Logical product of all the maxterms is equal to logical zero.

$$\prod_{i=0}^{2^n-1} M_i = 0$$

- Product = $(A' + B')(A' + B)(A + B')(A + B)$
 $= (A')(A)$
 $= 0$

Three variable

			Minterms		Maxterms	
A	B	C	Term	Designation	Term	Designation
0	0	0	$A'B'C'$	m_0	$A+B+C$	M_0
0	0	1	$A'B'C$	m_1	$A+B+C'$	M_1
0	1	0	$A'BC'$	m_2	$A+B'+C$	M_2
0	1	1	$A'BC$	m_3	$A+B'+C'$	M_3
1	0	0	$AB'C'$	m_4	$A'+B+C$	M_4
1	0	1	$AB'C$	m_5	$A'+B+C'$	M_5
1	1	0	ABC'	m_6	$A'+B'+C$	M_6
1	1	1	ABC	m_7	$A'+B'+C'$	M_7

Canonical SOP and POS

A	B	C	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

$$\begin{aligned}
 F(A,B,C) &= A'B'C + AB'C' + AB'C + ABC' + ABC \\
 &= m_1 + m_4 + m_5 + m_6 + m_7 \\
 &= \sum (m_1, m_4, m_5, m_6, m_7) \\
 &= \sum (1, 4, 5, 6, 7)
 \end{aligned}$$

$$\begin{aligned}
 F' &= A'B'C' + A'BC' + A'BC \\
 &= \sum (0, 2, 3)
 \end{aligned}
 \quad F' = m_0 + m_2 + m_3$$

$$\begin{aligned}
 F &= (A+B+C)(A+B'+C)(A'+B+C') \\
 &= M_0 M_2 M_3 \\
 &= \prod (M_0 M_2 M_3) \\
 &= \prod (0, 2, 3)
 \end{aligned}
 \quad
 \begin{aligned}
 F &= (m_0 + m_2 + m_3)' \\
 &= m_0' m_2' m_3' \\
 &= M_0 M_2 M_3
 \end{aligned}$$

Expressing a given equation in “Sum of minterms”

$$F = A + B'C$$

$$F = A + B'C(A + A')$$

$$= A + AB'C + A'B'C$$

$$= A(B + B') + AB'C + A'B'C$$

$$= AB + AB' + AB'C + A'B'C$$

$$= AB(C + C') + AB'(C + C') + AB'C + A'B'C$$

$$= ABC + ABC' + AB'C + AB'C' + A'B'C$$

$$= \sum (1, 4, 5, 6, 7)$$

x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1