

CS1102: Data Structures and Algorithms

Part 3

Linked List

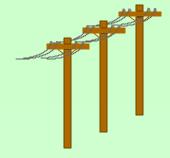
Zoltan KATO S16 06-12

Adopted from Chin Wei Ngan's cs1102 lecture notes



Linked-Lists

- Lists via Arrays
- Linked-List Approach
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- Linked-List Iteration ADT
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 - Doubly Linked-List
 - Implementation
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 - Implementation
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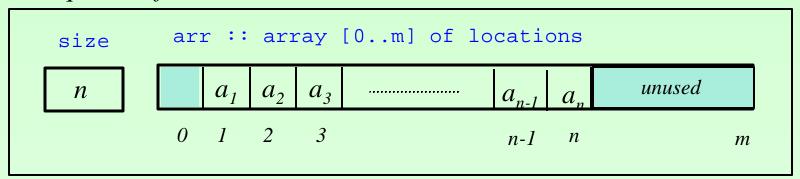


Lists via Arrays

Lists are very pervasive in computing, e.g. class list, list of chars, list of events

One very simple implementation is to use Java arrays

A sequence of n-elements.

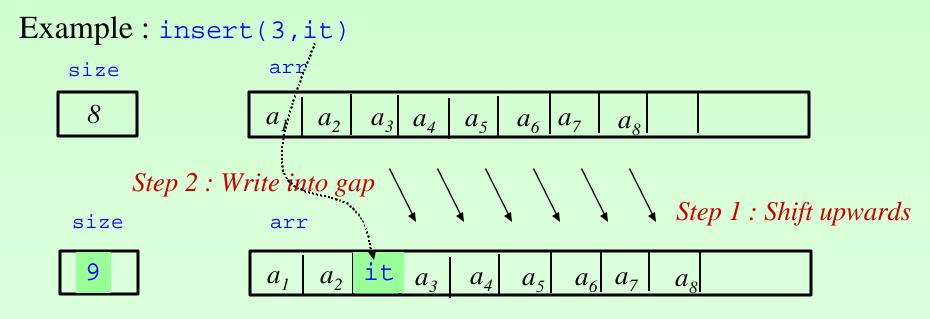


Note: In Java, arrays start from position 0, but here we use it starting from position 1. Position 0 is left unused.

Third Element: arr[3]

Inserting into an Array

While retrieval is very fast, insertion and deletion are slow. Insert has to shift upwards to create gap



Step 3 : Update Size

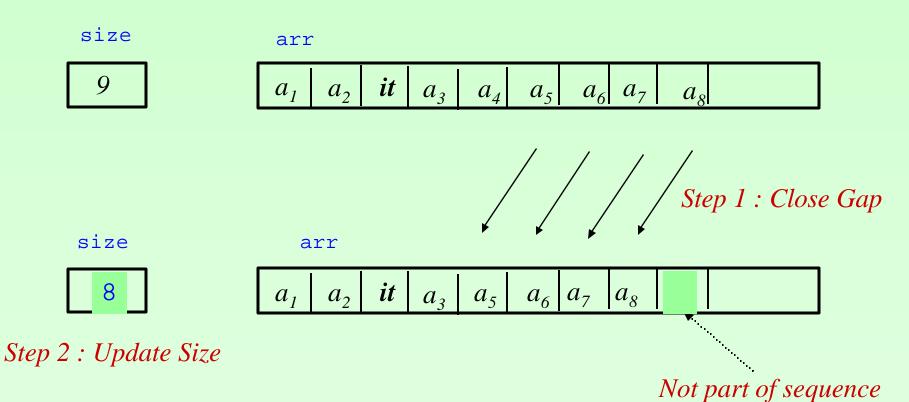
Inserting into an Array

```
class list {
  private int size
  private Object[] arr;
  public void insert(int j, Object it)
      { // pre : 1<=j<=size+1
       for (i=size; i>=j; i=i-1)
          { arr[i+1]=arr[i]; }; // Step 1 : Create gap
       arr[j]=it;
                               // Step 2 : Write to gap
       size = size + 1;  // Step 3 : Update size
```

Deleting from an Array

Delete has to shift downwards to close gap of deleted item,

Example: delete(5)



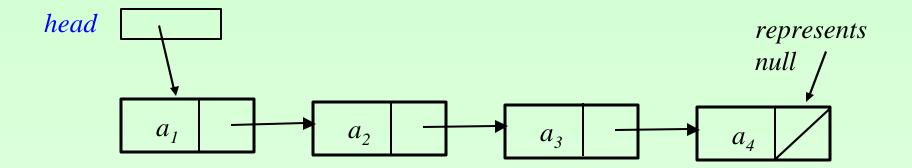
Deleting from an Array

Linked-List Approach

Main problem of array is deletion/insertion slow since it has to shift items in its *contiguous* memory.

Solution: linked list where items need *not be contiguous* with nodes of the form:

Sequence of four items $\langle a_1, a_2, a_3, a_4 \rangle$ can be represented by:



Sample

Linked-List Approach

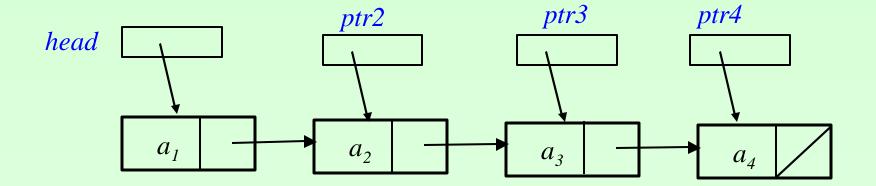
The earlier sequence can be built by:

```
ListNode ptr4 = new ListNode("a4",null);

ListNode ptr3 = new ListNode("a3",ptr4);

ListNode ptr2 = new ListNode("a2",ptr3);

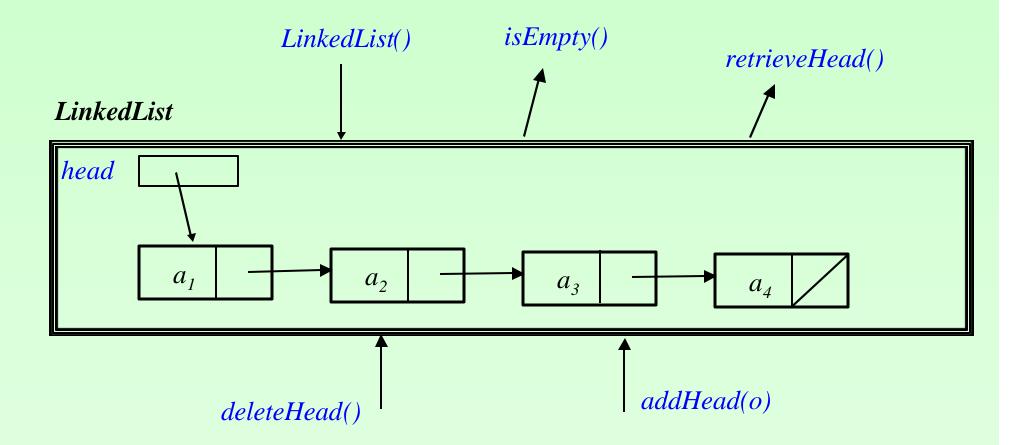
ListNode head = new ListNode("a1",ptr2);
```



Linked-List ADT

We can provide an ADT for linked-list.

This can help hide unnecessary internal details.

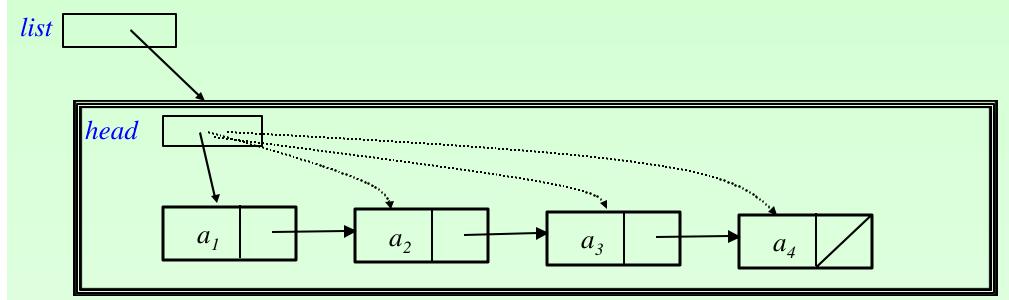


Sample

Linked-List ADT

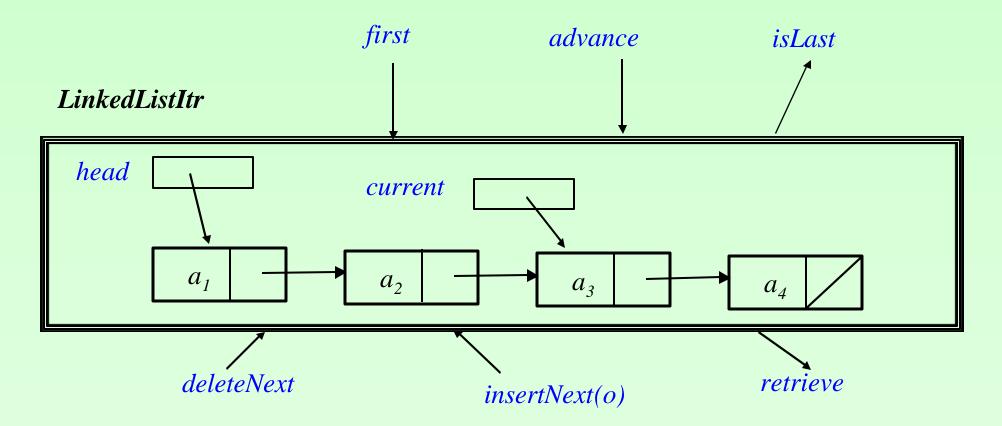
Sequence of four items $\langle a_1, a_2, a_3, a_4 \rangle$ can be built, as follows:

```
→LinkedList list = new LinkedList();
→list.addHead("a4");
→list.addHead("a3");
→list.addHead("a2");
→list.addHead("a1");
```



```
class LinkedList {
 protected ListNode head ;
 public LinkedList() {head = null; }
 public boolean isEmpty()
     {return (head == null); }
 public void addHead(Object o)
     {head = new ListNode(o,head); }
 public void deleteHead() throws ItemNotFound
     {if (head ==null)
           {throw new ItemNotFound("DeleteHead fails");}
       else head = head.next;}
  };
class ItemNotFound extends Exception {
 public ItemNotFound(String msg) {super(msg);}
```

To support better access to our linked-list, we propose to build a linked-list iteration ADT.



```
class LinkedListItr extends LinkedList{
                                                         data structure
   private ListNode current;
   private boolean zeroflag;
   public LinkedListItr() {... }
   public void zeroth() { ... }
                                                       access or change
   public void first() { ... }
                                                       current pointer
   public void advance() { ... }
   public boolean isLast() { ... }
   public boolean isInList() { ... }
   public Object retrieve() { ... }
public void insertNext(Object o) { ... }
                                                        access or change
   public void deleteNext() { ... }
```

Why zeroflag? – To distinguish "zeroth position" from "beyond list position".

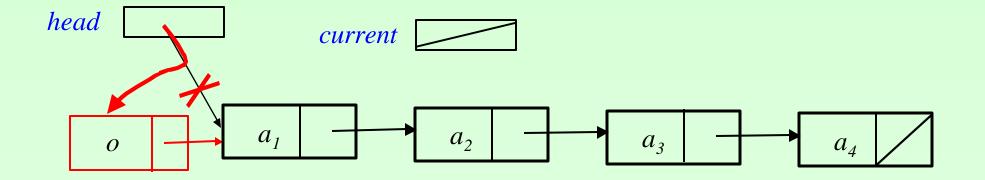
```
Zeroth Position
     (current == null) && (zeroflag==true)
InList
     (current != null)
Beyond List
     (current == null) && (zeroflag==false)
```

```
public void first() throws ItemNotFound
      // set position to first element
      { current = head;
        zeroflaq = false;
        if (current == null)
            {throw new ItemNotFound("No first element");};
public void advance() // advance to next item
    {if (current != null) {current = current.next }
     else {if (zeroflag) {current = head;
                         zeroflag = false;}
           else {}
          };
```

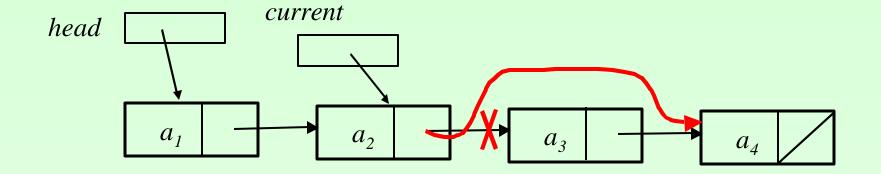
Insertion

```
public void insertNext (Object o) throws ItemNotFound
                     // insert after current position
   { ListNode temp;
     if (current!=null) {temp = new ListNode(o,current.next);
                         current.next = temp;}
     else if (zeroflag) { head = new ListNode(o,head); }
          else { throw new ItemNotFound("insert fails"); };
                                      temp
 head
                       current
          a_1
                        a_2
                                      a_3
```

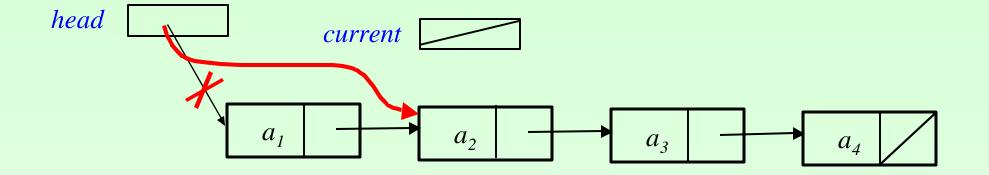
Insertion



Deletion

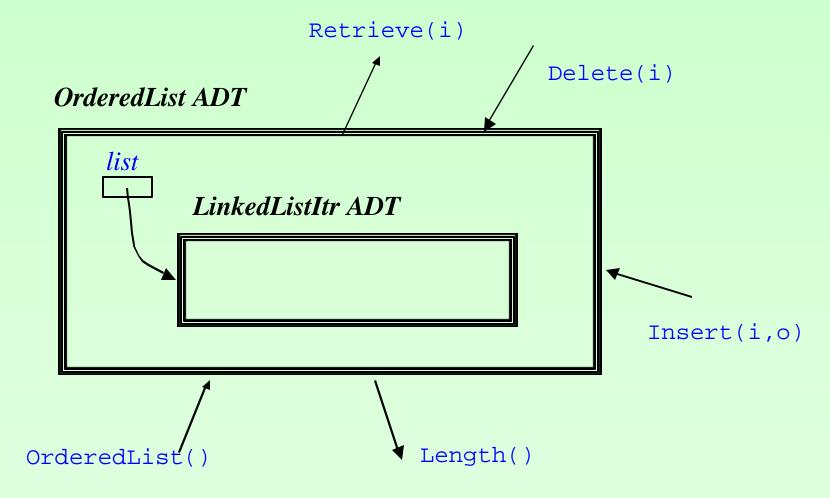


Deletion



Ordered List ADT

We can use LinkedListItr ADT to construct OrderedList ADT.



```
class OrderedList {
private LinkedListItr list;
public OrderedList()
       // pre : true
      // post : this = <>
  { list = new LinkedListItr(); }
public int Length()
       // pre : this = (a_1, ..., a_n)
       // post : returns n
   { int cnt = 0;
     try {
     for (list.first(); list.isInList(); list.advance())
         { cnt++; };
     } catch (ItemNotFound e) { cnt=0; };
     return cnt;
```

Ordered-List ADT

method to move the current pointer to the ith position of list. It returns false if the pointer goes beyond the end of the list.

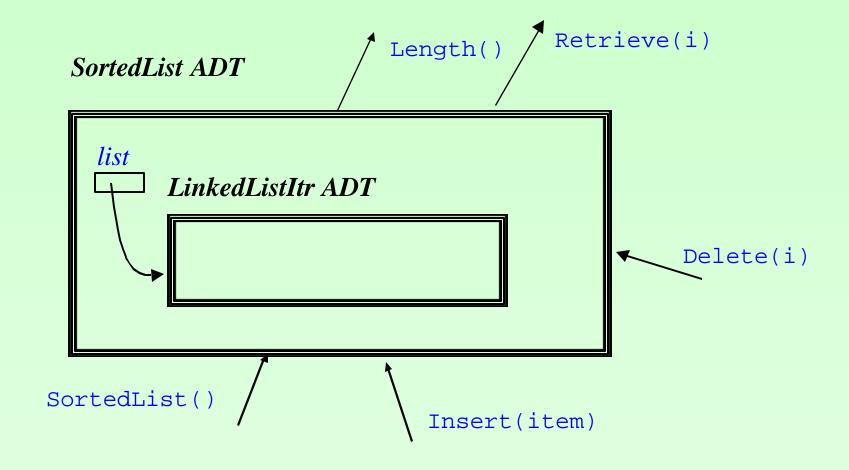
```
private boolean setPosn(int i)
{ // pre : i >= 0
  list.zeroth();
  for (j=1; j<=i; j++)
      {list.advance();
      if !(list.isInList()) return false;
    }
  return true;
}</pre>
```

Ordered-List ADT

```
public void Insert(int i, Object o) throws ItemNotFound
       // pre : this = < a_1, ..., a_n >
       // post : this = < a_1, ..., a_{i-1}, o, a_i, ..., a_n >
       // : or throws error if not(1<=i<=n+1)
 { if (this.setPosn(i-1))
     { list.insertNext(o); }
   else
     { throw new ItemNotFound("insert fails - bad position"); }
public void Delete(int i) throws ItemNotFound
       // pre : this = < a_1, ..., a_n >
       // post : this = < a_1, ..., a_{i-1}, a_{i+1}, ..., a_n >
       // : or throws error if not(1<=i<=n)</pre>
 { if (this.setPosn(i-1))
     { list.deleteNext(); }
   else
     { throw new ItemNotFound("delete fails - no item"); }
```

Sorted List ADT

We can also use *LinkedListItr* ADT to construct *SortedList* ADT.



Linked-List Implementation

Sorted-List ADT

The operations to retrieve and delete the ith item are similar to those given for *OrderedList* ADT.

The insert operation for SortedList ADT requires a lookahead capability. This requires LinkedList Iteration ADT to be supplied with an extra operation, called retrieveNext.

Sorted-List ADT

```
class SortedList {
private LinkedListItr list;
public void Insert(Comparable item)
       // inserts a new item into a sorted position
                                                         retrieve next object
 {try { list.zeroth;
                                                         without changes to
       boolean exitflag = false;
                                                         the current pointer
        do { Object nextobj = list.retrieveNext();
             if (nextobj.compareTo(item)<0) { list.advance }</pre>
             else {exitflag = true};
        while (!(exitflag))
      } catch (ItemNotFound e) { // do nothing
 try { list.insertNext(item);}
       catch (Exception e) {// should not be possible
```

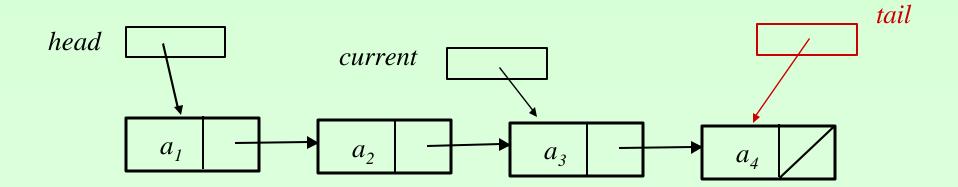
Variations of Linked-Lists

- with tail pointer
- with dummy node
- doubly linked-list
- circular linked-list

With Tail-Pointer

Allows more efficient access (i.e. insertion) at the end of linked list.

Useful for queue-like structures.



New Function

With Tail Pointer

```
class LinkedListItr {
    private ListNode head;
    private ListNode tail;
    private ListNode current;
    private boolean zeroflag;
    public void addTail (Object o) {
         if (tail!=null) { tail.next = new ListNode(o);
                           tail = tail.next;}
         else {tail = new ListNode(o);
               head = tail; }; }
                                                            tail
head
                     current
        a_1
                      a_2
```

With Tail Pointer

```
class LinkedListItr {
  private ListNode head;
  private ListNode tail;
  private ListNode current;
  private boolean iszeroth;
  public void addTail (Object o) {
        if (tail!=null) {tail.next = new ListNode(o);
                         tail = tail.next;}
       else {tail = new ListNode(o);
              head = tail; }; }
                              tail
                                   current
    head
```

Code Changes

With Tail Pointer

Functions, such as zeroth, first, advance, retrieve which do not change the nodes of the linked-list are unchanged.

However, functions that add or delete nodes may affect the tail-pointers. We must provide *extra code* to cater to this possibility.

With Tail Pointer

```
public void deleteNext() throws ItemNotFound
              // delete node after current position
{ if (current!=null)
      {if current.next!=null
             {current.next = current.next.next;
              if (current.next==null) {tail=current};
       else {throw new ItemNotFound("No Next Node to Delete")}
  else if (zeroflag && head!=null)
             {head=head.next;
               if (head==null) {tail=null};
      else {throw new ItemNotFound("No Next Node to Delete");}
```

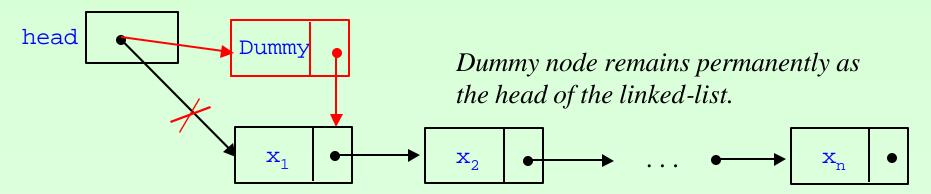
With Dummy Header

Need special codes for

- empty list
- inserting to the front of list

Solution: Use an extra dummy header node to simplify coding.

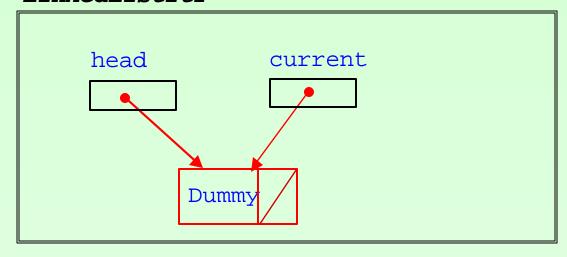
Linked List with Dummy Node



```
class LinkedListItr {
   private ListNode head;
   private ListNode current;

   public LinkedListItr()
        {head = new ListNode((Object) "Dummy", null);
        current = head;}
```

LinkedListItr



```
public void zeroth() // set position prior first element
 { current = head;
public void first() throws ItemNotFound
                       // set position to first element
 { current = head.next;
   if (current == null)
       throw new ItemNotFound("No first element");
public void advance() // advance to next item
 {if (current != null) {current = current.next }
 else (if (zeroflag) (current = head;
               ·····zeroflaq·=·false)··
       relse ()
```

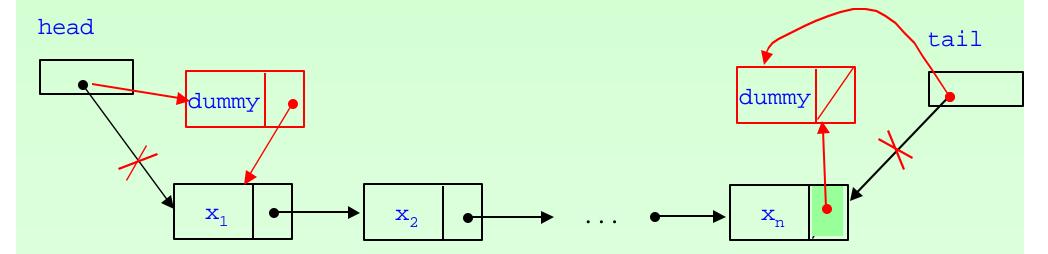
Simpler insertNext Code with Dummy Node

Trailer Node

With Dummy Node

If we have a tail-pointer, we may also wish to have a *dummy trailer* node too!

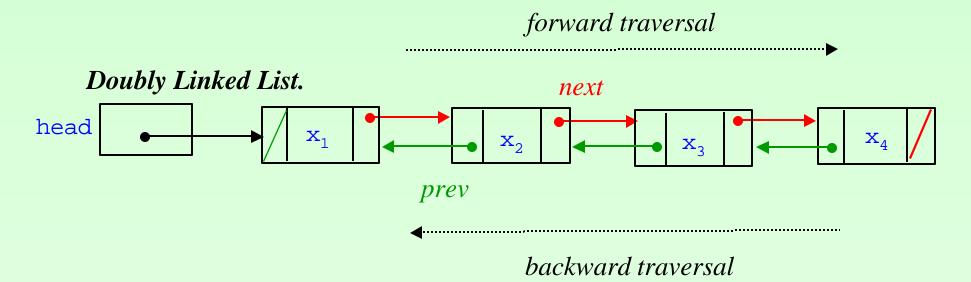
Linked List with Dummy Header & Trailer Node



Doubly Linked-List

Frequently, we need to traverse a sequence in BOTH directions efficiently.

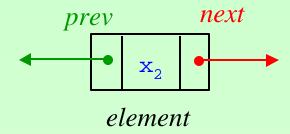
Solution: Use doubly-linked list where each node has two pointers.



Definition for Node

Doubly Linked-List

Each doubly-linked node has two pointers.



Class declaration for above node:

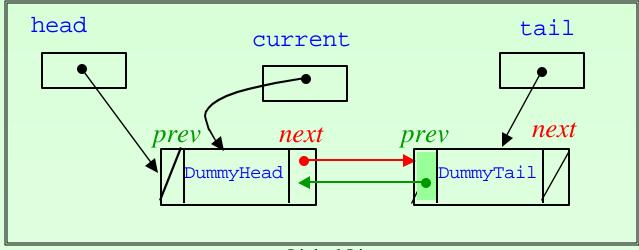
```
class DListNode {
   Object element;
   DListNode next;
   DListNode prev;

public DListNode(Object o, DListNode n, DListNode p)
   { element= o; prev=p; next = n; };
}
```

Implementation

Doubly Linked-List

LinkedListItr



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```
public void zeroth() // set to zeroth position
    current = head;
public void first() throws ItemNotFound
                       // set to zeroth position
    current = head.next;
     if (current==tail) {throw new ItemNotFound("Empty List");};
public boolean isInList() // checks if at a valid position
  { returns (current!=null) && (current!=head)
                   && (current!=tail);
public void advance() // move to next position
  { if (current!=null) {current=current.next;};
```

Insertion

```
public void insert(Object o) throws ItemNotFound
                        // insert at current position
   { if (current != null) && (current != head)
          DListNode temp = new DListNode(o,current,current.prev);
           current.prev.next = temp;
         current.prev = temp;
         current = temp;
     else {throw new ItemNotFound("insert fails");}
                                    temp
                                               current
                                                                  tail
head
 DummyHead
                                                               DummyTail
                X_1
                            \mathbf{X}_{2}
                                       X_2
```

Deletion

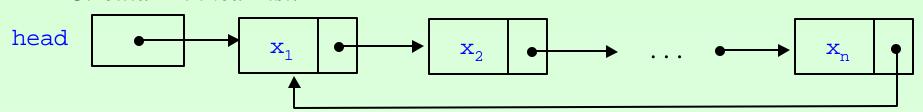
```
public void delete() throws ItemNotFound
                 // delete node at current position
   { if ((current!=null) && (current!=head) && (current != tail))
        current.next.prev = current.prev;
        current.prev.next = current.next;
          current=current.next;
     else { throw new ItemNotFound("No Node to Delete"); }
                                              current
                                                                  tail
head
 DummyHead
                                                                DummyTail
                X_1
                            \mathbf{X}_{2}
```

Circular Linked List

May need to cycle through a list repeatedly, e.g. round robin system for a shared resource.

Solution: Have the last node point to the first node.

Circular Linked List.

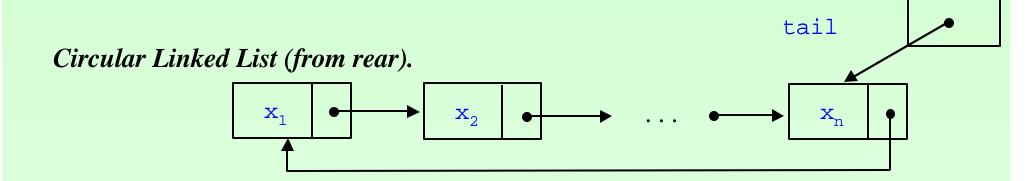


With Rear Pointer

Circular Linked-List

For singly circular linked-list, it may be better to have a pointer from the rear.

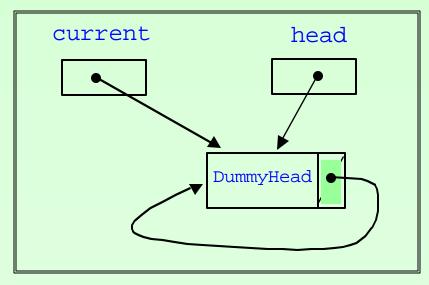
Two pointers for the price of one: (head == tail.next)



With Dummy Header

Circular Linked-List

LinkedListItr



More Codes

Circular Linked-List

```
public void first() throws ItemNotFound
     { current = head.next;
       if (current == head)
           throw new ItemNotFound("No first element");
 public void isLast()
     { if (current!=null && current!=head)
             {return current.next == head;}
      else {return false;};
 public void advance() // advance to next node
     { if (current!=null)
         {current = current.next;
          if (current==head) {current = null;};}
 public void advanceWrap() // supports wrap around
     { current = current.next;
       if (current==head) {current = current.next;};
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```

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Insertion

Circular Linked-List

```
public void insertNext(Object o) throws ItemNotFound
                      // insert after current position
   { ListNode temp;
     if (current!=null) { temp = new ListNode(o,current.next);
                          current.next = temp;}
     else { throw new ItemNotFound("insert fails"); };
                              temp
                                                            head
                                      0
               current
                                                        DummyTail
                a_2
                              a_3
                                             a_{1}
```

Deletion

Circular Linked-List

```
public void deleteNext() throws ItemNotFound
                     // delete node after current position
{ if (current!=null)
       {if (current.next!=head)
          {current.next = current.next.next;}
        else {throw new ItemNotFound("No Next Node to Delete");}
  else {throw new ItemNotFound("No Next Node to Delete");};
                                                           head
               current
                                                        DummyHead
                a_2
```