

# Sequences

Mairead Meagher

Waterford Institute of Technology,

2017



Waterford Institute *of* Technology  
INSTITIÚID TEICNEOLAÍOCHTA PHORT LÁIRGE

We will look at

- the idea of and need for a sequence
- how to construct a sequence, different kinds of sequences
- operations on sequences
- examples of use of sequences

# The need for sequences

- There are times when the properties of sets make them insufficient to model certain kinds of information.
- For example, sometimes order is important and it is important to include each occurrence of an element.
- In sets order is not important and duplicates are ignored i.e.

$$\{1, 2, 3\} = \{3, 2, 1\} = \{1, 3, 2, 1\}$$

- we introduce sequences to help model order (e.g. chronological, size, etc.)
- Sequences are ordered collections of objects.

- For example when talking about days of the week, their order is important. We write the sequence using the angle bracket notation.

$$\text{daysinweek} = \langle \text{sun}, \text{mon}, \text{tue}, \text{wed}, \text{thurs}, \text{fri}, \text{sat} \rangle$$

- This could be equivalently displayed without sequence notation as

$$\{1 \mapsto \text{sun}, 2 \mapsto \text{mon}, 3 \mapsto \text{tue}, 4 \mapsto \text{wed}, 5 \mapsto \text{thu}, 6 \mapsto \text{fri}, 7 \mapsto \text{sat}\}$$

# Sequences - notation

- Thus, a sequence holds not only the elements but also the position they occupy.
- The sequence denoted

$\langle mon, tue, wed \rangle$

contains information as follows:

1	mon
2	tue
3	wed

- Note from this, it follows that :

$\langle mon, tue, wed \rangle \neq \langle wed, tue, mon \rangle$

$\langle mon, tue, wed \rangle \neq \langle mon, tue, wed, mon \rangle$

# How to construct a sequence

Using the example as above:

$DAYS ::= sun \mid mon \mid tue \mid wed \mid thurs \mid fri \mid sat$

$daysGoneBy : seq\ DAYS$
$daysGoneBy = \langle mon, tue, wed \rangle$

More generally (and formally)

$seqX == \{f : \mathbb{N} \twoheadrightarrow X \mid domf = 1.. \# f \bullet f\}$

Note the use of the finite function symbol  $\twoheadrightarrow$  to indicate that sequences are finite functions. Infinite sequences are not part of standard Z.

- A non-empty sequence is a simple specialisation of sequences.
- It represents sequences with at least one element.
- It is written  $\text{seq}_1$  and can be defined as :

$$\text{seq}_1 X == \{s : \text{seq } X \mid \#s > 0 \bullet s\}$$

- Sometimes it is important to have an injective (one-to-one) sequence.
- This will not allow any element in the range to be used more than once in the sequence.
- The **daysinweek** sequence seen above is an example of an injective function.
- It is written `iseq` and is defined as:

$$\text{iseq } X == \text{seq } X \cap (\mathbb{N} \rightarrowtail X)$$

- It is precisely the set of finite of sequences over  $X$  which contain no repetitions.



# Example of use

$ACTIVITY ::= plan \mid code \mid test \mid review \mid deliver$

$Project : seq\ ACTIVITY$

$Project = \langle plan, code, deliver \rangle$

- All normal functions operations can be used on sequences.
- `dom` can be used to deliver the set of index numbers for the sequence and
- `ran` delivers the set of items in the sequence.

$Project(1) = plan$

$Project2 = code$

$ran\ Project = \{plan, code, deliver\}$

$dom\ Project = \{1, 2, 3\}$

$\# Project = 3$

- The main operator for constructing sequences is concatenation, written:



where both of its arguments must be sequences.

- Note that a simple union operators between the two sets containing sequences is not sufficient.

# Operations on Sequences - Concatenation

- Concatenation has the effect of joining two sequences by taking two sequences of the same type and producing a sequence numbered from one to the combined size as follows :

$$\langle \text{mon}, \text{tue}, \text{wed} \rangle \frown \langle \text{thu}, \text{fri}, \text{sat} \rangle = \\ \langle \text{mon}, \text{tue}, \text{wed}, \text{thu}, \text{fri}, \text{sat} \rangle$$

$$\{1 \mapsto \text{mon}, 2 \mapsto \text{tue}, 3 \mapsto \text{wed}\} \frown \\ \{1 \mapsto \text{thu}, 2 \mapsto \text{fri}, 3 \mapsto \text{sat}\}$$

$$= \{1 \mapsto \text{mon}, 2 \mapsto \text{tue}, 3 \mapsto \text{wed}, 4 \mapsto \text{thu}, 5 \mapsto \text{fri}, 6 \mapsto \text{sat}\}$$

# Operations on Sequences - Concatenation

The formal definition for the infix concatenation operation is as follows :

$$\begin{array}{l} \text{[X]} \\ \text{---} \wedge \text{---} : \text{seq } X \times \text{seq } X \rightarrow \text{seq } X \\ \forall s, t : \text{seq } X \bullet \\ \quad s \wedge t = s \cup \{n : 1.. \# t \bullet (n + \# s \mapsto t(n))\} \end{array}$$

The reverse operation takes a sequence and produces a sequence with the order of the elements reversed. For example,

$$\text{rev } \langle \text{mon}, \text{tue}, \text{wed} \rangle = \langle \text{wed}, \text{tue}, \text{mon} \rangle$$

Formally,

$$\begin{array}{l} \text{---} [X] \text{---} \\ \text{seq } X \rightarrow \text{seq } X \\ \hline \forall s : \text{seq } X \bullet \\ \quad \text{rev } s = \{ n : 1.. \# s \bullet n \mapsto s((\# s) - n + 1) \} \end{array}$$

# Splitting Sequences - Head and Tail

- The head of a sequence is its first element.
- The tail of a sequence the rest of the sequence when the head is removed.
- Example

$$\text{head}(\langle \text{mon}, \text{tue}, \text{wed} \rangle) = \text{mon}$$

Note that the head operation returns an element of the sequence, not a sequence.

- Example

$$\text{tail}(\langle \text{mon}, \text{tue}, \text{wed} \rangle) = \langle \text{tue}, \text{wed} \rangle$$

Note that the tail operation returns a sequence.

# Splitting Sequences - Front and Last

- These are similar to head and tail except that
- the last of a sequence is the last element in a sequence and
- the front of a sequence is the sequence containing everything except the last element.
- Example

*front*  $\langle \text{mon}, \text{tue}, \text{wed} \rangle = \langle \text{mon}, \text{tue} \rangle$

*last*  $\langle \text{mon}, \text{tue}, \text{wed} \rangle = \text{wed}$

# Formal definitions for Splitting Sequences

More formally, the operations are defined as follows:

$[X]$

$head, last : seq_1 X \rightarrow X$

$tail, front : seq_1 X \rightarrow seq X$

$\forall s : seq_1 X \bullet$

$head\ s = s(1) \wedge$

$s = \langle head\ s \rangle \hat{\ } tail\ s \wedge$

$last\ s = s(\#s) \wedge$

$front\ s = (1.. \#s - 1) \triangleleft s$



# Distributed Concatenation

- Analogous to the distributed union and intersection that we have already seen and used, there is a distributed concatenation operator.
- It takes a sequence of sequences and concatenates them all together. It is written  $\frown$

$$\frown : \text{seq}(\text{seq } X) \rightarrow \text{seq } X$$

- Example

$$ss = \langle \langle a, b \rangle, \langle c, d \rangle, \langle e \rangle, \langle f, g, h \rangle \rangle$$

$$\frown ss = \langle a, b, c, d, e, f, g, h \rangle$$

- Squash is a very important function. When we are working with sequences, we may be deleting elements etc. So for instance, if we wish to delete the third element of a sequence, e.g.

$$s : \text{seq } \mathbb{N}$$

$$s = \langle 3, 7, 10, 5, 16 \rangle$$

This can equivalently be written:

$$s = \{1 \mapsto 3, 2 \mapsto 7, 3 \mapsto 10, 4 \mapsto 5, 5 \mapsto 16\}$$

‘Delete the third element’

$$s' = \{1 \mapsto 3, 2 \mapsto 7, 4 \mapsto 5, 5 \mapsto 16\}$$

$s'$  is not a sequence, but if we apply squash to it:

$$\begin{aligned} \text{squash}(\{1 \mapsto 3, 2 \mapsto 7, 4 \mapsto 5, 5 \mapsto 16\}) &= \{1 \mapsto 3, 2 \mapsto 7, 3 \mapsto 5, \\ &= \langle 3, 7, 5, 16 \rangle \end{aligned}$$

This is again a sequence.

- Squash takes a finite function defined strictly on the positive integers and compacts it into a sequence,
- i.e. there are no 'holes' left.
- 

$[X]$

$squash : (\mathbb{N}_1 \multimap X) \rightarrow seq\ X$

# Using squash

- 1 Let  $s'$  be the sequence  $s$  with the  $n^{th}$  element of it deleted.  $s'$  should be a sequence.

$$\begin{aligned}s &: \text{seq } \mathbb{N} \\ s' &= \text{squash}(\{3\} \triangleleft s)\end{aligned}$$

$s'$  is a sequence with the third element of  $s$  removed.

- 2 Let  $s'$  be the sequence  $s$  with all occurrences of an element deleted.  $s'$  should be a sequence.

$$\begin{aligned}s &: \text{seq } \mathbb{N} \\ s' &= \text{squash}(s \triangleright \{23\})\end{aligned}$$

$s'$  is a sequence with all the occurrences of the number 23 removed from  $s$ .

- We revisit the concept of sets being disjoint and partitioning another set.
- Simply, a sequence of sets is disjoint if they have no elements in common( i.e. the distributed intersection operator yields the empty set).

$$\begin{aligned} \text{disjoint\_} &: \mathbb{P}(\text{seq } \mathbb{P} X) \\ \text{disjoint}\langle A, B \rangle &\Leftrightarrow A \cap B = \emptyset \end{aligned}$$

- We model the employees in a factory.

*management, personnel, production, employees :  $\mathbb{P}$  PERSON*

**management** is the set of all employees in management roles.

**personnel** is the set of all employees in personnel roles.

**production** is the set of all employees in production roles.

**employees** is the set of all employees in the factory.

- We model the following thus:
  - 1 No employee has more than one role (e.g. they cannot be in management and production), this can be affected by:

*disjoint <management, personnel, production>*

- 2 Furthermore, if we wish to state that each employee must in one of the three roles(sets), this can be affected by:

*<management, personnel, production> partition employees*

- Formally,

$$| \quad \_ \text{partition} \_ : (\text{seq}(\mathbb{P} X)) \leftrightarrow \mathbb{P} X$$

- A sequence of sets partition a set  $S$  if
  - they are disjoint and
  - the distributed union of the sequence of sets is  $S$ .



Any questions?

Irgendwelche Fragen?

Aon céisteanna?

